# 15-442/15-642: Machine Learning Systems

# **Memory Optimizations**

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Tianqi Chen
Carnegie Mellon University

### Outline

Activation Checkpointing and Rematerialization

**Mixed Precision** 

Fully Sharded Data Parallelism

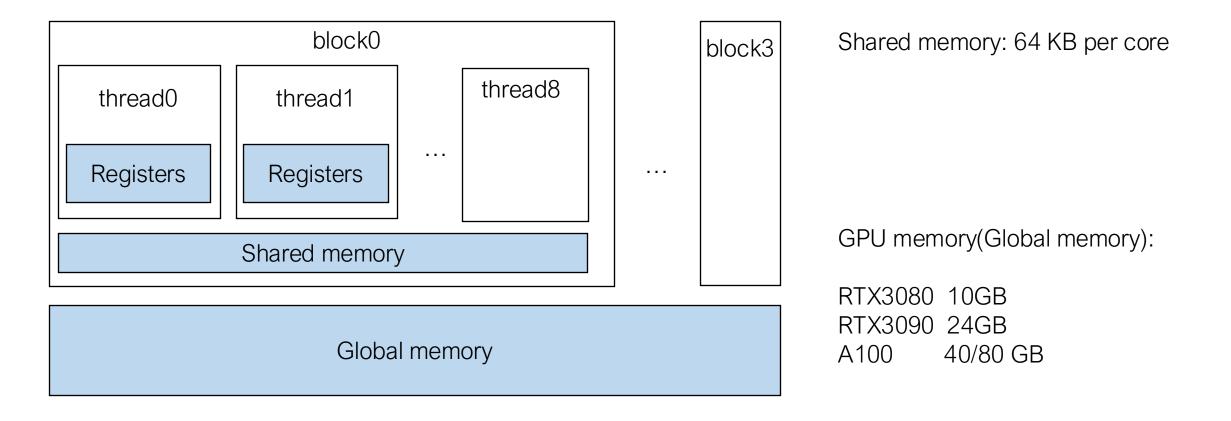
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Activation Checkpointing and Rematerialization

Mixed Precision

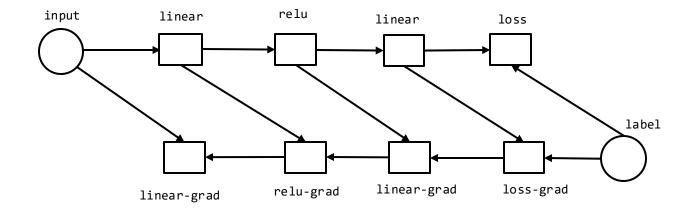
Fully Sharded Data Parallelism

## Recap: GPU memory hierarchy



## Sources of memory consumption

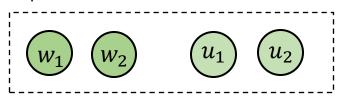
A simplified view of a typical computational graph for training, weights are omitted and implied in the grad steps.



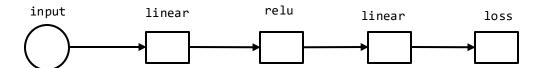
Sources of memory consumption

- Model weights
- Optimizer states
- Intermediate activation values

#### Optimizer states

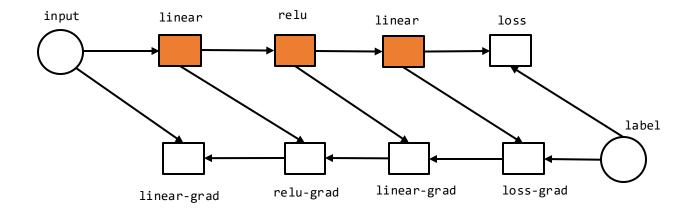


## Techniques for Memory Saving, Inference Only



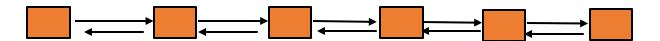
We only need O(1) memory for computing the final output of a N layer deep network by cycling through two buffers

## **Activation Memory Cost for Training**

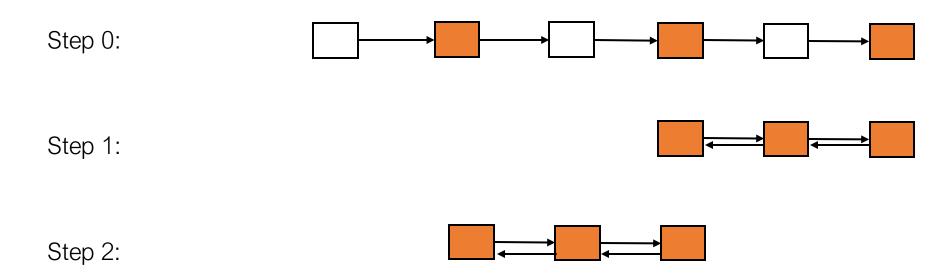


Because the need to keep intermediate value around (checkpoint) for the gradient steps. Training a N-layer neural network would require O(N) memory.

We will use the following simplified view to combine gradient and forward computation



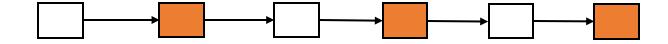
## Checkpointing Techniques in AD



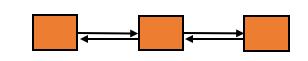
- Only checkpoint colored nodes (step 0)
- Recompute the missing intermediate nodes in small segments (step 1, 2)

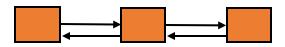
## **Sublinear Memory Cost**

Forward computation



Gradient per segment with re-computation





For a *N* layer neural network, if we checkpoint every *K* layers

$$Memory\ cost = O\left(\frac{N}{K}\right) + O(K) \qquad \text{Pick}\ K = \sqrt{N}$$
 Checkpoint cost   
 Re-computation cost

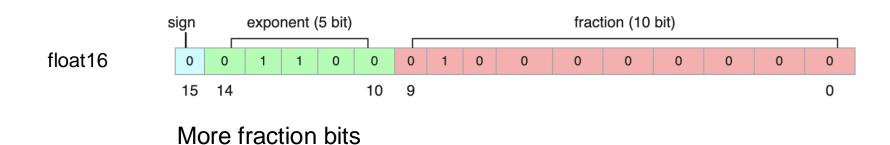
### Outline

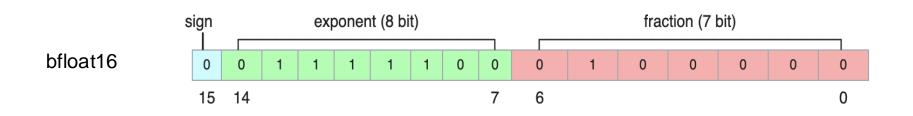
Rematerialization

**Mixed Precision** 

Fully Sharded Data Parallelism

## **16bit Floating Points**



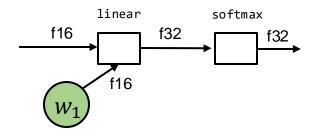


Less easy to overflow

source: wikipedia

### **Mixed Precision**

- Some layers are more sensitive to dynamic range
- Common issues: aggregation of a lot of entries
- Mixed precision: different input/output/accumulation types



### Outline

Activation Checkpointing and Rematerialization

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## Recap: AllReduce Abstraction

Interface result = allreduce(float buffer[size])

#### Running Example

#### Worker 0

#### comm = communicator.create()

a = [1, 2, 3]

b = comm.allreduce(a, op=sum)

#### Worker 1

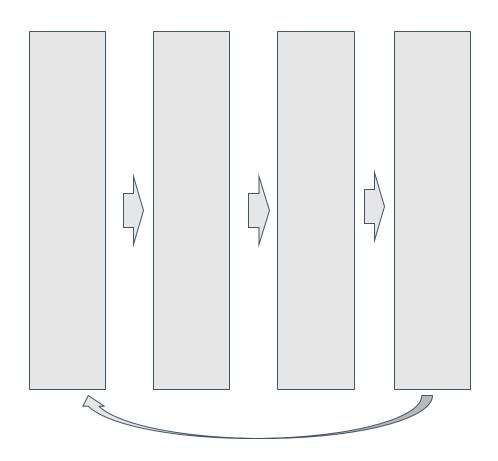
comm = communicator.create()

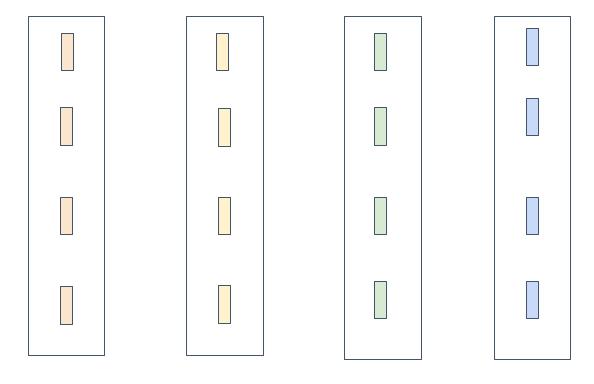
a = [1, 0, 1]

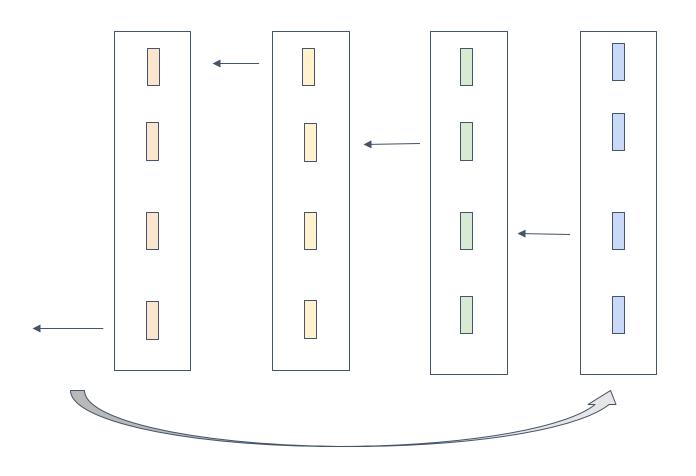
b = comm.allreduce(a, op=sum)

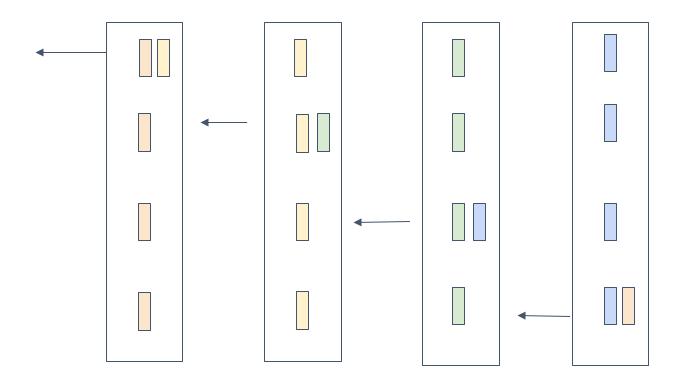
assert b == [2, 2, 4] assert b == [2, 2, 4]

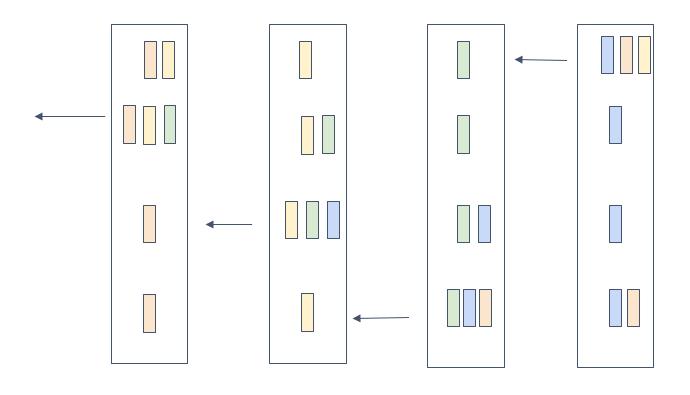
- Form a logical ring between nodes
- Streaming aggregation

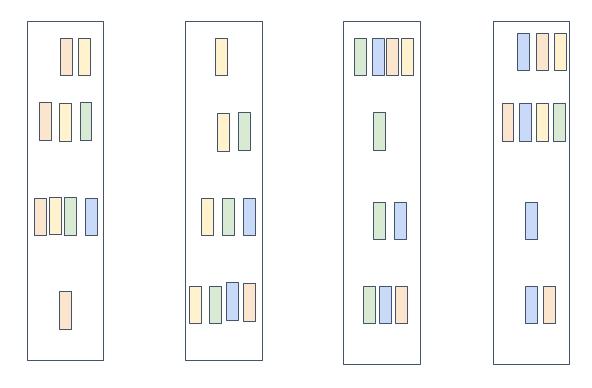












Each node have correctly reduced result of one segment! This is called *reduce\_scatter* 

#### Reduce Scatter Abstraction

Interface result = reduce\_scatter(float buffer[size])

#### Running Example

#### Worker 0

## comm = communicator.create()

$$a = [1, 2, 3, 4]$$

b = comm.allreduce(a, op=sum)

#### Worker 1

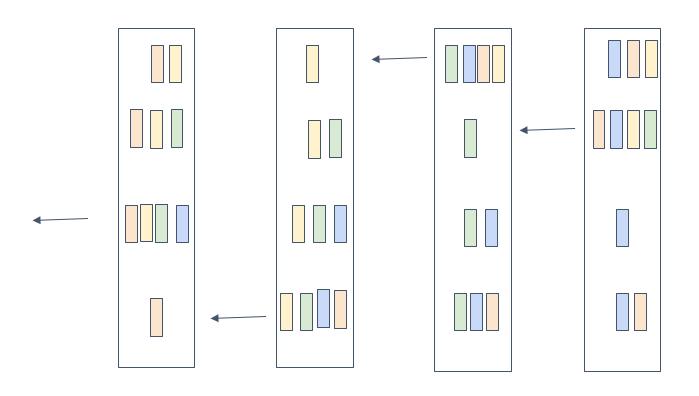
comm = communicator.create()

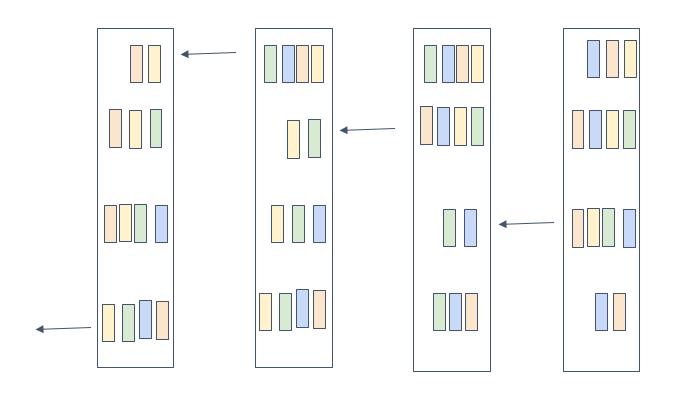
a = [1, 0, 1, 1]

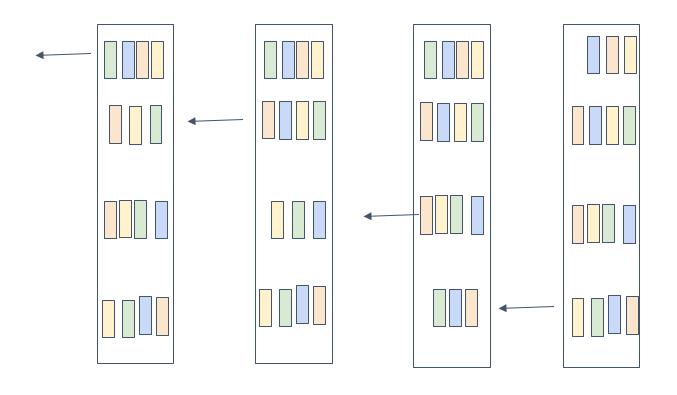
b = comm.allreduce(a, op=sum)

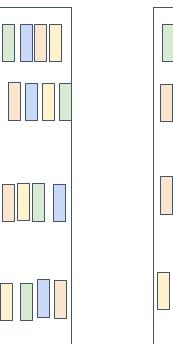
assert b == [2, 2]

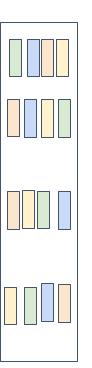
assert b == [4, 5]

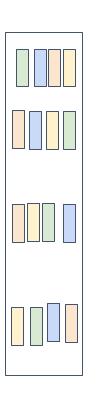


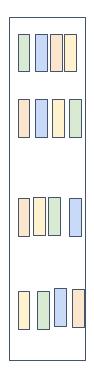












Question: What is Time Complexity of Ring based Reduction

## Allgather abstraction

Interface result = allgather(float buffer[size])

#### Running Example

#### Worker 0

### comm = communicator.create()

$$a = [1, 2]$$

#### Worker 1

comm = communicator.create()

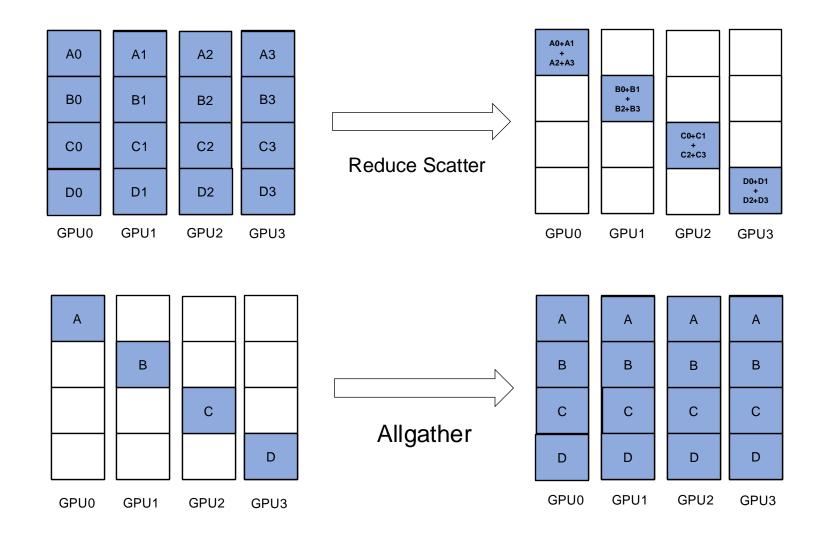
a = [3, 4]

b = comm.allgather(a)

assert b == [1, 2, 3, 4]

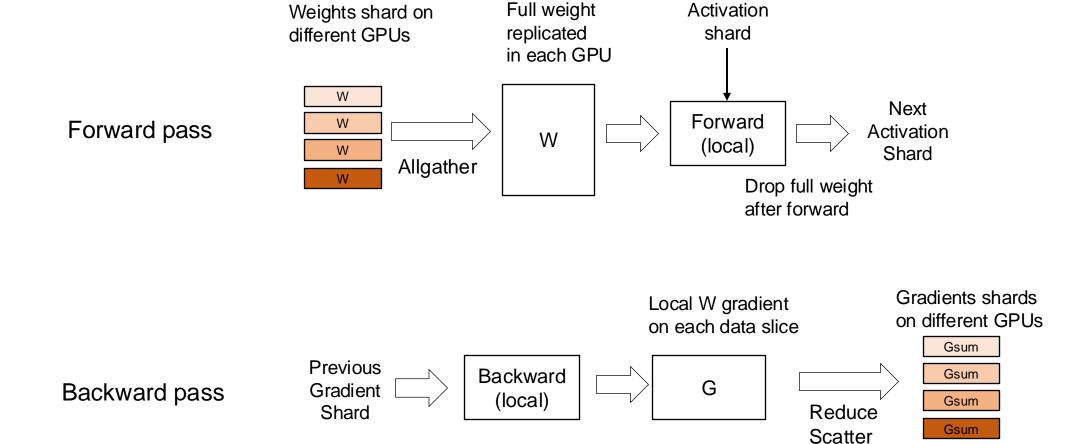
assert b == [1, 2, 3, 4]

### **Overall Relations**



Combine both we get Allreduce

## FSDP: Fully Sharded Data Parallel



Core idea is equivalent to ZeRO3