

**Theorem 3:** *Leveraging multiple subgraphs enhances entropy and improves label discrimination. Specifically,*

$$P_{\text{multi}}^{(t)} \geq P_{\text{single}}^{(t)}, \quad H_{\text{multi}}^{(t)} \geq H_{\text{single}}^{(t)}, \quad (1)$$

where  $P_{\text{single}}^{(t)}$  and  $H_{\text{single}}^{(t)}$  represent the probability and entropy of label conflicts for a single subgraph, while  $P_{\text{multi}}^{(t)}$  and  $H_{\text{multi}}^{(t)}$  represent the corresponding quantities for multiple subgraphs.

## Proof

The probability of label conflict for a single subgraph at the  $t$ -th update is:

$$P_{\text{single}}^{(t)} = \frac{\text{Conflicting pairs}_{\text{single}}^{(t)}}{\binom{n}{2}}, \quad (2)$$

where  $\text{Conflicting pairs}_{\text{single}}^{(t)}$  is the number of conflicting pairs in a single subgraph.

Let there be  $k$  subgraphs. The aggregated label conflict probability across multiple subgraphs is:

$$P_{\text{multi}}^{(t)} = \frac{\text{Conflicting pairs}_{\text{multi}}^{(t)}}{\binom{n}{2}}, \quad (3)$$

where  $\text{Conflicting pairs}_{\text{multi}}^{(t)}$  accounts for the combined label diversity across all subgraphs.

It is easy to find that multiple subgraphs increase label diversity through diverse sampling strategies, reducing the chance of label conflicts:

$$P_{\text{multi}}^{(t)} \geq P_{\text{single}}^{(t)}. \quad (4)$$

The entropy for a single subgraph is given by:

$$H_{\text{single}}^{(t)} = - \sum_{x \in \mathcal{H}} P_{\text{single}}^{(t)}(x) \log P_{\text{single}}^{(t)}(x), \quad (5)$$

where  $P_{\text{single}}^{(t)}(x)$  is the probability distribution of labels in the single subgraph.

The entropy for multiple subgraphs combines the diversity across  $k$  subgraphs are defined as follows:

$$H_{\text{multi}}^{(t)} = - \sum_{x \in \mathcal{H}} P_{\text{multi}}^{(t)}(x) \log P_{\text{multi}}^{(t)}(x), \quad (6)$$

where  $P_{\text{multi}}^{(t)}(x)$  represents the aggregated label distribution across all subgraphs.

As multiple subgraphs reduce label conflicts and increase diversity, the entropy satisfies:

$$H_{\text{multi}}^{(t)} > H_{\text{single}}^{(t)}. \quad (7)$$

From (4) and (7), we establish:

$$P_{\text{multi}}^{(t)} \geq P_{\text{single}}^{(t)}, \quad H_{\text{multi}}^{(t)} > H_{\text{single}}^{(t)}. \quad (8)$$