Allocating Memory





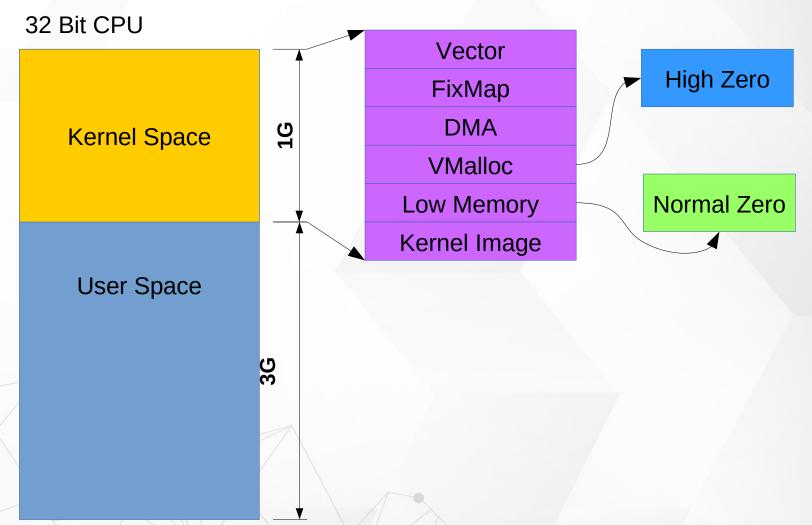
Kernel Allocate Memory API

- > kmalloc()
- kfree
- __get_free_page()
- __free_page()
- __get_free_pages()
- __free_pages()
- > vmalloc()
- vfree()
- ioremap()
- iounmap()





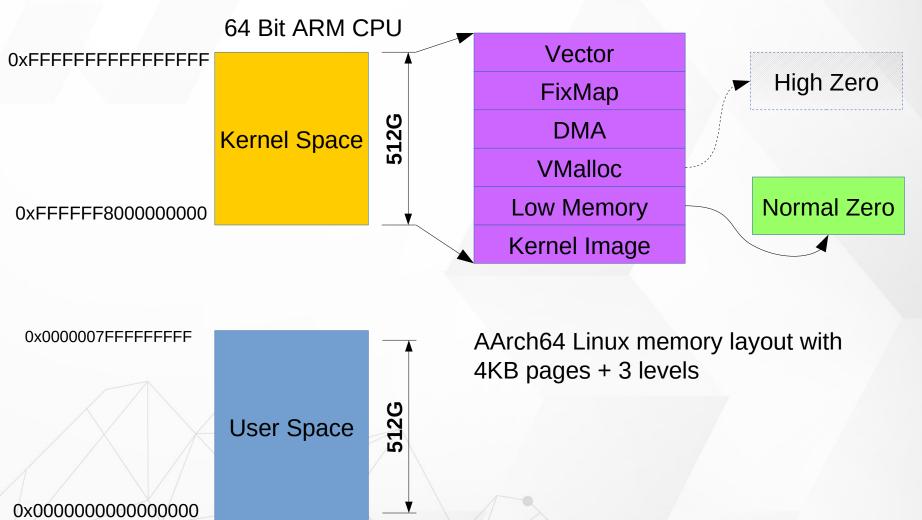
Kernel Memory Configuration







Kernel Memory Configuration





The Flags Argument

> kmalloc prototype

```
#include <linux/slab_def.h>
void *kmalloc(size_t size, int flags);
```

- Flags
 GFP_KERNEL
 - the most commonly used flag
 - Cannot be used in atomic context





kmalloc

- >> kmalloc(size_t size, gfp_t flags);
 - >> Allocates consecutive virtual/physical memory pages
 - Offset by PAGE_OFFSET
 - PAGE_OFFSET 2^12 (4K)
 - No changes to page tables
 - Marian Allocate memory limit: 128K







- Normal memory
 - GFP_KERNEL
- DMA-capable memory
 - GFP_DMA
 - Platform dependent
 - >> First 16MB of RAM on the x86 for ISA devices
 - PCI devices have no such limit
- Migh memory
 - GFP_HIGHMEM
 - Platform dependent
 - > 32-bit addressable range





Driver SHT21

```
static int sht21 probe(struct i2c client *client,
    const struct i2c device id *id)
    struct device *dev = &client->dev;
    struct device *hwmon dev;
    struct sht21 *sht21;
    if (!i2c check functionality(client->adapter,
                     I2C FUNC SMBUS WORD DATA)) {
        dev err(&client->dev,
            "adapter does not support SMBus word transactions\n");
        return - ENODEV:
                                                kmalloc and flag
   sht21 = devm kzalloc(dev, sizeof(*sht21), GFP KERNEL);
   if (!sht21)
        return - ENOMEM:
    sht21->client = client;
   mutex init(&sht21->lock);
    hwmon dev = devm hwmon device register with groups(dev, client->name,
                               sht21, sht21 groups);
    return PTR ERR OR ZERO(hwmon dev);
```





The Size Argument

- Xernel manages physical memory in pages
 - Needs special management to allocate small memory chunks
- Linux creates pools of memory objects in fixed sizes
 - 32-byte
 - 3 64-byte,
 - 128-byte memory objects
- kmalloc create smallest allocation unit
 - 32 or 64 bytes
- Largest portable allocation unit is 128KB



Supplement





get_free_page and Friends

- For allocating big chunks of memory, it is more efficient to use a page-oriented allocator
- To allocate pages, call

Kernel API

```
/* returns a pointer to a zeroed page */
get_zeroed_page(unsigned int flags);

/* does not clear the page */
__get_free_page(unsigned int flags);

/* allocates multiple physically contiguous pages */
__get_free_pages(unsigned int flags, unsigned int order);
```





get_free_page and Friends

- flags
 - >> Same as flags for kmalloc
- order
 - → Allocate 2order pages
 - order = 0 for 1 page
 - order = 3 for 8 pages
 - Can use get_order(size) to find out order
 - Maximum allowed value is about 10 or 11
 - See /proc/buddyinfo statistics





Page and order

Order		
2 ^ 10	1024 Pages	
2 ^ 9	512 Pages	
2 ^ 8	256 Pages	
2 ^ 7	128 Pages	
2 ^ 6	64 Pages	
2 ^ 5	32 Pages	
2 ^ 4	16 Pages	
2 ^ 3	8 Pages	
2 ^ 2	4 Pages	
2 ^ 1	2 Pages	
2^0	1 Page	



get_free_page and Friends

- Subject to the same rules as kmalloc
- To free pages, call

```
void free_page(unsigned long addr);
void free_pages(unsigned long addr, unsigned long order);
```

- Make sure to free the same number of pages
 - Or the memory map becomes corrupted





Sample – get_free_page (1)

3-1-2-buddy

```
static int buddy init(void)
  printk("Buddy driver installed...\n");
  order = 4;
  ptr = (void *) __get_free_pages(GFP_KERNEL, order);
  if (ptr)
    printk("2^%d * %ld bytes allocated\n", order, PAGE_SIZE);
  order2 = get_order(PAGE_SIZE * 4);
  ptr2 = (void *) __get_free_pages(GFP_KERNEL, order2);
  if (ptr2)
    printk("2^%d * %ld bytes allocated\n", order2, PAGE_SIZE);
  return 0;
```





Sample – get_free_page (2)

```
static void buddy_exit(void)
{
    printk("buddy driver removed...\n");
    printk("buddy driver removed order %d ...\n", order);
    printk("buddy driver removed order2 %d ...\n", order2);

    free_pages((unsigned long)ptr, order);
    free_pages((unsigned long)ptr2, order2);
}
```





vmalloc and Friends

- Allocates a virtually contiguous memory region
 - Not consecutive pages in physical memory
 - Each page retrieved with a separate alloc_page call
 - Less efficient
 - Can sleep (cannot be used in atomic context)





>> vmalloc-related prototypes

```
#include unx/vmalloc.h>

void *vmalloc(unsigned long size);

void vfree(void * addr);

void *ioremap(unsigned long offset, unsigned long size);

void iounmap(void * addr);
```





vmalloc and Friends

- ioremap builds page tables
 - Does not allocate memory
 - Takes a physical address (offset) and return a virtual address
 - Useful to map the address of a Hardware buffer to kernel space
 - >> Should use **readb** and other relation functions to access remapped memory

