Types	A, B, C, D	$::=$	α	Type variable
		$\mid$	$\top$	Top type
		$\mid$	$A \to B$	Function type
		$\mid$	$\forall \alpha. A$	Universal quantification
		$\mid$	$A \cap B$	Intersection type
		$\mid$	$(A * B) \Rightarrow C$	Disjoint constraint
Expressions	e	$::=$	x	Variable
		$\mid$	$\top$	Top
		$\mid$	$\lambda(x{:}A).\, e$	Lambda
		$\mid$	$e_1\ e_2$	Application
		$\mid$	$\Lambda \alpha. e$	Big lambda
		$\mid$	$e\ A$	Type application
		$\mid$	$e_1 ,, e_2$	Merge
Contexts	Γ	$::=$	ϵ	
		$\mid$	Γ, α	
		$\mid$	$\Gamma, x{:}A$	
		$\mid$	$\Gamma, A * B$	

Figure 1. Syntax.

$$\boxed{A <: B \hookrightarrow F}$$

$$\frac{}{\alpha <: \alpha \hookrightarrow \lambda(x{:}|\alpha|).\, x}\ \textsc{SubVar}$$

$$\frac{}{A <: \top \hookrightarrow \lambda(x{:}|A|).\, ()}\ \textsc{SubTop}$$

$$\frac{\tau_3 <: \tau_1 \hookrightarrow C_1 \qquad \tau_2 <: \tau_4 \hookrightarrow C_2}{\tau_1 \to \tau_2 <: \tau_3 \to \tau_4 \hookrightarrow \lambda(f{:}|\tau_1 \to \tau_2|).\, \lambda(x{:}|\tau_3|).\, C_2\ (f\ (C_1\ x))}\ \textsc{SubFun}$$

$$\frac{\tau_1 <: [\alpha_1/\alpha_2]\tau_2 \hookrightarrow C}{\forall \alpha_1. \tau_1 <: \forall \alpha_2. \tau_2 \hookrightarrow \lambda(f{:}|\forall \alpha_1. \tau_1|).\, \Lambda \alpha.\, C\ (f\ \alpha)}\ \textsc{SubForall}$$

$$\frac{\tau_1 <: \tau_2 \hookrightarrow C_1 \qquad \tau_1 <: \tau_3 \hookrightarrow C_2}{\tau_1 <: \tau_2 \cap \tau_3 \hookrightarrow \lambda(x{:}|\tau_1|).\, (C_1\ x, C_2\ x)}\ \textsc{SubAnd}$$

$$\frac{\tau_1 <: \tau_3 \hookrightarrow C}{\tau_1 \cap \tau_2 <: \tau_3 \hookrightarrow \lambda(x{:}|\tau_1 \cap \tau_2|).\, C\ (\mathsf{proj}_1 x)}\ \textsc{SubAnd}_1$$

$$\frac{\tau_2 <: \tau_3 \hookrightarrow C}{\tau_1 \cap \tau_2 <: \tau_3 \hookrightarrow \lambda(x{:}|\tau_1 \cap \tau_2|).\, C\ (\mathsf{proj}_2 x)}\ \textsc{SubAnd}_2$$

Figure 2. Subtyping.

Coherence for well-typed terms.

$$\epsilon \vdash 1 ,, 2 : (\mathsf{Int} * \mathsf{Int}) \Rightarrow \mathsf{Int} \cap \mathsf{Int}$$

$$\boxed{\Gamma \vdash e : A \hookrightarrow E}$$

$$\frac{x{:}A \in \Gamma}{\Gamma \vdash x : A \hookrightarrow x}\ \textsc{Ty/Var} \qquad \frac{}{\Gamma \vdash \top : \top \hookrightarrow ()}\ \textsc{Ty/Top}$$

$$\frac{\Gamma, x{:}A \vdash e : B \hookrightarrow E \qquad \Gamma \vdash A}{\Gamma \vdash \lambda(x{:}A).\, e : A \to B \hookrightarrow \lambda(x{:}|A|).\, E}\ \textsc{Ty/Lam}$$

$$\frac{\Gamma \vdash e_1 : A_1 \to A_2 \hookrightarrow E_1 \qquad \Gamma \vdash e_2 : A_3 \hookrightarrow E_2 \qquad A_3 <: A_1 \hookrightarrow C}{\Gamma \vdash e_1\ e_2 : A_2 \hookrightarrow E_1\ (C\ E_2)}\ \textsc{Ty/App}$$

$$\frac{\Gamma, \alpha \vdash e : A \hookrightarrow E}{\Gamma \vdash \Lambda \alpha. e : \forall \alpha. A \hookrightarrow \Lambda \alpha. E}\ \textsc{Ty/BLam}$$

$$\frac{\Gamma \vdash e : \forall \alpha. B \hookrightarrow E \qquad \Gamma \vdash A}{\Gamma \vdash e\ A : [A/\alpha]B \hookrightarrow E\ |A|}\ \textsc{Ty/TApp}$$

$$\frac{\Gamma \vdash e_1 : A \hookrightarrow E_1 \qquad \Gamma \vdash e_2 : B \hookrightarrow E_2}{\Gamma \vdash e_1 ,, e_2 : (A * B) \Rightarrow A \cap B \hookrightarrow (E_1, E_2)}\ \textsc{Ty/Merge}$$

$$\frac{\Gamma \vdash e : (A_1 * A_2) \Rightarrow B \hookrightarrow E \qquad \Gamma \vdash A_1 * A_2}{\Gamma \vdash e : B \hookrightarrow E}\ \textsc{Ty/DisjointCheck}$$

Figure 3. Typing.