

$$(\lambda M)_{i,j} = \lambda(M)_{i,j}$$

↳ 2 cells hidden

ENR145 Computational Methods:

Hamming Numpy 301: a detour from traffic control so

Or this:

we can do one step operation for many things

$$P = \begin{bmatrix} Q_{11}R_{11} + Q_{12}R_{21} + \cdots + Q_{1n}R_{n1} & Q_{11}R_{12} + Q_{12}R_{22} + \cdots + Q_{1n}R_{n2} & \cdots & Q_{11}R_{1q} + Q_{12}R_{2q} + \cdots + Q_{1n}R_{nq} \\ Q_{21}R_{11} + Q_{22}R_{21} + \cdots + Q_{2n}R_{n1} & Q_{21}R_{12} + Q_{22}R_{22} + \cdots + Q_{2n}R_{n2} & \cdots & Q_{21}R_{1q} + Q_{22}R_{2q} + \cdots + Q_{2n}R_{nq} \\ \vdots & \vdots & \ddots & \vdots \\ Q_{m1}R_{11} + Q_{m2}R_{21} + \cdots + Q_{mn}R_{n1} & Q_{m1}R_{12} + Q_{m2}R_{22} + \cdots + Q_{mn}R_{n2} & \cdots & Q_{m1}R_{1q} + Q_{m2}R_{2q} + \cdots + Q_{mn}R_{nq} \end{bmatrix}$$

I think using an example is better:

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Spring 2026

$$E = AD = \begin{bmatrix} 10 & 20 & 30 \\ 40 & 50 & 60 \end{bmatrix} \begin{bmatrix} 2 & 3 & 5 & 7 \\ 11 & 13 & 17 & 19 \\ 23 & 29 & 31 & 37 \end{bmatrix} = \begin{bmatrix} 150 & 150 & 1520 & 1560 \\ 2010 & 2510 & 2910 & 3450 \end{bmatrix}$$

So, a $[2 \times 3] \times [3 \times 4] = [2, 4]$

And a $[1 \times 3] \times [2 \times 4]$ is not mathematically legal.

↳ 5 cells hidden



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Download the in-class notebook here:

 **ENR 145: Computational Methods for Physicists and Engineers**
Department of Engineering Physics, Coe College | Cedar Rapids, Iowa

[Download Syllabus](#) [Upload to Moodle](#)

Week 5:

[Assignment #5](#) [Example of coding quests](#)
[Slides #6](#) [Hamming numpy in-class](#)

```
[1]  # Let's test our theory to do hamming (16,11) with numpy
✓ 0s  # Run this block first to call numpy in.
      import numpy as np # this is how you call a python library
```

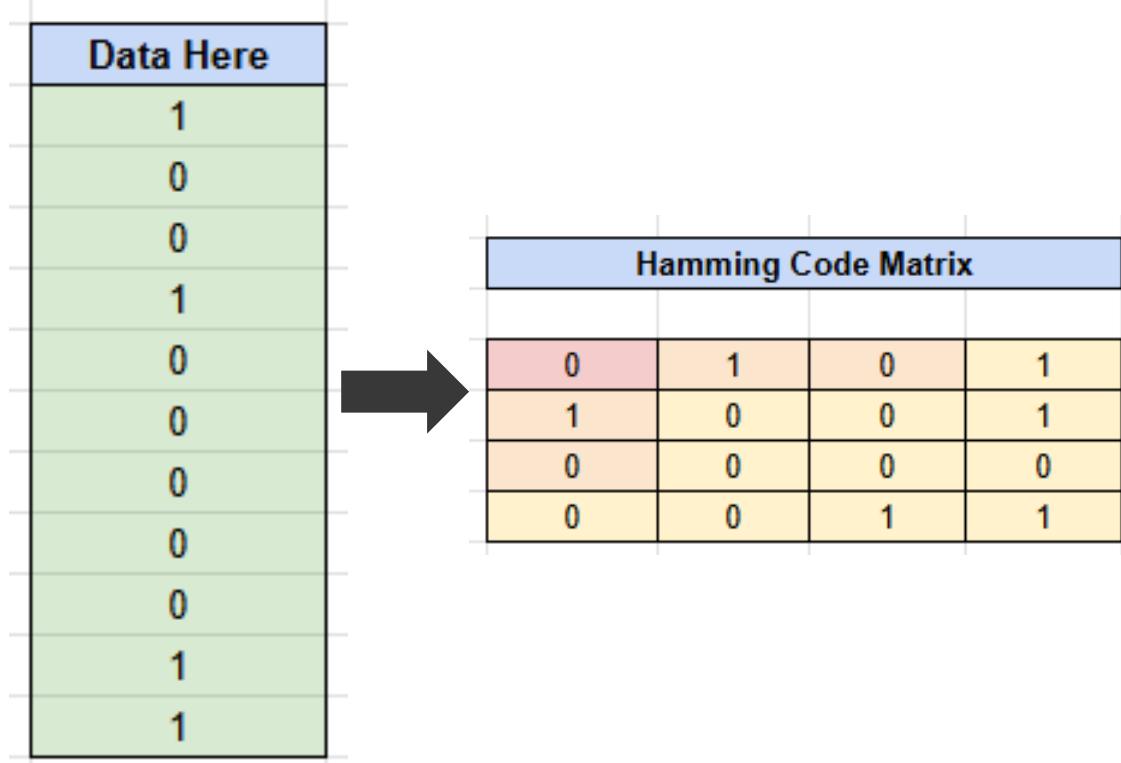
Numpy array quick intro

- > Build numpy array and check its spec
 - ↳ 4 cells hidden
- > Create an empty numpy array or a special array:
 - ↳ 2 cells hidden
- > Matrix multiply and numpy
 - ↳ 3 cells hidden
- > How it's done in numpy?
 - ↳ 2 cells hidden
- > That's "good enough" math and matrix for ya. Let's reconsider Hamming (16,11):
 - ↳ 11 cells hidden

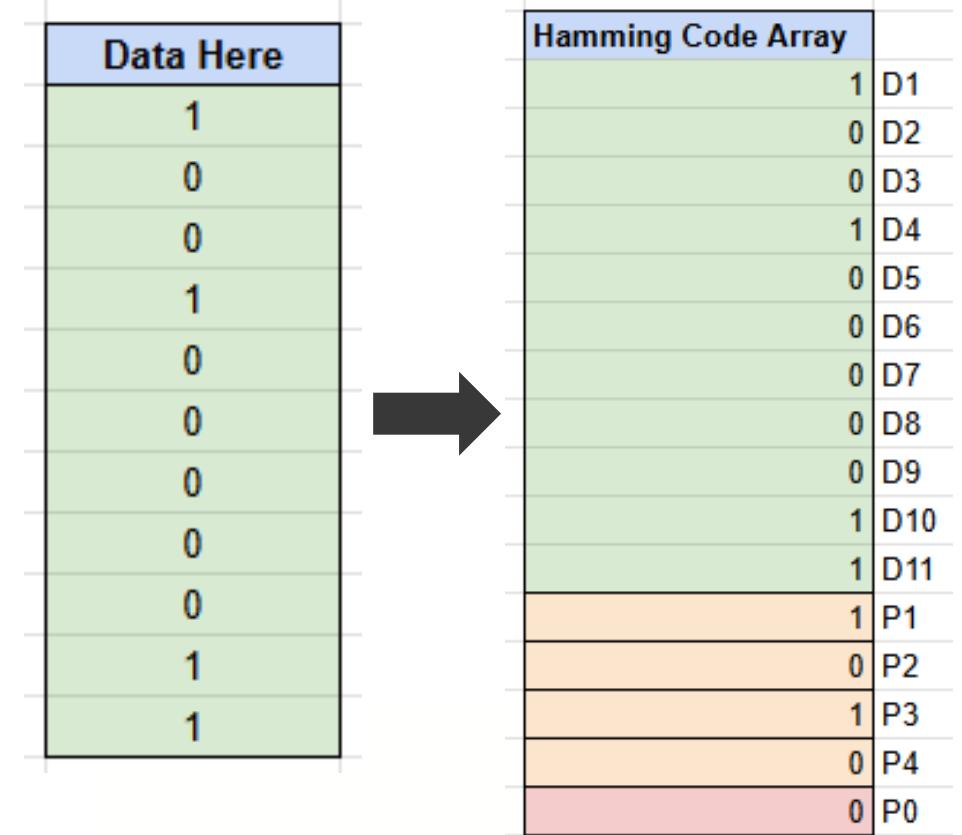
Two things about hamming (1/2):

1. Hamming code doesn't' need to be a 4x4 matrix

This can work:



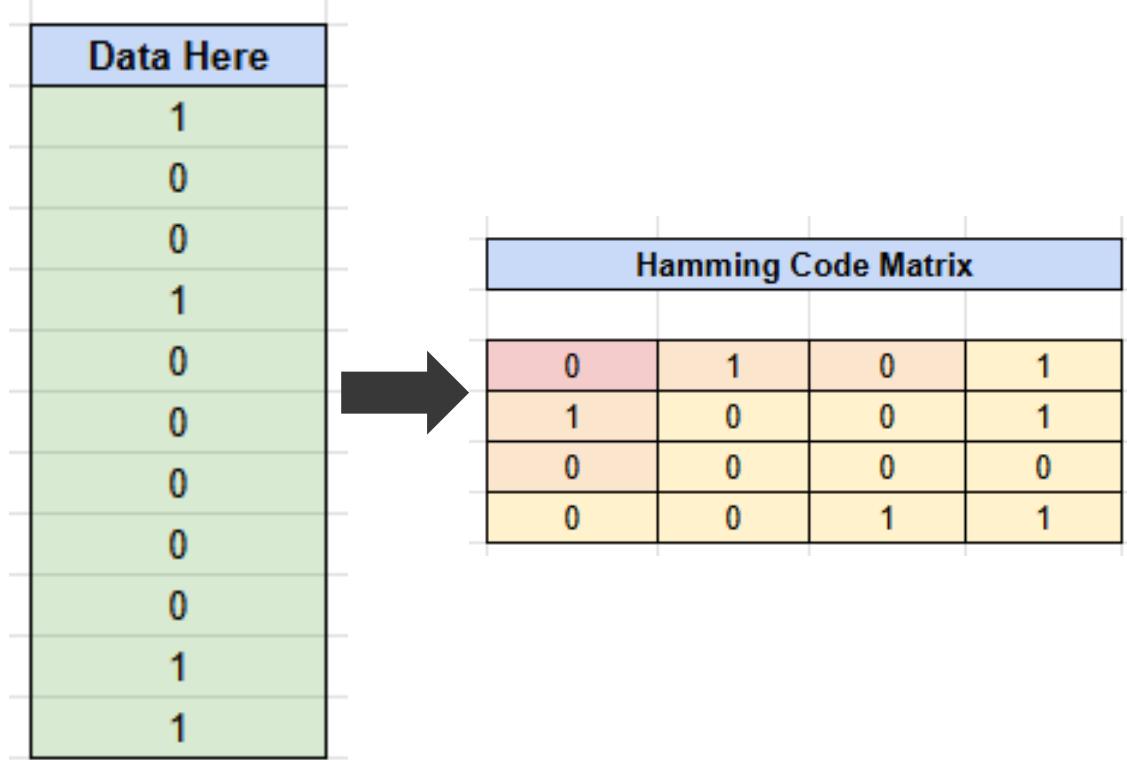
Can this work?



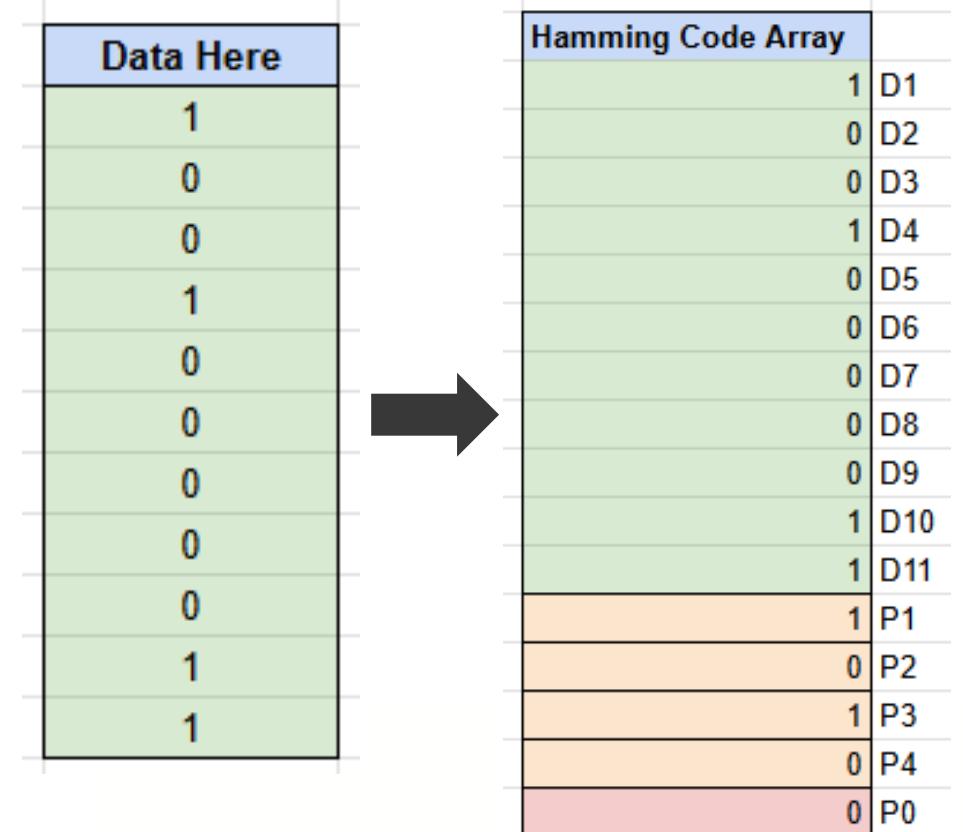
Two things about hamming (1/2):

1. Hamming code doesn't' need to be a 4x4 matrix

This can work:



Can this work?



Might have error code mapping issue, maybe?

Two things about hamming (2/2):

2. Array to array conversion will be so easy:

This is math:

$$P_{i,j} = \sum_{k=1}^n Q_{i,k} \times R_{k,j}$$

This is math application:

$$[2,3] \times [3,4] = [2,4]$$

This is Engineering application:

11 bits in

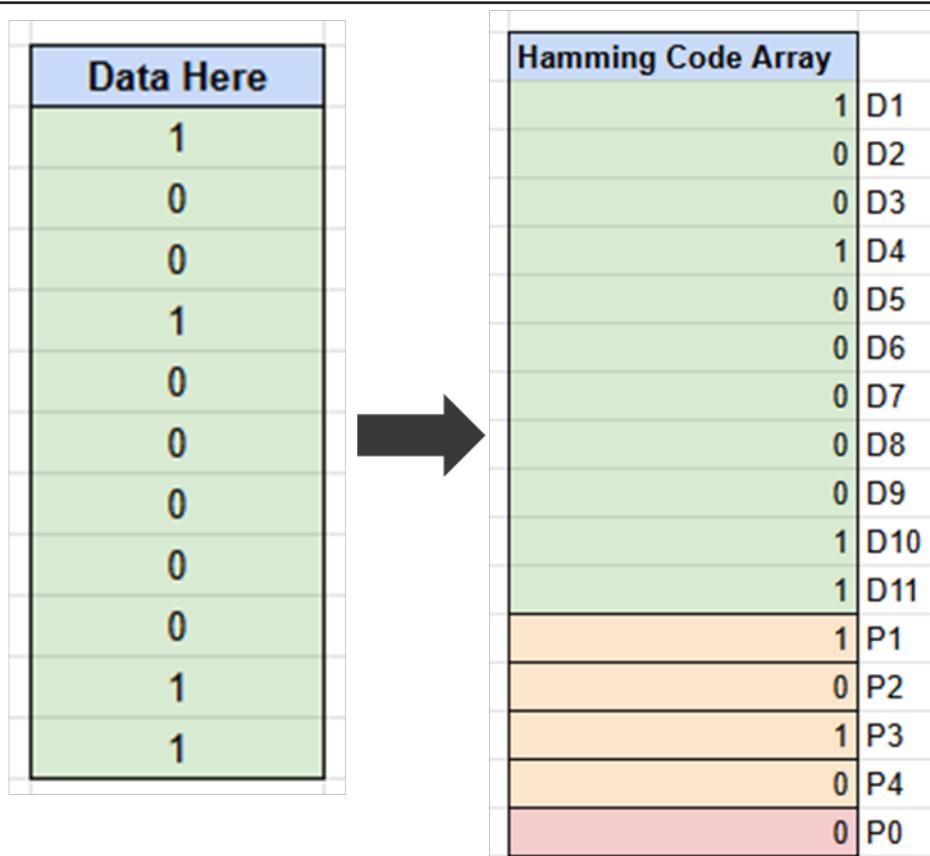
16 bits out

So instead of double for loops, we can do matrix operation to get it in one go?

$$[1,11] \times [?] = [1,16]?$$

Would it mathematically work?

Can this work?



I can show it could theoretically work in 3 steps.



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Step 1: how does matrix multiply works:

When “math” gives an example:

$$P_{i,j} = \sum_{k=1}^n Q_{i,k} \times R_{k,j}$$

$$P = \begin{bmatrix} Q_{11}R_{11} + Q_{12}R_{21} + \cdots + Q_{1n}R_{n1} & Q_{11}R_{12} + Q_{12}R_{22} + \cdots + Q_{1n}R_{n2} & \cdots & Q_{11}R_{1q} + Q_{12}R_{2q} + \cdots + Q_{1n}R_{nq} \\ Q_{21}R_{11} + Q_{22}R_{21} + \cdots + Q_{2n}R_{n1} & Q_{21}R_{12} + Q_{22}R_{22} + \cdots + Q_{2n}R_{n2} & \cdots & Q_{21}R_{1q} + Q_{22}R_{2q} + \cdots + Q_{2n}R_{nq} \\ \vdots & \vdots & \ddots & \vdots \\ Q_{m1}R_{11} + Q_{m2}R_{21} + \cdots + Q_{mn}R_{n1} & Q_{m1}R_{12} + Q_{m2}R_{22} + \cdots + Q_{mn}R_{n2} & \cdots & Q_{m1}R_{1q} + Q_{m2}R_{2q} + \cdots + Q_{mn}R_{nq} \end{bmatrix}$$

When “CS” gives an example:

$$E = AD = \begin{bmatrix} 10 & 20 & 30 \\ 40 & 50 & 60 \end{bmatrix} \times \begin{bmatrix} 2 & 3 & 5 & 7 \\ 11 & 13 & 17 & 19 \\ 23 & 29 & 31 & 37 \end{bmatrix} = \begin{bmatrix} 930 & 1160 & 1320 & 1560 \\ 2010 & 2510 & 2910 & 3450 \end{bmatrix}$$



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Step 1: how does matrix multiply works:

When “engineering” gives an example:

Indexing of M3

$$M3 = M1 \times M2 = \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix} \times \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix} = \begin{bmatrix} a_1 \times a_2 + b_1 \times c_2 & a_1 \times b_2 + b_1 \times d_2 \\ c_1 \times a_2 + d_1 \times c_2 & c_1 \times b_2 + d_1 \times d_2 \end{bmatrix} \quad [1,] \times [,1] \quad [1,1]$$

$$M3 = M1 \times M2 = \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix} \times \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix} = \begin{bmatrix} a_1 \times a_2 + b_1 \times c_2 & a_1 \times b_2 + b_1 \times d_2 \\ c_1 \times a_2 + d_1 \times c_2 & c_1 \times b_2 + d_1 \times d_2 \end{bmatrix} \quad [1,] \times [,2] \quad [1,2]$$

$$M3 = M1 \times M2 = \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix} \times \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix} = \begin{bmatrix} a_1 \times a_2 + b_1 \times c_2 & a_1 \times b_2 + b_1 \times d_2 \\ c_1 \times a_2 + d_1 \times c_2 & c_1 \times b_2 + d_1 \times d_2 \end{bmatrix} \quad [2,] \times [,1] \quad [2,1]$$

$$M3 = M1 \times M2 = \begin{bmatrix} a_1 & b_1 \\ c_1 & d_1 \end{bmatrix} \times \begin{bmatrix} a_2 & b_2 \\ c_2 & d_2 \end{bmatrix} = \begin{bmatrix} a_1 \times a_2 + b_1 \times c_2 & a_1 \times b_2 + b_1 \times d_2 \\ c_1 \times a_2 + d_1 \times c_2 & c_1 \times b_2 + d_1 \times d_2 \end{bmatrix} \quad [2,] \times [,2] \quad [2,2]$$

Matrix indexing: [i, j]

Row #

Column #



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Step 2: how to do parity checks

Recall:

Two ways to do parity check:

Logical operation

Math operation



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Step 2: how to do parity checks

Recall:

Two ways to do parity check:

Logical operation

i) XOR (the superior way, imo)
“exclusive OR”

XOR gate truth table

Input		Output
A	B	A XOR B
0	0	0
0	1	1
1	0	1
1	1	0

Math operation

ii) Modulo 2 (mod 2, %2)

0 modulo 2 =	0
2 modulo 2 =	0
4 modulo 2 =	0
1 modulo 2 =	1
3 modulo 2 =	1
5 modulo 2 =	1



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Step 3: how to preserve and operate numbers in matrix

```
## TODO: understand this matrix operation
input = np.array([
    [1, 2, 3, 4]])
Operation_I = np.array([
    [1,0,0,0,0],
    [0,1,0,0,1],
    [0,0,1,0,1],
    [0,0,0,1,1],
])
output = input @ Operation_I
```

```
shape check of input:
(1, 4)
shape check of operation matrix:
(4, 5)
shape check of output matrix:
(1, 5)
this is the output:
[[1 2 3 4 9]]
```

Since we can do SUM, let's do %2 this time for parity check.
Then we just need to “SUM” the right bits for P1, P2, ..., P4 and P0.



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Step 3: how to preserve and operate numbers in matrix

We already had a road map for P1-P4 and P0 from “classic hamming”.

0 1 0000	1 1 0001	2 0 0010	3 0 0011
4 0 0100	5 0 0101	6 1 0110	7 1 0111
8 1 1000	9 1 1001	10 1 1010	11 1 1011
12 0 1100	13 0 1101	14 0 1110	15 0 1111

<https://harryli0088.github.io/hamming-code/>

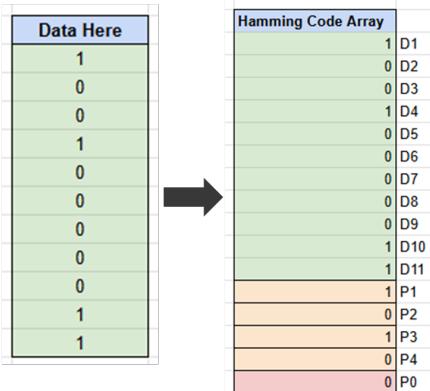


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Step 3: how to preserve and operate numbers in matrix

This can work...

```
## TODO: understand this matrix operation
input = np.array([
    [1, 2, 3, 4]])
Operation_I = np.array([
    [1,0,0,0,0],
    [0,1,0,0,1],
    [0,0,1,0,1],
    [0,0,0,1,1],
])
output = input @ Operation_I
```



As it becomes the combination of two simple matrix operation problem:

1.) To preserve D1-D13: an “Identity matrix”

(11,11)

```
[[1 0 0 0 0 0 0 0 0 0 0 0 0 0 0]
 [0 1 0 0 0 0 0 0 0 0 0 0 0 0 0]
 [0 0 1 0 0 0 0 0 0 0 0 0 0 0 0]
 [0 0 0 1 0 0 0 0 0 0 0 0 0 0 0]
 [0 0 0 0 1 0 0 0 0 0 0 0 0 0 0]
 [0 0 0 0 0 1 0 0 0 0 0 0 0 0 0]
 [0 0 0 0 0 0 1 0 0 0 0 0 0 0 0]
 [0 0 0 0 0 0 0 1 0 0 0 0 0 0 0]
 [0 0 0 0 0 0 0 0 1 0 0 0 0 0 0]
 [0 0 0 0 0 0 0 0 0 1 0 0 0 0 0]
 [0 0 0 0 0 0 0 0 0 0 1 0 0 0 0]
 [0 0 0 0 0 0 0 0 0 0 0 1 0 0 0]]
```

2.) To generate sum of the some: a “SoS matrix ”

(11,5)

```
[[1 1 1 0 1]
 [1 1 0 1 1]
 [1 0 1 1 1]
 [0 1 1 1 1]
 [1 1 0 0 1]
 [1 0 1 0 1]
 [1 0 0 1 1]
 [0 1 1 0 1]
 [0 1 0 1 1]
 [0 0 1 1 1]
 [1 1 1 1 1]]
```

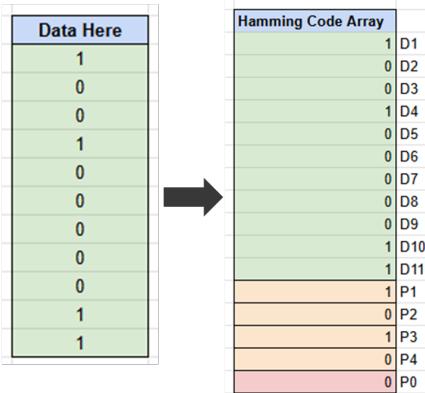


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Step 3: how to preserve and operate numbers in matrix

This can work...

```
## TODO: understand this matrix operation
input = np.array([
    [1, 2, 3, 4]])
Operation_I = np.array([
    [1,0,0,0,0],
    [0,1,0,0,1],
    [0,0,1,0,1],
    [0,0,0,1,1],
])
output = input @ Operation_I
```



As it becomes the combination of two simple matrix operation problem:

3.1) AKA a generator matrix: (11,11+5)

$$\begin{bmatrix} [1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0] & | & 1 & 1 & 1 & 0 & 1 \\ [0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0] & | & 1 & 1 & 0 & 1 & 1 \\ [0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0] & | & 1 & 0 & 1 & 1 & 1 \\ [0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0] & | & 0 & 1 & 1 & 1 & 1 \\ [0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0] & | & 1 & 1 & 0 & 0 & 1 \\ [0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0] & | & 1 & 0 & 1 & 0 & 1 \\ [0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0] & | & 1 & 0 & 0 & 1 & 1 \\ [0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0] & | & 0 & 1 & 1 & 0 & 1 \\ [0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0] & | & 0 & 1 & 0 & 1 & 1 \\ [0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0] & | & 1 & 0 & 0 & 0 & 1 \\ [0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1] & | & 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

3.2) Then do a %2 of the output, duh

Time to test our theory out in Google Colab!



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