# 22. Searching a List

#### Topics:

Linear Search

Binary Search

Measuring Execution Time

The Divide and Conquer Framework

### Search

#### Examples:

Is this song in that playlist?

Is this number in that phone book?

Is this name in that phone book?

Is this fingerprint in that archive of fingerprints?

Is this photo in that yearbook?

# More on Using Phone Books

The Manhatten phone book has 1,000,000+ entries.

How is it possible to locate a name by examining just a tiny, tiny fraction of those entries?



There must be a great search algorithm behind the scenes.

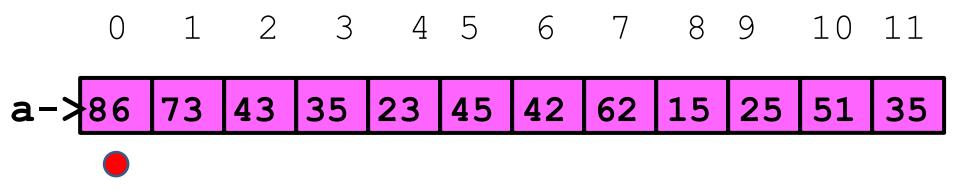
### LinSearch: The Spec

```
def LinSearch(x,a):
        Returns an int k with the
    property that a[k] == x is True.
    If no such k exists, then
    k = -1.
    PreC: a is a nonempty list of
    ints and x is an int.
    ** ** **
```

0 1 2 3 4 5 6 7 8 9 10 11 a-86 73 43 35 23 45 42 62 15 25 51 35

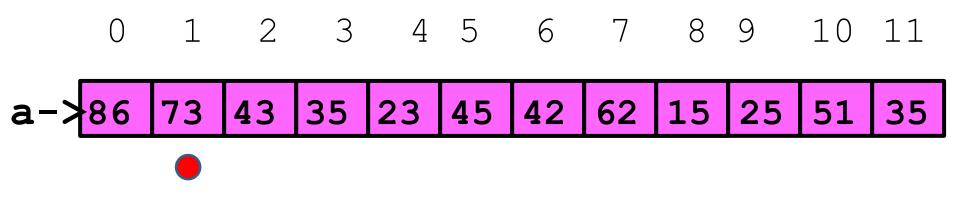
```
k-> 23
```

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]:
        return k
    return -1
```



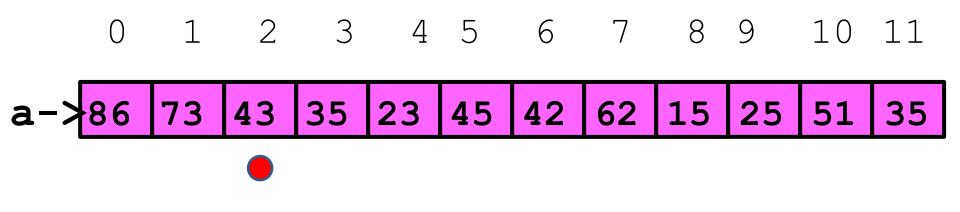
```
k-> 23
```

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```



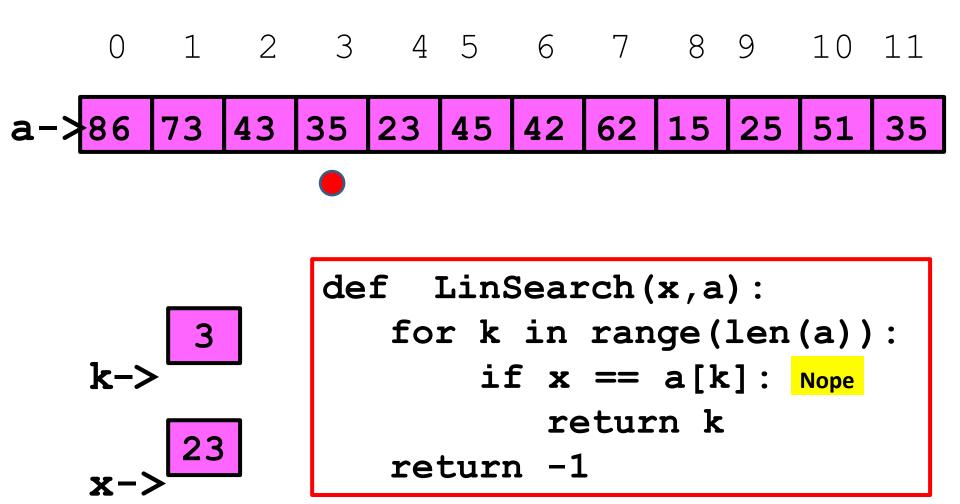
```
k-> 23
```

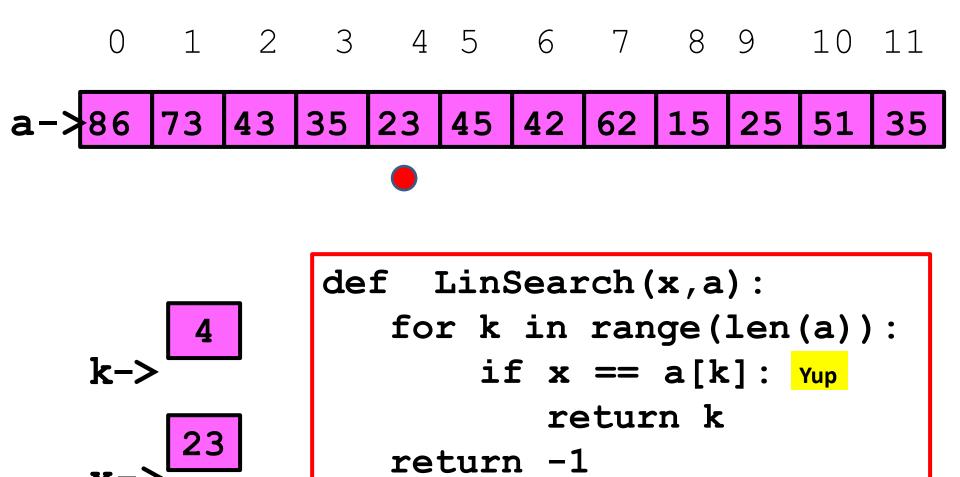
```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```

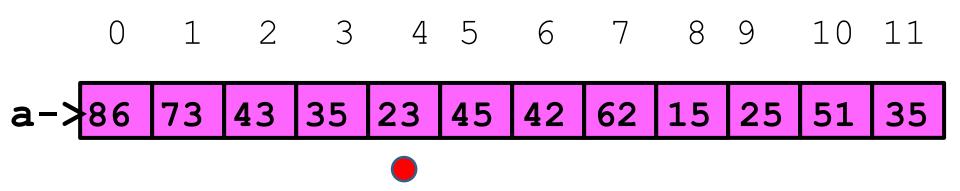


```
k-> 23
```

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```

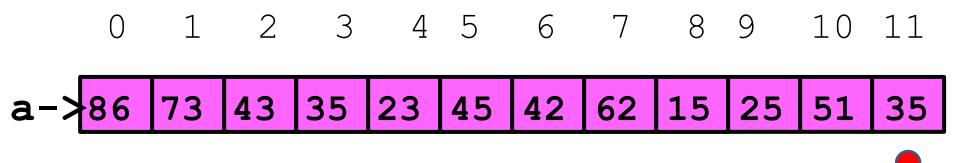


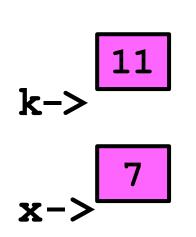




```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]:
            return k Alldone
    return -1
```

### Linear Search: No Match Case





```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```

### Linear Search: No Match Case

```
a - > 86 73 43 35 23 45 42 62 15 25 51 35
```

```
x-> 7
```

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]:
            return k
    return -1 Yup
```

# Linear Search: While Implementation

```
def LinSearchW(x,a):
    k=0
    while k<len(a) and a[k]!=x:
        k+=1
    if k==len(a):
        return -1
    else:
        return k</pre>
```

# Binary Search

Now we assume that the list to be searched is sorted from little to big.

```
a = [10,20,40,60,90]
a = ['brown','dog','fox','lazy','quick','the']
```

# Back to Using Phone Books

The Ithaca phone book has 10,000+ entries.

The Manhatten
phone book has

1,000,000+ entries. But it does not take 100 x longer to look something up. Why?

wide at	SuperPages.com	195	Car U
nack lent	Cartage New England Inc	Carter F 24 Hillion Ros 02131	Carter Nella E 333 Maschalls Av Bos 07115
17 566-1282	26 Allen La Ignavich 01938	Faye & Ricky 157 Calumbus Av Box 07116	Nicholas S F
81 447-4101	18 Jewett Rus 02131	Francis S 114 Terrole W Res 02132 617 323-6781	115 Rangelob Av MII 07160
01 40 4101	Cartagena Avith	Franklin & Anne	Nick 21 Fairfield 8os 02116
00 257-9981	9 Saxcroft Ray 02119	223 Mt Auburn Cam 02138	Nick & Debbi
	8 14/0 02136	Fred 47 Haverlord Jan 02130	196 Herrick Rd Newton 02459617 527-0480
17 566-1282	Jessica 50 Decater Chia 07129	Fred 96 Hindley Rt MI 02186	Nicole
17 364-5188	Lucilla 174 Harvard Cam 02139 617 491-5621	G & R # Verdum Der 102124	38 Oxfortawbut Dor 07177
	M 95 Rowe Ras 00131	Gayle 25 Frontenac Der 00124	P 94 Cristoward Ps Rox 82121 617 427-4754
361-0380	Carte Nicholas	Geo S 115 Mars Hill Rd Jan 02130 617 522-3215	P E 501 E Serb 5 Box 02127
17 566-4548	18 Appleton Seston 07116	George 125 Nashua Bas 02114617 367-9548	P L 44 Philippings Rox 07171
27 300 1310	Cartegena O 4 Millord Box 07118617 338-8219	Carter Halliday Associate	P R 91 Bysner Jam 02130617 983-8692
17 628-8248	Carten Thos J Sr & Claire	107 S Street Bos 02111	Paul & Constance
	1 Peradise Rd Mii 071M		114 Angeware Av W Rox 02137
17 445-5116	Thomas & Kathleen	26 Rurrig Brit Rd W Rox 02112 617 325-5465 Carter Hide Co Inc	Paul M 27 Unton Bri 02135
17 000 0000	50 Thompson Ln Mil 02186	146 Suntrear des 07110	Carter Pile Driving Inc 17 Beaver Ct
17 822-2982 17 427-5712	A Rosbury	Carter Hilary 61 Harvey Cam 07140 617 876-2750	Framingtons 01/02 Wellesley TelNo-781 235-8488
17 569-2698	A 31 Rethune Wy Roxbury 02119 617 442-1219	Horace	Carter Prudence
III JUT LUTE	A 260 Putram Av Cambridge 07139 617 492-4174	241 Walnut Air Roobury 02119617 442-5307	46 Franklin Watertown 021/2
17 667-5190	A M 255 Maschets Av Bos 02115 617 266-7153	Howard Jr 36 Notre Drue Rox 02119.617 445-5552	Prudence
	Adams 361 Centre St Mil 02186 617 698-9074	J Cam	46 Franklin Watertown 02172 617 926-7063 Reginald
17 569-1417	Alice 108 Kimurnock Bos 07215 617 425-0193	J 15 Chatham fire 02446	106 Brunswick Dorchester 02122617 541 - 2843
17 338-9110	Alice 45 Market Cambridge 02139 617 945-2711 Andrew F 62 Vinal Av Son 02143 617 625-7623	J 775 VW Play West Rosbury 02132 617 323-5574	Renee & Andrew
17 825-9195		Carter J Jacques MD	10 Walnut Ses 02108
11, 053-4142	1101 Beacon Bro 02446	1 Brookline Pl Bro 0246	Carter Rice Dowd
17 296-1593		Carter J M	Buildey Dunton Publishing 163 Main Wilmington 61887 Tall Free-Cial '2' & Then
	272 Newbury Baston 02116	1410 Columbia Rd S See 02127 617 464-1040	Toli Free-Dial '1' & Then
17 670-2078	B E 68 Gladeside Av Mat 02126 617 296-6911	Carter J M Ornamental Ironworks	Tail Free-Dail 'I' & Then
17 623-9001		Carter J Veal Co	Cust Svc-Printing 613 Main Wilmington
17 296-4725	Tufts-New England Medical Center Box 02111 Call	48 Newmarket So Rox 07118	Tot Free-Oat '1' & Then
17 270-4725	Carter Becky 8cs 02114	Carter James	Headquarters 613 Main Wilmington (1)897 Call 978 988-7447
17 542-1521	Bernard J	1573 Cambridge St Cam 02138617 492-1214	Ingalls Cronin 163 Main Wilmington 01887
	112 Gladstone E Bos 02128617 567-3430	James 182 Fisher Av Rosbury 02120617 739-2193	Toll Free-Dial '1' & Then
17 364-5232	Bithiah 25 Medway Dor 02124617 298-8713	James (17 074 0041	Carter Richard
17 541-5649	Blake 25 Mt Verson But 02108	37 Gold Star Rd Cambridge 82140617 876-8841 Jas L 14 Roseberry Rd Mat 02126617 361-0773	1679 Commwith Av Brighton 02215 617 987-0836 Richard A 57 Mt Vernan fon 02108617 566-7293
		Jane 114 Adena Rd Newton 02465617 964-0435	Richard A 57 Mt Vernor for 02108617 306-7293 Carter Richard A MD
17 739-2662	Carter & Buruess Consultants Inc	Jeffrey 41 Warren As Box 02116	170 Commwith Av Sos 02116
17 879-0030	23 East St Cam 02141	John 11 Manufield Bri 07134	Carter Richard K
17 541-3948	Carter C 2000 Commetts Av Bri 02135 617 782-2118	John 327 Summer Box 02210	15 Mercar 5 Box 07127617 268-0448
17 436-1513	C 228 Faywood Av East Boston 02326617 569-1545	John 40 Westwind Rd Der 00125 617 282-1235	Robert L 175 Richarle Av Care 02140. 617 864-1535
517 569-4119	C 359 Harvard Cam 82138	June 0 329 A Summit As Bri 02135 617 734-6109	Roger 150 St Butosph Bas 07115617 424-6148
ton 02120	C 610 Walk Hill Mart 02126	K 38 Browning Av Dorchester 02124	Roy 44 Concord Av Cart 02138
300 569-8782	C & M 43 Burroughs Jam 02130 617 524-9558	K 17 Esmond Dorchester 02121	Royce 18 Seninary Cha 02129

# Key Idea: Repeated Halving

To find Derek Jeter's number...

```
B = phone book
while (B is longer than 1 page):
   1. P = middle page of B
   2. Let Q be the first name on P
   3. if 'Jeter" comes before Q:
         Rip away the 2nd half of B
      else:
         Rip away the 1st half of B.
Scan remaining page P line-by-line for 'Jeter'
```

# What Happens to Phone Book Length?

```
Original: 3000 pages
```

After 1 rip: 1500 pages

After 2 rips: 750 pages

After 3 rips: 375 pages

After 4 rips: 188 pages

After 5 rips: 94 pages

After 12 rips: 1 page

# Binary Search

The idea of repeatedly halving the size of the "search space" is the main idea behind the method of binary search.

An item in a sorted array of length n can be located with approximately log2 n comparisons.

## What is log2(n)?

n	ceil(log2(n))	
10	4	
100	7	
1000	10	
10000	14	
100000	17	
1000000	20	

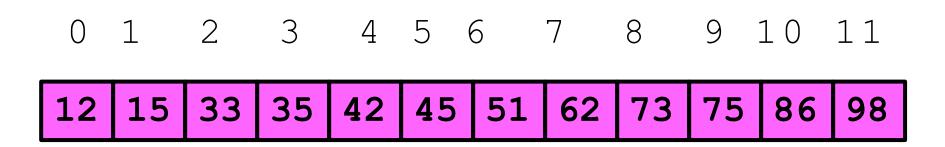
### BinSearch: The Spec

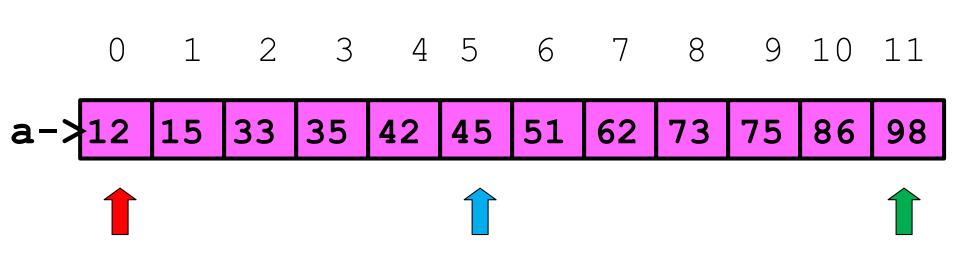
```
def BinSearch(x,a):
    """ Returns an int k with the
    property that a[k]==x is True.
    If no such k exists, then
    k==-1.
```

PreC: a is a nonempty list of ints that is sorted from smallest to largest. x is an int.

\*\* \*\* \*\*

# Example: Does this List have an Element With Value Equal to 70?





L: 0

 $a[Mid] \le x ????$ 

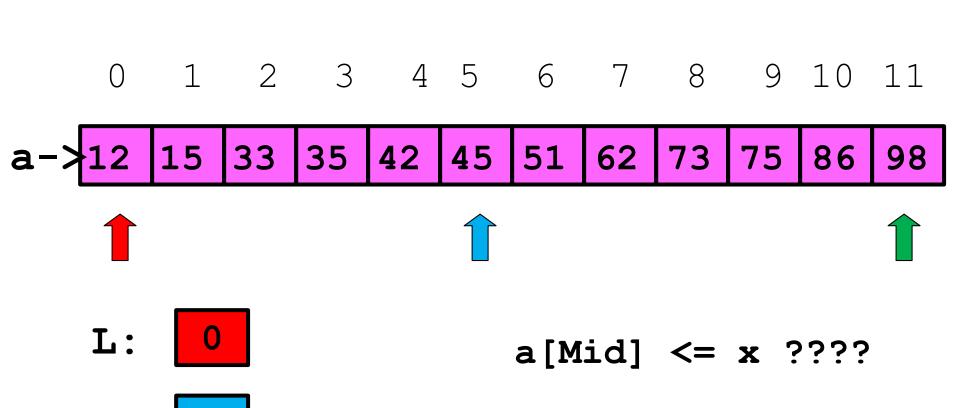
Mid: 5

R: 11 x: 70

Mid = (L+R)/2

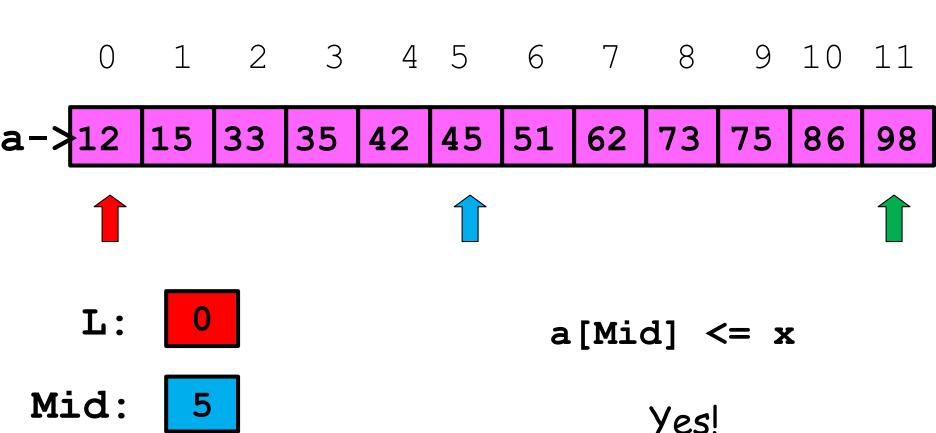
# The Midpoint Computations

L	R	(L+R)/2
0	11	5
2	6	4
1	100	50



Mid: 5

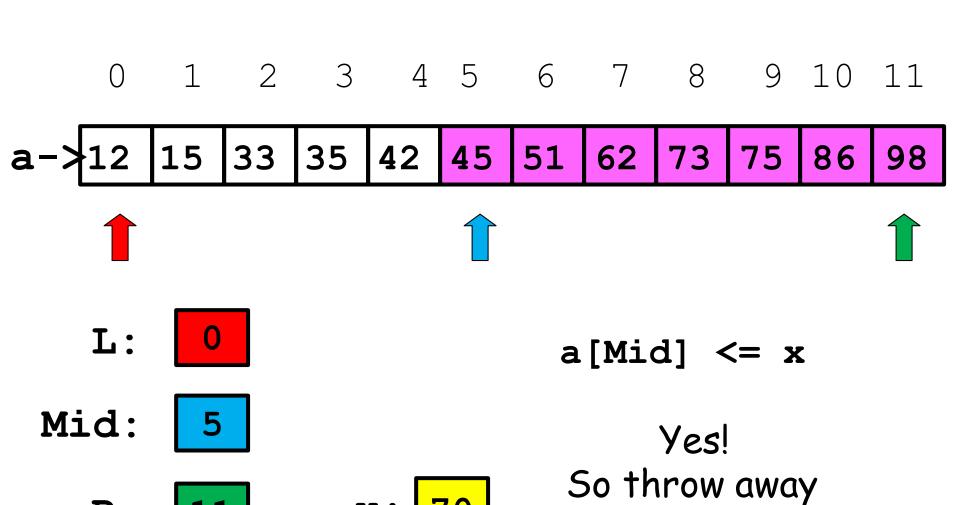
R: 11 x: 70



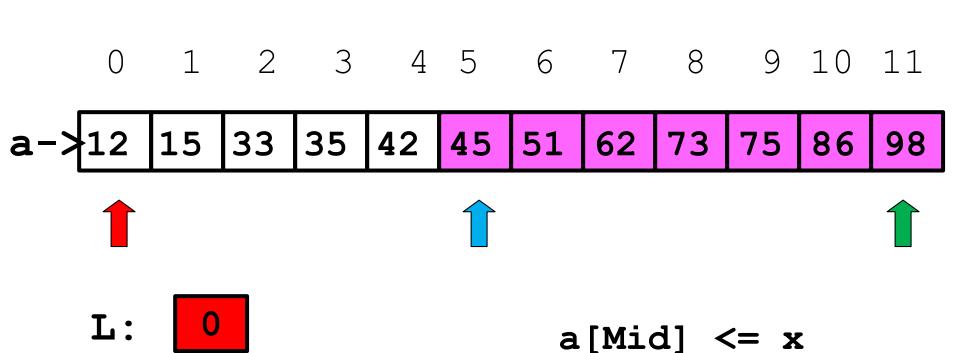
R: 11

x: 70

So throw away
The "left half"



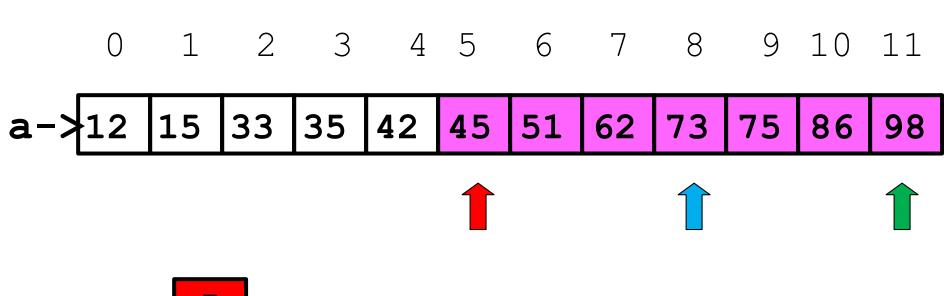
The "left half"



Mid: 5

Revise L and Mid

R: 11 x: 70

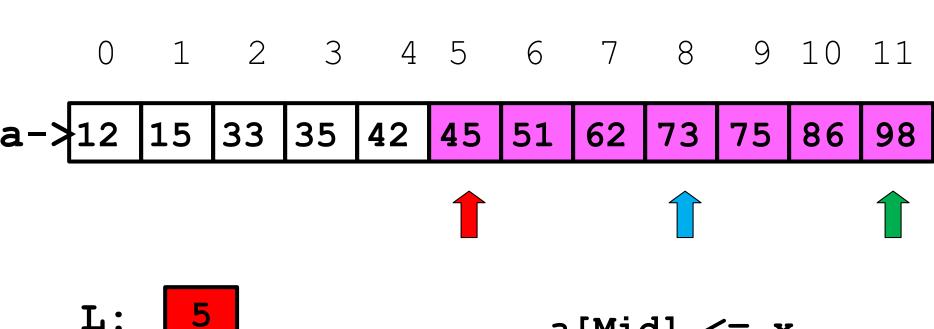


L: 5

a[Mid] <= x ???

Mid: 8

R: **11** x: 70



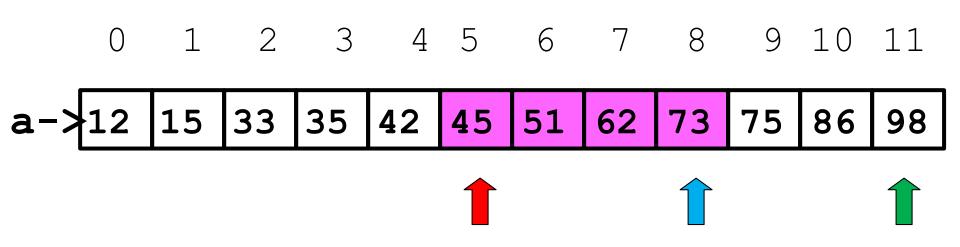
 $a[Mid] \le x$ 

No

x: 70

So throw away the "right half"

### Let's Look For x = 70



L: 5

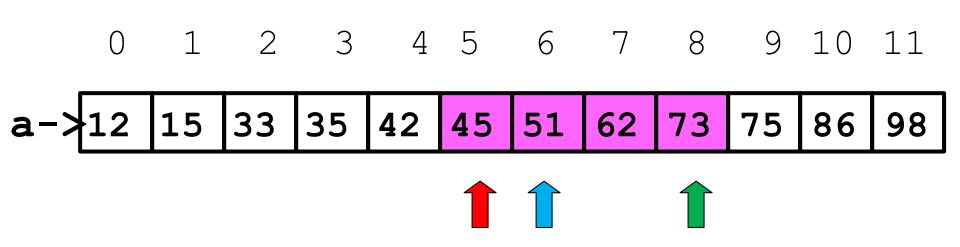
 $a[Mid] \le x$ 

Mid: 8

Revise R and Mid

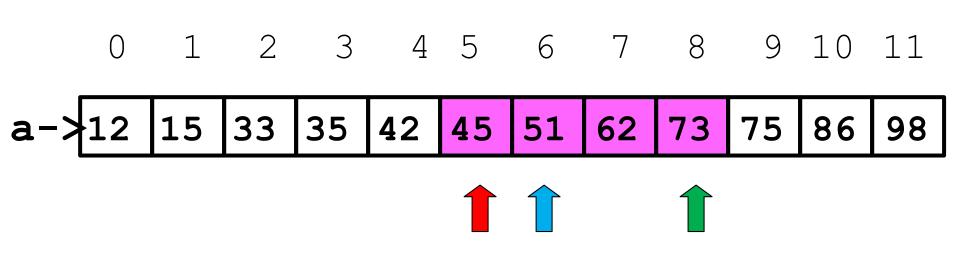
R: **11** x: 70

### Let's Look For x = 70



 $a[Mid] \le x$ 

Revise R and Mid

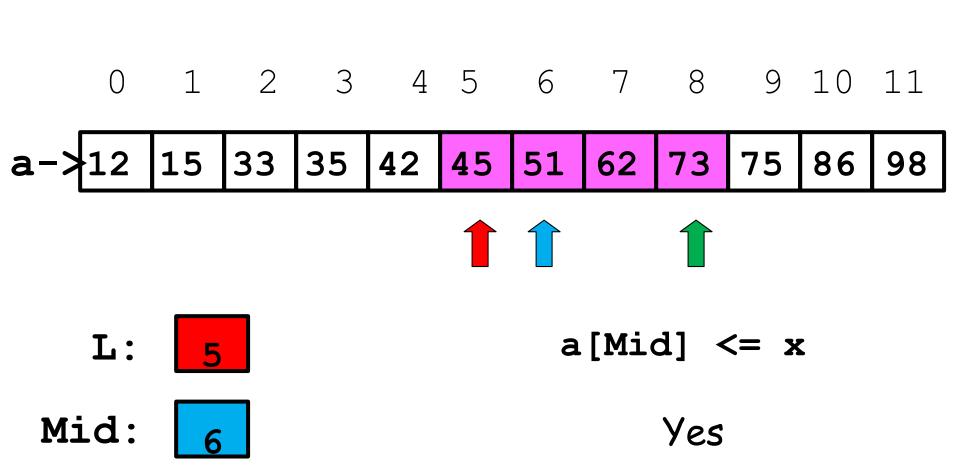


L: 5

 $a[Mid] \le x ????$ 

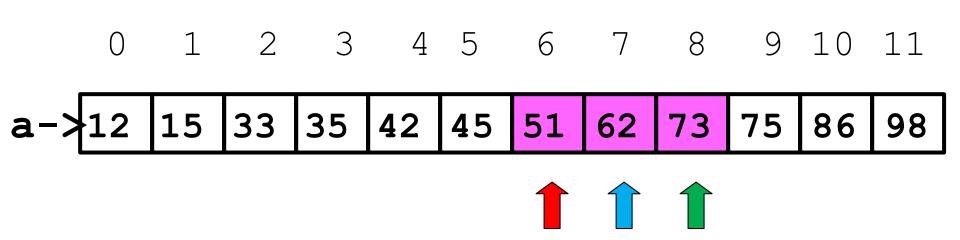
Mid: 6

R: 8 x: 70



R: 8 x: 70

Throw away the Left half



L: 6

 $a[Mid] \le x$ 

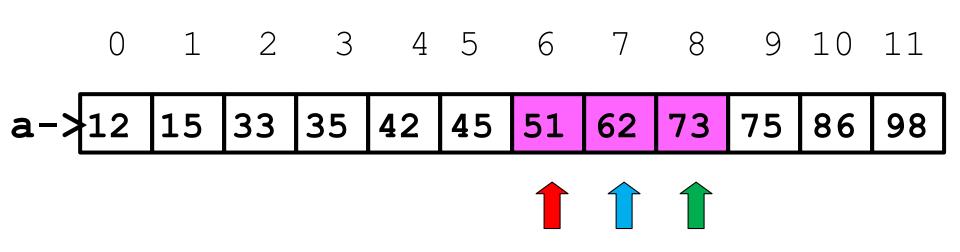
Mid: 7

Yes

R: 8

**x**: 70

### Let's Look For x in a



Mid: 7

R: 8

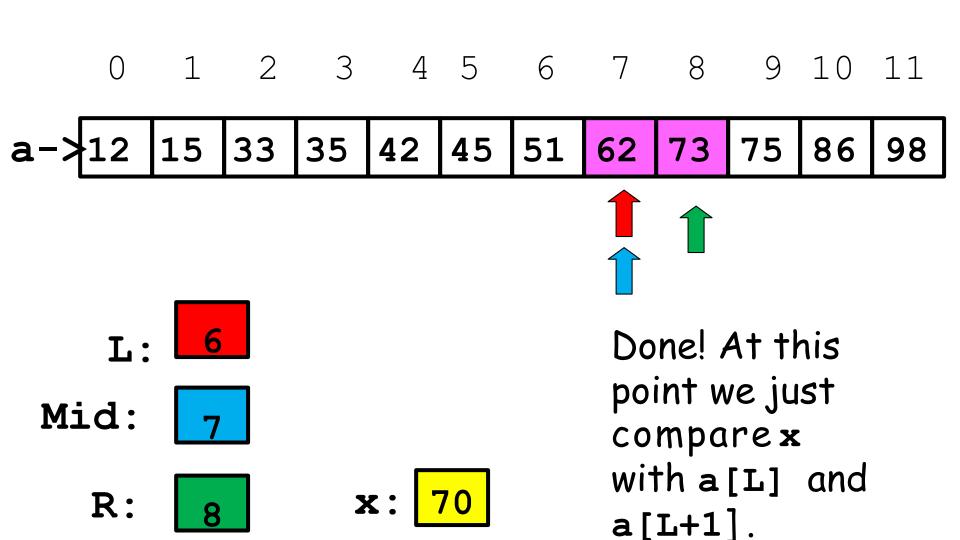
7

x: 70

$$a[Mid] \le x$$

Throw away the left half

### Let's Look For x in a



### What We Just Did

```
R = len(a)-1
while R-L > 1:
                         ALoop
    # a[L]<=x<=a[R]
                         Invariant
     Mid = (L+R)/2
     if x \le a[mid]:
          R = Mid
     else:
          L = Mid
```

### What We Just Did

```
R = len(a)-1
while R-L > 1:
    \# a[L] <= x <= a[R]
     Mid = (L+R)/2
     if x \le a[mid]:
          R = Mid
     else:
          L = Mid
```

### What We Just Did

```
L = 0
R = len(a)-1
while R-L > 1:
    \# a[L] <= x <= a[R]
     Mid = (L+R)/2
     if x \le a[mid]:
          R = Mid
     else:
          L = Mid
```

## After the Loop Ends



This is True:  $a[L] \le x \le a[L+1]$ 

## After the Loop Ends

```
a[L+1]
a[L]
if x==a[L]:
   return L
elif x==a[L+1]:
   return L+1
else:
   return -1
```

## Measuring Execution Time

We now have two ways to search a list:

LinSearch(x,a)
BinSearch(x,a)

Intuition: BinSearch much faster.

Can we quantify this with a "stop watch"?

### The timeit Module

This module can be used to time how long it takes to execute a chunk of code.

Typical chunk = some function of interest.

This is called benchmarking.

## Benchmarking

Let's benchmark LinSearch (x,a) and BinSearch (x,a).

Compare how long it takes when len(a) equals 1000, 10000, 100000, and 100000.

Our intuition tells us that as len(a) increases, BinSearch will be dramatically faster.

#### BinSearch vs LinSearch

n	tBin	tLin	tLinW
1000	0.0007	0.0064	0.0119
10000	0.0009	0.0668 0.8296	0.1203 1.2082
1000000	0.0015	17.7388	13.9341

```
tBin = time for BinSearch
tLin = time for LinSearch (for loop version)
tLinW = time for LinSearch (while-loop version)
```

#### BinSearch vs LinSearch

n	tLin/tBin	
1000	9	
10000	74	
100000	754	
1000000	7095	

Reporting ratios is more illuminating since we do not really care about the time units in this informal comparison

## Using the timeit Module

We show how this module was used to get the results on the previous slides.

Our LinSearch vs BinSearch example is very typical: is one function faster than another?

```
from timeit import *
S = """
   Set-up code
                               Yes, these are doc
                               strings.
// // //
    11 11 11
   Code to Benchmark
// // //
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

## The Set-Up and Bench Codes

```
from random import randint as randi
from ShowSearch import BinSearch
n = 10000
s = [randi(0,10*n) for i in range(n)]
s.sort()
x = s[n/2]
```

k=BinSearch(x,s)

The set-up code is run once.

It is not timed.

It just sets up the code to be timed.

```
from timeit import *
     11 11 11
   Set-up code
                               m times.
// // //
   Code to Benchmark
// // //
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

An "experiment" consists of running the blue code

The stopwatch will time how long it takes to do one experiment

Larger values necessary if the blue code executes very quickly

```
from timeit import *
     11 11 11
    Set-up code
                                   Timer returns
// // //
                                   alength-p
      11 11 11
                                   list. Each
                                   element is
    Code to Benchmark
                                   the stopwatch
                                   time for 1
// // //
                                   experiment
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

This helps control for other stuff that may be running on your computer.

```
from timeit import *
      11 11 11
    Set-up code
                                    In general, it is
// // //
                                     best to take
      11 11 11
                                    the mininum as
                                    the most reliable.
    Code to Benchmark
                                     The benchmark
                                    time is assigned
// // //
                                    to t
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

This helps control for other stuff that may be running on your computer.

# Why Benchmarking is Important

Confirms/refutes what our intuition might say about efficiency.

Makes us sensitive to the various issues that affect efficiency.

Steers us away from simplistic comparisons of different methods that can be used on the same problem.

#### Demonstration!

 Download and run the python script ShowBenchmarking.py

#### TODO:

- Run the above script for the following values of "n" (line 55)
  - 1000, 10000, 100000, 500000, 1000000, 10000000
- Of the three search methods—binary search, linear search with for loop, linear search with while—which seems the fastest?
- Which is faster of the following two—linear search with for loop, linear search with while loop?