# Programming Assignment Lecture I

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### Basic Info

- Course Content
  - ► PA1-PA??
- ► Time
  - ▶ ???
- Assessment methods
  - Submit project and report before deadline.
  - An act of plagiarism is absolutely forbidden.
  - ▶ For more details, see SUBMIT REQUIREMENT WEBPAGE.
- Instruction
  - https://nju-ics.gitbooks.io/ics2017-programming-assignment/
  - Please check the "news" module of website at least once each day.

### Teacher and T.A.s

emmmmm

### Resources

$$\begin{aligned} & Plantform \ and \ tools \ IA-32 + GNU/Linux + gcc + C \\ & Guidebook \ https://nju-ics.gitbooks.io/ics2017-programming-assignment/ \\ & Skeleton \ https://github.com/NJU-ProjectN/ics-pa \end{aligned}$$

### Tip

You can download the PDF or epub version of guide in github.

### Programming Assignment Lecture I

- 1 From Computer System, to ICS, to PA
- $ext{@}$  How we emulate a computer?  $\longrightarrow$  The story of computer
- $\bigcirc$  Help you do it!  $\longrightarrow$  Brand new PA based on AM

## Why we need learn ICS?

### Motivation

```
Question
int main()
{
    printf("Hello World");
    return 0;
}
```

What the computer are doing when you execute the program above

### Tip

This may appear in exams.

### System Stack

Application
Algorithm
Programming Language
Operating System/Virtual Machines
Instruction Set Architecture
Micro-architecture
Register-Transfer Level
Gates
Circuits
Devices
Physics

### What you'll get after finishing PA

#### You'll

- Get Systems thinking
- ▶ Understand how program run on a computer
- Enhance coding ability
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- Way Complete a tiny but entire computer system and run program on it.
  - PA NEMU (i.e. NJU Emulator)

#### Question

### What is an *emulator*?

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### What is an *emulator*?

### Emulator -Wikipedia

In computing, an **emulator** is hardware or software that enables one computer system (called the *host*) to behave like another computer system (called the *guest*).

An emulator typically enables the host system to run software or use peripheral devices designed for the guest system.

### What is an emulator?









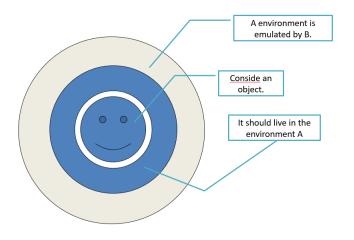
### What is an emulator?











### Question

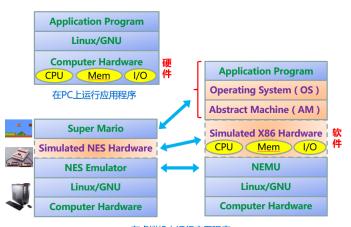
Does the object know he is living in a virtualized environment?

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### Task of PA

#### Task of PA

To realize **NEMU**, a simplified x86 system-wide **emulator**.



### **Transition**

So how we emulate a computer?

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### **DLC**

- Most of you have completed the Digital Logic Circuit Course. (Taught by zzs whj)
  - Adder
  - Register
  - Multiplexer
  - etc.

### DLC

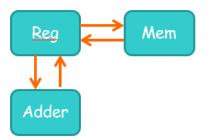
- Most of you have completed the Digital Logic Circuit Course. (Taught by zzs whj)
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  - etc.
- ▶ Logic Gate → digital logic device → Computer

### The simplest computer- Turing machine

Ultimate goal of computer

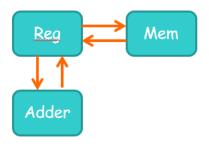
Run programs on it.

To place programs Memory
To process data Adder
To store temporary results efficiently Reg



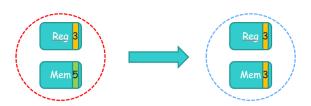
### The simplest computer- Turing machine

- ightharpoonup TRM = Reg + Adder + Mem
  - These are all what we learned in DLC!
- Computers read data from Mem, store them to Reg, do some calculate and write the result to Mem.



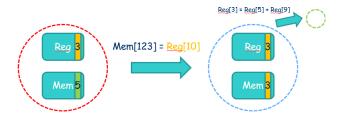
### State of TRM

- ▶ Reg and Mem are all Sequential Logic Device. They can store value (i.e. state).
- ► The process of computer working = These sequential logical devices transfer one state to another state



### Instructions of TRM

- This large state machine does have abundant states.
- ▶ But how one state transfer to another is not at liberty.
- ▶ Instructions = Commands that guide the state transition of computer
  - ▶ R/W Mem ,Computing
- ► Program = A lot of instructions



### Instruction Set

- ► Instruction Set = All things computer can do
  - mov, inc,jmp

### Question

What can TRM do with just these instructions?

### Instruction Set

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  - ▶ mov, inc,jmp

#### Question

What can TRM do with just these instructions?

Computability theory shows that TRM can do anything!

```
loop jmp
add inc,inc,inc,inc,inc...
multiply add,add,add,add...
```

function call mov (transfer of parameters),jmp(transfer of control)

### Instruction Set

- However, the speed of TRM is so slow.
- Why do we add more powerful instructions to the computer?
- Modern computer Instruction Set(x86, mips)
  - mul,div, cmp, bit operation, string manipulation
  - ▶ Adder → ALU



### Input & Output

#### Question

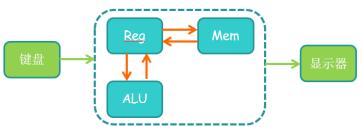
Now, it seems that we have a powerful computer. Are there any unsatisfactory drawbacks?

### Input & Output

#### Question

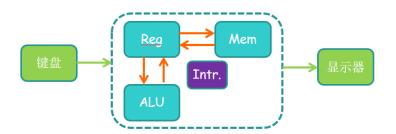
Now, it seems that we have a powerful computer. Are there any unsatisfactory drawbacks?

- It can only do computing!
- ► To interact with the outside world, we need to add I/O devices
- plus with some relative instructions.



### Interrupt & Exception

- Besides running regular program, computer should deal with emergencies whenever possible.
- Such as,
  - ▶ Inner exception: division by 0, access violation, trap etc.
  - Outer interruption: Keyboard, device ready, etc.
- ▶ We need add asynchronous processing unit to deal with emergencies.



### Multitask

### Question again

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- ▶ Monotask → Mutitask
- ► Time-division multiplexing: Improve the efficiency of resources usages.
  - ► CPU, I/O Device: Take turns!

#### Question

Can tasks use memory exclusively by turns? What else should we do in practice?

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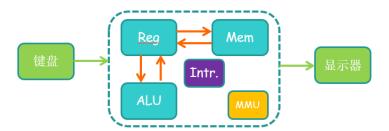
#### Answer

Certainly not. Exclusive memory using will cause inefficiency. We should add MMU for memory protection.

### AM(Abstract Machine)

$$\mathsf{AM} = \mathsf{TRM} + \mathsf{IOE} + \mathsf{ASYE} + \mathsf{PTE} + \mathsf{MPE}$$

- TRM(Turing Machine)
- ► IOE(I/O Extension)
- ASYE(Asynchronous Extension)
- ► PTE(Protection Extension)
- MPE(Multi-Processor Extension) (not included in PA)



### How to realize NEMU?

By now, we have known what a modern computer should look like, then how we emulate it? Or what is the constituent of NEMU?

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#### What tools we already have?

- Variables, data structure
- Algorithm

Devices C variables, arrays, link list, etc.

Logic functions Algorithmic functions

Runtime support Loop

## E.g. cpu

```
struct CPU_State
{
    rtlreg_t eax, ecx, edx, ebx, esp, ebp, esi, edi;
    vaddr_t eip;
    EFLAGS;
    CRs;
    ......
}
```



## E.g. memory

uint8\_t pmem[PMEM\_SIZE];



## The simplest computer- Turing machine

Architecture of the simplest computer

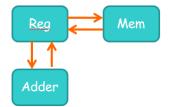
To place programs Memory

To process data Adder

To store temporary results efficiently Reg

Working mode of the simplest computer

- Fetch instruction from Mem using PC.
- Execute instruction.
- Update PC.





### TRM in NEMU

#### Architecture of the simplest computer

Memory **pmem array**(128MB nemu/src/memory/memory.c)

Adder(ALU) functions

Reg CPU\_state struct (nemu/include/cpu/reg.h)

#### Working mode of the simplest computer

Function cpu\_exec() (nemu/src/monitor/cpu-exec.c)

- Fetch instruction from Mem using PC.
- Execute instruction.
- ▶ Update PC.

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- $\textbf{ 3} \ \mathsf{Help} \ \mathsf{you} \ \mathsf{do} \ \mathsf{it}! \ \longrightarrow \ \mathsf{Brand} \ \mathsf{new} \ \mathsf{PA} \ \mathsf{based} \ \mathsf{on} \ \mathsf{AM}$

### Structure of PA-1

```
ics2017
|---nanos-lite  # mini operating system kernel
|---navy-apps  # apps
|---nemu  # NEMU
|---nexus-am  # abstract machine
```

### Structure of PA-2

Navy-apps

libos, newlib

Nanos-lite

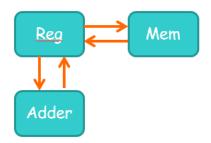
n86 Nexus-am

NEMU

Application			
Algorithm			
Programming Language			
Operating System/Virtual Machines			
Instruction Set Architecture			
Micro-architecture			
Register-Transfer Level			
Gates			
Circuits			
Devices			
Physics			

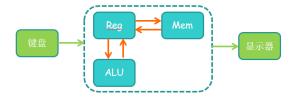
### PA0 and PA1

- ▶ PA0 Environment configuration
- ▶ PA1 Simple monitor



## PA2-Von Neumann Computer System (TRM+IOE)

- ▶ Instructions system + I/O
- Run interesting apps on NEMU

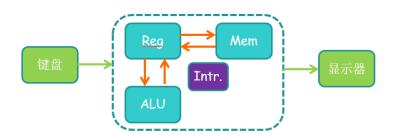






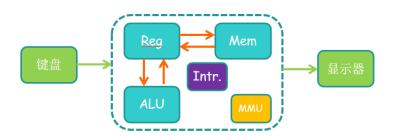
## PA3- Asynchronous Extension

- Run tiny operating system Nanos-lite on NEMU.
  - ramdisk, fs, raw img loader
- ▶ Implement some library functions which packing the system call .
- ▶ App calls library functions and run PAL.



## PA4- Multitask(Protection Extension)

- Virtual storage.
- Interruption.
- Run pal,typing game,clock at the same time on Nanos-lite.



### PA5- Performance

- ► Elf32 loader
- cache
- ► TLB
- profiler, performance optimization
- ▶ JIT (i.e just-in-time) compilation

实验内容(括号中为新方案)	持续时间/周	预计耗时/小时	代码量/行
PAO - 开发环境配置(不变)	1	10	无
PA1 - 简易调试器(不变)	3	30	400
PA2 - 指令系统 (PA2 - 冯诺依曼计算机系统)	6 (4)	60 (30)	800 (300)
课时不足可选择完成到PA2 [小计]	(8)	(70)	(700)
PA3 - 存储管理 (PA3 - 异常控制流)	4 (3)	50 (30)	500 (200)
PA4 - 中断与I/O (PA4 - 分时多任务)	3 (3)	30 (30)	300 (200)
无 (PA5 - 程序性能)[可设置为选做]	无	无 -	无 -
总计	17 (14)	180 (130)	2000 (1100)

### The End

