

# CSCE 451/851 Homework 3

Assigned: Apr 5, 2020

Due: Apr 19, 2020 23:59:59

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100 points total

## Problem 1 (40 points)

Consider the following set of processes:

| Process Name | Arrival Time | Processing Time |
|--------------|--------------|-----------------|
| A            | 0            | 3               |
| B            | 1            | 5               |
| C            | 3            | 2               |
| D            | 9            | 5               |
| E            | 12           | 5               |

Perform FCFS, RR ( $q = 4$ ), SPN, SRT, on them and get the Finish Time and Turnaround Time for each process. For reference:

- FCFS: First Come First Served
- RR: Round Robin
- SPN: Shortest Process Next
- SRT: Shortest Remaining Time Next

|      | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
|------|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|----|----|----|----|----|
| FCFS |   |   |   |   |   |   |   |   |   |    |    |    |    |    |    |    |    |    |    |    |
| RR   |   |   |   |   |   |   |   |   |   |    |    |    |    |    |    |    |    |    |    |    |
| SPN  |   |   |   |   |   |   |   |   |   |    |    |    |    |    |    |    |    |    |    |    |
| SRT  |   |   |   |   |   |   |   |   |   |    |    |    |    |    |    |    |    |    |    |    |

## Problem 2 (20 points)

Consider a computer with two processes,  $H$ , with high priority, and  $L$ , with low priority. The scheduling rules are such that  $H$  is run whenever it is in ready state. At a certain moment, with  $L$  in its critical region,  $H$  becomes ready to run (e.g., an I/O operation completes).  $H$  now begins busy waiting, but since  $L$  is never scheduled with  $H$  is running,  $L$  never gets the chance to leave its critical region, so  $H$  loops forever. This situation is sometimes referred to as the priority inversion problem.

If instead of priority scheduling, we use round-robin scheduling, will we encounter the same problem? Please explain your answer in detail.

## Problem 3 (40 points)

Based on measurements, we know for a certain system the average process runs for a time  $X$  before blocking on I/O. It takes a time  $Y$  to do a process switch, which is effectively wasted (overhead). For round-robin scheduling with quantum  $Q$ , give a formula for the CPU efficiency for each of the following:

- $Q = \infty$

- $Q > X$
- $Y < Q < X$
- $Q = Y$
- $Q \rightarrow 0, Q \neq 0$

*NOTE: The CPU efficiency is the useful CPU time divided by the total CPU time.*