



**L-Università ta' Malta**  
**Faculty of Information &  
Communication Technology**

## Assignment Report

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# 1 Introduction

This assignment is implemented in c++.

## 2 Question 1

### 2.1 How the problem was tackled

First off, an FSA was dedigned to accept numbers. Then the required structures where initialized in code. The FSA was translated to code in the form of transition table. Then, code was developed to successfully read the numbers with a stack and rollback capabilities. This was done according to the pseudo code provided in the text book. Following that, more and more elements where added, testing the lexer as it grew.

### 2.2 Data Structures

The following structures are defined in "transitions.h". Three enumerations were done to keep track of classifiers, states and token types. The enum for token types is divided into two parts (shown via the comment). The second part is used exclusively by the parser.

An array of accepting states was declared to keep track of the states which were final. The transisition table was implemented as a 2d array with classifiers as columns and states as rows.

The char\_cat and type tables were implemented as functions.

### 2.3 FSA

The following are parts of the FSA in the order they where added to the transition table.

## 3 Question 2

### 3.1 How the problem was tackled

For this question, a single path was taken through the EBNF. This path was implemented as boolean functions to test the system.

Listing 1: "Initial boolean functions"

```
bool parse_statement() {
```

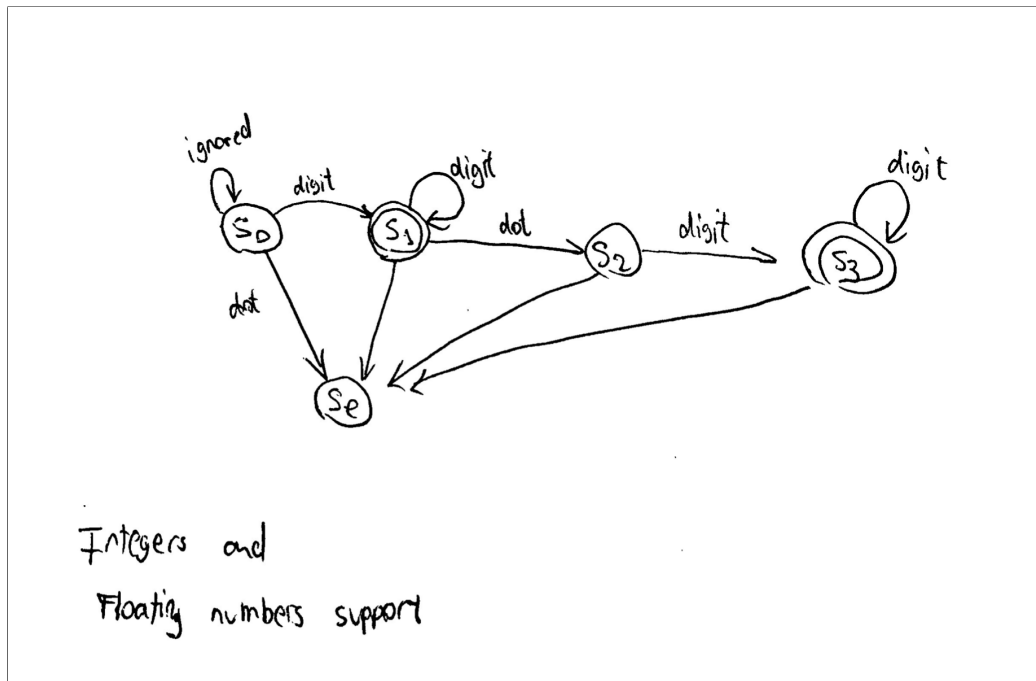


Figure 1: FSA (part concerned with numbers)

```

    if ( parse_rtrn_statement() ) {
        if ( parse_semicolon() ) { return true; }
        else { return false; }
    }
    /* TODO: Add other types of statements */

    else { return false; }
}

bool parse_semicolon() {
    if ( tok.type == s_colon ) { tell("semicolon");
        advance(); return true; }
    else { return true; }
}

bool parse_rtrn_statement() {
    if ( !parse_return() ) { return false; }
    if ( !parse_expression() ) { return false; }
    return true;
}

```

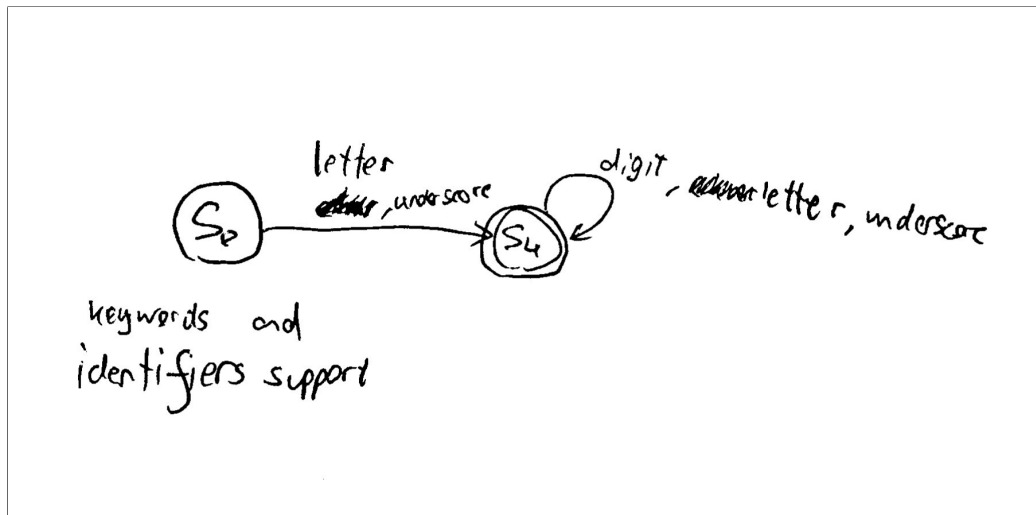


Figure 2: FSA (part concerned with identifiers and keywords)

```

bool parse_return() {
    if ( tok.value.compare("return") == 0 ) { tell("
        return"); advance(); return true; }
    else { return false; }
}
  
```

Later on, a node was declared with binary tree properties. The node holds a token and points to a left node and a right node. The boolean functions were translated to return a pointer to the node created instead of true and a nullptr instead of false.

Listing 2: "The change to return node pointer"

```

AST* parse_rtrn_statement() {
    tell("rtrn_stmnt");
    AST* rtn_node = parse_return();
    if ( rtn_node == nullptr ) {
        return nullptr;
    }

    AST* expression_node = parse_expression();
    if ( expression_node == nullptr ) {
        return nullptr;
    }
}
  
```

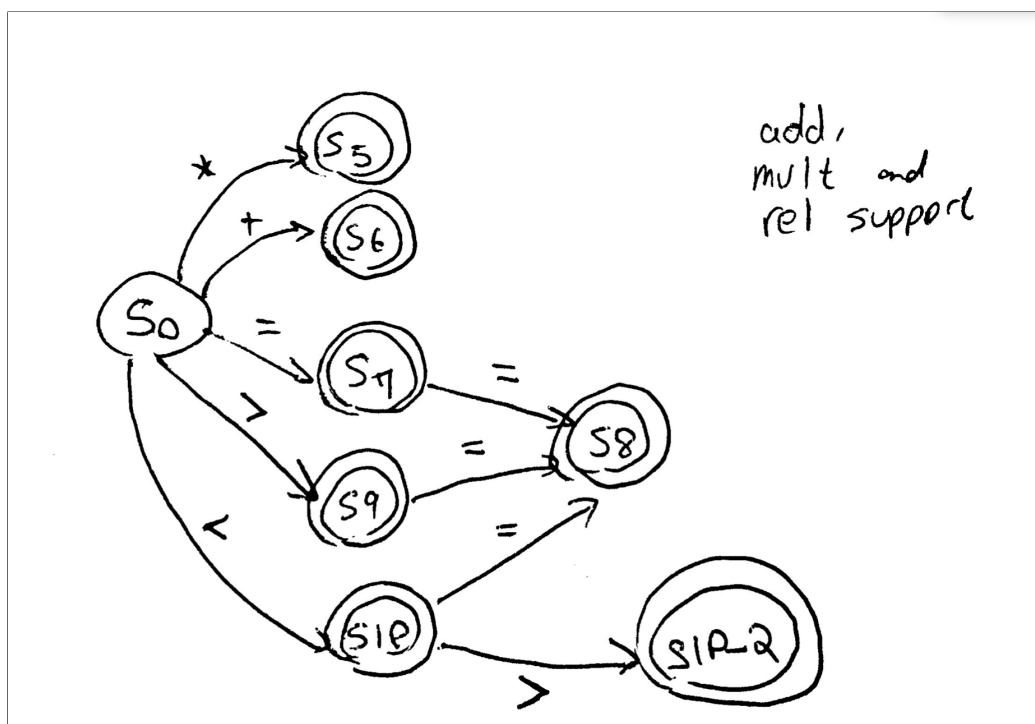


Figure 3: FSA (part concerned with operators)

```

token new_tok;
new_tok.type = rtn_statement;
return make_node(new_tok, rtn_node,
    expression_node);
}

```

An example of tree building would be the following:

## 3.2 Types of parsing functions

The section explains how the control forms were implemented within the parsing functions.

### 3.2.1 Sequence

An example from the EBNF that uses sequence is FormalParam. This was implemented by using if statements to and return nullptr if the required node is not found.

Listing 3: "Sequence example"

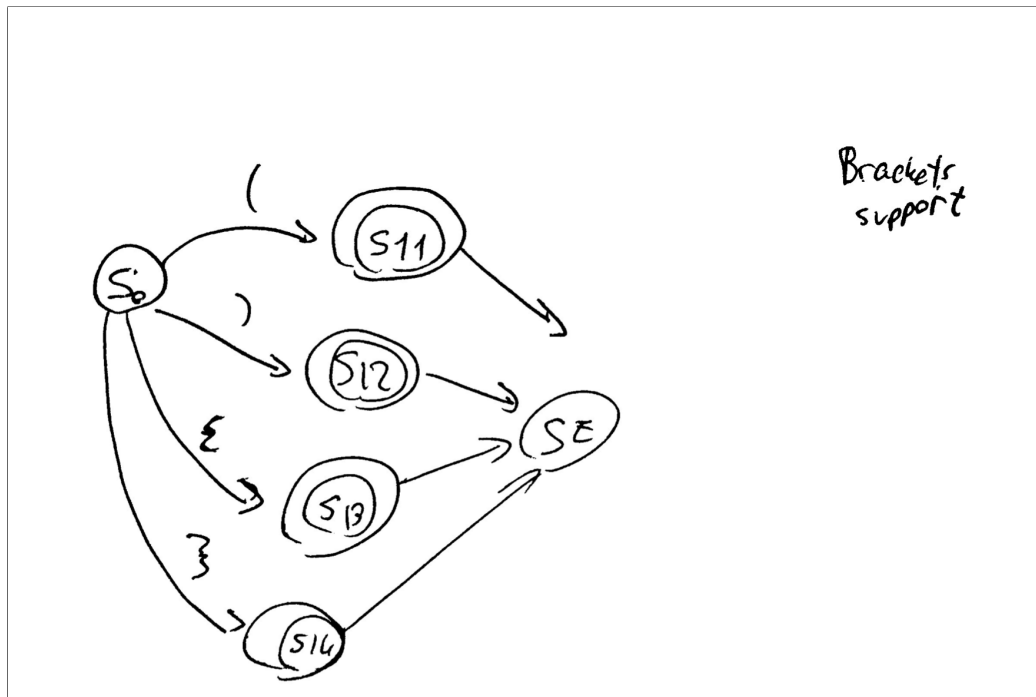


Figure 4: FSA (part concerned with brackets)

```

AST* parse_rtrn_statement() {
    tell("rtrn_stmnt");
    AST* rtn_node = parse_return();
    if ( rtn_node == nullptr ) {
        return nullptr;
    }

    AST* expression_node = parse_expression();
    if ( expression_node == nullptr ) {
        return nullptr;
    }

    token new_tok;
    new_tok.type = rtn_statement;
    return make_node(new_tok, rtn_node,
        expression_node);
}

```

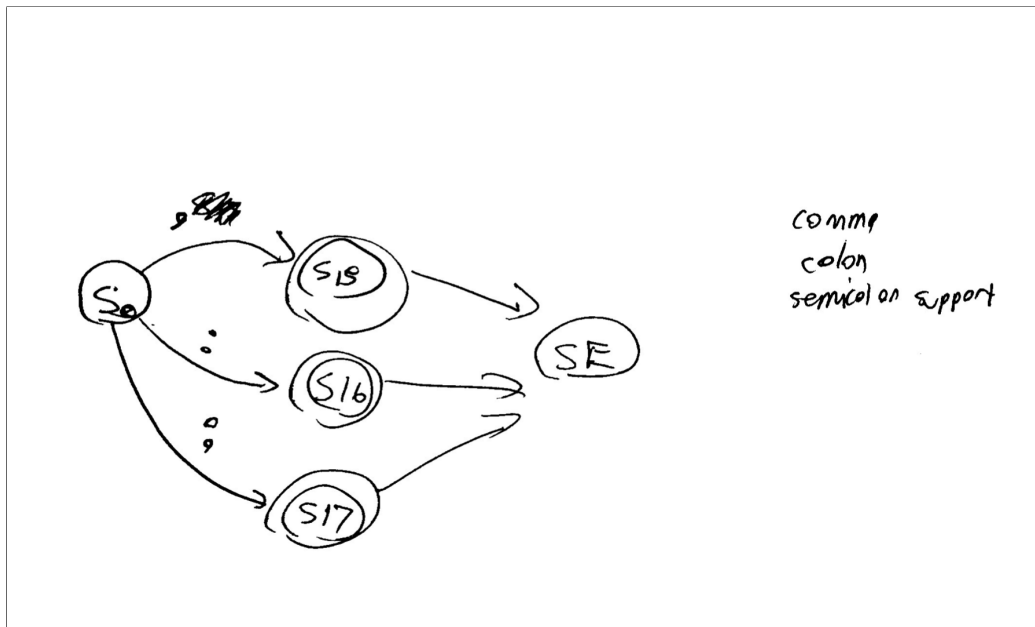


Figure 5: FSA (part concerned with colons and comma)

### 3.2.2 Choice

Statement is an example that uses choices within the EBNF. These types are implemented by trying to get a valid (non nullptr) node. If the function successfully gets a node, it returns the said node in the correct form. If not (the node was a nullptr) it tries again with the next choice.

Listing 4: "Choice example"

```
AST* parse_statement() {
    tell("statement");
    AST* rtn_statement_node = parse_rtrn_statement();
    if ( rtn_statement_node != nullptr ) {
        AST* semicolon_node = parse_semicolon();
        if ( semicolon_node != nullptr ) {
            token new_tok;
            new_tok.type = statement;
            return make_node(new_tok,
                            rtn_statement_node, semicolon_node);
        }
        else { return nullptr; }
    }
}
```

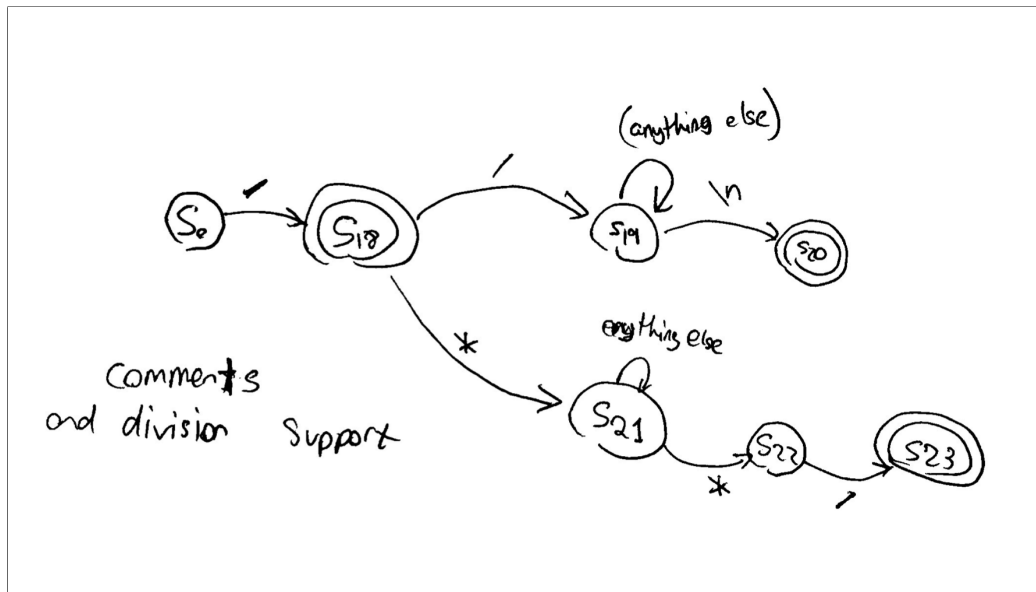


Figure 6: FSA (part concerned with comments and division operator)

```

AST* print_statement_node = parse_print_statement()
;
if ( print_statement_node != nullptr ) {
    AST* semicolon_node = parse_semicolon();
    if ( semicolon_node != nullptr ) {
        token new_tok;
        new_tok.type = statement;
        return make_node(new_tok,
            print_statement_node, semicolon_node);
    }
    else { return nullptr; }
}

AST* var_decl_node = parse_var_decl();
if ( var_decl_node != nullptr ) {
    AST* semicolon_node = parse_semicolon();
    if ( semicolon_node != nullptr ) {
        token new_tok;
        new_tok.type = statement;
        return make_node(new_tok, var_decl_node,
            semicolon_node);
    }
}

```



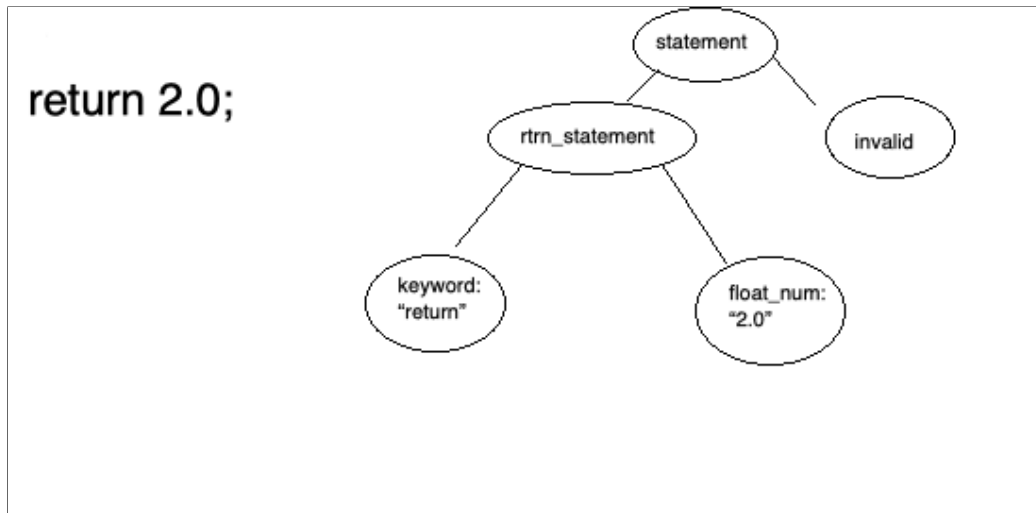


Figure 7: Example of AST

```

    else { return nullptr; }
}

AST* assignment_node = parse_assignment();
if ( assignment_node != nullptr ) {
    AST* semicolon_node = parse_semicolon();
    if ( semicolon_node != nullptr ) {
        token new_tok;
        new_tok.type = statement;
        return make_node(new_tok, assignment_node,
            semicolon_node);
    }
    else { return nullptr; }
}

AST* block_node = parse_block();
if ( block_node != nullptr ) {
    token new_tok;
    new_tok.type = statement;
    return make_node(new_tok, block_node, nullptr);
}

AST* while_statement_node = parse_while_statement()
;

```

```

if ( while_statement_node != nullptr ) {
    token new_tok;
    new_tok.type = statement;
    return make_node(new_tok, while_statement_node,
        nullptr);
}

AST* if_statement_node = parse_if_statement();
if ( if_statement_node != nullptr ) {
    token new_tok;
    new_tok.type = statement;
    return make_node(new_tok, if_statement_node,
        nullptr);
}

AST* for_statement_node = parse_for_statement();
if ( for_statement_node != nullptr ) {
    token new_tok;
    new_tok.type = statement;
    return make_node(new_tok, if_statement_node,
        nullptr);
}

AST* func_decl_node = parse_func_decl();
if ( func_decl_node != nullptr ) {
    AST* semicolon_node = parse_semicolon();
    if ( semicolon_node != nullptr ) {
        token new_tok;
        new_tok.type = statement;
        return make_node(new_tok, func_decl_node,
            semicolon_node);
    }
    else { return nullptr; }
}

else { return nullptr; }
}

```

### 3.2.3 Options

Options were implemented by declaring a node but the node doesn't have to be filled. This means the node can be `nullptr` without the function quitting.

An example of this is the 'else' part of the if statement.

Listing 5: "Option example"

```
AST* parse_if_statement() {
    tell("if_statement");
    AST* if_node = parse_if();
    if ( if_node == nullptr)
        {
            return nullptr; }

    AST* l_par_node = parse_l_par();
    if ( l_par_node == nullptr)
        { return
            nullptr; }

    AST* expression_node = parse_expression();
    if ( expression_node == nullptr )
        { return
            nullptr; }

    AST* r_par_node = parse_r_par();
    if ( r_par_node == nullptr )
        { return
            nullptr; }

    AST* block_node = parse_block();
    if ( block_node == nullptr )
        { return
            nullptr; }

    AST* else_node = parse_else();
    AST* else_block_node = parse_block();

    if ( (else_node == nullptr && else_block_node !=
        nullptr)
        || (else_node != nullptr && else_block_node ==
            nullptr) )
        { return nullptr; }
```

```

AST* block = nullptr;
token new_tok_0;
new_tok_0.type = if_block;
block = make_node(new_tok_0, block_node,
    else_block_node);

token new_tok_1;
new_tok_1.type = if_head;
return make_node(new_tok_1, expression_node, block)
    ;
}

```

### 3.2.4 Repetition

Repetition is the most complex of the control forms. A recursive function was developed for the tree to add a node to the rightmost. This enables parsing functions to add elements to the rightmost node of the tree from an iterative loop. An example of repetition is the simple expression.

Listing 6: "Repetition example"

```

AST* parse_simp_expression() {
    tell("simp");
    AST* term_node_l = parse_term();
    if ( term_node_l == nullptr ) {
        return nullptr; }

    AST* add_op_node      = nullptr;
    AST* term_node_r      = nullptr;
    AST* temp_add_op_node = nullptr;
    AST* temp_term_node_r = nullptr;

    for(int i=0;; ++i) {
        temp_add_op_node = parse_add_op();
        if ( temp_add_op_node == nullptr ) {
            break; }

        temp_term_node_r = parse_term();
        if ( temp_term_node_r == nullptr ) {
            return nullptr; }
    }
}

```

```

        if ( i==0 ) {
            add_op_node = temp_add_op_node;
        }
        add_to_rightmost( term_node_r , temp_add_op_node ,
                           temp_term_node_r );
    }

    if ( add_op_node == nullptr ) {
        return term_node_l;
    }

    return make_node( add_op_node->data , term_node_l ,
                      term_node_r );
}

```

### 3.3 Types of nodes

The way the node was implemented, only a binary tree can be built. Therefore a problem arose for EBNF statements like FuncDecl or ForStatements since they contain too much information to be stored in just two child nodes. The workaround used is by creating a sub tree as a node for the statement. The tree structure of FuncDecl, IfStatement and ForStatement is included.

## 4 Appendix

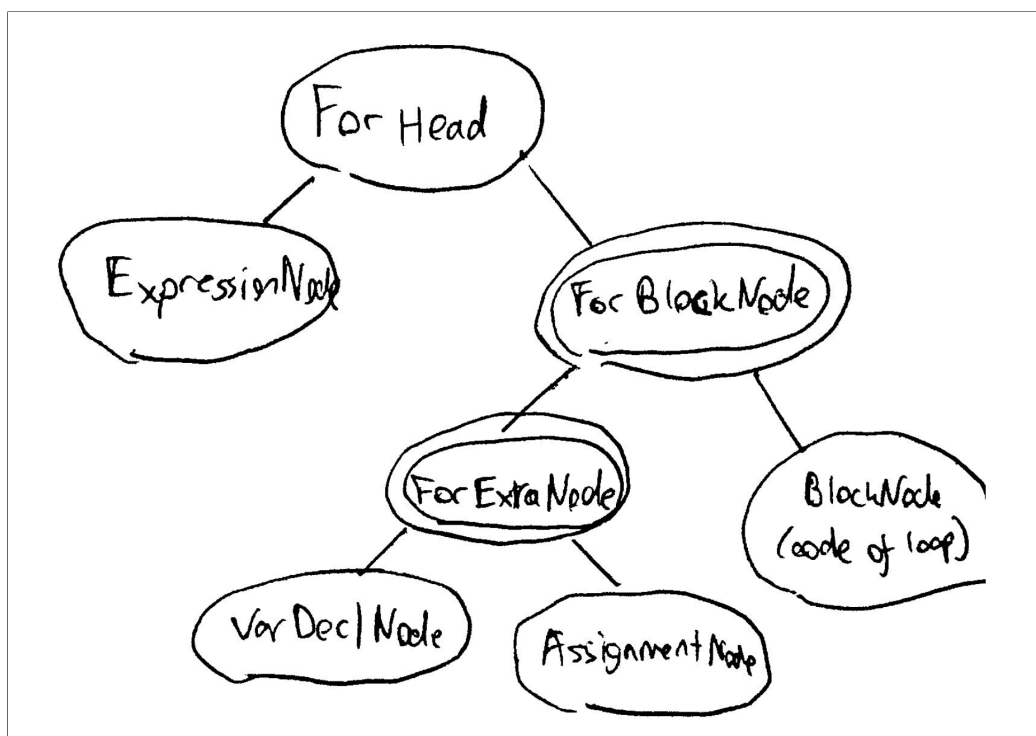


Figure 8: Tree for ForStatement

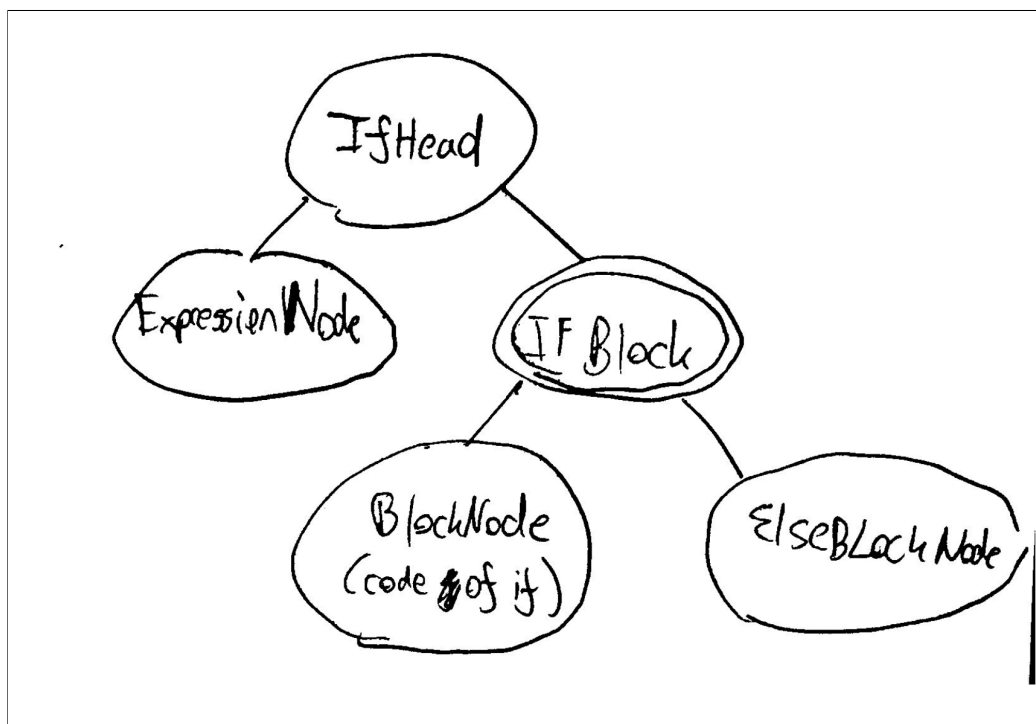


Figure 9: Tree for IfStatement

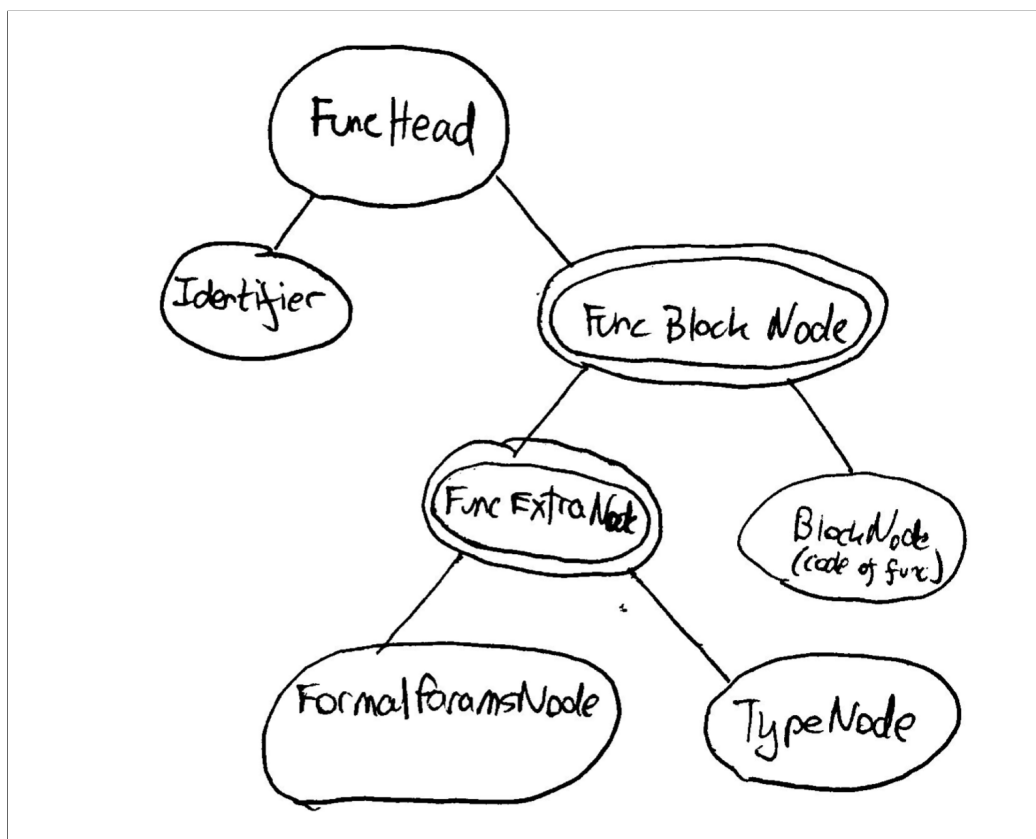


Figure 10: Tree for `FuncDeclStatement`