

# System Programming

## *13. Network*

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# Network Protocols for Communications

## ■ Protocol

- A pre-defined communication step for **error-free** communications on an erroneous data link.
- Usually a protocol is implemented in a multiple-layered architecture. Why?
  - too big,
  - various level of abstraction,
  - various level of service,
  - variable media & communications types



# Network Protocols for Communications

- OSI 7 layers (OSI reference model)
  - **Physical layer** : electrical signaling system, (wired, wireless)
  - **Data link layer** : error free communications between adjacent nodes;
    - MAC layer (Medium Access Control) : multiple shared accesses to a link/bus (ethernet, RF)
    - Point-to-point (a private link between nodes)
  - **Network layer** : Routing (which must be the next node to deliver the received packet to the final target?), **IP in the Internet Protocol**
  - **Transport layer** : host APIs for end-to-end communications
    - **TCP/UDP in the Internet Protocol**
  - **Session layer** : session management, error recovery
  - **Presentation layer** : Encryption/decryption, network standard data format, other libraries/utilities (address translation, etc.)
  - **Application layer** : ftp, email, rlogin, telnet, web server/browser (http), etc



# Network Protocols for Communications

- A message from a user process can be split into multiple segments in each protocol layer. (**Fragmentation**)
- Messages from a user can be merged into a segment in each protocol layer.
- Each layer attaches a layer's **packet header** to the segment. → a packet frame
  - *TCP segment : TCP Header + Data segment*
  - *IP : IP header + TCP segment + CRC checksum*
  - ...
- Each layer has its **MTU (Maximum Transfer Unit)** : max. sized data segment in each layer.



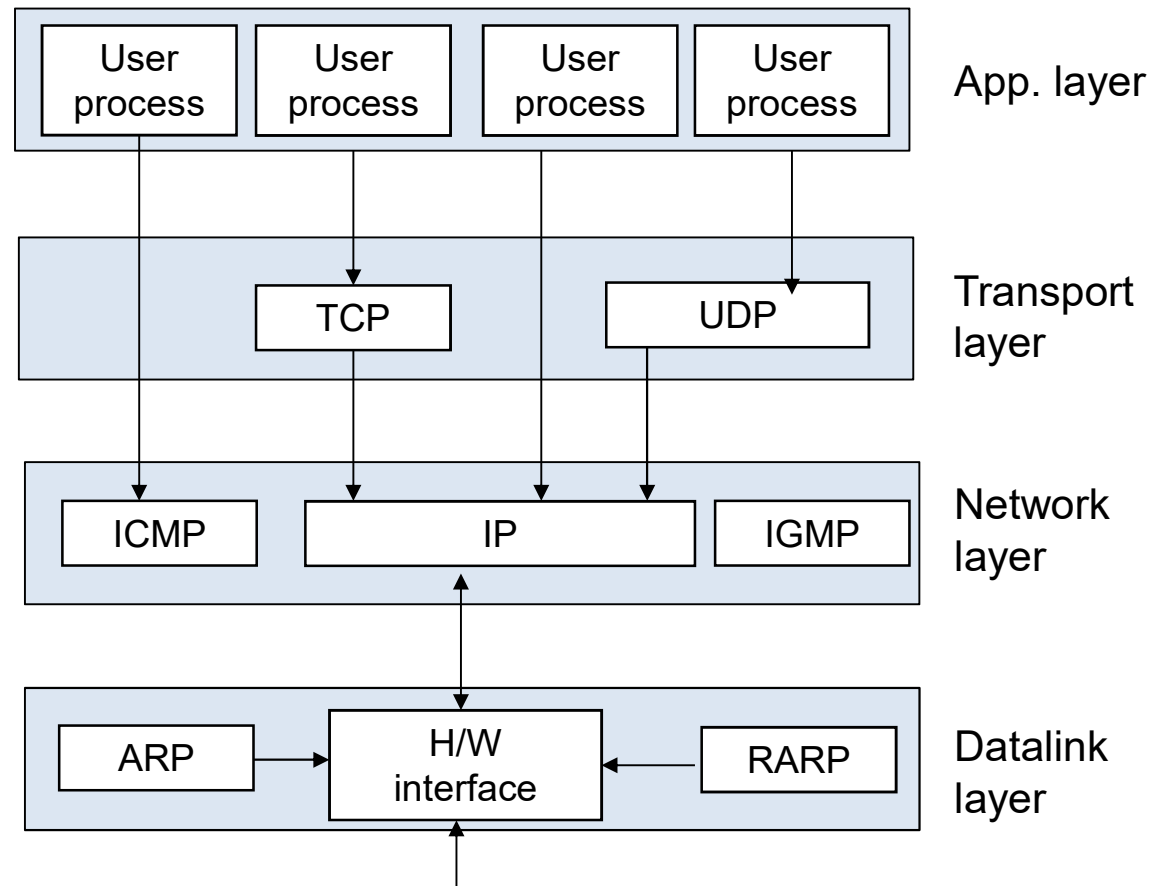
# Protocol Examples

<b>Application layer</b>	<i>Telnet, ftp, SNMP (Simple Network Management Protocol), etc.</i>
<b>Transport layer</b>	<i>TCP (Transmission Control Protocol), UDP (User Datagram Protocol), etc.</i>
<b>Network layer</b>	<i>IP, ICMP, IGMP, etc.</i>
<b>Data Link layer</b>	<i>Network device drivers, Interface/controller cards, etc.</i>

- ICMP (Internet Control Message Protocol)
  - A protocol used by the “ping” service
- IGMP (Internet Group Message Protocol)
  - A router’s protocol for multicasting



# Protocol Usages



**ARP** (Address Resolution Protocol): IP to a physical MAC address

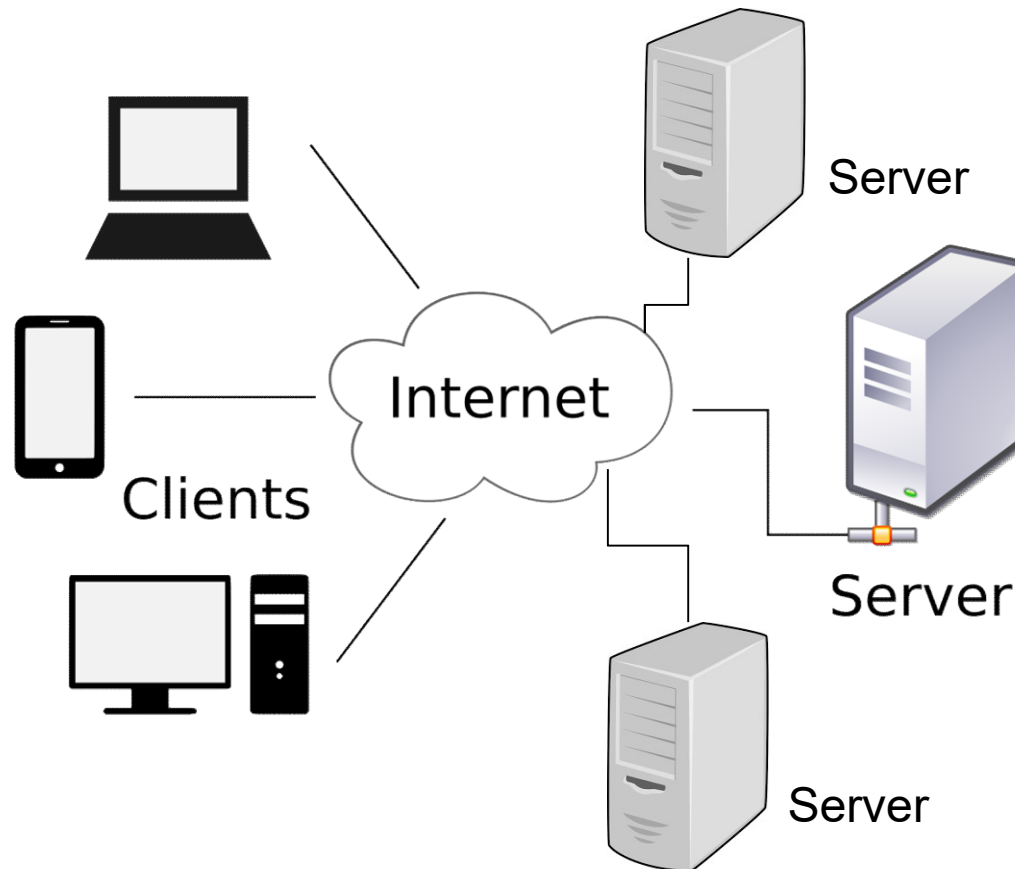
**RARP** (Reverse Address Resolution Protocol) : Physical MAC address to an IP



## ***App. Layer Services on Transport Protocols***

<i><b>Transport Protocols</b></i>	<i><b>Application Services</b></i>
<b>TCP</b>	<b>FTP (File Transfer Protocol)</b> <b>TELNET</b> <b>SMTP (Simple Mail Transfer Protocol)</b> <b>HTTP (Hyper Text Transfer Protocol) of WWW</b>
<b>UDP</b>	<b>SNMP (Simple Network Management Protocol)</b> <b>TFTP (Trivial FTP)</b>

# Client/Server Model (1)



## ■ *Types of network application services*

- **Client/server** : a client program  $\leftrightarrow$  a server program
- **Web-based** : web-browser(client)  $\leftrightarrow$  web-server(server)
- **Peer-to-peer(P2P)** : a node can be a server or can be a client



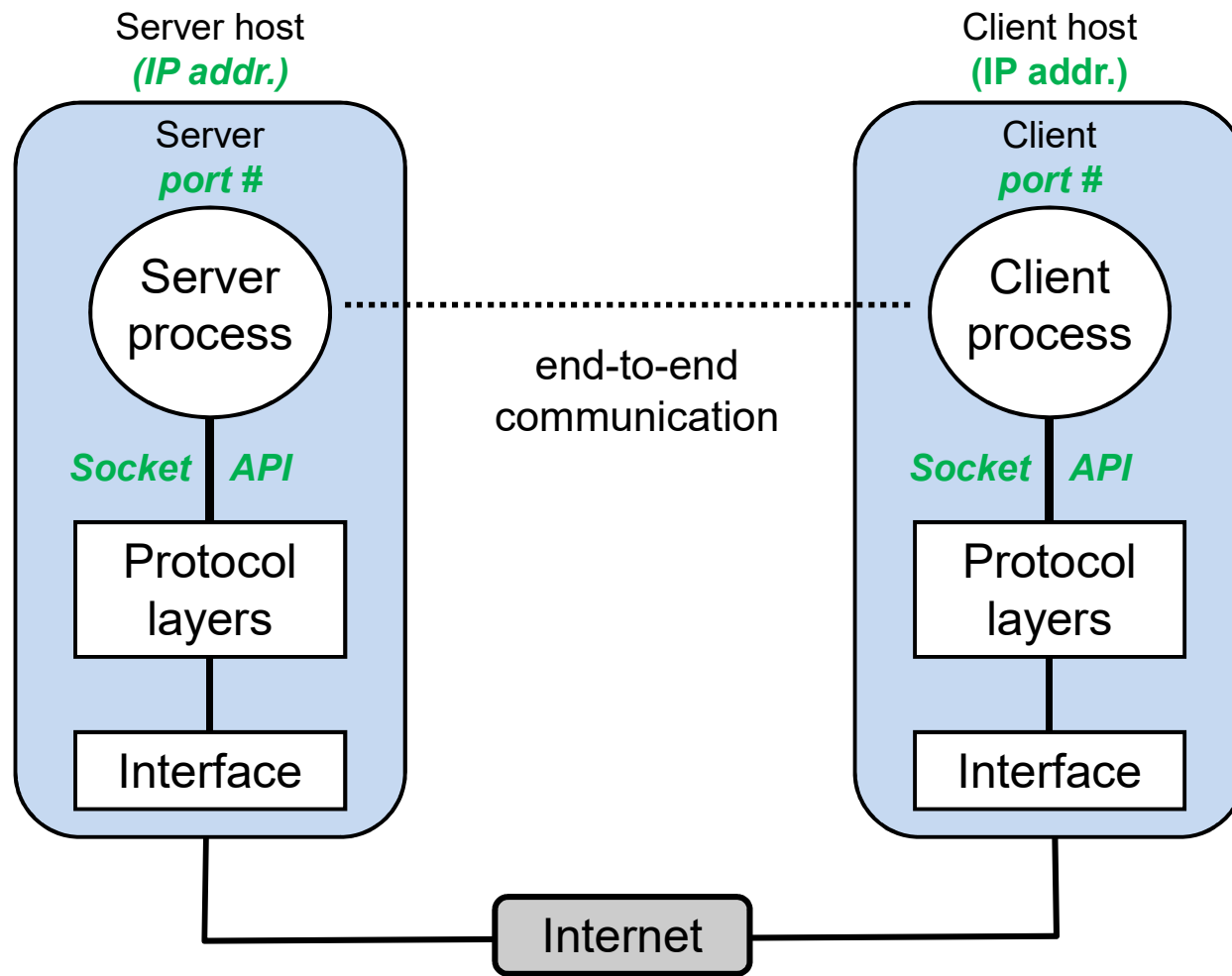


# Client/Server Model (2)

- Both of client & server processes (programs) are necessary.
  - Client & server processes use a set of protocol APIs.
  - A server demon is always waiting for connections of clients.
    - Upon a connection, the server spawns a child (or a thread) to serve the client.
  - A client tries to connect to a server.
    - Target (server/client) addresses are designated by the (*IP address, Port #*) couples.
    - Packets are delivered to a node with the IP address.
    - In the target node, a network protocol delivers the packet to the server/client process/thread using the port #.



# Client/Server Model (3)



# *Connection-oriented vs. Connectionless Communications*

<i><b>Types</b></i>	<i><b>Description</b></i>
<i><b>Connection-oriented</b></i>	Use the <b>TCP</b> (Transmission Control Protocol) protocol, Reliable data transfer is guaranteed, Must set up a link to each client, More clients, higher pressure on the server. <b>Think telephone communications! (Virtual Circuit)</b>
<i><b>Connectionless</b></i>	Use the <b>UDP</b> (User Datagram Protocol) protocol, Suitable for a single message transmission, No link for clients: low pressure to the server, Can be used for broadcast or multicast services. <b>Think datagram of post offices!</b>



# Server Types

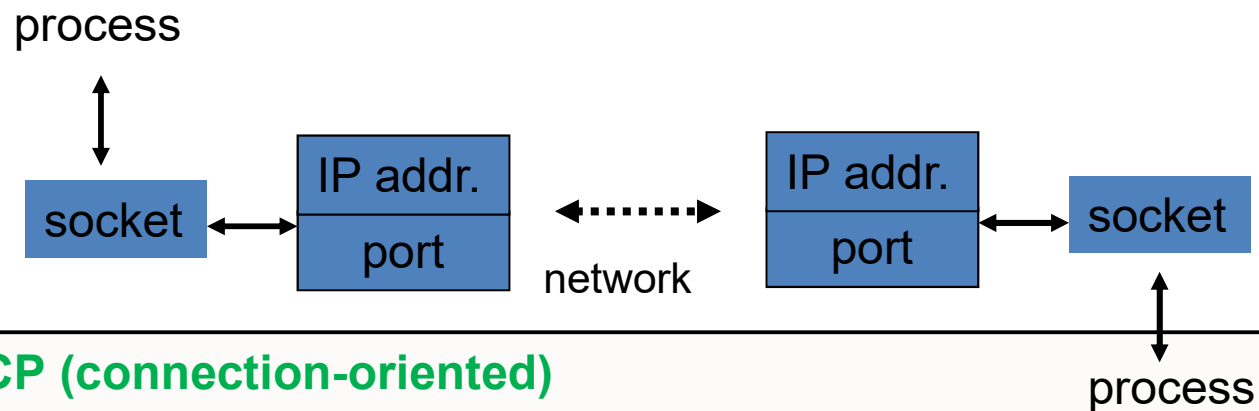
<i>Server types</i>	<i>Description</i>
Repetitive server type ( <i>single-threaded server</i> )	A single server process handles the requests of multiple clients one by one. (e.g. FIFO) Slow response.
Concurrent server type ( <i>multi-threaded server</i> )	A server process consists of several concurrent threads. For each client, a server thread/process is created at the connection time. (or pre-created at the time of server-launch) High performance. Concurrency problems (e.g. Mutex)



## Programs/APIs of Each Layer

<i>Layer</i>	<i>APIs, Programs</i>	<i>Description</i>
App. layer	<i>http, ftp, email, rsh, RPC</i>	Application services for easy use.
Transport layer	<i>Socket (Berkley), Winsock, TLI</i>	For <i>end-to-end</i> communications (message, stream). Supports <i>TCP/IP, UDP/IP,</i>
Device driver layer	Packet Driver, NDIS(Network Driver Interface Spec, window), ODI (Open DataLink Interface)	Handle MAC frame transmission on a LAN. Support various MAC protocols. Error control, flow control.

# Socket Communications (end-to-end: TCP, UDP)



## ■ TCP (connection-oriented)

- Binding : assign (IP addresses, port #) to a socket
- Every transmission of packets use the same link
- Stream I/O, reliable, flow control, error control.

## ■ UDP (connectionless)

- For each transmission of a message, IP addr. & port # of the target are necessary.
- Useful for one-time small message transmission.
- Message-based, unreliable (a message can be lost), order of message delivery can be reversed.
- Message size must be smaller than the UDP packet size
- No flow control, restricted error control, low overhead.

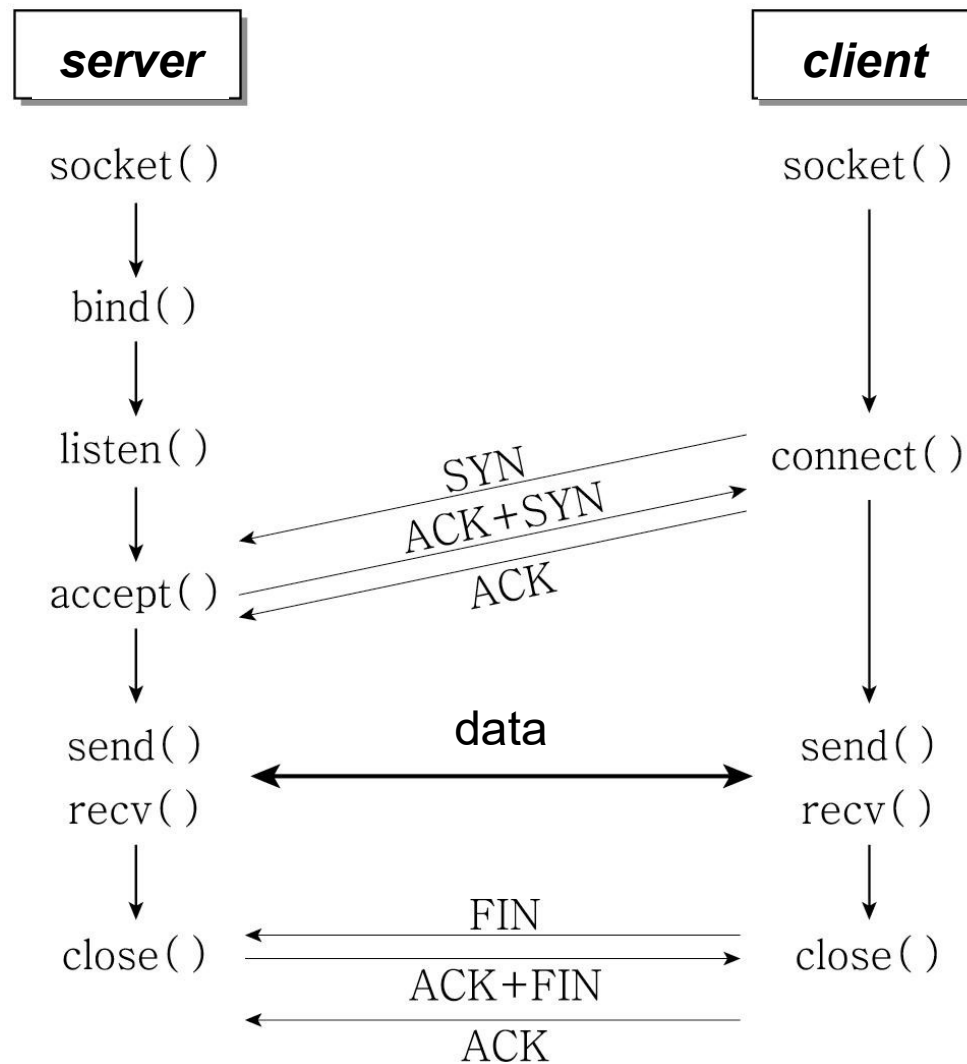


# Ports

- A network application process uses a port in the local host.
  - So **(IP addr. + port #)** can designate the **peer** process of the remote host.
- Port number usage (0 ~ 65535), **IANA** allocation
  - **Well-known ports** (0 ~ 1023) : already assigned to exiting network services. (ex: 23: DNS)
    - `$cat /etc/services | grep tftp → tftp 69/tcp tftp 69/udp`
  - **Registered ports** (1024 ~ 49151), **Dynamic ports** (49152 ~ 65535)
    - All ports can be used by a user, but a registered port can be registered to IANA.
    - Dynamic ports: free use by users.

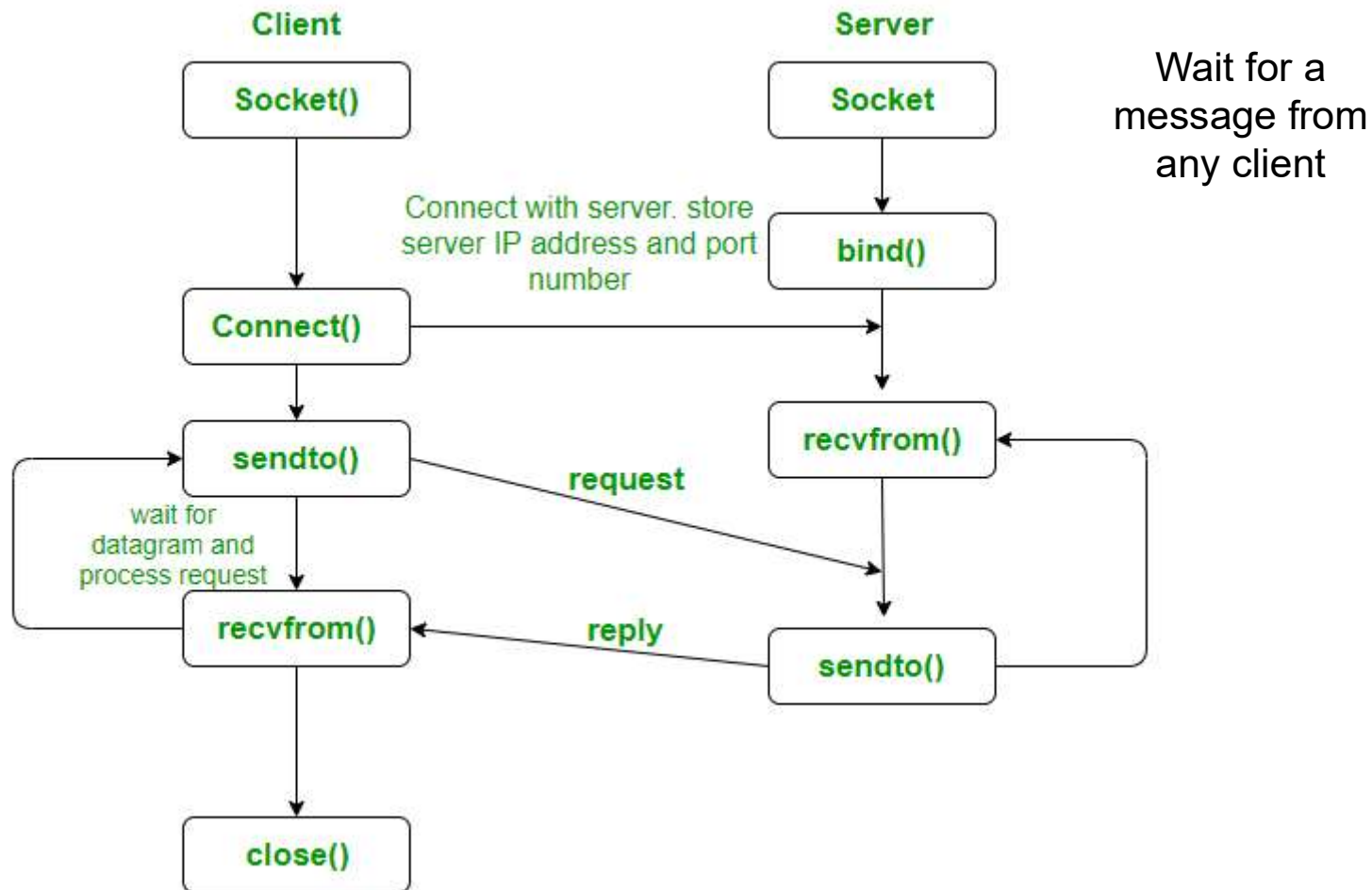


# Connection-oriented Communication (TCP)





# Connectionless Communication (UDP)



# Socket Address Structure

- socket-addr structure (a representative structure for all types of networks)

```
#include <sys/socket.h>
```

```
struct sockaddr {  
    u_char sa_len;    /* address structure length */  
    u_char sa_family /* address type */  
    char sa_data[14]; /* 14 byte-address */  
};
```



# Socket Address Structure of the Internet

```
#include <netinet/in.h>
#include <sys/types.h>
// sockaddr internet (practical address structure used by internet TCP/UDP)
// when use this structure in relevant syscalls, address casting to sockaddr is needed

struct sockaddr_in {
    u_char sin_len;           // address structure length
    u_char sin_family;        // address type
    u_short sin_port;         // 16 bit port #
    struct in_addr sin_addr // 32 bit IP address
    char sin_zero[8];         // not used, for the further use: must set to be all zeros
};

struct in_addr
{
    u_long    s_addr;         // 32 bit IP address
}
```



# Socket Address Structure

## ■ **sin\_family:**

- **AF\_INET:** internet IP address
- **AF\_UNIX:** UNIX or used for local communications
  - The server & client are all in the local host,
  - But use the whole protocol stack.
  - Usually be used for testing
- **AF\_NS:** XEROX network address



# socket()

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int socket (int domain, int type, int protocol);
```

input

- **domain** : address types

- *PF\_INET*: internet protocol
- *PF\_INET6*: IPv6 protocol
- *PF\_UNIX*: UNIX, local communications
- *PF\_NS*: XEROX: Xerox network address
- *PF\_IMPLINK*: IMP link layer address

- **type** : socket type

- **protocol** : protocol for use  
(packet header format)

return

- normal : socket id
- error : -1

<i>type</i>	<i>protocol</i>	<i>actual protocol</i>
<b>SOCK_DGRAM</b>	<b>IPPROTO_UDP</b>	<b>UDP</b>
<b>SOCK_STREAM</b>	<b>IPPROTO_TCP</b>	<b>TCP</b>
<b>SOCK_RAW</b> (user creates a packet header)	<b>IPPROTO_ICMP</b>	<b>ICMP</b>
<b>SOCK_RAW</b> (user creates a packet header)	<b>IPPROTO_RAW</b>	<b>RAW</b>



# bind()

- Connect my host IP addr. & port # to my socket

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int bind (int sockfd, struct sockaddr *myaddr, int addrlen);
```

input

- sockfd : socket descriptor (file descriptor)
- myaddr : address of the socket address structure
- addrlen : size of the “myaddr” structure

return

- normal : 0
- error : -1



# connect()

- A client waits for the connection to a server by sending a “connection-request message”. : telephoning
- Auto-binding to the socket will be done, no bind() is necessary!

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int connect (int sockfd, struct sockaddr *servaddr, int addrlen);
```

input:

- sockfd : socket descriptor
- myaddr : socket addr. structure that contains the IP & port of the server.
- addrlen : size of the “servaddr” structure

return:

- normal : 0, error : -1



# listen()

- A server declares the max. queue-length of client-requests.

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int listen (int socket, int queuesize);
```

input

- socket : bound socket descriptor
- queuesize : max. number of client connection requests

return

- normal : 0
- error : -1





# accept()

- A server waits for (being blocked) a connection request from a client.
- When a connection is established, returns a new socket descriptor for communication with the client.
- The old socket will be used for other client-requests for further connections.

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int accept (int sockfd, struct sockaddr *peer, int *addrlen);
```

input:

- sockfd : socket descriptor
- peer : the socket address structure that contains the IP & port of the connected client
- addrlen : size of the “peer” structure

return:

- normal : a **new** socket descriptor to communicate the client
- error : -1



# send() for TCP

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int send (int sockfd, char * buf, int bytes, int flag);
```

input

- sockfd : socket descriptor
- buf : data buffer holding the data to be sent
- bytes : size of the buffer
- flag : options
  - » MSG\_OOB: OOB(out of bound) data: used for urgent data sending
  - » MSG\_PEEK: keep the data in the buffer
  - » MSG\_DONTROUTE: ignore the usual routing

return

- normal : the size of the data that actually be sent
- error : -1



# send() for TCP

- Completion of send() // successful return
  - This means that the data message has been stored in the sender's protocol buffer successfully.
  - Does not mean that the delivery has been completed.
- The data stored in the protocol buffer will be removed when an ACK message is arrived from the receiver's protocol.
- If **ACK timeout or NACK**, the sender's protocol resends the data packet.  
(ARP: Automatic Repeat Request)
- TCP send/recv : **stream I/O**
  - A message can be split or messages can be merged at the time of arrival.
  - So if a message splitting is required, the user must handle this by buffering, using of header/length.



# recv() for TCP

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int recv (int sockfd, char * buf, int bytes, int flag);
```

input

- sockfd : socket descriptor
- buf : user's receive-buffer
- bytes : size of the buffer
- flag : options (maybe NULL)

return

- normal : the size of data actually received
- error : -1



# sendto() for UDP

- Send a message to a server using the UDP protocol
- At every sendto(), the IP addr. & port # of the receiver must be given because this is a datagram.

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int sendto (int sockfd, char * buf, int bytes, int flag, struct sockaddr *to, int addrlen);
```

input: - sockfd : socket descriptor

- buf : buffer address

- bytes : size of the buffer

- flag : options

- to : the address structure containing the receiver's IP address and port #

- addrlen : size of the "to" structure

return: - normal : the size of the data message actually transferred

- error : -1



# recvfrom() for UDP

```
#include <sys/socket.h>
```

```
#include <sys/types.h>
```

```
int recvfrom (int sockfd, char * buf, int bytes, int flag,  
              struct sockaddr *from, int *addrlen);
```

input

- sockfd : socket descriptor
- buf : buffer for data reception
- bytes : size of the buffer
- flag : options
- from : sender's address structure
- addrlen : size of the "from" structure

return

- normal: the size of the message actually received
- error : -1



# close()

```
#include <unistd.h>
```

```
int close (int sockfd);
```

input

- sockfd : socket descriptor

return

- normal : 0

- error : -1



# Client of an echo program (TCP)

client.c

```
#include <sys/socket.h>
#include <sys/types.h>
#include <netinet/in.h>
#include <arpa/inet.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

#define MAX 1024           // max KB input length
#define PORT 30000         // server port #
#define HOSTADDR "xxx.xxx.xxx.xxx" // server IP

int main (int argc, char *argv[])
{
    int sd, send_bytes, n, recv_bytes;
```





# (con't)

```
struct sockaddr_in servaddr;  
char snddata[MAX], rcvdata[MAX];  
  
bzero ((char*) &servaddr, sizeof(servaddr)); // prepare server address, port  
servaddr.sin_family = AF_INET;  
servaddr.sin_addr.s_addr = inet_addr (HOSTADDR);  
servaddr.sin_port = htons (PORT);  
  
if (( sd = socket (AF_INET, SOCK_STREAM, 0)) < 0) {  
    fprintf( stderr, "can't open socket.\n");  
    exit(1);  
}  
if (connect (sd, (struct sockaddr *) &servaddr, sizeof(servaddr)) < 0) {  
    fprintf (stderr, "can't connect to server.\n");  
    exit(1);  
} // auto binding for this client and connect request
```



# (con't)

```
while (fgets (snddata, MAX, stdin) != NULL) { // get a string from KB
    send_bytes = strlen (snddata);
    if (send (sd, snddata, send_bytes, 0) != send_bytes) { // to server
        fprintf( stderr, "can't send data.\n");
        exit(1);
    }
    recv_bytes = 0;
    while (recv_bytes < send_bytes) { // while loop for stream I/O !
        if ((n = recv (sd, rcvdata + recv_bytes, MAX, 0)) < 0) { // from server
            fprintf (stderr, "can't receive data.\n");
            exit(1);
        }
        recv_bytes += n;
    }
    rcvdata[recv_bytes] = 0; // NULL char for string
    fputs (rcvdata, stdout); // display
}
close (sd);
return 0;
}
```



# Server of an Echo Program (TCP)

server.c

```
#include <sys/socket.h>
#include <sys/types.h>
#include <netinet/in.h>
#include <arpa/inet.h>
#include <unistd.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

#define MAX 1024           // max client data length
#define PORT 30000         // server port #
#define BACKLOG 5          // queue length

int main (int argc, char *argv[])
{
    int sd, nsd, pid, bytes, cliaddrsz;
```



# (con't)

```
struct sockaddr_in cliaddr, servaddr;  
char data[MAX];  
  
if (( sd = socket (AF_INET, SOCK_STREAM, 0)) < 0) {  
    fprintf( stderr, "can't open socket.\n");  
    exit(1);  
}  
  
// to bind the server itself to the socket  
bzero ((char*) &servaddr, sizeof( servaddr));  
servaddr.sin_family = AF_INET;  
servaddr.sin_addr.s_addr = htonl (INADDR_ANY);  
servaddr.sin_port = htons (PORT);
```

- Why **INADDR\_ANY** was used instead of the IP address?  
**INADDR\_ANY** means the all of IP address of a multi-home host.



# (con't)

```
if (bind (sd, (struct sockaddr *) &servaddr, sizeof(servaddr)) < 0) {
    fprintf (stderr, "can't bind to socket.\n");
    exit(1);
} // bind itself to the socket

listen (sd, BACKLOG); // declare the client-queue length

while (1) { // a typical server waiting loop
    cliaddrsz = sizeof (cliaddr);
    if ((nsd = accept (sd, (struct sockaddr *) &cliaddr, &cliaddrsz)) < 0) {
        fprintf (stderr, "can't accept connection.\n");
        exit(1);
    } // upon return: client addr. is known and a new socket is created
    if ((pid = fork()) < 0) { // fork error, a new thread may be used!
        fprintf (stderr, "can't fork process.\n"); exit(1);
    }
}
```



# (con't)

child  
server



```
if (pid == 0) {           // the new child server for the connected client
    close (sd);           // old socket is not necessary for me
    while(1) {
        bytes = recv (nsd, data, MAX, 0); // from client
        if (bytes == 0)      // client quit
            break;
        else if (bytes < 0) { // error
            fprintf (stderr, "can't receive data.\n"); exit(1);
        }
        if (send (nsd, data, bytes, 0) != bytes) { // echo back
            fprintf (stderr, "can't send data.\n"); exit(1);
        }
    } // end while, client quits
    return 0;              // child server exit.
}

else                      //: parent
    close (nsd);           // parent: close the new socket
                          // end while: parent goes to the client waiting-loop again
}
} /* main */
```



# Run & Results

```
$ ./server&                // run the server as a background process
```

```
[1] 25345
```

```
$ ./client                // run the client
```

```
It's client-server test
```

```
It's client-server test
```

```
Linux programming
```

```
Linux programming
```

```
^C
```

```
$
```



# TCP's stream I/O

- TCP does not guarantee to receive a whole message sent.
  - Messages can be split or merged.
- When the message size is fixed and known.

- Method 1

```
while(1) {  
    len = recv (sd, data, MAX, MSG_PEEK); // do not remove from TCP buffer.  
    if (len >= desired_length)  
        break;  
}  
recv (sd, data, MAX, 0); // read a message & remove from the TCP buffer
```

- Method 2

```
len=0;  
size = sizeof (struct message); // the fixed known message size  
while(1) {  
    p = (char*)message_buf + len;  
    len += recv (sd, (void*)p, size-len, 0);  
    if ( len >= size)  
        break;  
}
```





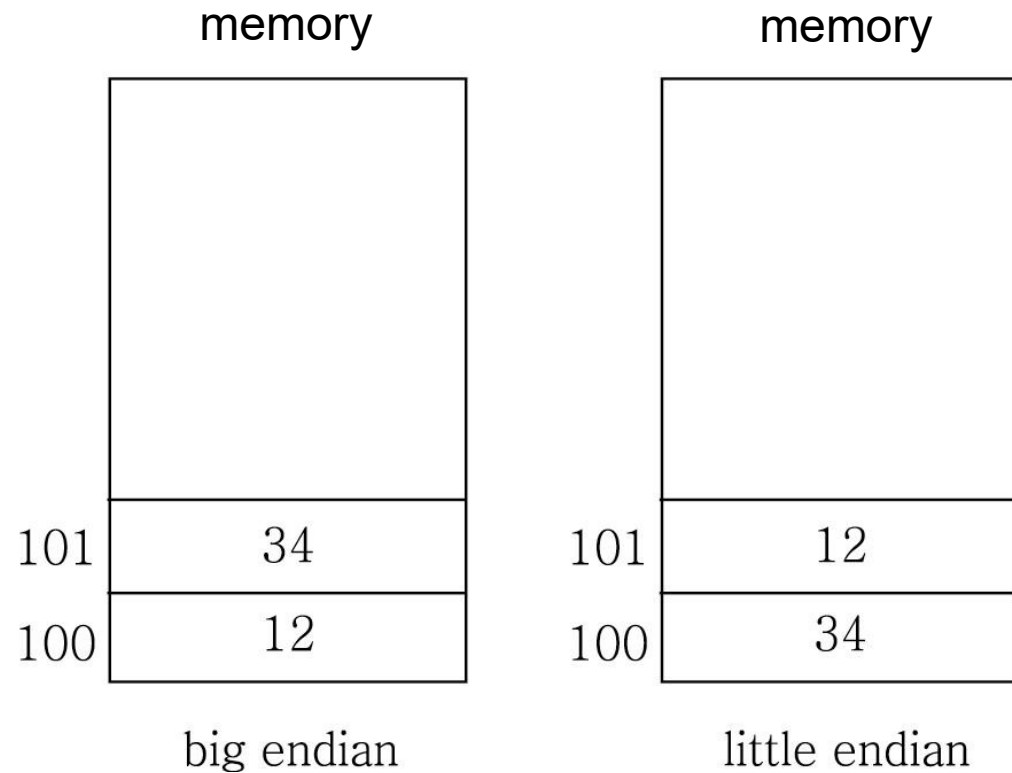
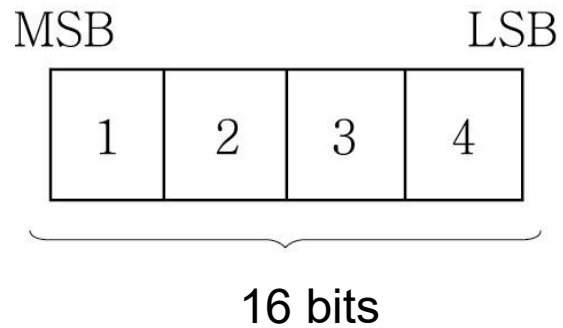
# Data Format Conversion

Functions	Description
u_long <b>htonl</b> (u_long hostlong);	Convert host long format to the network standard
u_short <b>htons</b> (u_short hostshort);	Host to network short
u_long <b>ntohl</b> (u_long netlong);	Network to host long
u_short <b>ntohs</b> (u_short netshort);	Network to host short

- The above functions are necessary because data formats are machine dependent.
- *Little endian vs. big endian*
- Different float/double format (exponent, mantissa)
  - Float/double must be handled by a user.

Functions	Description
void <b>bcopy</b> (char *src, char *dst, int bytes);	Bytes copy from src to dst
void <b>bzero</b> (char *dst, int bytes);	Clear the bytes in dst
int <b>bcmp</b> (char *ogrn, char *tgt, int bytes);	Compare two byte strings, If same, return 0, else otherwise

# *Little Endian & Big Endian*



Byte-orderings are different.  
This is CPU dependent.

# Byte ordering, Little endian, Big endian

- Variable's bytes-ordering is reversed in some other machines.
  - For integer, short, ***ntohl()***, ***htonl()*** can be used.
    - Big endian, Little endian, network standard
  - But for float or double, the ordering must be cared by users.
  - Endian may differs from each other.