

بِسْمِ اللّٰهِ الرَّحْمٰنِ الرَّحِيْمِ

مبانی رایانش امن

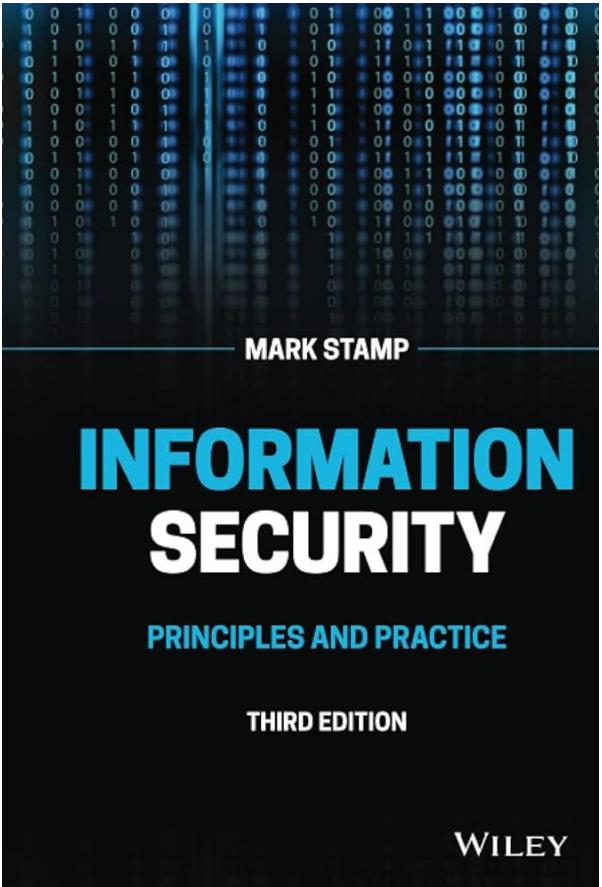
جلسه ۱۳

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IUT-ECE

فصل ؟ استمپ ◀



Blockchain

Preliminaries: Work

- How to measure (digital) *work* ?
- Our unit of work will be 1 hash
- Suppose that we have a hash function $h(x)$ that generates an N -bit output
- Then randomly chosen input generates one of 2^N equally likely outputs
 - For any input R , have, $0 \leq h(R) < 2^N$
 - Different R yield uncorrelated hashes

Work and Hashing

- For 16-bit hash, how many hashes until
 $h(R) = y = (0000000000y_5y_4y_3y_2y_1y_0)$?
- For random R, we have a 1/2 chance that
 $y = (0y_{14}y_{13}y_{12}y_{11}y_{10}y_9y_8y_7y_6y_5y_4y_3y_2y_1y_0)$
- And 1/4 chance that
 $y = (00y_{13}y_{12}y_{11}y_{10}y_9y_8y_7y_6y_5y_4y_3y_2y_1y_0)$
- And 1/8 chance that
 $y = (000y_{12}y_{11}y_{10}y_9y_8y_7y_6y_5y_4y_3y_2y_1y_0)$
- And so on...

Work and Hashing

- For 16-bit hash, if someone gives us an R such that $h(R) < 64$
 - Then expected number of hashes computed is 2^{10} ("expected" means average case)
 - That is, they have done 1,000 units of work
- We use hashing to show work was done
- Why this obsession with work?
 - That will become clear later...

Work and Hashing

- We can adjust parameter so more work (or less) is required
 - For N-bit hash, if we require $h(R) < 2^n$ then expected work is 2^{N-n} hashes
- **Note:** We can easily verify that the expected amount of work was done
 - Only requires one single hash
 - No matter how much work to find R

Distributed Ledger and Work

- Every ledger will have some amount of work associated with it
- Ledger with most work always “wins”
 - That is, everyone accepts ledger that has the most work put into it
- Recall, work is measured in hashes
- So, more hashes is “more better”

Blocks and Hashes

- Each transaction is signed
- Transactions grouped into *blocks*
 - Let B be one such block
- Find (nonce) R so that $h(B,R) < 2^n$
 - Equivalent to saying $h(B,R)$ starts with a specified number of 0s
- Work required to find R ?
 - On average 2^{N-n} hashes for N -bit hash

Chain

- Don't want to revalidate each block, want to order blocks, and so on
- We'll *chain* blocks together
 - Put hash of previous block in header of current block before computing hash
- So, must find R so that $h(Y, B, R) < 2^n$
 - Where Y is hash of previous block

Blockchain

- We now have

$$Y_{i+1} = h(Y_i, B_i, R_i) < 2^n$$

$$Y_{i+2} = h(Y_{i+1}, B_{i+1}, R_{i+1}) < 2^n$$

$$Y_{i+3} = h(Y_{i+2}, B_{i+2}, R_{i+2}) < 2^n$$

- Each B is a block
 - Block is a group of signed transactions
- Each R is chosen so inequality holds
 - Lot of work to find R, easy to verify $Y < 2^n$

Mining?

- Anyone can create a new block
- But lots of work to find a valid hash
- So what is the incentive to do work?
- “Free” money!
 - Get (new) money for doing work, say, § 1
 - Put this info at start of block, does not need to be signed (since new money)

One Block

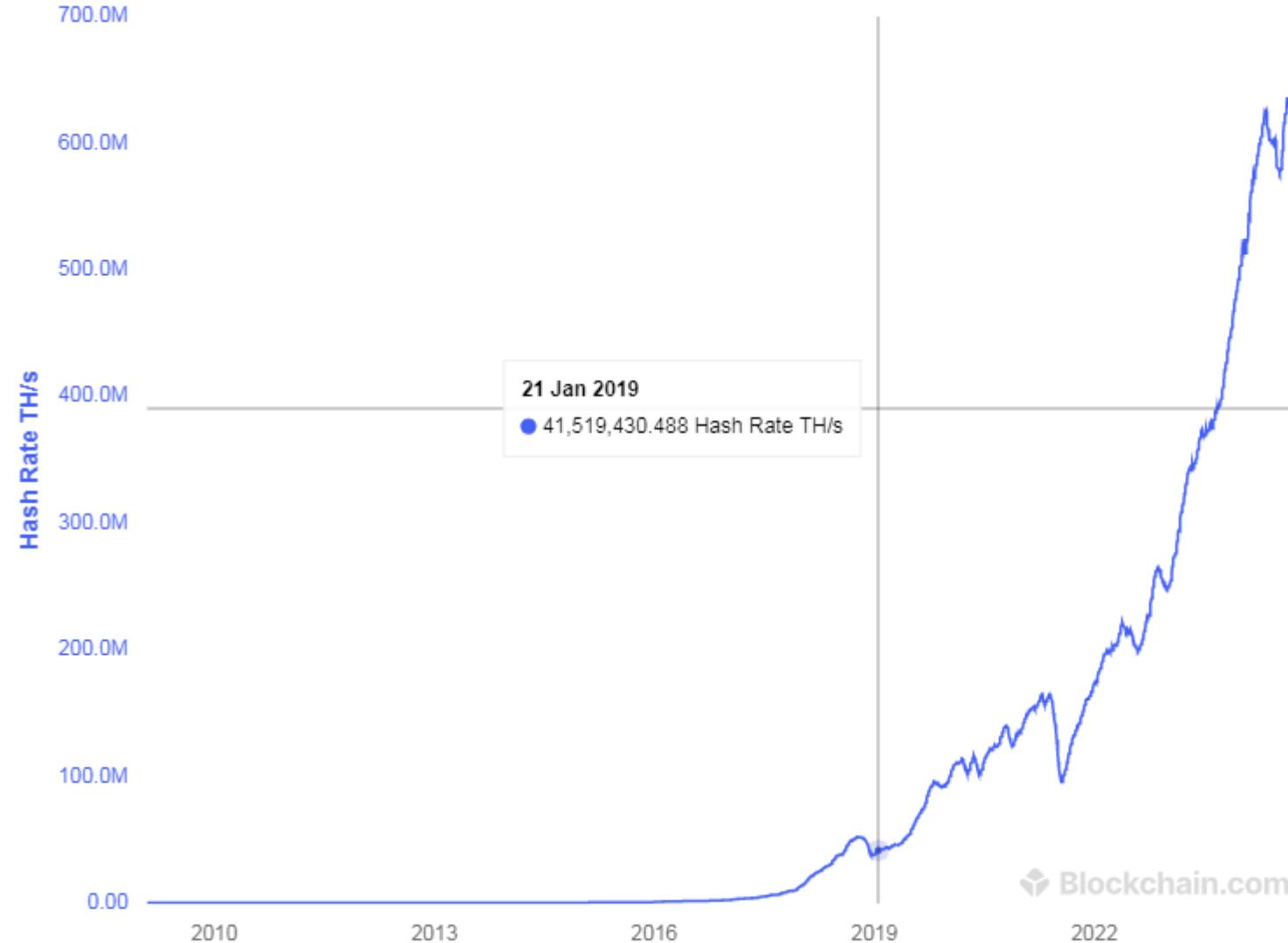
- Block B_i looks like...

$Y_i = h(B_{i-1})$
Miner _x gets §1
[1, Bob owes Alice §10] _{Bob}
[2, Charlie owes Trudy §30] _{Charlie}
[3, Trudy owes Alice §25] _{Trudy}
⋮
R_i

Block B_i

Mining

- ❑ Free money, so miners are in a race to find hashes that yield valid blocks
- ❑ The more computing power a miner has, the better chance to win race
- ❑ Once a valid hash is found, miner sends the block out to everybody
- ❑ Again, easy to verify hash is correct



Mining Hardware



GPU



FPGA



ASIC

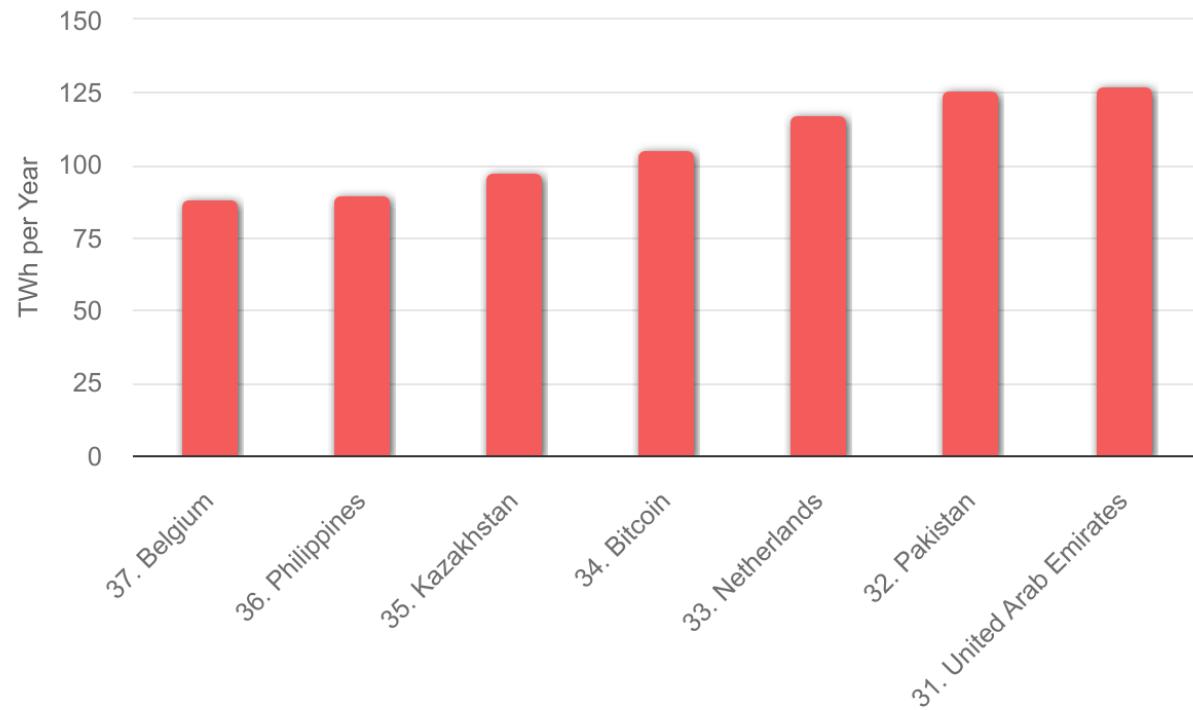
Mining Hardware



Mining Hardware

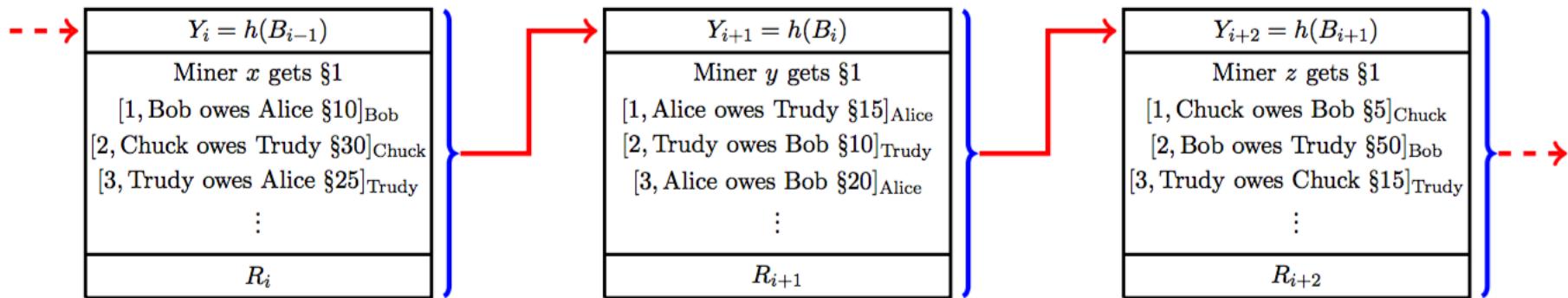


Energy Consumption by Country



Blockchain

- ❑ Blockchain looks like...



- ❑ Require that $h(Y_i, B_i, R_i) < 2^n$ and so on

Mining

- Why is “mining” called mining ?
 - Really, just finding a valid block hash
- Miner is doing work, and creating new money that did not previously exist
 - In a sense, this is comparable to mining gold or silver (for example)
- This may be the most misunderstood part of cryptocurrency protocols