

Name1
Affiliation1
Email1

Name2 Name3
Affiliation2/3
Email2/3

Abstract

This is the text of the abstract.

Categories and Subject Descriptors CR-number [*subcategory*]: third-level

General Terms Design, Languages, Theory

Keywords Intersection Types, Polymorphism, Type System
bruno: Make all figures fit and be well-formatted!

1. Introduction

bruno: The following are just some loose notes for the paper.
Modularity has always been a hot topic on OOP research.

Looking for new programming language models with increased expressiveness in terms of modularity.

Two previous lines of work:

- development of new languages with more expressive mechanisms. Examples: family polymorphism and virtual classes.
- Use of existing languages to express more modular designs with existing features. In particular Object Algebras.

The use of existing languages has shown that we can get quite far, but it has also exposed some limitations. For example, work on object algebras has shown that a delegation-based OO model would be better suited for OAs. Moreover, Object Algebra composition can benefit from intersection types. Additionally Object Algebras suffer from the "Inheritance is not Subtyping" problem. Although there are very few examples in the literature illustrating programs where the separating inheritance from subtyping is desirable, ob-

ject algebras illustrate a particularly compelling example of this.

Motivated by the insights gained in previous work, this paper presents a minimal core calculus that addresses current limitations and provides a better foundational model for statically typed delegation-based OOP? We show that Object Algebras fit nicely in this model.

Lots of work on developing new languages
finitary overloading: yes but have other merits of intersection been explored?

```
-- Compare Scala: -- merge[A,B] = new A with B
-- type IEval = eval : Int -- type IPrint = print : String
-- F[]
```

We present a polymorphic calculus containing intersection types and records, and show how this language can be used to solve various common tasks in functional programming in a nicer way.

Intersection types provides a power mechanism for functional programming, in particular for extensibility and allowing new forms of composition.

Prototype-based programming is one of the two major styles of object-oriented programming, the other being class-based programming which is featured in languages such as Java and C#. It has gained increasing popularity recently with the prominence of JavaScript in web applications. Prototype-based programming supports highly dynamic behaviors at run time that are not possible with traditional class-based programming. However, despite its flexibility, prototype-based programming is often criticized over concerns of correctness and safety. Furthermore, almost all prototype-based systems rely on the fact that the language is dynamically typed and interpreted.

In summary, the contributions of this paper are:

- elaboration typing rules which given a source expression with intersection types, typecheck and translate it into an ordinary F term. Prove a type preservation result: if a term e has type τ in the source language, then the translated term $\llbracket e \rrbracket$ is well-typed and has type $\llbracket \tau \rrbracket$ in the target language.
- present an algorithm for detecting incoherence which can be very important in practice.

- explores the connection between intersection types and object algebra by showing various examples of encoding object algebra with intersection types.

2. An Overview of **fi**

george: Change the examples later to something very simple.

This section provides the reader with the intuition of **fi**, while we postpone the presentation of the details in later sections.

In short, **fi** generalizes System F by adding intersection polymorphism. **fi** terms are elaborated into **f**, a variant of System F . System F , or polymorphic lambda calculus lays the foundation of functional programming languages such as Haskell.

The type system of **fi** permits a subtyping relation naturally and enables prototype-based inheritance. **bruno: Don't mention prototype-based inheritance in the summary, unless you are going to mention it later in the section. More generally don't mention things that you don't talk about later!** We will explore the usefulness of such a type system in practice by showing various examples.

2.1 Intersection Types

The central addition to the type system of **f** in **fi** is intersection types. **bruno: mention (and cite) that there are a number of OO languages supporting intersection types** What is an intersection type? One classic view is from set-theoretic interpretation of types: $A \wedge B$ stands for the intersection of the set of values of A and B . The other view, adopted in this paper, regards types as a kind of interface: a value of type $A \wedge B$ satisfies both of the interfaces of A and B . For example, $\text{eval} : \text{Int}$ is the interface that supports evaluation to integers, while $\text{eval} : \text{Int} \ \& \ \text{print} : \text{String}$ supports both evaluation and pretty printing. Those interfaces are akin to interfaces in Java or traits in Scala. But one key difference is that they are unnamed in **fi**.

Intersection Types and Overloading Intersection types provide a simple mechanism for ad-hoc polymorphism, similar to what type classes in Haskell achieve. The key constructs are the "merge" operator, denoted by \wedge , at the value level and the corresponding type intersection operator, denoted by \wedge at the type level. **bruno: We take a particular approach to intersection types, namely that of Dundfield. Some other approaches are different. The key feature in Dunfield's approach is the \wedge operator, which allows for run-time value/object composition.**

For example, we can define an (ad-hoc)-polymorphic `show` function that is able to convert integers and booleans to strings. In **fi** such function can be given the type

```
(Int -> String) & (Bool -> String)
```

and be defined using the merge operator \wedge , as

```
let show = showInt  $\wedge$  showBool
```

From another viewpoint, one can regard intersections in our system as "implicit pairs" whose introduction is explicit by using the merge operator and elimination is implicit (with no source-level construct for elimination).

bruno: This is not a proper sentence. Polish text. where `showString` and `showBool` are ordinary monomorphic functions. Later suppose the integer 1 is applied to the `show` function, the first component `showInt` will be picked because the type of `showInt` is compatible with 1 while `showBool` is not.

2.2 Labels and a Record Encoding

bruno: The title of this section needs adjusting. Moreover we do not want to talk about record elimination and update. fi supports a generalization of record elimination and update that works for any type. That's what we want to mention instead.

In addition to introduction of record literals using the usual notation, **fi** support two more operations on records: record elimination and record update.

A record type of the form $\{l : \tau\}$ can be thought as a normal type τ tagged by the label l . **bruno: I would call this a labelled type. Record types can be viewed as intersections of labelled types.**

bruno: Not a proper sentence! e_1 and e_2 are two expressions that support both evaluation and pretty printing and each has type $\text{eval} : \text{Int}$, $\text{print} : \text{String}$. `add` takes two expressions and computes their sum. **bruno: Don't inline examples in the text. Give proper separate code examples..** Note that in order to compute a sum, `add` only requires that the two expressions support evaluation and hence the type of the parameter $\text{eval} : \text{Int}$. As a result, the type of e_1 and e_2 are not exactly the same with that of the parameters of `add`. However, under a structural type system, this program should typecheck anyway because the arguments being passed has more information than required. In other words, $\text{eval} : \text{Int}$, $\text{print} : \text{String}$ is a subtype of $\text{eval} : \text{Int}$.

How is this subtyping relation derived? In **fi**, multi-field record types are excluded from the type system because $\text{eval} : \text{Int}$, $\text{print} : \text{String}$ can be encoded as $\text{eval} : \text{Int} \ \& \ \text{print} : \text{String}$. And by one of subtyping rules derives that $\text{eval} : \text{Int} \ \& \ \text{print} : \text{String}$ is a subtype of $\text{eval} : \text{Int}$.

bruno: Examples illustrating update and elimination of intersections are missing!

2.3 Intersection Types and Parametric Polymorphism

The presence of both parametric polymorphism and intersection is critical, as we shall see in the next section, in solving modularity problems. Here is a code snippet from the next section (The reader is not required to understand the purpose of this code at this stage; just recognizing the two types of polymorphism is enough.) **bruno: use the join example here.**

```
type SubExpAlg E = (ExpAlg E) \& { sub : E -> E -> E };
let e2 E (f : SubExpAlg E) = f.sub (exp1 E f) (f.lit 2);
```

SubExpAlg is a type synonym (a la Haskell) defined as the intersection of ExpAlg E and $\text{sub} : E \rightarrow E \rightarrow E$, parametrized by a type parameter E . $e2$ exhibits parametric polymorphism as it takes a type argument E .

3. Application

Structural subtyping facilitates reuse [18]. **bruno: orphan sentence!**

This section shows that the System F plus intersection types are enough for encoding extensible designs, and even beat the designs in languages with a much more sophisticated type system. In particular, **fi** has two main advantages over existing languages:

1. It supports dynamic composition of intersecting values.
2. It supports contravariant parameter types in the subtyping relation.

Various solutions have been proposed to deal with the extensibility problems and many rely on heavyweight language features such as abstract methods and classes in Java.

bruno: I would like to see a story about Church Encodings in fi. Can you look at Pierce's papers and try to write something along those lines? That will be a good intro for object algebras and visitors!

These two features can be used to improve existing designs of modular programs.

The expression problem refers to the difficulty of adding a new operations and a new data variant without changing or duplicating existing code.

There has been recently a lightweight solution to the expression problem that takes advantage of covariant return types in Java. We show that FI is able to solve the expression problem in the same spirit. The A)

```
trait Expr {
  def eval: Int
}

class Lit(n: Int) extends Expr {
  def eval: Int = n
}

class Add(n: Int) extends Expr {
  def eval: Int = e1.eval + e2.eval
}
```

3.1 Overloading

bruno: You already talk about overloading in the previous section. Need to decide where to put the text!

Dunfield [14] notes that using merges as a mechanism of overloading is not as powerful as type classes.

3.2 Encoding bounded polymorphism

bruno: IMPORTANT: Always capitalize the relevant words in a title!. This title should be "Encoding Bounded Polymorphism"

As **fi** extends System F with intersection types, F_{\leq} extends System F with bounded polymorphism. F_{\leq} [23] allows giving an upper bound to the type variable in type abstractions. The idea of bounded universal quantification was discussed in the seminal paper by Cardelli and Wegner [5]. They show that bounded quantifiers are useful because it is able to solve the "loss of information" problem.

In fact, the extension of System F in the other direction, i.e., with intersection types, is able to address the same problem effectively. Suppose we have the following definitions:

```
def user = { name = "George", admin = true }
def id(user: {name: String}) = user
```

Under a structural type system, programmers would expect that passing the user to the function is allowed. They are correct. Note that in the source language multi-field records are just syntactic sugars for merges of single-field records. Therefore, user is of a subtype of the parameter due to subtyping introduced by intersection types. So far so good. But there is a problem: what if programmers wants to access the admin field later, like:

```
(id user).admin
```

They cannot do so as the above will not typecheck. After going through the function, the user now has only the type

```
{name: String}
```

This is rather undesired because it indeed has an admin field!

Bounded polymorphism enable the function to return the exact type of the argument so that we have no problem in accessing the admin field later. Consider the example below:

```
def id [A <: {name: String}] (user: A) = user
(id [{name: String, admin: Bool}] user).admin
```

We do not have bounded polymorphism in the source language. But we can encode that via intersection types:

```
def id [A] (user: A & {name: String}) = user in
(id [{admin: Bool}] user).admin
```

Polymorphism plus intersection types is as powerful as bounded polymorphism.

[22]**bruno: I don't think this has been shown. What we can say is: we can encode a form of bounded polymorphism with intersection types.**

3.3 Object algebras

Object algebras provide an alternative to *algebraic data types* (ADT).**bruno: We are targeting an OO crowd. Mentioning algebraic datatypes is not going to be very useful there.** For example, the following Haskell definition of the type of simple expressions

```
data Exp where
  Lit :: Int -> Exp
  Add :: Exp -> Exp -> Exp
```

can be expressed by the *interface* of an object algebra of simple expressions:

```
trait ExpAlg[E] {
  def lit(x: Int): E
  def add(e1: E, e2: E): E
}
```

Similar to ADT, data constructors in object algebras are represented by functions such as `lit` and `add` inside an interface `ExpAlg`. Different with ADT, the type of the expression itself is abstracted by a type parameter `E`.

which can be expressed similarly in **fi** as:

```
type ExpAlg E = {
  lit : Int -> E,
  add : E -> E -> E
}
```

Scala supports intersection types via the `with` keyword. The type `A with B` expresses the combined interface of `A` and `B`. The idea is similar to

```
interface AwithB extends A, B {}
```

in Java.¹

The value level counterpart are functions of the type `A =>`

```
B => A with B.2
```

Our type system is a simple extension of System *F*; yet surprisingly, it is able to solve the limitations of using object algebras in languages such as Java and Scala. We will illustrate this point with an step-by-step of solving the expression problem using a source language built on top of **fi**.

Oliveira noted that composition of object algebras can be cumbersome and intersection types provides a solution to that problem.

We first define an interface that supports the evaluation operation:

```
type IEval = { eval : Int };
type ExpAlg E = { lit : Int -> E, add : E -> E -> E };
let evalAlg = {
  lit = \x : Int. { eval = x },
  add = \x : IEval. \y : IEval. { eval = x.eval
    + y.eval }
};
```

The interface is just a type synonym `IEval`. In **fi**, record types are structural and hence any value that satisfies this interface is of type `IEval` or of a subtype of `IEval`.³

In the following, `ExpAlg` is an object algebra interface of expressions with literal and addition case. And `evalAlg` is an

object algebra for evaluation of those expressions, which has type `ExpAlg Int`

```
type SubExpAlg E = (ExpAlg E) & { sub : E -> E -> E };
let subEvalAlg = evalAlg ,, { sub = \ (x : IEval).
  \ (y : IEval). { eval = x.eval - y.eval } };
```

Next, we define an interface that supports pretty printing.

```
type IPrint = { print : String };
let printAlg = {
  lit = \x : Int. { print = x.toString() },
  add = \x : IPrint. \y : IPrint. { print = x.
    print.concat(" + ").concat(y.print) },
  sub = \x : IPrint. \y : IPrint. { print = x.
    print.concat(" - ").concat(y.print) }
};
```

Provided with the definitions above, we can then create values using the appropriate algebras. For example: defines two expressions.

The expressions are unusual in the sense that they are functions that take an extra argument `f`, the object algebras, and use the data constructors provided by the object algebra (factory) `f` such as `lit`, `add` and `sub` to create values. Moreover, The algebras themselves are abstracted over the allowed operations such as evaluation and pretty printing by requiring the expression functions to take an extra argument `E`.

```
let merge A B (f : ExpAlg A) (g : ExpAlg B) = {
  lit = \x : Int. f.lit x ,, g.lit x,
  add = \x : A & B. \y : A & B.
    f.add x y ,, g.add x y
};
```

If we would like to have an expression that supports both evaluation and pretty printing, we will need a mechanism to combine the evaluation and printing algebras. Intersection types allows such composition: the `merge` function, which takes two expression algebras to create a combined algebra. It does so by constructing a new expression algebra, a record whose each field is a function that delegates the input to the two algebras taken.

```
let newAlg = merge IEval IPrint subEvalAlg
  printAlg in
let o1 = e1 (IEval & IPrint) newAlg in
o1.print
```

bruno: Don't start a sentence with `o1`. `o1` is a single object created that supports both evaluation and printing, thus achieving full feature-oriented programming.

3.4 Visitors

Constructing instances seems clumsy!

The visitor pattern allows adding new operations to existing structures without modifying those structures. The type of expressions are defined as follows:

```
trait Exp[A] {
  def accept(f: ExpAlg[A]): A
```

¹ However, Java would require the `A` and `B` to be concrete types, whereas in Scala, there is no such restriction.

² FIXME

³ Should be mentioned in S2.

```
}

trait SubExp[A] extends Exp[A] {
  override def accept(f: SubExpAlg[A]): A
}
```

The body of `Exp` and `SubExp` are almost the same: they both contain an `accept` method that takes an algebra `f` and returns a value of the carrier type `A`. The only difference is at `f` --- `SubExpAlg[A]` is a subtype of `ExpAlg[A]`. Since `f` appear in parameter position of `accept` and function parameters are contravariant, naturally we would hope that `SubExp[A]` is a supertype of `Exp[A]`. However, such subtyping relation does not fit well in Scala because inheritance implies subtyping in such languages⁴. As `SubExp[A]` extends `Exp[A]`, the former becomes a subtype of the latter.

Such limitation does not exist in **fi**. For example, we can define the similar interfaces `Exp` and `SubExp`:

```
type Exp A = { accept: forall A. ExpAlg A -> A };
type SubExp A = { accept: forall A. SubExpAlg A ->
  A };
```

Then by the typing judgment it holds that `SubExp` is a supertype of `Exp`. This relation gives desired results. To give a concrete example:

`A` is called is the *interpretation*. It works for any interpretation you want.

First we define two data constructors for simple expressions:

```
let lit (n : Int): Exp A = {
  accept = /\A. \ (f : ExpAlg A). f.lit n
};

let add (e1 : Exp) (e2 : Exp): Exp A = {
  accept = /\A. \ (f : ExpAlg A).
    f.add (e1.accept A f) (e2.accept A f)
};
```

Suppose later we decide to augment the expressions with subtraction:

```
let sub (e1 : SubExp) (e2 : SubExp): SubExp A =
  { accept = /\A. \ (f : SubExpAlg A).
    f.sub (e1.accept A f) (e2.accept A f)
  };
```

One big benefit of using the visitor pattern is that programmers is able to write in the same way that would do in Haskell. For example, `e2 = sub (lit 2)(lit 3)` defines an expression.

Another important property that does not exist in Scala is that programmer is able to pass `lit 2`, which is of type `Exp A`, to `sub`, which expects a `SubExp A` because of the subtyping relation we have. After all, it is known statically that `lit 2` can be passed into `sub` and nothing will go wrong. **bruno: Subtyping needs to be much more emphasized! See Modular Visitor Components!**

⁴ It is still possible to encode contravariant parameter types in Scala but doing so would require some technique. **bruno: what technique?**

3.5 Yanlin stuff

bruno: This can be dropped.

This subsection presents yet another lightweight solution to the Expression Problem, inspired by the recent work by Wang. It has been shown that contravariant return types allows refinement of the types of extended expressions.

First, we define the type of expressions that support evaluation and implement two constructors:

```
type Exp = { eval: Int }
let lit (n: Int) = { eval = n }
let add (e1: Exp) (e2: Exp)
  = { eval = e1.eval + e2.eval }
```

If we would like to add a new operation, say pretty printing, it is nothing more than refining the original `Exp` interface by *intersecting* the original type with the new `print` interface using the `&` primitive and *merging* the original data constructors using the `,,` primitive.

```
type ExpExt = Exp & { print: String }
let litExt (n: Int) = lit n ,, { print = n.
  toString() }
let addExt (e1: ExpExt) (e2: ExpExt)
  = add e1 e2 ,,
    { print = e1.print.concat(" + ").concat(e2.
      print) }
```

Now we can construct expressions using the constructors defined above:

```
let e1: ExpExt = addExt (litExt 2) (litExt 3)
let e2: Exp = add (lit 2) (lit 4)
```

`e1` is an expression capable of both evaluation and printing, while `e2` supports evaluation only.

We can also add a new variant to our expression:

```
let sub (e1: Exp) (e2: Exp) = { eval = e1.eval -
  e2.eval }
let subExt (e1: ExpExt) (e2: ExpExt)
  = sub e1 e2 ,, { print = e1.print.concat(" - ").
    concat(e2.print) }
```

Finally we are able to manipulate our expressions with the power of both subtraction and pretty printing.

```
(subExt e1 e1).print
```

3.6 Mixins

Mixins are useful programming technique wildly adopted in dynamic programming languages such as JavaScript and Ruby. But obviously it is the programmers' responsibility to make sure that the mixin does not try to access methods or fields that are not present in the base class.

In Haskell, one is also able to write programs in mixin style using records. However, this approach has a serious drawback: since there is no subtyping in Haskell, it is not possible to refine the mixin by adding more fields to the records. This means that the type of the family of the mixins has to be determined upfront, which undermines extensibility.

fi is able to overcome both of the problems: it allows composing mixins that (1) extends the base behavior, (2) while ensuring type safety.

The figure defines a mini mixin library. The apostrophe in front of types denotes call-by-name arguments similar to the \Rightarrow notation in the Scala language.

```
type Mixin S = 'S -> 'S -> S;
let zero S (super : 'S) (this : 'S) : S = super;
let rec mixin S (f : Mixin S) : S
  = let m = mixin S in f (\ (_ : Unit). m f) (\ (_
    : Unit). m f);
let extends S (f : Mixin S) (g : Mixin S) : Mixin
  S
  = \ (super : 'S). \ (this : 'S). f (\ (d : Unit).
    g super this) this;
```

We define a factorial function in mixin style and make a noisy mixin that prints "Hello" and delegates to its superclass. Then the two functions are composed using the `mixin` and `extends` combinators. The result is the `noisyFact` function that prints "Hello" every time it is called and computes factorial.

```
let fact (super : 'Int -> Int) (this : 'Int -> Int)
  : Int -> Int
  = \ (n : Int). if n == 0 then 1 else n * this (n
    - 1)
let noisy (super : 'Int -> Int) (this : 'Int ->
  Int) : Int -> Int
  = \ (n : Int). { println("Hello"); super n }
let noisyFact = mixin (Int -> Int) (extends (Int
  -> Int) foolish fact)
noisy 5
```

4. The fi calculus

Why no union types?

Why no empty records? Limitations of encoding records as intersections Local type inference⁵

This section formalizes the syntax, subtyping, and typing of **fi**. In the next section, we will go through the type-directed translation from **fi** to System F.
bruno: Mention Dundfield

4.1 Syntax

Figure 1 specifies the syntax of **fi**. It extends the syntax of System *F* by adding the two features: intersection types and records.
bruno: labelled types (single field records are fine too). The formalization includes only single records because single record types as the multi-records can be desugared into the merge of multiple single records.

Types τ have five constructs. The first three are standard (present in System *F*): type variable α , function types $\tau \rightarrow \tau$, and type abstraction $\forall \alpha. \tau$; while the last two, intersection types $\tau \wedge \tau$
bruno: be consistent: in the source language we use &, and in the formalization we use \wedge and record types

Types

$\tau ::=$	$\alpha \mid \tau \rightarrow \tau \mid \forall \alpha. \tau$	
	$\mid \tau \wedge \tau$	Intersection type
	$\mid \{l : \tau\}$	Record type

Terms

$e ::=$	x	
	$\Lambda \alpha. e \mid e \tau$	Type abstraction/application
	$\lambda(x:\tau). e \mid e e$	Term abstraction/application
	e, e	Merge
	$\{l = e\}$	Record construction
	$e.l$	Record projection
	$e \text{ with } \{l = e\}$	Record update

Contexts

$\Gamma ::= \epsilon \mid \Gamma, \alpha \mid \Gamma, x : \tau$

Labels

l

Figure 1: Syntax of **fi**

$\{l : \tau\}$, are novel in **fi**. In record types, l is the label and τ the type.

The first five constructs of expressions are also standard: variables x and two abstraction-elimination pairs: $\lambda(x:\tau). e$ abstracts expression e over values of type τ and is eliminated by application $e e$; $\Lambda \alpha. e$ abstracts expression e over types and is eliminated by type application $e \tau$.

The last four constructs are novel. e, e is the *merge* of two terms.
bruno: explain further, after all this is new.. $\{l = e\}$ introduces a record literal having l as the label for field containing expression e . $e.l$ access the field with label l in e . Finally, $e \text{ with } \{l = e_1\}$ is a *new* record which is exactly the same as e except the field labelled l is updated to become e_1 . For simplicity, we omit other constructs in order to focus on the essence of the calculus. For example, fixpoints can be added in standard ways.

The field F is non-standard and introduced to deal with records. It is an associative list. Each item is a pair whose first item is either empty or a label and the second the types.
bruno: What field F ? I don't know what this is referring to!

4.2 Subtyping

Subtyping is transitive and forms a lattice? Do we have a subtyping lattice?
bruno: Intersection types require a simple form of subtyping. The subtyping relation is reflexive and transitive. Our subtyping relation follows previous work. Need to cite! The main difference is the interaction with parametric polymorphism

bruno: Fix the overflow. For example by using 2 lines in the conclusion.
bruno: Subtyping figure is doing the translation already?

⁵ Joshua Dunfield

$\tau < \tau \hookrightarrow C$	
$\alpha < \alpha \hookrightarrow \lambda(x : \llbracket \alpha \rrbracket). x$	(SVar)
$\tau_3 < \tau_1 \hookrightarrow C_1 \quad \tau_2 < \tau_4 \hookrightarrow C_2$	
$\tau_1 \rightarrow \tau_2 < \tau_3 \rightarrow \tau_4 \hookrightarrow \lambda(f : \llbracket \tau_1 \rightarrow \tau_2 \rrbracket). \lambda(x : \llbracket \tau_3 \rrbracket). C_2 (f (C_1 x))$	(SFun)
$\tau_1 < [\alpha_2 \mapsto \alpha_1] \tau_2 \hookrightarrow C$	
$\forall \alpha_1. \tau_1 < \forall \alpha_2. \tau_2 \hookrightarrow \lambda(f : \llbracket \forall \alpha. \tau_1 \rrbracket). \Lambda \alpha. C (f \alpha)$	(SForall)
$\tau_1 < \tau_3 \hookrightarrow C$	
$\tau_1 \wedge \tau_2 < \tau_3 \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C (\text{fst } x)$	(SAnd1)
$\tau_2 < \tau_3 \hookrightarrow C$	
$\tau_1 \wedge \tau_2 < \tau_3 \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C (\text{snd } x)$	(SAnd2)
$\tau_1 < \tau_2 \hookrightarrow C_1 \quad \tau_1 < \tau_3 \hookrightarrow C_2$	
$\tau_1 < \tau_2 \wedge \tau_3 \hookrightarrow \lambda(x : \llbracket \tau_1 \rrbracket). \langle C_1 x, C_2 x \rangle$	(SAnd3)
$\tau_1 < \tau_2 \hookrightarrow C$	
$\{l : \tau_1\} < \{l : \tau_2\} \hookrightarrow \lambda(x : \llbracket \{l : \tau_1\} \rrbracket). C x$	(SRcd)

Figure 2: Subtyping

Intersection types introduce natural subtyping relations among types. For example, $\text{Int} \wedge \text{Bool}$ should be a subtype of Int , since the former can be viewed as either Int or Bool . The subtyping rules are standard except for three points listed below:

1. $\tau_1 \wedge \tau_2$ is a subtype of τ_3 , if *either* τ_1 or τ_2 are subtypes of τ_3 ,
2. τ_1 is a subtype of $\tau_2 \wedge \tau_3$, if τ_1 is a subtype of both τ_2 and τ_3 .
3. $\{l_1 : \tau_1\}$ is a subtype of $\{l_2 : \tau_2\}$, if l_1 and l_2 are identical and τ_1 is a subtype of τ_2 .

The first point is captured by two rules (S-And-1) and (S-And-2), whereas the second point by (S-And-3). Note that the last point means that record types are covariant in the type of the fields.

4.3 Typing

The typing judgment for **fi** is of the form: $\Gamma \vdash e : \tau$. This judgment uses the context Γ . The typing rules for our core languages are mostly standard ones for System F . **bruno:** Sort the rules by those, which are exactly the same as System F , following by those which are different (application) and finally the rules for the new constructs. In particular we introduce a (T-Merge) rule that applies to *merge* constructs. **bruno:**

$\Gamma \vdash t$	
$\frac{\alpha \in \Gamma}{\Gamma \vdash \alpha}$	(WF-Var)
$\frac{\Gamma \vdash \tau_1 \quad \Gamma \vdash \tau_2}{\Gamma \vdash \tau_1 \rightarrow \tau_2}$	(WF-Fun)
$\frac{\Gamma, \alpha \vdash \tau}{\Gamma \vdash \forall \alpha. \tau}$	(WF-Forall)
$\frac{\Gamma \vdash \tau_1 \quad \Gamma \vdash \tau_2}{\Gamma \vdash \tau_1 \wedge \tau_2}$	(WF-And)
$\frac{\Gamma \vdash \tau}{\Gamma \vdash \{l : \tau\}}$	(WF-Red)

Figure 3: Well-formedness

A lot more explanation needed here! You want to explain: 1) the rules which are different, and why they are different; 2) the new rules and the intuition for the new rules.

dom reads: "the domain of". $F(l)$ means the result of lookup for l inside the associative list F . The order of lookup can be either from left to right or from right to left but has to be consistent inside one implementation. We prefer the order from the right to the left because it make possible record overriding. For example, $(\{count = 1\}, \{count = 2\}).count$ will evaluate to 2 in this case.

5. Type-directed translation to System F

In this section we define the semantics of **fi** by means of a type-directed translation to System F . This translation removes the labels of records and turns intersections into products, much like Dunfield's elaboration. But our translation also deals with parametric polymorphism and records.

5.1 Informal discussion

To help the reader have a high-level understanding of how the translation works, in this subsection we present the translation informally. Take the **fi** expression for example:

```
{ eval = 4, print = ``4'' }.eval
```

First, multi-field record literals are desugared into merges of single-field record literals. Therefore $\{eval = 4, print = ``4''\}$ becomes $\{eval = 4\}, \{print = ``4''\}$. Merges of two values are translated into just a pair of them by (Merge) and single-field record literals lose their field labels by (RedIntro). Hence $\{eval = 4\}, \{print = ``4''\}$ becomes $(4, ``4'')$. **bruno:** Don't abuse inlining of examples in the text!

Finally, e_1 and e_2 are both coerced by a projection function

$$\boxed{\Gamma \vdash e : \tau \hookrightarrow E}$$

$$\frac{(x, \tau) \in \Gamma}{\Gamma \vdash x : \tau \hookrightarrow x} \quad (\text{Var})$$

$$\frac{\Gamma, x : \tau \vdash e : \tau_1 \hookrightarrow E \quad \Gamma \vdash \tau}{\Gamma \vdash \lambda(x : \tau). e : \tau \rightarrow \tau_1 \hookrightarrow \lambda(x : \llbracket \tau \rrbracket). E} \quad (\text{Abs})$$

$$\frac{\Gamma, \alpha \vdash e : \tau \hookrightarrow E}{\Gamma \vdash \Lambda \alpha. e : \forall \alpha. \tau \hookrightarrow \Lambda \alpha. E} \quad (\text{TAbs})$$

$$\frac{\begin{array}{c} \Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \hookrightarrow E_1 \\ \Gamma \vdash e_2 : \tau_3 \hookrightarrow E_2 \\ \tau_3 \leq \tau_1 \hookrightarrow C \end{array}}{\Gamma \vdash e_1 e_2 : \tau_2 \hookrightarrow E_1 (C E_2)} \quad (\text{App})$$

$$\frac{\Gamma \vdash e : \forall \alpha. \tau_1 \hookrightarrow E \quad \Gamma \vdash \tau}{\Gamma \vdash e \tau : [\alpha \mapsto \tau] \tau_1 \hookrightarrow E \llbracket \tau \rrbracket} \quad (\text{TApp})$$

$$\frac{\begin{array}{c} \Gamma \vdash e_1 : \tau_1 \hookrightarrow E_1 \\ \Gamma \vdash e_2 : \tau_2 \hookrightarrow E_2 \end{array}}{\Gamma \vdash e_1, e_2 : \tau_1 \wedge \tau_2 \hookrightarrow \langle E_1, E_2 \rangle} \quad (\text{Merge})$$

$$\frac{\Gamma \vdash e : \tau \hookrightarrow E}{\Gamma \vdash \{l = e\} : \{l : \tau\} \hookrightarrow E} \quad (\text{RcdCon})$$

$$\frac{\Gamma \vdash e : \tau \hookrightarrow E \quad \Gamma \vdash_{\text{get}} (\tau; l) : \tau_1 \hookrightarrow C}{\Gamma \vdash e.l : \tau_1 \hookrightarrow C E} \quad (\text{RcdProj})$$

$$\frac{\begin{array}{c} \Gamma \vdash e : \tau \hookrightarrow E \quad \Gamma \vdash e_1 : \tau_1 \hookrightarrow E_1 \\ \vdash_{\text{put}} (t; l; e_1 : \tau_1 \hookrightarrow E_1) : (\tau'_1, \tau') \hookrightarrow C \quad \tau_1 \leq \tau'_1 \end{array}}{\Gamma \vdash e \text{ with } \{l = e_1\} : \tau' \hookrightarrow C E} \quad (\text{RcdUpd})$$

$$\boxed{\vdash_{\text{get}} (\tau_1; l) : \tau_2 \hookrightarrow C}$$

The field with label l inside τ_1 is of type τ_2 .

$$\vdash_{\text{get}} (\{l : \tau\}; l) : \tau \hookrightarrow \lambda(x : \llbracket \{l : \tau\} \rrbracket). x \quad (\text{GetBase})$$

$$\frac{\vdash_{\text{get}} (\tau_1; l) : \tau \hookrightarrow C}{\vdash_{\text{get}} (\tau_1 \wedge \tau_2; l) : \tau \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{fst } x)} \quad (\text{GetLeft})$$

$$\frac{\vdash_{\text{get}} (\tau_2; l) : \tau \hookrightarrow C}{\vdash_{\text{get}} (\tau_1 \wedge \tau_2; l) : \tau \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{snd } x)} \quad (\text{GetRight})$$

$$\boxed{\vdash_{\text{put}} (\tau; l; e : \tau \hookrightarrow E) : (\tau, \tau) \hookrightarrow C}$$

$$\vdash_{\text{put}} (\{l : \tau\}; l; e : \tau' \hookrightarrow E) : (\tau, \{l : \tau'\}) \hookrightarrow \lambda(x : \llbracket \{l : \tau\} \rrbracket). E \quad (\text{PutBase})$$

$$\frac{\text{fi} \quad \vdash_{\text{put}} (\tau_1; l; e : \tau \hookrightarrow E) : (\tau', \tau'_1) \hookrightarrow C}{\vdash_{\text{put}} (\tau_1 \wedge \tau_2; l; e : \tau \hookrightarrow E) : (\tau', \tau'_1 \wedge \tau_2) \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{fst } x)} \quad (\text{PutLeft})$$

$(x : (Int, String)).\text{fst } x$. **bruno: Show the source program, and the program that it gets translated to. Then explain how that translation works.**

5.2 Target language

The target language is System F extended with pairs. The syntax and typing is completely standard. The syntax of System F is as follows:

$$\begin{array}{ll} \text{Types} & T ::= \alpha \mid T \rightarrow T \mid \forall \alpha. T \mid \langle T, T \rangle \\ \text{Terms} & E, C ::= x \mid \lambda(x : T). E \mid \Lambda \alpha. E \mid E E \mid E T \\ & \quad \mid \langle E, E \rangle \mid \text{fst } E \mid \text{snd } E \end{array}$$

The dynamic semantics of System F can be found in ...**bruno: fill!**

Lemma 1. *If*

$$\Gamma \vdash \tau_1 \leq \tau_2 \hookrightarrow C$$

then

$$\llbracket \Gamma \rrbracket \vdash C : \llbracket \tau_1 \rrbracket \rightarrow \llbracket \tau_2 \rrbracket$$

bruno: Why is this lemma placed here?

bruno: Generally Speaking this text seems out of place. Move to 5.4, maybe? The main translation judgment is $\Gamma \vdash e : \tau \hookrightarrow E$ which states that with respect to the environment Γ , the **fi** expression e is of a **fi** type τ and its translation is a System F expression E .

We also define the type translation function $\llbracket \cdot \rrbracket$ from **fi** types τ to System F types T .

The first three rules of the translation is standard. For the last two, the intersection of two types are translated into a product of them, and the label of record types are erased.

The translation consists of four sets of rules, which are explained below:

5.3 Subtyping (coercion)

Talk about η -expansion.

The coercion judgment $\Gamma \vdash \tau_1 \leq \tau_2 \hookrightarrow C$ extends the subtyping judgment with a coercion on the right hand side of \hookrightarrow . A coercion C is an expression in the target language and has type $\tau_1 \rightarrow \tau_2$, as proved by Lemma 1. It is read "In the environment Γ , τ_1 is a subtype of τ_2 ; and if any expression e has a type τ_1 that is a subtype of the type of τ_2 , the elaborated e , when applied to the corresponding coercion C , has exactly type $\llbracket \tau_2 \rrbracket$ ". For example, $\Gamma \vdash Int \& Bool \leq Bool \hookrightarrow \text{fst}$, where fst is the projection of a tuple on the first element. The coercion judgment is only used in the (App) case. As (SFun) supports contravariant parameter type and covariant return type, the coercion of the parameter types and that of the return types are used to create a coercion for the function type. (SAnd1), (SAnd2), and (SAnd3) deal with intersection types. The first two are complementary to each other. Take (SAnd1) for example, if we know τ_1 is a subtype of τ_3 and C is a coercion from τ_1 to τ_3 , then we can conclude that $\tau_1 \wedge \tau_2$ is also a subtype of τ_3 and the new coercion is a function that

$$\llbracket \tau \rrbracket = T$$

$$\begin{aligned} \llbracket \alpha \rrbracket &= \alpha \\ \llbracket \tau_1 \rightarrow \tau_2 \rrbracket &= \llbracket \tau_1 \rrbracket \rightarrow \llbracket \tau_2 \rrbracket \\ \llbracket \forall \alpha. \tau \rrbracket &= \forall \alpha. \llbracket \tau \rrbracket \\ \llbracket \tau_1 \wedge \tau_2 \rrbracket &= \langle \llbracket \tau_1 \rrbracket, \llbracket \tau_2 \rrbracket \rangle \\ \llbracket \{l : \tau\} \rrbracket &= \llbracket t \rrbracket \end{aligned}$$

Figure 5: Type translation

$$\tau < \tau \hookrightarrow C$$

$$\alpha < \alpha \hookrightarrow \lambda(x : \llbracket \alpha \rrbracket). x \quad (\text{SVar})$$

$$\frac{\tau_3 < \tau_1 \hookrightarrow C_1 \quad \tau_2 < \tau_4 \hookrightarrow C_2}{\tau_1 \rightarrow \tau_2 < \tau_3 \rightarrow \tau_4 \hookrightarrow \lambda(f : \llbracket \tau_1 \rightarrow \tau_2 \rrbracket). \lambda(x : \llbracket \tau_3 \rrbracket). C_2 (f (C_1 x))} \quad (\text{SFun})$$

$$\frac{\tau_1 < [\alpha_2 \mapsto \alpha_1] \tau_2 \hookrightarrow C}{\forall \alpha_1. \tau_1 < \forall \alpha_2. \tau_2 \hookrightarrow \lambda(f : \llbracket \forall \alpha. \tau_1 \rrbracket). \Lambda \alpha. C (f \alpha)} \quad (\text{SForall})$$

$$\frac{\tau_1 < \tau_3 \hookrightarrow C}{\tau_1 \wedge \tau_2 < \tau_3 \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C (\text{fst } x)} \quad (\text{SAnd1})$$

$$\frac{\tau_2 < \tau_3 \hookrightarrow C}{\tau_1 \wedge \tau_2 < \tau_3 \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C (\text{snd } x)} \quad (\text{SAnd2})$$

$$\frac{\tau_1 < \tau_2 \hookrightarrow C_1 \quad \tau_1 < \tau_3 \hookrightarrow C_2}{\tau_1 < \tau_2 \wedge \tau_3 \hookrightarrow \lambda(x : \llbracket \tau_1 \rrbracket). \langle C_1 x, C_2 x \rangle} \quad (\text{SAnd3})$$

$$\frac{\tau_1 < \tau_2 \hookrightarrow C}{\{l : \tau_1\} < \{l : \tau_2\} \hookrightarrow \lambda(x : \llbracket \{l : \tau_1\} \rrbracket). C x} \quad (\text{SRcd})$$

Figure 6: Coercion

takes a value x of type $\tau_1 \wedge \tau_2$, project x on the first item, and apply C to it. (SAnd3) uses both of two coercions and constructs a pair. **bruno: Give a couple of concrete examples when explaining the rules.**

5.4 Typing (translation)

In this subsection we now present formally the translation rules that convert **fi** expressions into System F ones. This set of rules essentially extends those in the previous section with the light-blue part for the translation.

bruno: Badly structured. Don't mention Coercion here, as it was already explained in the previous section. bruno: Don't use itemize and items. Use paragraphs instead!

- **Coercion**

Explained in the previous subsection.

- **Translation**

The elaboration judgment $\Gamma \vdash e : \tau \hookrightarrow E$ extends the typing judgment with an elaborated expression on the right hand side of \hookrightarrow . The translation ensures that E has type $\llbracket \tau \rrbracket$. It is also standard, except for the case of (App), in which a coercion from the inferred type of the argument, e_2 , to the expected type of the parameter, τ_1 , is inserted before the argument; (Merge) translates merges into pairs. (RcdIntro) uses the same System F expression E for e as for $\{l = e\}$. And in (RcdElim) and (RcdUpd) the coercions generated by the ``get" and ``put" rules will be used to coerce the main **fi** expression.

(RcdElim) typechecks e and use the ``get" rule to return the type of the field τ_1 and the coercion C . The type of the whole expression is τ_1 and its translation of $C E$.

(RcdUpd) is similar to (RcdElim) in that it uses the auxiliary ``put" rule. This rule typechecks e and e_1 , and uses

$\Gamma \vdash e : \tau \hookrightarrow E$

$$\frac{(x, \tau) \in \Gamma}{\Gamma \vdash x : \tau \hookrightarrow x} \quad (\text{Var})$$

$$\frac{\Gamma, x : \tau \vdash e : \tau_1 \hookrightarrow E \quad \Gamma \vdash \tau}{\Gamma \vdash \lambda(x : \tau). e : \tau \rightarrow \tau_1 \hookrightarrow \lambda(x : \llbracket \tau \rrbracket). E} \quad (\text{Abs})$$

$$\frac{\Gamma, \alpha \vdash e : \tau \hookrightarrow E}{\Gamma \vdash \Lambda \alpha. e : \forall \alpha. \tau \hookrightarrow \Lambda \alpha. E} \quad (\text{TAbs})$$

$$\frac{\begin{array}{c} \Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \hookrightarrow E_1 \\ \Gamma \vdash e_2 : \tau_3 \hookrightarrow E_2 \\ \tau_3 \leq \tau_1 \hookrightarrow C \end{array}}{\Gamma \vdash e_1 e_2 : \tau_2 \hookrightarrow E_1 (C E_2)} \quad (\text{App})$$

$$\frac{\Gamma \vdash e : \forall \alpha. \tau_1 \hookrightarrow E \quad \Gamma \vdash \tau}{\Gamma \vdash e \tau : [\alpha \mapsto \tau] \tau_1 \hookrightarrow E \llbracket \tau \rrbracket} \quad (\text{TApp})$$

$$\frac{\begin{array}{c} \Gamma \vdash e_1 : \tau_1 \hookrightarrow E_1 \\ \Gamma \vdash e_2 : \tau_2 \hookrightarrow E_2 \end{array}}{\Gamma \vdash e_1, e_2 : \tau_1 \wedge \tau_2 \hookrightarrow \langle E_1, E_2 \rangle} \quad (\text{Merge})$$

$$\frac{\Gamma \vdash e : \tau \hookrightarrow E}{\Gamma \vdash \{l = e\} : \{l : \tau\} \hookrightarrow E} \quad (\text{RcdCon})$$

$$\frac{\Gamma \vdash e : \tau \hookrightarrow E \quad \Gamma \vdash_{\text{get}} (\tau; l) : \tau_1 \hookrightarrow C}{\Gamma \vdash e.l : \tau_1 \hookrightarrow C E} \quad (\text{RcdProj})$$

$$\frac{\begin{array}{c} \Gamma \vdash e : \tau \hookrightarrow E \quad \Gamma \vdash e_1 : \tau_1 \hookrightarrow E_1 \\ \vdash_{\text{put}} (t; l; e_1 : \tau_1 \hookrightarrow E_1) : (\tau'_1, \tau') \hookrightarrow C \quad \tau_1 \leq \tau'_1 \end{array}}{\Gamma \vdash e \text{ with } \{l = e_1\} : \tau' \hookrightarrow C E} \quad (\text{RcdUpd})$$

$\vdash_{\text{get}} (\tau_1; l) : \tau_2 \hookrightarrow C$

The field with label l inside τ_1 is of type τ_2 .

$$\vdash_{\text{get}} (\{l : \tau\}; l) : \tau \hookrightarrow \lambda(x : \llbracket \{l : \tau\} \rrbracket). x \quad (\text{GetBase})$$

$$\frac{\vdash_{\text{get}} (\tau_1; l) : \tau \hookrightarrow C}{\vdash_{\text{get}} (\tau_1 \wedge \tau_2; l) : \tau \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{fst } x)} \quad (\text{GetLeft})$$

$$\frac{\vdash_{\text{get}} (\tau_2; l) : \tau \hookrightarrow C}{\vdash_{\text{get}} (\tau_1 \wedge \tau_2; l) : \tau \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{snd } x)} \quad (\text{GetRight})$$

$\vdash_{\text{put}} (\tau; l; e : \tau \hookrightarrow E) : (\tau, \tau) \hookrightarrow C$

$$\vdash_{\text{put}} (\{l : \tau\}; l; e : \tau' \hookrightarrow E) : (\tau, \{l : \tau'\}) \hookrightarrow \lambda(x : \llbracket \{l : \tau\} \rrbracket). E \quad (\text{PutBase})$$

$$\frac{\vdash_{\text{put}} (\tau_1; l; e : \tau \hookrightarrow E) : (\tau', \tau'_1) \hookrightarrow C}{\vdash_{\text{put}} (\tau_1 \wedge \tau_2; l; e : \tau \hookrightarrow E) : (\tau', \tau'_1 \wedge \tau_2) \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{fst } x)} \quad (\text{PutLeft})$$

$$\text{fi} \quad \frac{\vdash_{\text{put}} (\tau_2; l; e : \tau \hookrightarrow E) : (\tau', \tau'_2) \hookrightarrow C}{\vdash_{\text{put}} (\tau_1 \wedge \tau_2; l; e : \tau \hookrightarrow E) : (\tau', \tau'_1 \wedge \tau_2) \hookrightarrow \lambda(x : \llbracket \tau_1 \wedge \tau_2 \rrbracket). C(\text{snd } x)} \quad (\text{PutRight})$$

2015/3/8

the ``put" rule. Note that it allows refining of types by an e_1 that is of a subtype of τ'_1 , which is the type of the field l in e . The type of the updated expression then takes the type τ' returned by the ``put" rule, while its translation is E , applied to the coercion generated by the ``put" rule, C . The two set of rules are explained below.

• ``get" rules

The ``get" judgment deals specifically with record elimination and yields a coercion can be thought as a field accessor. For example: **bruno: Still not showing the derivations!**

$\Gamma \vdash_{\text{get}} (\{eval : Int\}, eval) : \{eval : Int\} \hookrightarrow \lambda(x : Int). x$

The lambda is the field accessor and when applied to a translated expression of type $\{eval : Int\}$, it is able to give the desired field. (GetBase) is the base case: the type of the field labelled l in a $\{l : \tau\}$ is just τ and the coercion is an identity function specialized to type $\llbracket \{l : \tau\} \rrbracket$ (GetLeft) and (GetRight) are complementary to each other.

Consider the source program:

```
{ name = ``Isaac'', age = 10 }.name
```

Multi-field records are desugared into merge of single-field records:

```
{ name = ``Isaac'' } ,, { age = 10 }.name
```

By (GetBase),

$\vdash_{\text{get}} (\{name : String\}; name) : String$

we have the coercion

$\lambda(x : \llbracket \{name : String\} \rrbracket). x$

which is just $\lambda(x : String). x$ according to type translation.

By (GetLeft),

$\vdash_{\text{get}} (\{name : String\} \wedge \{age : Int\}; name) : String$

By typing rules, the translation of the program is

$(\text{``Isaac''}, 10)$

. If we apply the coercion to it, we get

``Isaac''

• ``put" rules **bruno: Missing example (and derivation)**

The ``put" judgment deals specifically with record update can be thought as producing a field updater. Compared to the ``get" rules, the ``put" rules take an extra input e , which is the desired expression to replace the field labelled l in values of type τ . (PutBase) is the base case. This rule allows refinement of record fields in the sense

that the type of e can be a subtype of the type of the field labelled by l . The resulting type is $\{l : \tau'\}$ and the generated coercion is a constant function that always returns E . (PutLeft) and (PutRight) are complementary to each other: the idea is exactly the same as (GetLeft) and (GetRight) except that the refined type τ'_1 and τ'_2 is used.

5.5 Meta-theory

Lemma 2. *Subtyping is reflexive Given a type τ , $\tau <: \tau$.*

Lemma 3. *Subtyping is transitive If $\tau_1 <: \tau_2$ and $\tau_2 <: \tau_3$, then $\tau_1 <: \tau_3$.*

Lemma 4 (Get rules produce the type-correct coercion). *If*

$\Gamma \vdash_{\text{get}} \tau; l = C; \tau_1$

then

$\llbracket \Gamma \rrbracket \vdash C : \llbracket \tau \rrbracket \rightarrow \llbracket \tau_1 \rrbracket$

Proof. By induction on the given derivation. \square

Lemma 5 (Put rules produce the type-correct coercion). *If*

$\Gamma \vdash_{\text{put}} \tau; l; E = C; \tau_1$

then

$\llbracket \Gamma \rrbracket \vdash C : \llbracket \tau \rrbracket \rightarrow \llbracket \tau_1 \rrbracket$

Proof. By induction on the given derivation. \square

Lemma 6 (Translation preserves well-formedness). *If*

$\Gamma \vdash \tau$

then

$\llbracket \Gamma \rrbracket \vdash \llbracket \tau \rrbracket$

Proof. By induction on the given derivation. \square

Theorem 1 (Type preserving translation). *If*

$\Gamma \vdash e : \tau \hookrightarrow E$

then

$\llbracket \Gamma \rrbracket \vdash E : \llbracket \tau \rrbracket$

Proof. (Sketch) By structural induction on the expression and the corresponding inference rule. The full proof can be found in the appendix. \square

Type-Directed Translation to System F . Main results: type-preservation + coherence.

6. Implementation

Our implementation currently supports multi-field records.

6.1 Type Synonyms

We extend the implementation of the type system extended with type synonyms and lazy arguments.

```
type T A1 A2 = ... in
```

6.2 Optimization

7. Related work

http://homepages.inf.ed.ac.uk/gdp/publications/Sub_Par.pdf
[24]

Also discussed intersection types! [18].

Intersection types with polymorphism. Intersection types date back to as early as Coppo et al. [9]. Recently, some form of intersection types have been adopted in object-oriented languages such as Scala, Ceylon, and Grace. One defining difference, among others, is that all those languages only allow intersections of concrete types (classes), whereas our language allows intersections of type variables, such as $A \ \& \ B$. Without that vehicle, we would not be able to define the generic merge function (below) for all interpretations of a given algebra, and would incur boilerplate code:

```
let merge [A, B] (f: ExpAlg A) (g: ExpAlg B) = {  
  lit (x : Int) = f.lit x ,, g.lit x,  
  add (x : A & B) (y : A & B) =  
    f.add x y ,, g.add x y  
}
```

In Scala community, there have been attempts to provide a foundational calculus for Scala that incorporates intersection types [1, 2].

Our type system combines intersection types and polymorphism. The closest to ours is Pierce's work [20] on a prototype compiler for a language with both intersection types, union types, and parametric polymorphism. The important difference with our system is that in his language there is no explicit introduction construct like our merge operator. As shown in Section 3, this feature is pivotal in supporting modularity and extensibility because it allows dynamic composition of values. Pierce has also studied a system where both intersection types and bounded polymorphism are present in his Ph.D dissertation [21] and a 1997 report [22]. Going in the direction of higher kinds, Compagnoni and Pierce [8] add intersection types to System F^ω and use the new calculus, F^ω_\wedge , to model multiple inheritance. In their system, types include the construct of intersection of types of the same kind K . Compared to our work, they do not have a term-level construct for intersection introduction. Davies and Pfenning [12] study the interactions between intersection types and effects in call-by-value languages. And they propose a "value restriction" for intersection types, similar to value restriction on parametric polymorphism.

Other type systems with intersection types. Dunfield [14] describes a similar approach to ours: compiling a system with

intersection types into ordinary λ -calculus terms. The major difference is that his system does not include parametric polymorphism, while ours does not include unions. Besides, our rules are algorithmic.

Reynolds invented Forsythe [25] in the 1980s. Our merge operator is analogous to his p_1, p_2 . Castagna, and Dunfield describe elaborating multi-fields records into merge of single-field records. As Dunfield has noted, in Forsythe merges can be only used unambiguously.⁶ For instance, it is not allowed in Forsythe to merge two functions.

Refinement intersection [11, 13, 15] is the more conservative approach of adopting intersection types. It increases only the expressiveness of types but not terms. But without a term-level construct like "merge", it is not possible to encode various language features. As an alternative to syntactic subtyping described in this paper, Frisch et al. [16] study semantic subtyping.

Type systems for modularity. bruno: fill me in!

Goto <http://i.cs.hku.hk/bruno/papers/ecoop2012.pdf> Language-based

* <http://www.informatik.uni-marburg.de/kos/papers/pop106-virtualclasses.pdf> * Poly variants

Intersection types have been shown to be useful in designing languages that support modularity. [19]

Extensible records. <http://elm-lang.org/learn/Records.elm>

Understanding records is important for understanding object-oriented languages. And we are the first to elaborate records to System F . Encoding records using intersection types appear in Reynolds [25] and Castagna et al. [6]. Although Dunfield also discusses this idea in his paper [14], he only provides an implementation but not formalization. Very similar to our treatment of elaborating records is Cardelli's work [3] on translating a calculus, named $F_{<,p}$, with extensible records to a simpler calculus that without records primitives (in which case is $F_{<}$). But he does not consider encoding multi-field records as intersections; hence his translation is more heavyweight. Crary [10] uses intersection types and existential types to address the problem that arises when interpreting method dispatch as self-application. But in his paper, intersection types are not used to encode multi-field records.

Wand [26] started the work on extensible records and proposes row types [27] for records. Cardelli and Mitchell [4] define three primitive operations on records that are different from ours: *selection*, *restriction*, and *extension*. Following this approach, Leijen [17] define record update in terms of restriction and extension, while in our record system record update is a primitive operation. Both Leijen's system and ours allows records that contain duplicate labels. Arguably Leijen's system is stronger. For example, it supports passing record labels as arguments to functions. He also shows en-

⁶ Why the restriction?

coding an intersection type using first-class labels. Chlipala's Ur [7] explains record as type level constructs.

Indeed, our system can be adapted to simulate systems that support extensible records but not intersection of ordinary types like `Int` and `Float` by allowing only intersection of record types.

$\vdash_{\text{rec}} \tau$ states that τ is a record type, or the intersection of record types, and so forth.

$$\vdash_{\text{rec}} \{l : \tau\} \quad (\text{RecBase})$$

$$\frac{\vdash_{\text{rec}} \tau_1 \quad \vdash_{\text{rec}} \tau_2}{\vdash_{\text{rec}} \tau_1 \wedge \tau_2} \quad (\text{RecStep})$$

$$\frac{\begin{array}{c} \Gamma \vdash e_1 : \tau_1 \hookrightarrow E_1 \quad \vdash_{\text{rec}} \tau_1 \\ \Gamma \vdash e_2 : \tau_2 \hookrightarrow E_2 \quad \vdash_{\text{rec}} \tau_2 \end{array}}{\Gamma \vdash e_1, e_2 : \tau_1 \wedge \tau_2 \hookrightarrow \langle E_1, E_2 \rangle} \quad (\text{Merge'})$$

Of course our approach has its limitation as duplicated labels in a record are allowed. This has been discussed in a larger issue by Dunfield [14].

8. Conclusions and Further Work

Acknowledgments

Acknowledgments, if needed.

References

- [1] N. Amin, A. Moors, and M. Odersky. Dependent object types. In *19th International Workshop on Foundations of Object-Oriented Languages*, number EPFL-CONF-183030, 2012.
- [2] N. Amin, T. Rompf, and M. Odersky. Foundations of path-dependent types. In *Proceedings of the 2014 ACM International Conference on Object Oriented Programming Systems Languages & Applications*, pages 233--249. ACM, 2014.
- [3] L. Cardelli. *Extensible records in a pure calculus of subtyping*. Digital. Systems Research Center, 1992.
- [4] L. Cardelli and J. C. Mitchell. Operations on records. In *Mathematical foundations of programming semantics*, pages 22--52. Springer, 1990.
- [5] L. Cardelli and P. Wegner. On understanding types, data abstraction, and polymorphism. *ACM Computing Surveys (CSUR)*, 17(4):471--523, 1985.
- [6] G. Castagna, G. Ghelli, and G. Longo. A calculus for overloaded functions with subtyping. *Information and Computation*, 117(1):115--135, 1995.
- [7] A. Chlipala. Ur: statically-typed metaprogramming with type-level record computation. In *ACM Sigplan Notices*, volume 45, pages 122--133. ACM, 2010.
- [8] A. B. Compagnoni and B. C. Pierce. Higher-order intersection types and multiple inheritance. *Mathematical Structures in Computer Science*, 6(5):469--501, 1996.
- [9] M. Coppo, M. Dezani-Ciancaglini, and B. Venneri. Functional characters of solvable terms. *Mathematical Logic Quarterly*, 27(2-6):45--58, 1981.
- [10] K. Cray. Simple, efficient object encoding using intersection types. Technical report, Cornell University, 1998.
- [11] R. Davies. *Practical refinement-type checking*. PhD thesis, University of Western Australia, 2005.
- [12] R. Davies and F. Pfenning. Intersection types and computational effects. In *ACM Sigplan Notices*, volume 35, pages 198--208. ACM, 2000.
- [13] J. Dunfield. Refined typechecking with stardust. In *Proceedings of the 2007 workshop on Programming languages meets program verification*, pages 21--32. ACM, 2007.
- [14] J. Dunfield. Elaborating intersection and union types. *Journal of Functional Programming*, 24(2-3):133--165, 2014.
- [15] T. Freeman and F. Pfenning. *Refinement types for ML*, volume 26. ACM, 1991.
- [16] A. Frisch, G. Castagna, and V. Benzaken. Semantic subtyping: Dealing set-theoretically with function, union, intersection, and negation types. *Journal of the ACM (JACM)*, 55(4):19, 2008.
- [17] D. Leijen. Extensible records with scoped labels. *Trends in Functional Programming*, 5:297--312, 2005.
- [18] D. Malayeri and J. Aldrich. Integrating nominal and structural subtyping. In *ECOOP 2008--Object-Oriented Programming*, pages 260--284. Springer, 2008.
- [19] N. Nystrom, X. Qi, and A. C. Myers. J&: nested intersection for scalable software composition. In *ACM SIGPLAN Notices*, volume 41, pages 21--36. ACM, 2006.
- [20] B. C. Pierce. Programming with intersection types, union types, and polymorphism. 1991.
- [21] B. C. Pierce. *Programming with intersection types and bounded polymorphism*. PhD thesis, Carnegie Mellon University Pittsburgh, PA, 1991.
- [22] B. C. Pierce. Intersection types and bounded polymorphism. *Mathematical Structures in Computer Science*, 7(02):129--193, 1997.
- [23] B. C. Pierce. *Types and programming languages*. MIT press, 2002.
- [24] G. Plotkin, M. Abadi, and L. Cardelli. Subtyping and parametricity. In *Logic in Computer Science, 1994. LICS'94. Proceedings., Symposium on*, pages 310--319. IEEE, 1994.
- [25] J. C. Reynolds. *Design of the programming language Forsythe*. Springer, 1997.
- [26] M. Wand. Complete type inference for simple objects. In *LICS*, volume 87, pages 37--44, 1987.
- [27] M. Wand. Type inference for record concatenation and multiple inheritance. In *Logic in Computer Science, 1989. LICS'89, Proceedings., Fourth Annual Symposium on*, pages 92--97. IEEE, 1989.