

AMS526: Numerical Analysis I (Numerical Linear Algebra for Computational and Data Sciences)

Lecture 13: Other Methods for Least Squares Problems;
Rank-Deficient Least Squares Problems;
Linear Algebra Software

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Outline

- 1 Solution of Least Squares Problems
- 2 Rank-Deficient Least Squares Problems
- 3 Software for Linear Algebra

Algorithms for Solving Least Squares Problems

- There are many variants of algorithms for solving least squares problems
 - ▶ Householder QR (with/without pivoting, explicit or implicit Q): **Backward stable**
 - ▶ Classical Gram-Schmidt: **Unstable**
 - ▶ Modified Gram-Schmidt with explicit Q : **Unstable**
 - ▶ Modified Gram-Schmidt with augmented system of equations with implicit Q : **Backward stable**
 - ▶ Normal equations (solve $A^T A x = A^T b$): **Unstable**
 - ▶ Singular value decomposition: **Stable and most accurate**

Stability of Gram-Schmidt Orthogonalization

- Gram-Schmidt QR is unstable, due to loss of orthogonality
- Gram-Schmidt can be stabilized using augmented system of equations
 - 1 Compute QR factorization of augmented matrix: $[Q,R1]=mgs([A,b])$
 - 2 Extract R and $\hat{Q}^T b$ from $R1$: $R=R1(1:n,1:n)$; $Qb=R1(1:n,n+1)$
 - 3 Back solve: $x=R \backslash Qb$

Theorem

The solution of the full-rank least squares problem by Gram-Schmidt orthogonality is backward stable in the sense that the computed solution \tilde{x} has the property

$$\|(A + \delta A)\tilde{x} - b\| = \min, \quad \frac{\|\delta A\|}{\|A\|} = O(\epsilon_{machine})$$

for some $\delta A \in \mathbb{R}^{m \times n}$, provided that $\hat{Q}^T b$ is formed implicitly.

Other Methods

- The method of *normal equation* solves $x = (A^T A)^{-1} A^T b$, due to squaring of condition number of A

Theorem

The solution of the full-rank least squares problem via normal equation is unstable. Stability can be achieved, however, by restriction to a class of problems in which $\kappa(A)$ is uniformly bounded above.

- Another method is to SVD

Solution by SVD

- Using $A = \hat{U}\hat{\Sigma}V^T$, b can be projected onto $\text{range}(A)$ by $P = \hat{U}\hat{U}^T$, and therefore $\hat{U}\hat{\Sigma}V^Tx = \hat{U}\hat{U}^Tb$
- Left-multiply by \hat{U}^T and we get $\hat{\Sigma}V^Tx = \hat{U}^Tb$

Least squares via SVD

Compute reduced SVD factorization $A = \hat{U}\hat{\Sigma}V^T$

Compute vector $c = \hat{U}^Tb$

Solve diagonal system $\hat{\Sigma}w = c$ for w

Set $x = Vw$

- Work is dominated by SVD, which is $\sim 2mn^2 + 11n^3$ flops, very expensive if $m \approx n$
- Question: If A is rank deficient, how to solve $Ax \approx b$?

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Rank-Deficient Least Squares Problems

- Least squares problems $Ax \approx b$ is the most challenging if A is (nearly) rank deficient
- If A is rank deficient, there are an infinite number of x that minimizes $\|b - Ax\|$. This is because if $y \in \text{null}(A)$, for any x that minimizes $\|b - Ax\|$, $x + y$ also minimizes $\|b - Ax\|$
- “Uniqueness” is recovered by requiring $x \perp \text{null}(A)$. Or equivalently, minimize $\|x\|$ subject to $(b - Ax) \perp \text{range}(A)$
- In practice, however, we often have near rank deficiency instead of exact rank deficiency
- For rank deficiency, (left or right) null space is the space span by (left or right) singular vectors corresponding to zero singular values
- For nearly rank deficient least squares problem, define “numerical null space” to be singular vectors corresponding to smallest singular values

Solving Rank-Deficient Least Squares Problems by SVD

- If A is full rank, $A = \hat{U} \hat{\Sigma} V^T = \sum_{j=1}^{\min\{m,n\}} \sigma_j u_j v_j^T$, and $A^+ = \sum_{j=1}^{\min\{m,n\}} \frac{1}{\sigma_j} v_j u_j^T$
- If A is rank deficient, $A^+ = \sum_{j=1}^r \frac{1}{\sigma_j} v_j u_j^T$, where r is rank of A
- If A is nearly rank deficient, $\tilde{A}^+ = \sum_{j=1}^r \frac{1}{\sigma_j} v_j u_j^T$, where r is *numerical rank* of A , i.e., largest j such that $\sigma_j \geq \epsilon \sigma_1$ for some small ϵ . This is called *truncated SVD*
- $\tilde{A} = \sum_{j=1}^r \sigma_j u_j v_j^T$ is a low-rank approximation to A

Rank-deficient least squares via truncated SVD

Compute reduced SVD factorization $A = \hat{U} \hat{\Sigma} V^T$ and estimate r

Compute vector $c = \left(\hat{U}_{:,1:r} \right)^T b$

Solve diagonal system $\hat{\Sigma}_{1:r,1:r} w = c$ for w

Set $x = V_{1:m,1:r} w$

A Note on Pseudoinverse

- If $A \in \mathbb{R}^{m \times n}$ is rank deficient, the pseudoinverse of A is defined as

$$A^+ = \sum_{j=1}^r \frac{1}{\sigma_j} v_j u_j^T,$$

where r is rank of A

- It is unique minimum Frobenius norm solution to

$$\min_{X \in \mathbb{R}^{n \times m}} \|AX - I_m\|_F$$

- It is also unique matrix $X \in \mathbb{R}^{n \times m}$ that satisfies four *Moore-Penrose conditions*:
 - 1 $AXA = A$
 - 2 $XAX = X$
 - 3 $(AX)^T = AX$
 - 4 $(XA)^T = XA$

QR with Column Pivoting

- Another approach is to use QR with column pivoting, or QRCP
- Suppose $A \in \mathbb{R}^{m \times n}$, and r be its rank. In **exact arithmetic**, QR with column pivoting is rank revealing if

$$Q^T A \Pi = \begin{bmatrix} R_{11} & R_{12} \\ 0 & 0 \end{bmatrix} \begin{matrix} r \\ m-r \\ r & n-r \end{matrix}$$

where Π is a permutation matrix. $\text{range}(A) = \text{span}\{q_1, \dots, q_r\}$

- In exact arithmetic, a rank-revealing QRCP is obtained by permuting columns such that diagonal entry in R is maximized at each step
- In particular, at k th step,

$$(Q_{k-1} \cdots Q_1) A (\Pi_1 \cdots \Pi_{k-1}) = R^{(k-1)} = \begin{bmatrix} R_{11}^{(k-1)} & R_{12}^{(k-1)} \\ 0 & R_{22}^{(k-1)} \end{bmatrix} \begin{matrix} k-1 \\ m-k+1 \\ k-1 & n-k+1 \end{matrix}$$

permute column with maximum 2-norm in $R_{22}^{(k-1)}$ to k th column

Solving Rank-Deficient Least Squares Problems by QRCP

- With rounding errors, one terminates if the computed $R_{22}^{(k-1)}$ ($\tilde{R}_{22}^{(k-1)}$) has a sufficient small 2-norm compared to that of A
 - ▶ If $\tilde{R}_{22}^{(k-1)}$ is small, then A is (numerically) rank deficient
 - ▶ However, if $\text{rank}(A) = k$, it does not follow that $\tilde{R}_{22}^{(k-1)}$ is small, so it may not reveal rank deficiency (and still lead to instability)
- In practice, QRCP needs to be coupled with a condition number estimator to help reveal the rank

Rank-deficient least squares via truncated QRCP

Compute QRCP $AP = QR$ and estimate r

Compute vector $c = (Q_{:,1:r})^T b$

Solve triangular system $R_{1:r,1:r}y = c$ for y

Set $x = P_{1:m,1:r}y$

- Truncated QRCP is far less expensive than truncated SVD, and is robust with a good condition number estimator
- Unlike SVD, QRCP uses a subset of columns of A

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Software for Linear Algebra

- LAPACK: Linear Algebra PACKage (www.netlib.org/lapack/lug)
 - ▶ Standard library for solving linear systems and eigenvalue problems
 - ▶ Successor of LINPACK (www.netlib.org/linpack) and EISPACK (www.netlib.org/eispack)
 - ▶ Depends on BLAS (Basic Linear Algebra Subprograms)
 - ▶ Parallel extensions include ScaLAPACK and PLAPACK (with MPI)
 - ▶ Note: Uses Fortran conventions for matrix arrangements
- MATLAB
 - ▶ Factorization A : $\text{lu}(A)$ and $\text{chol}(A)$
 - ▶ Solve $Ax = b$: $x = A \backslash b$
 - ★ Uses back/forward substitution for triangular matrices
 - ★ Uses Cholesky factorization for positive-definite matrices
 - ★ Uses LU factorization with column pivoting for nonsymmetric matrices
 - ★ Uses Householder QR for least squares problems
 - ★ Uses some special routines for matrices with special sparsity patterns
 - ▶ Uses LAPACK and other packages internally
- Solvers for sparse matrices (e.g., SuperLU, TAUCS, SuiteSparse)

Some Commonly Used Functions

Example BLAS routines: Matrix-vector multip.: dgemv; Matrix-matrix multip: dgemm

	LU Factorization		Solve linear system		Est. cond
	General	Symmetric	General	Symmetric	
LAPACK	dgetrf	dpotrf/dsytrf	dgesv	dposv/dposvx	dgecon
LINPACK	dgefa	dpofa/dsifa	dgesl	dposl/dsisl	dgeco
MATLAB	lu	chol	\	\	rcond

	Linear least squares			Eigenvalue/vector		SVD
	QR	Solve	Rank-deficient	General	Sym.	
LAPACK	dgeqrf	dgels	dgelsy/s/d	dgeev	dsyev	dgesvd
LINPACK	dqrdc	dqrsl	dqrst	-	-	dsvdc
MATLAB	qr	\	\	eig	eig	svd

For BLAS, LINPACK, and LAPACK, first letter s stands for single-precision real, d for double-precision real, c for single-precision complex, and z for double-precision complex. Boldface LAPACK routines are **driver** routines; others are **computational** routines.

Using LAPACK Routines in C Programs

- LAPACK was written in Fortran 77. Special attention is required when calling from C.
- Key differences between C and Fortran
 - ① Storage of matrices: column major (Fortran) versus row major (C/C++)
 - ② Argument passing for subroutines in C and Fortran: pass by reference (Fortran) and pass by value (C/C++)
- Example C code (example.c) for solving linear system using sgesv
 - ▶ See class website for sample code
 - ▶ To compile, issue command “cc -o example example.c -llapack -lblas”
- Hint: To find a function name, refer to LAPACK Users' Guide
- To find out arguments for a given function, search on netlib.org