The chromatic number of a random lift of a regular graph

Xavier Pérez-Giménez joint work with JD Nir

University of Nebraska-Lincoln



Moscow Institute of Physics and Technology, Dec 2020

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Surjective graph homomorphism $\Pi: L \to G$ that is a bijection between edges incident with ν and edges incident with $\Pi(\nu)$.

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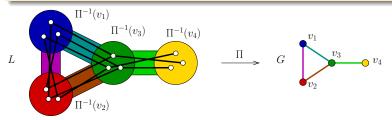
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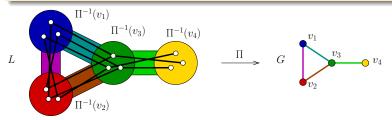


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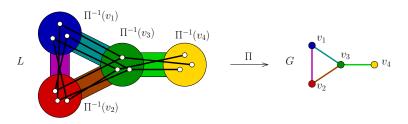
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Note: G may have loops and multiple edges.

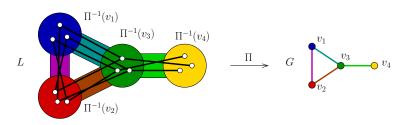
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Fact 1: G connected \Longrightarrow all fibers $\Pi^{-1}(v)$ have same cardinality.

L is an *n***-lift** if all $|\Pi^{-1}(v)| = n$.

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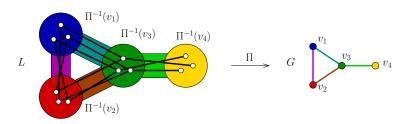


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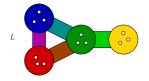
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Fact 3: $\chi(L) \leq \chi(G)$



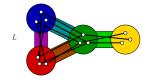
Random *n*-lift model (Amit, Linial 2002):

- Replace $v \in V(G)$ by bin with n vertices.
- Replace $e \in E(G)$ by random perfect matching. (slightly different for loops.)



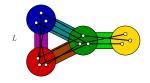
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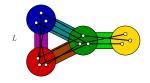
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Facts (for d-regular G):

- $G = B_{d/2} \Longrightarrow L$ contiguous to uniform d-regular multigraph.
- Not true for $G = K_{d+1}$



(Amit, Linial, Matoušek 2002)

Problem 1

For any
$$G$$
: is $\chi(L) = \Omega\left(\frac{\chi(G)}{\log \chi(G)}\right)$ a.a.s.?

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Conjecture

For any G, there is k_G with $\chi(L) = k_G$ a.a.s.

...and many more open questions!



Our results

Thm (P-G, Nir 2019++):

Let $d \ge 3$ and $k_d = \min\{k \in \mathbb{N} : d < 2k \log k\}$. $(k_d \approx \frac{d}{2 \log d})$

Let *L* be random *n*-lift of $G = K_{d+1}$.

Then a.a.s. $\chi(L) \in \{k_d, k_d + 1\}.$

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Thm (Kemkes, P-G, Wormald 2010):

Analogous result holds for uniform d-regular graphs.

(Improved by Coja-Oghlan, Efthymiou, Hetterich 2016.)

Main tools

- Small subgraph conditioning method (Robinson & Wormald 1992)
- Optimization over stochastic matrices (Achlioptas, Naor 2005)
- Laplace summation method (Greenhill, Janson, Ruciński 2010)
- Saddle-point method
- Algebraic graph theory
 - Kirchhoff Matrix-Tree Thm
 - Counting non-backtracking closed walks

Proof structure

Lower bound on $\chi(L)$:

X = # k-colourings of L.

Thm: If $k < k_d$, then $\mathbf{E}X = o(1)$

Then $P(X > 0) \le EX = o(1)$.

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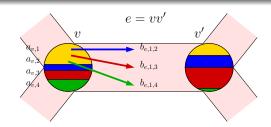
Thm: If $k > k_d$, then $\mathbf{E}Y^2 = \Theta((\mathbf{E}Y)^2)$

Then $P(Y > 0) \ge \frac{(EY)^2}{EY^2} \sim C > 0$ (Paley-Zygmund)

Unfortunately, C < 1 due to the influence of short cycles in L.

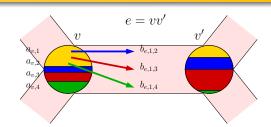


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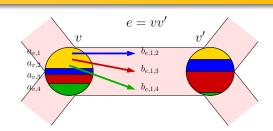
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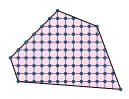
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Claim:

Max contribution is from $a_{v,i} = 1/k$, $b_{e,i,i'} = 1/k(k-1)$.

(We extend result by Achlioptas, Naor 2005.) 990

Summation domain:



$$\begin{cases} \mathbf{x} \in \mathcal{P} \\ B\mathbf{x} = \mathbf{y} \\ \mathbf{x} \in \left(\frac{1}{n}\mathbb{Z}\right)^D \end{cases}$$

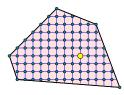
where:

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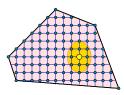
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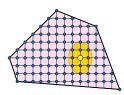
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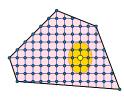
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• About C: It depends on

 $\begin{cases} \text{Hessian of } f \\ \text{Volume of fundamental cell in lattice} \end{cases}$

(Greenhill, Janson, Ruciński 2010)

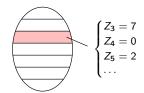
Instead, we count maximal forests in Γ with incidence matrix B.

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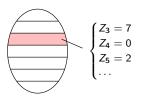


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Rough idea:

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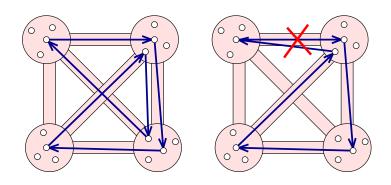
•
$$\frac{\mathsf{E}(YZ_i)}{\mathsf{E}Y} = 1 + \delta_i + o(1)$$
 (& joint factorial moments)
(i.e. $Z_i \sim \mathsf{Poi}(1 + \delta_i)$ in space "weighted" by Y).

$$\bullet \ \frac{\mathsf{E} Y^2}{(\mathsf{E} Y)^2} = \exp(\sum_i \lambda_i \delta_i^2) + o(1).$$

Then $P(Y > 0) \rightarrow 1$ (+ contiguity [...]).

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Counting non-backtracking closed walks



Some algebraic tools: (Friedman 2008)

(Amit, Linial, Matoušek 2002)

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