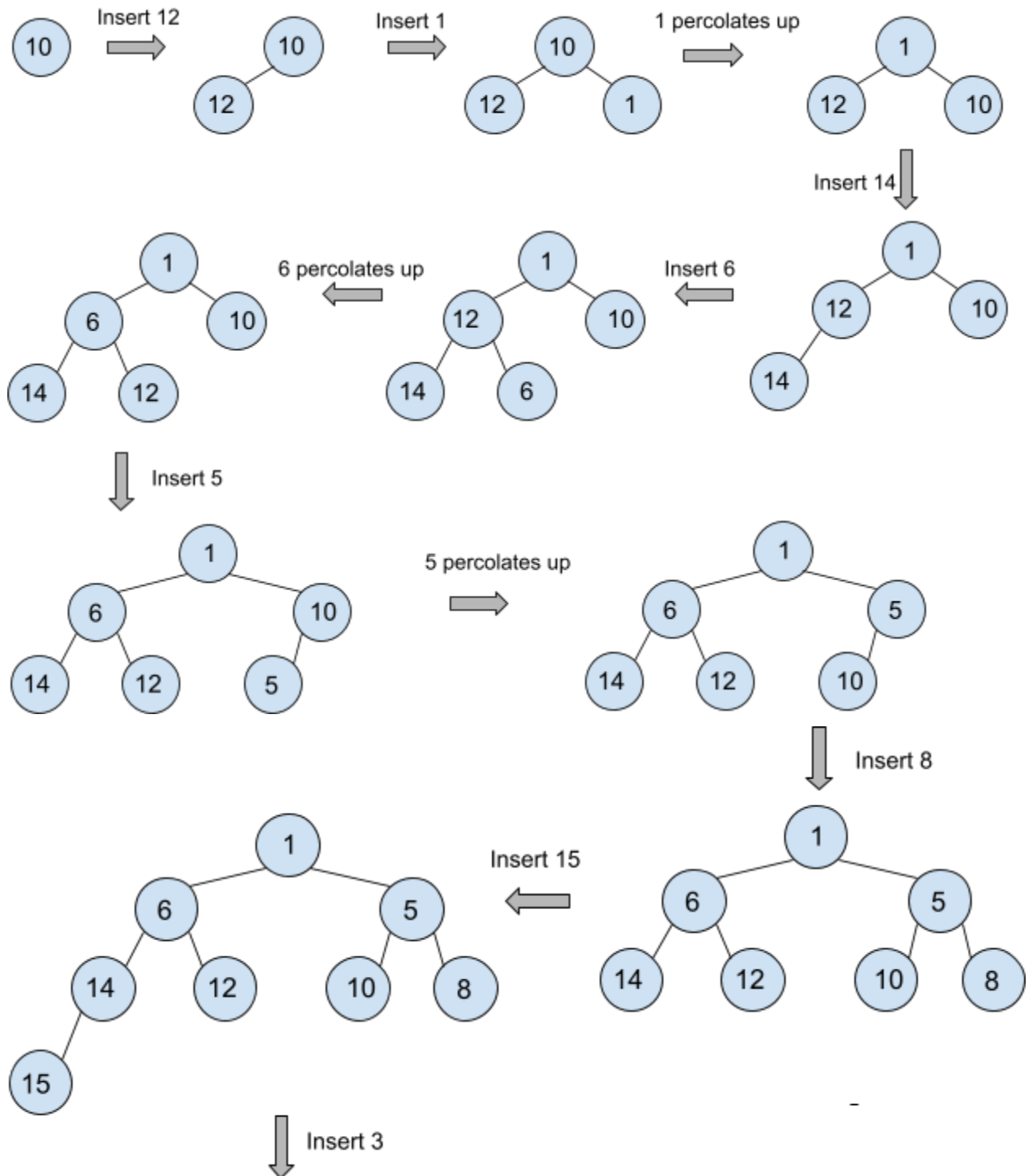
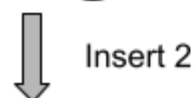
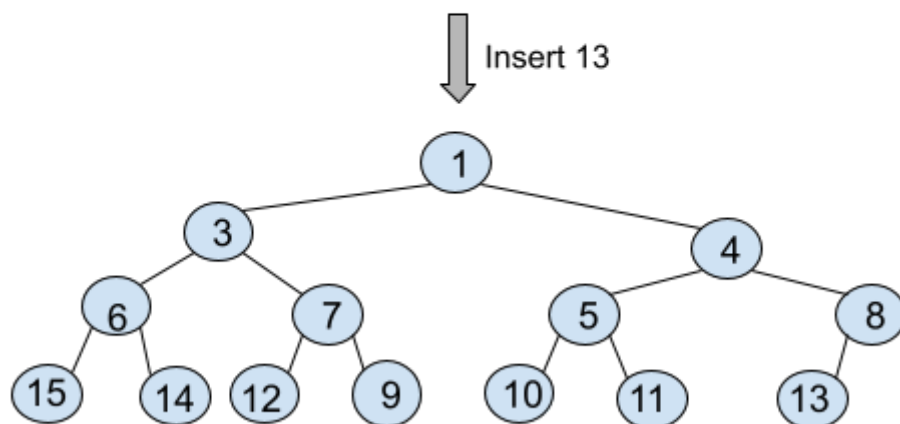
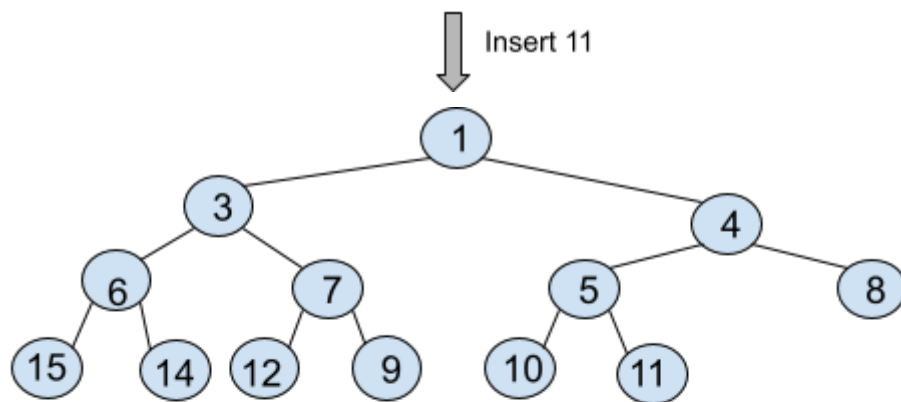
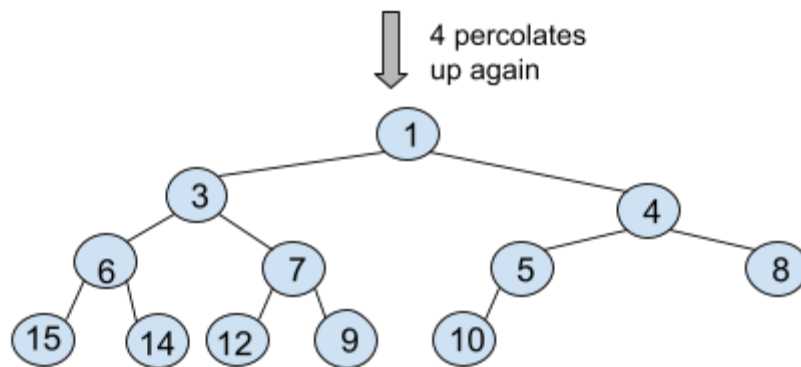
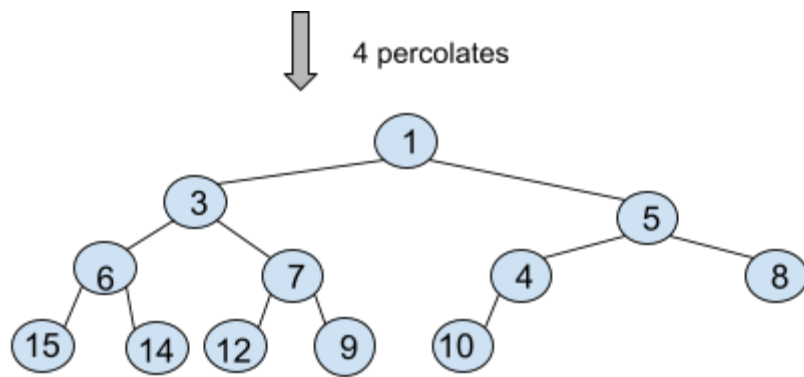
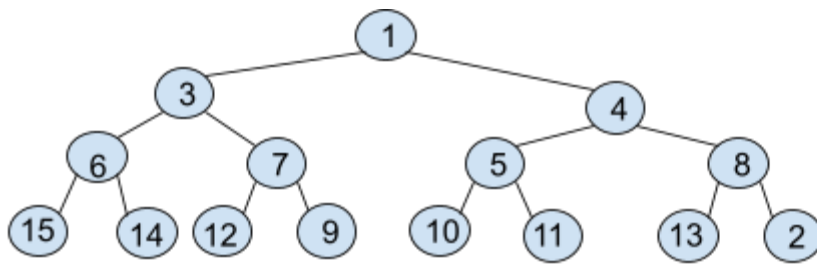


Problem 2: (45 pts)

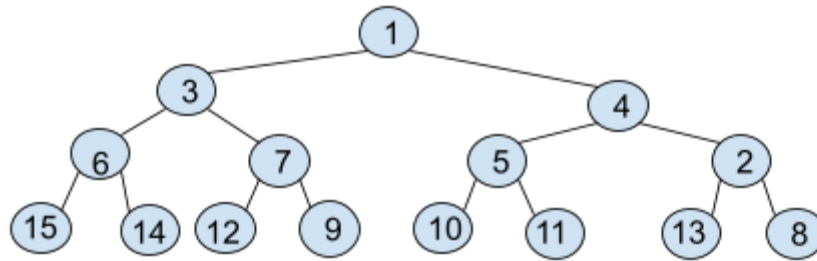
A) Show the result of inserting 10, 12, 1, 14, 6, 5, 8, 15, 3, 9, 7, 4, 11, 13, and 2 at one time and in the given order, into an initially empty **binary min heap**.



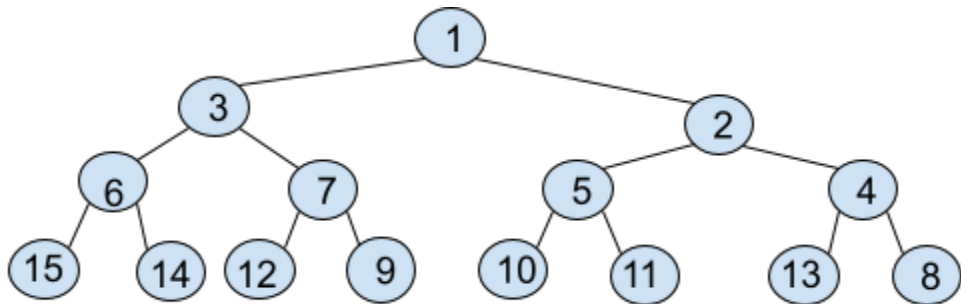




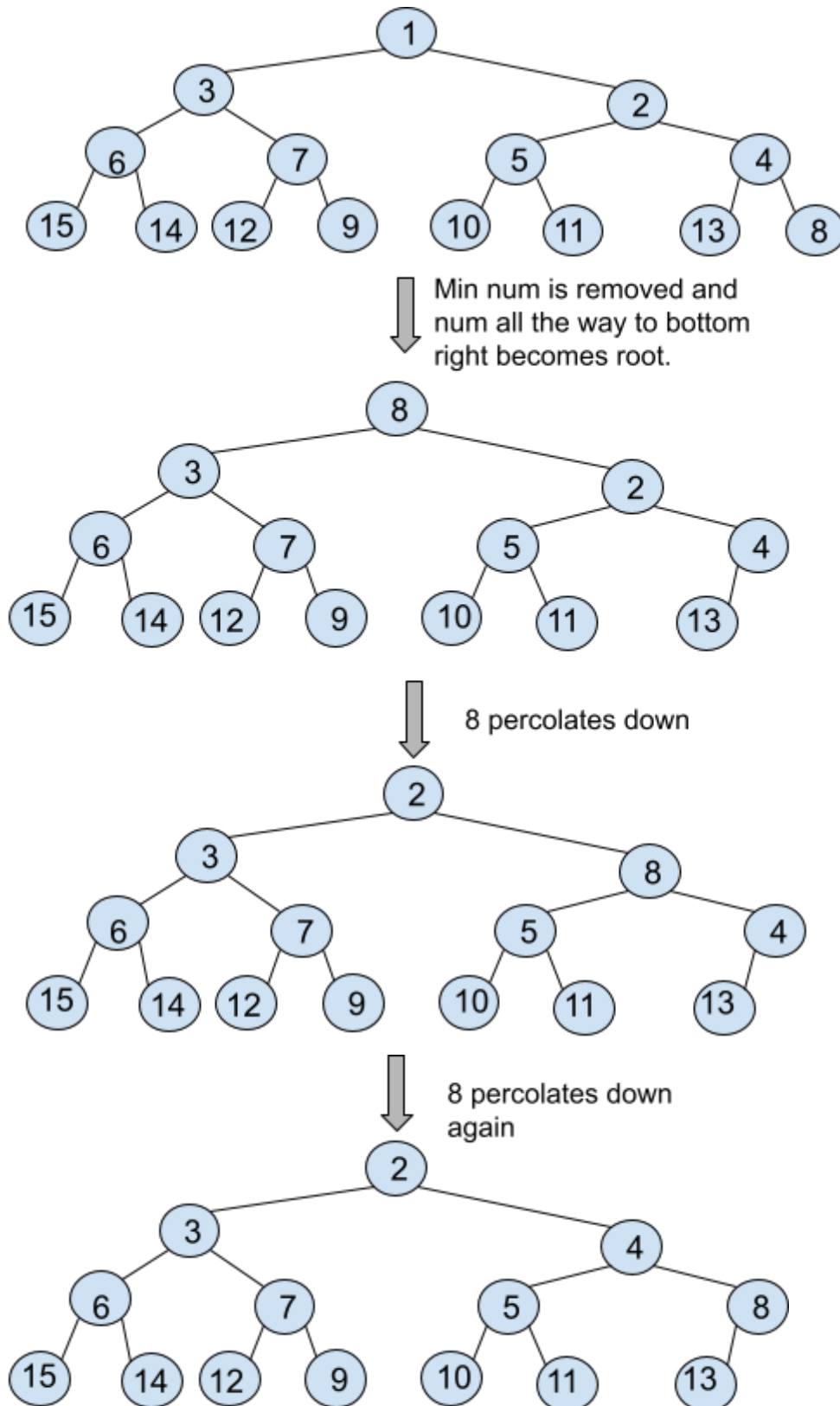
2 percolates up

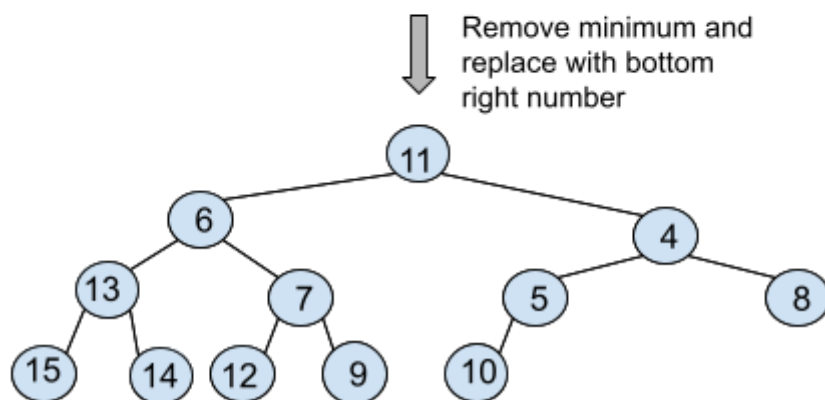
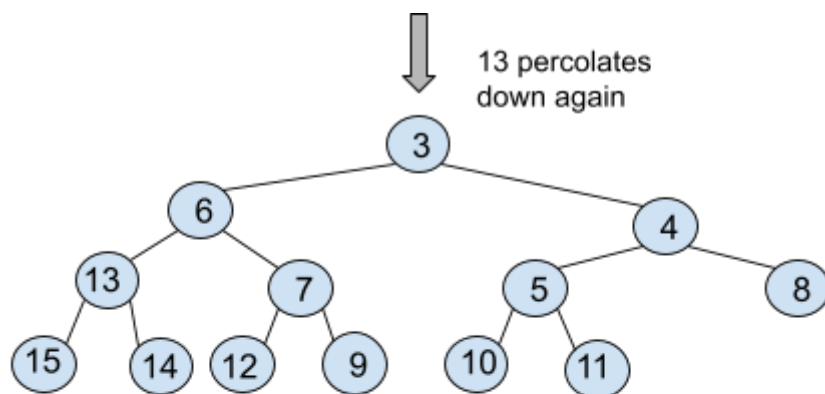
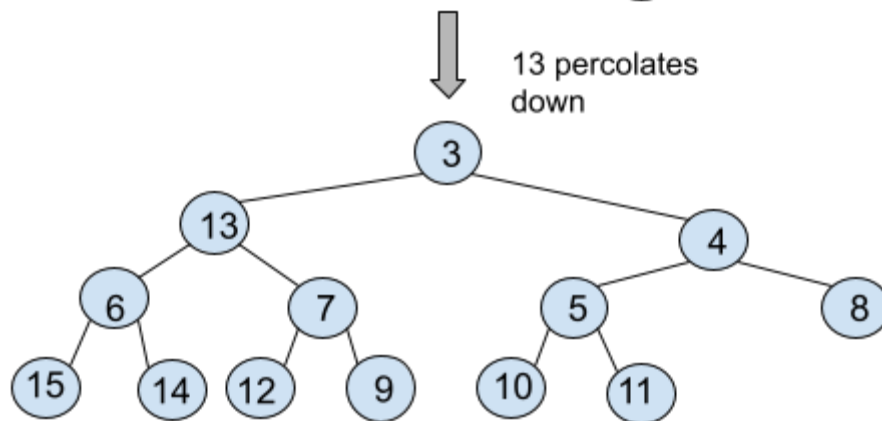
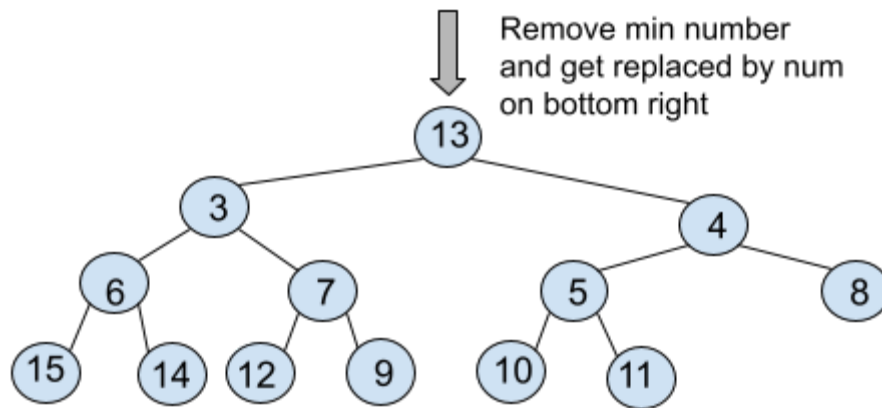


2 percolates up again

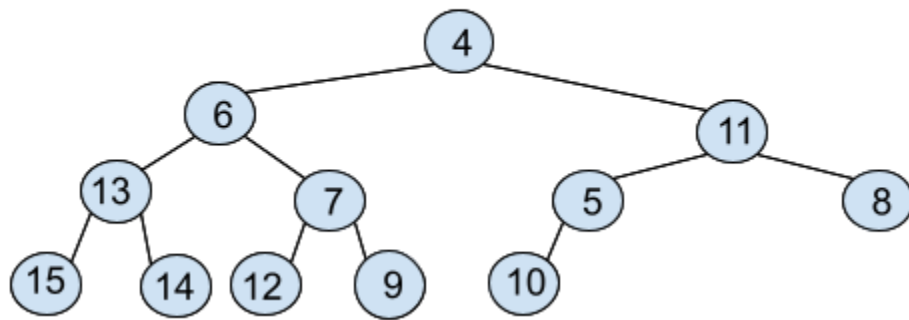


B) Show the result of performing three **deleteMin** operations in the heap of the previous **binary min heap**.

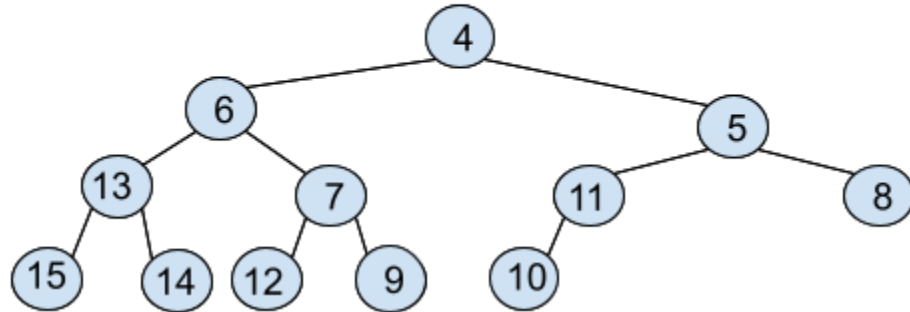




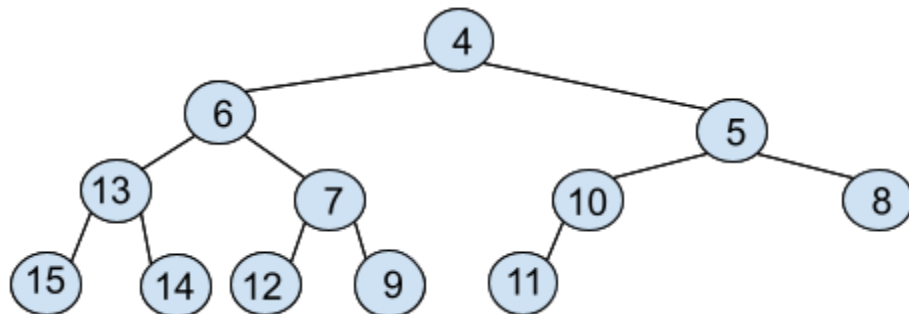
↓ 11 percolates down
and 4 percolates up



↓ 11 percolates down,
and 5 percolates up



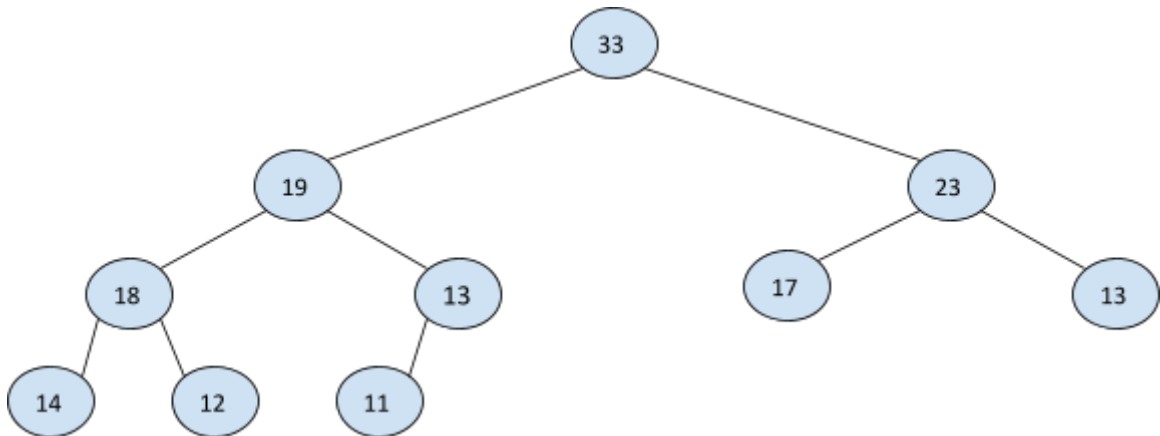
↓ 11 percolates down,
and 10 percolates up



D) What is the running time complexity of your replace key method?

- Worst case scenario is $O(n)$ because you're looking for a key using a for loop iterating 1 by one and you don't know where the key is.

E) Consider an initially empty **max-heap**, where the following keys are to be inserted one at a time: 11, 19, 23, 12, 13, 14, 18, and 33. Draw a **tree** that results after building this **max-heap**.



F) Is it possible to find the maximum in a min heap in $O(\log(n))$ time? Justify.

- It is not possible to find the maximum in a minimum heap at $O(\log(n))$ time. Since the maximum number is stored at the leaf node, we would need to check each leaf node in the worst case scenario making it an $O(n)$ time complexity.