

## Ex-Ante versus Ex-Post Compromise

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# Introducing the problem

**Setting:** Several voters express their preferences over a set of alternatives

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**Goal:** Find a procedure determining a collective choice that promote a notion of compromise

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- **Majoritarian Compromise**: MVR and ties are broken according to the quantity of support these receive
- **Fallback Bargaining**: bargainers fall back to less and less preferred alternatives until they reach a unanimous agreement
- **q-approval FB**: picks the alternatives which receive the support of  $q$  voters at the highest possible quality, breaking ties according to the quantity of support



# Example 1

$$|N| = 100, A = \{a, b, c\}$$

|           |          |          |          |
|-----------|----------|----------|----------|
| <b>51</b> | <i>a</i> | <i>b</i> | <i>c</i> |
| <b>49</b> | <i>c</i> | <i>b</i> | <i>a</i> |

- Plurality:  $\{a\}$

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- Plurality:  $\{a\}$
- MVR:  $\{a\}$

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- Plurality:  $\{a\}$
- MVR:  $\{a\}$
- MC:  $\{a\}$

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- Plurality:  $\{a\}$
- MVR:  $\{a\}$
- MC:  $\{a\}$
- FB:  $\{b\}$
- q-approval FB  $q \in \{1, \dots, \frac{n}{2} + 1\}$ :  $\{a\}$

## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
|-----------------------------|----------|----------|----------|----------|
| <b>26</b>                   | <i>a</i> | <i>b</i> | <i>c</i> | <i>d</i> |
| <b>25</b>                   | <i>c</i> | <i>b</i> | <i>a</i> | <i>d</i> |
| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$

## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
|-----------------------------|----------|----------|----------|----------|
| <b>26</b>                   | <i>a</i> | <i>b</i> | <i>c</i> | <i>d</i> |
| <b>25</b>                   | <i>c</i> | <i>b</i> | <i>a</i> | <i>d</i> |
| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$
- MVR: for  $z < 76$   $\{a, b\}$

## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
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| <b>26</b>                   | <i>a</i> | <i>b</i> | <i>c</i> | <i>d</i> |
| <b>25</b>                   | <i>c</i> | <i>b</i> | <i>a</i> | <i>d</i> |
| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$
- MVR: for  $z < 76$   $\{a, b\}$

|          |           |           |
|----------|-----------|-----------|
|          | $1^\circ$ | $2^\circ$ |
| <b>a</b> | 26        | $126 - z$ |
| <b>b</b> | 0         | $z$       |
| <b>c</b> | 25        | 25        |
| <b>d</b> | 49        | 49        |



## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
|-----------------------------|----------|----------|----------|----------|
| <b>26</b>                   | <i>a</i> | <i>b</i> | <i>c</i> | <i>d</i> |
| <b>25</b>                   | <i>c</i> | <i>b</i> | <i>a</i> | <i>d</i> |
| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

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|          |           |           |
|----------|-----------|-----------|
|          | $1^\circ$ | $2^\circ$ |
| <b>a</b> | 26        | 51        |
| <b>b</b> | 0         | 75        |
| <b>c</b> | 25        | 25        |
| <b>d</b> | 49        | 49        |

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$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
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| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$
- MVR: for  $z < 76$   $\{a, b\}$  , for  $z \geq 76$   $\{b\}$

|          |           |           |
|----------|-----------|-----------|
|          | $1^\circ$ | $2^\circ$ |
| <b>a</b> | 26        | 50        |
| <b>b</b> | 0         | 76        |
| <b>c</b> | 25        | 25        |
| <b>d</b> | 49        | 49        |

## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

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| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$
- MVR: for  $z < 76$   $\{a, b\}$  , for  $z \geq 76$   $\{b\}$
- MC:  $\{b\}$

## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
|-----------------------------|----------|----------|----------|----------|
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| <b>25</b>                   | <i>c</i> | <i>b</i> | <i>a</i> | <i>d</i> |
| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$
- MVR: for  $z < 76$   $\{a, b\}$  , for  $z \geq 76$   $\{b\}$
- MC:  $\{b\}$
- FB:  $\{a\}$

## Example 2

$$|N| = 100, A = \{a, b, c, d\}, z \in \{64, \dots, 99\}$$

|                             |          |          |          |          |
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| <b><math>z - 51</math></b>  | <i>d</i> | <i>b</i> | <i>a</i> | <i>c</i> |
| <b><math>100 - z</math></b> | <i>d</i> | <i>a</i> | <i>c</i> | <i>b</i> |

- Plurality:  $\{d\}$
- MVR: for  $z < 76$   $\{a, b\}$  , for  $z \geq 76$   $\{b\}$
- MC:  $\{b\}$
- FB:  $\{a\}$
- q-approval FB  $q \in \{\frac{n}{2}, \dots, z\}$ :  $\{b\}$

# Motivation

$$|N| = 2, |A| = 2k + 2$$

| $\mathbf{i_1}$ | $\mathbf{i_2}$ |
|----------------|----------------|
| $x$            | $b_1$          |
| $a_1$          | $\cdot$        |
| $\cdot$        | $\cdot$        |
| $\cdot$        | $b_{k-1}$      |
| $\cdot$        | $y$            |
| $a_k$          | $x$            |
| $y$            | $b_k$          |
| $b_1$          | $a_1$          |
| $\cdot$        | $\cdot$        |
| $\cdot$        | $\cdot$        |
| $b_k$          | $a_k$          |

# Ex-Ante versus Ex-Post Perspective

## ex-ante compromise

imposes over individuals a willingness to compromise but it does not ensure an outcome where everyone has effectively compromised

## ex-post compromise

favors an outcome where every voter gives up her most preferred positions if this increases equality

# Cardinal Compromise

## Setting

$A$  alternatives

$N$  voters

$u_i \in \mathbb{R}^A$  utility function depending on the rank of voter  $i$



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## Setting

$A$  alternatives

$N$  voters

$u_i \in \mathbb{R}^A$  utility function depending on the rank of voter  $i$

$\lambda_i^u(x) = \max_{a \in A} u_i(a) - u_i(x)$  represents the loss of utility for the voter  $i$  if the alternative  $x$  is elected instead of her favorite one; and  $\lambda^u(x)$  represents the vector of these losses

# Spread Measure

$$\sigma : \mathbb{R}_+^N \longrightarrow \mathbb{R}_+$$

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## Pure Equality Recognition

$$r_i = r_j \quad \forall i, j \in N \Rightarrow \sigma(r) = 0 \quad r \in \mathbb{R}_+^N$$

# Spread Measure

$$\sigma : \mathbb{R}_+^N \longrightarrow \mathbb{R}_+$$

## Pure Equality Recognition

$$r_i = r_j \ \forall i, j \in N \Rightarrow \sigma(r) = 0 \quad r \in \mathbb{R}_+^N$$

## Pairwise Pareto Dominance

$$[|r_i - r_j| \leq |s_i - s_j| \ \forall i, j \in N] \Rightarrow \sigma(r) \leq \sigma(s) \quad r, s \in \mathbb{R}_+^N$$

# Spread Measure

## Example

$$\bar{r} = \frac{\sum_{i=1}^n r_i}{n}$$

$$\sigma_{avg}(r) = \sum_{i=1}^n |\bar{r} - r_i|$$

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$$\sigma_{avg}(r) = \sum_{i=1}^n |\bar{r} - r_i|$$

Examples:

$$s = (3, 3, 3, 3)$$

$$\sigma_{avg}(s) = \sum_{i=1}^4 (3-3) = 0$$

$$t = (1, 2, 3, 4)$$

$$\sigma_{avg}(t) = |2.5-1|+|2.5-2|+|2.5-3|+|2.5-4| = 4$$

$$w = (1, 3, 5, 7)$$

$$\sigma_{avg}(w) = |4-1|+|4-3|+|4-5|+|4-7| = 8$$

# Cardinal Compromise

$\mathcal{U}$  set of injective utility functions defined over  $A$

$PO(u)$  set of Pareto optimal alternatives at  $u \in \mathcal{U}$

$\lambda^u(x)$  losses vector when electing the alternative  $x$

$\sigma$  spread measure

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$\lambda^u(x)$  losses vector when electing the alternative  $x$

$\sigma$  spread measure

$$C^\sigma(u) = \{x \in PO(u) : \sigma(\lambda^u(x)) \leq \sigma(\lambda^u(y)), \forall y \in A\}$$



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$$|N| = 100, A = \{a, b, c\}$$

|           |          |          |          |
|-----------|----------|----------|----------|
| <b>51</b> | <i>a</i> | <i>b</i> | <i>c</i> |
| <b>49</b> | <i>c</i> | <i>b</i> | <i>a</i> |

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## Example

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|          |  |                       |           |                        |                        |          |                         |
|----------|--|-----------------------|-----------|------------------------|------------------------|----------|-------------------------|
|          |  |                       | <b>51</b> | <i>a</i>               | <i>b</i>               | <i>c</i> |                         |
|          |  |                       | <b>49</b> | <i>c</i>               | <i>b</i>               | <i>a</i> |                         |
|          |  | <i>i</i> <sub>1</sub> | ...       | <i>i</i> <sub>51</sub> | <i>i</i> <sub>52</sub> | ...      | <i>i</i> <sub>100</sub> |
| <b>a</b> |  | 2                     | ...       | 2                      | 0                      | ...      | 0                       |
| <b>b</b> |  | 1                     | ...       | 1                      | 1                      | ...      | 1                       |
| <b>c</b> |  | 0                     | ...       | 0                      | 2                      | ...      | 2                       |

# Cardinal Compromise

## Example

$$|N| = 100, A = \{a, b, c\}$$

$$\begin{array}{cc} \mathbf{51} & a & b & c \\ \mathbf{49} & c & b & a \end{array}$$

$$\begin{array}{l} \lambda(a) = (i_1, \dots, i_{51}, i_{52}, \dots, i_{100}) \\ \lambda(b) = (0, \dots, 0, 2, \dots, 2) \\ \lambda(c) = (1, \dots, 1, 1, \dots, 1) \\ \lambda(c) = (2, \dots, 2, 0, \dots, 0) \end{array}$$

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## Example

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$$\begin{array}{cc} \mathbf{51} & a & b & c \\ \mathbf{49} & c & b & a \end{array}$$

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$$\sigma_{avg}(\lambda(a)) = 99.96, \quad \sigma_{avg}(\lambda(b)) = 0, \quad \sigma_{avg}(\lambda(c)) = 100.04$$

# Cardinal Compromise

## Example

$$|N| = 100, A = \{a, b, c\}$$

$$\begin{array}{cc} \mathbf{51} & a & b & c \\ \mathbf{49} & c & b & a \end{array}$$

$$\begin{array}{l} \lambda(a) = (i_1, \dots, i_{51}, i_{52}, \dots, i_{100}) \\ \lambda(b) = (0, \dots, 0, 2, \dots, 2) \\ \lambda(c) = (1, \dots, 1, 1, \dots, 1) \\ \lambda(c) = (2, \dots, 2, 0, \dots, 0) \end{array}$$

$$\sigma_{avg}(\lambda(a)) = 99.96, \quad \sigma_{avg}(\lambda(b)) = 0, \quad \sigma_{avg}(\lambda(c)) = 100.04$$

$$C^{\sigma_{avg}}(u) = b$$

# Ordinal Compromise

$P_i \in L(A)$  linear order over  $A$  which represents the preference of  $i \in N$

$v : \{1, \dots, m\} \rightarrow \mathbb{R}$  utility assignment

$v_{P_i} \in \mathbb{R}^A$  utility function for  $P_i \in L(A)$  induced by  $v$

$\mathbf{v} : L(A)^N \rightarrow \mathcal{U}$  function mapping a profile of ordinal preferences to a utility profile

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$$(C^\sigma \circ \mathbf{v})(\{P_i, i \in N\}) = C^\sigma(\{v_{P_i}, i \in N\})$$

# Ordinal Compromise

$\Sigma^{\text{All}} = \mathbb{R}_+^N$  the set of all spread measures

## UA-independence

A class of spread measure  $\Sigma \subseteq \Sigma^{\text{All}}$  is UA-independent iff, given any  $\sigma \in \Sigma$  and any two UAs  $v$  and  $v'$ , there exists a  $\sigma' \in \Sigma$  such that

$$C^\sigma \circ v = C^{\sigma'} \circ v'$$



# Ordinal Compromise

## UA-independence

$\Sigma^{\text{PPd}} \subseteq \Sigma^{\text{All}}$  the class of spread measures that satisfy PPd

### Proposition 1:

$\Sigma^{\text{PPd}}$  is not UA-independent

# Ordinal Compromise

## UA-independence

$\Sigma^{\text{PPd}} \subseteq \Sigma^{\text{All}}$  the class of spread measures that satisfy PPd

### Proposition 1:

$\Sigma^{\text{PPd}}$  is not UA-independent

|       |       |       |       |       |       |       |       |       |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| $i_1$ | $x$   | $a_1$ | $a_2$ | $a_3$ | $y$   | $b_1$ | $b_2$ | $b_3$ |
| $i_2$ | $b_1$ | $b_2$ | $y$   | $x$   | $b_3$ | $a_1$ | $a_2$ | $a_3$ |

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|       |       |       |       |       |       |       |       |       |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| $i_1$ | $x$   | $a_1$ | $a_2$ | $a_3$ | $y$   | $b_1$ | $b_2$ | $b_3$ |
| $i_2$ | $b_1$ | $b_2$ | $y$   | $x$   | $b_3$ | $a_1$ | $a_2$ | $a_3$ |

|         |      |     |     |     |   |   |   |   |
|---------|------|-----|-----|-----|---|---|---|---|
| $k =$   | 1    | 2   | 3   | 4   | 5 | 6 | 7 | 8 |
| $v(k)$  | 7    | 6   | 5   | 4   | 3 | 2 | 1 | 0 |
| $v'(k)$ | 1000 | 999 | 998 | 997 | 3 | 2 | 1 | 0 |

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### Proposition 1:

$\Sigma^{\text{PPd}}$  is not UA-independent

|       | $v_{P_1}(\cdot)$ | $v_{P_2}(\cdot)$ | $\lambda^P(\cdot)$ |
|-------|------------------|------------------|--------------------|
| $x$   | 7                | 4                | (0, 3)             |
| $y$   | 3                | 5                | (4, 2)             |
| $a_1$ | 6                | 2                | (1, 5)             |
| $a_2$ | 5                | 1                | (2, 6)             |
| $a_3$ | 4                | 0                | (3, 7)             |
| $b_1$ | 2                | 7                | (5, 0)             |
| $b_2$ | 1                | 6                | (6, 1)             |
| $b_3$ | 0                | 3                | (7, 4)             |

$$C^\sigma(v) \in \{y\}$$

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$\Sigma^{\text{PPd}}$  is not UA-independent

|       | $v'_{P_1}(\cdot)$ | $v'_{P_2}(\cdot)$ | $\lambda'^P(\cdot)$ |
|-------|-------------------|-------------------|---------------------|
| $x$   | 1000              | 997               | (0, 3)              |
| $y$   | 3                 | 998               | (997, 2)            |
| $a_1$ | 999               | 2                 | (1, 998)            |
| $a_2$ | 998               | 1                 | (2, 999)            |
| $a_3$ | 997               | 0                 | (3, 1000)           |
| $b_1$ | 2                 | 1000              | (998, 0)            |
| $b_2$ | 1                 | 999               | (999, 1)            |
| $b_3$ | 0                 | 3                 | (1000, 997)         |

$$C^\sigma(v') \in \{x, b_3\}$$

# Ordinal Compromise

## UA-independence

$\Sigma_{\text{threshold}} = \{\sigma^k, k \in \mathbb{R}\}$  where  $\sigma^k(\lambda) = \#\{i \in N \mid \lambda_i \geq k\}$

Proposition 2:

$\Sigma_{\text{threshold}}$  is UA-independent

# Ordinal Compromise

UA-independence

## Proposition 3:

$\sigma^k \in \Sigma_{\text{threshold}}$  fails Pure Equality Recognition and Pairwise Pareto Dominance

# Ordinal Compromise

UA-independence

## Proposition 3:

$\sigma^k \in \Sigma_{\text{threshold}}$  fails Pure Equality Recognition and Pairwise Pareto Dominance

$$\lambda(x) = (4, 4, 4, 4) \quad \sigma^3(\lambda(x)) = 4$$



# Ordinal Compromise

## UA-independence

### Proposition 3:

$\sigma^k \in \Sigma_{\text{threshold}}$  fails Pure Equality Recognition and Pairwise Pareto Dominance

|       |     |     |     |     |
|-------|-----|-----|-----|-----|
| $i_1$ | $a$ | $d$ | $b$ | $c$ |
| $i_2$ | $b$ | $c$ | $d$ | $a$ |

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$v$  assigns the utility values 10, 2, 1, 0 respectively to the ranks 1, 2, 3, 4

$$\begin{array}{rcl}
 & & \sigma^1(\lambda^P(\cdot)) \\
 \lambda^P(a) & = & \begin{pmatrix} 0 & 10 \end{pmatrix} & 1 \\
 \lambda^P(b) & = & \begin{pmatrix} 9 & 0 \end{pmatrix} & 1 \\
 \lambda^P(c) & = & \begin{pmatrix} 10 & 8 \end{pmatrix} & 2 \\
 \lambda^P(d) & = & \begin{pmatrix} 8 & 9 \end{pmatrix} & 2
 \end{array}$$

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- ...

# Appendix I

## Minimal Liberty

A social planner must choose between a world  $x$  where individuals may sell their organs, and a world  $y$  where they do not

|     | $u_1$ | $u_2$ |
|-----|-------|-------|
| $x$ | 1     | 100   |
| $y$ | 0     | 0     |

Even though  $y$  is Pareto dominated, the social planner might prefer  $y$  to  $x$

Thank You!



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