Assignment #B: 贪心、矩阵和动态规划

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2023 fall, Complied by Xinjie Song, Phy

说明:

本周作业留点难题,期中考试结束了,需要学习计算概论了。这次不分必做选做题目了,如果耗时太长,直接找答案看。两个题解,经常更新。所以最好从这个链接下载最新的,https://github.com/GMyhf/2020fall-cs101。

- 1)请把每个题目解题思路(可选),源码Python,或者C++(已经在Codeforces/Openjudge上AC),截图(包含Accepted, 学号),填写到下面作业模版中(推荐使用 typora https://typoraio.cn,或者用word)。AC或者没有AC,都请标上每个题目大致花费时间。
- 2) 提交时候先提交pdf文件,再把md或者doc文件上传到右侧"作业评论"。Canvas需要有同学清晰头像、提交文件有pdf、作业评论有md或者doc。
- 3) 如果不能在截止前提交作业,请写明原因。

编程环境

操作系统: Windows 11 22H2

Python编程环境: PyCharm 2023.2 (Community Edition)

C/C++编程环境: g++ (x86_64-win32-seh-rev0, Built by MinGW-W64 project) 8.1.0

1. 题目

如果耗时太长,直接看解题思路,或者源码

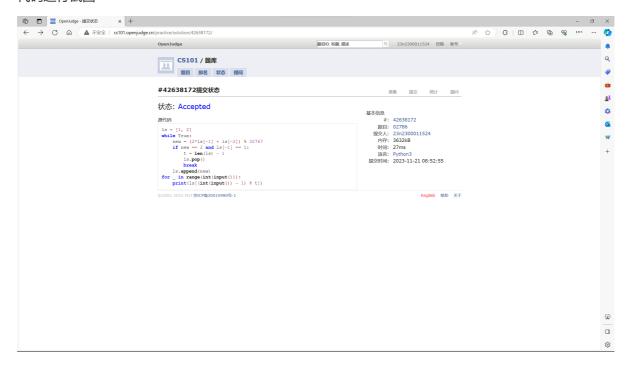
02786:Pell数列

http://cs101.openjudge.cn/practice/02786/

思路:找到周期,直接输出

```
ls = [1, 2]
while True:
    new = (2*ls[-1] + ls[-2]) % 32767
    if new == 2 and ls[-1] == 1:
        t = len(ls) - 1
        ls.pop()
        break
    ls.append(new)
for _ in range(int(input())):
    print(ls[(int(input()) - 1) % t])
```

代码运行截图



04133:垃圾炸弹

matrices, http://cs101.openjudge.cn/practice/04133/

思路:将某个垃圾能被清除的范围内的矩阵值加一,最后整个矩阵值最大的点便是所求点,统计并输出 即可

代码

```
maps = [[None for j in range(1025)] for i in range(1025)]
d, n = int(input()), int(input())
max_num = 0
count = 0
for i in range(n):
    x, y, num = map(int, input().split())
```

```
for j in range(max(x - d, 0), min(x + d + 1, 1025)):
    for k in range(max(y - d, 0), min(y + d + 1, 1025)):
        if maps[j][k]:
            maps[j][k] += num
        else:
            maps[j][k] = num
        if maps[j][k] > max_num:
            max_num = maps[j][k]
            count = 1
        elif maps[j][k] == max_num:
            count += 1

print(count, max_num)
```

代码运行截图



26971:分发糖果

greedy, http://cs101.openjudge.cn/routine/26971/

思路:分别正向、反向遍历数据,标记单增和单减部分即可

代码

```
n = int(input())
ratings = [int(i) for i in input().split()]
candies = [0 for i in range(n)]
i = 0
while i < n:
    j = i + 1
    while j < n:
    if ratings[j] <= ratings[j - 1]:</pre>
```

```
break
       candies[j-1] = j-i
       if j == n - 1:
           candies[j] = j - i + 1
       j += 1
   i = j
i = 0
while i < n:
   j = i + 1
   while j < n:
       if ratings[n - 1 - j] \leftarrow ratings[n - j]:
       candies[n - j] = j - i
       if j == n - 1:
           candies[n-1-j]=j-i+1
       j += 1
   i = j
candies[0] = [candies[0], 1][candies[0] == 0]
candies[-1] = [candies[-1], 1][candies[-1] == 0]
for i in range(1, n - 1):
   if candies[i] == 0:
        candies[i] = 1 + [max(candies[i - 1], candies[i + 1]), 0][ratings[i] ==
ratings[i - 1] == ratings[i + 1]]
print(sum(candies))
```

26976:摆动序列

greedy, http://cs101.openjudge.cn/routine/26976/

思路: 统计单增单减片段数量

代码

```
n = int(input())
ls = list(map(int, input().split()))
count, i = 1, 0
try:
    while i < n - 1:
        j = i + 1
        count += 1
        if ls[i + 1] > ls[i]:
            while ls[j] >= ls[j - 1]:
                 j += 1
        elif ls[i + 1] < ls[i]:
            while ls[j] \leftarrow ls[j - 1]:
                 j += 1
        else:
            count -= 1
            j += 1
        i = j - 1
except:
    pass
print(count)
```

27104:世界杯只因

http://cs101.openjudge.cn/practice/27104/

思路:经hrc同学启发,首先用列表存储每个摄像头的左右端点,根据自定义规则排序:左端点更小的在前,左端点相同的右端点靠后的在前;然后依次遍历排序后的列表,记录右侧可以达到的最大值,当索引超出当前右端时,数量加1,当前右端更改为此时右端点的最大值。

代码

```
n = int(input())
ls = list(map(int, input().split()))
for i in range(n):
    ls[i] = [max(0, i - ls[i]), min(n - 1, i + ls[i])]
ls.sort(key=lambda t: (t[0], -t[1]))
count = 1
left = right = ls[0][1]
for i in range(1, n):
    right = max(right, ls[i][1])
    if ls[i][0] >= left + 1:
        left = right
        count += 1
print(count)
```



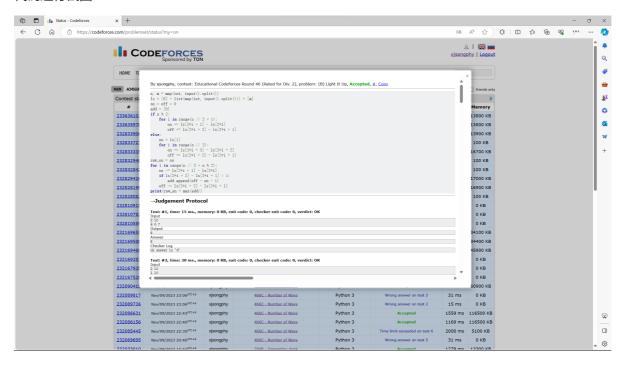
CF1000B: Light It Up

greedy, 1500, https://codeforces.com/problemset/problem/1000/B

思路:统计在每个关灯初后一个位置插入(如果可以插入的话)能够增加的开灯时间,然后找到增量的最大值即可。

代码

```
n, m = map(int, input().split())
ls = [0] + list(map(int, input().split())) + [m]
on = off = 0
add = [0]
if n % 2:
    for i in range(n // 2 + 1):
        on += 1s[2*i + 1] - 1s[2*i]
        off += 1s[2*i + 2] - 1s[2*i + 1]
else:
    on = ls[1]
    for i in range(n // 2):
        on += 1s[2*i + 3] - 1s[2*i + 2]
        off += ls[2*i + 2] - ls[2*i + 1]
raw_on = on
for i in range(n // 2 + n % 2):
    on -= ls[2*i + 1] - ls[2*i]
    if ls[2*i + 2] - ls[2*i + 1] > 1:
        add.append(off - on - 1)
    off -= 1s[2*i + 2] - 1s[2*i + 1]
print(raw_on + max(add))
```



2. 学习总结和收获

部分题目以前做过,剩下的题目让我了解了一些新的类型的问题,学会了一些新的贪心/动态规划思路。

截至2023年11月21日,OJ完成题目110道,CF完成题目47道。