Testing Framework for Zero Knowledge Programs

Motivation

- Zero knowledge programming now is practical
- Developers want to verify or test their programs as they want this to work in privacy or financial applications with a lot at stake

Background: Making ZK Available to Developers

- ullet People developed DSL to let developers describe a function in the family ${\mathscr F}$
- What exactly are the statements can the system prove?

Let's look at a concrete system called zkSNARK

Background: Rank-1 Constraint System

• We say $(x_1, x_2, ..., x_n) \in F_p^n$ satisfied a Rank-1 Constraint System (R1CS) iff the following are all satisfied

$$(a_{11}x_1 + a_{12}x_2 + \dots, a_{1n}x_n) \cdot (b_{11}x_1 + b_{12}x_2 + \dots, b_{1n}x_n) = c_{11}x_1 + c_{12}x_2 + \dots, c_{1n}x_n$$

$$(a_{21}x_1 + a_{22}x_2 + \dots, a_{2n}x_n) \cdot (b_{21}x_1 + b_{22}x_2 + \dots, b_{2n}x_n) = c_{21}x_1 + c_{22}x_2 + \dots, c_{2n}x_n$$

$$(a_{31}x_1 + a_{32}x_2 + \dots, a_{3n}x_n) \cdot (b_{31}x_1 + b_{32}x_2 + \dots, b_{3n}x_n) = c_{31}x_1 + c_{32}x_2 + \dots, c_{3n}x_n$$

$$\dots$$

$$(a_{m1}x_1 + a_{m2}x_2 + \dots, a_{mn}x_n) \cdot (b_{m1}x_1 + b_{m2}x_2 + \dots, b_{mn}x_n) = c_{m1}x_1 + b_{m2}x_2 + \dots, b_{mn}x_n$$
where $a_{ij}, b_{i,j}, c_i$ are all in F_p

Background: Compiler

- People have developed some high level DSL and can compile to R1CS form
 - xJsnark by Elaine Shi et al. in S&P 2018
 - Viaduct by Andrew Myers et al. in PLDI 2021
 - Cairo by Starkware in 2021
 - CirC by in 2021
 - Leo in 2021

Problem: Developers are still writing R1CS Form!

- Two main reason:
 - Developers and tooling still not catch up yet
 - Manually optimize R1CS form is sometimes necessary to get efficient program because proving time is proportional to the number of constraints
 - Writing R1CS is like writing assembly level code!

- Given an integer x compute the multiplicative inverse y corresponding in the underlying field F_p and assume 0 <= x < p
 - If we Implement this in a general programming language without field arithmetic: extended GCD algorithm or using Fermat's little theorem
 - R1CS form

$$x \cdot y = 1$$

•Given a number x in [0, 15] return its binary representation

Input: x of type integer

Output: four boolean variables y0, y1, y2, y3 corresponding to the binary representation of x

$$x := x / 2$$

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$$x := x/2$$

- Given a number in [0, 15] return its binary representation
 - Optimized R1CS form

$$x = y_0 + 2 \cdot y_1 + 4 \cdot y_2 + 8 \cdot y_3$$

$$y_0 \cdot (y_0 - 1) = 0$$

$$y_1 \cdot (y_1 - 1) = 0$$

$$y_2 \cdot (y_2 - 1) = 0$$

$$y_3 \cdot (y_3 - 1) = 0$$

• Input an integer ${\bf x}$ in ${\cal F}_p$

Output an integer
$$y = \begin{cases} 1, & \text{if } x \text{ is zero} \\ 0, & \text{otherwise} \end{cases}$$

• R1CS form

$$x \cdot y = 0$$
$$y + x \cdot z = 1$$

(z is also a variable, and it's there to help to make it work)

Problem: Verifying or Testing R1CS form is hard

- People now are still writing this kind of low level R1CS form
- Given a spec, people want to verifying if their handwritten R1CS form is correct

Given a function $f: F_p \to F_p$ as our spec and

an R1CS
$$CS: F_p \times F_p^m \times F_p \rightarrow \{\text{True}, \text{False}\}$$

• (Existence) $\forall x . \exists z_1, z_2, ..., z_m$.

$$CS(x, z_1, z_2, ..., z_m, f(x))$$

• (Uniqueness) $\forall x . \forall z_1, z_2, ..., z_m . \forall y$.

$$CS(x, z_1, z_2, ..., z_m, y) \Rightarrow y = f(x)$$

Given x as a test case input

• (Testing Existence) $\exists z_1, z_2, ..., z_m$.

$$CS(x_1, z_1, z_2, ..., z_m, f(x_1, x_2, ..., x_n))$$

• (Testing Uniqueness) $\forall z_1, z_2, ..., z_m . \forall y$.

$$CS(x, z_1, z_2, ..., z_m, y) \Rightarrow y = f(x_1, x_2, ..., x_n)$$

Given x as a test case input

• (Testing Existence) $\exists z_1, z_2, ..., z_m$.

$$CS(x_1, z_1, z_2, ..., z_m, f(x_1, x_2, ..., x_n))$$

People have written tests for this

• (Testing Uniqueness) $\forall z_1, z_2, ..., z_m . \forall y$.

$$CS(x, z_1, z_2, ..., z_m, y) \Rightarrow y = f(x_1, x_2, ..., x_n)$$

People don't have tests for this

$$CS(x_1, x_2, ..., x_n, z_1, z_2, ..., z_m, y)$$

Substitute $x_1, x_2, ..., x_n$ for the test case input

$$MQ(z_1, z_2, \ldots, z_m, y)$$

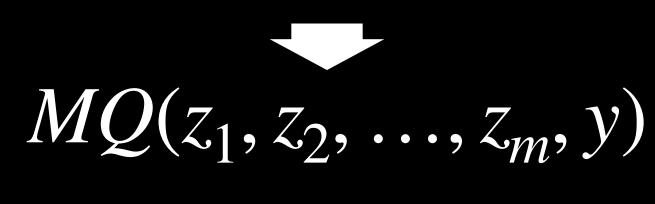
Grobner Basis Solver

Solutions of
$$(z_1, z_2, ..., z_m, y)$$

Check if $y = f(x_1, x_2, ..., x_n)$ for all solutions

$$CS(x_1, x_2, ..., x_n, z_1, z_2, ..., z_m, y)$$

Substitute $x_1, x_2, ..., x_n$ for the test case input





Grobner Basis Solver



Assume $(z_1, z_2, ..., z_m)$ should also be unique



Check if the solution is unique

Demo

```
pragma circom 2.0.2;
template Num2FourBits() {
    signal input x;
    signal output bits[4];
    for(var i = 0; i < 4; i++){
        bits[i] \leftarrow (x >> i) & 1;
    x === bits[0] + bits[1] * 2 + bits[2] * 4 + bits[3] * 8;
    bits[0] * (bits[0]-1) === 0;
    bits[1] * (bits[1]-1) === 0;
    bits[2] * (bits[2]-1) === 0;
    bits[3] * (bits[3]-1) === 0;
}
```

Demo

Demo: Buggy Code

```
pragma circom 2.0.2;
template Num2FourBits() {
    signal input x;
    signal output bits[4];
    for(var i = 0; i < 4; i++){
        bits[i] <-- (x >> i) & 1;
    x === bits[0] + bits[1] * 2 + bits[2] * 4 + bits[3] * 8;
    bits[0] * (bits[0]-1) === 0;
    bits[1] * (bits[1]-1) === 0;
    bits[2] * (bits[2]-1) === 0;
    //bits[3] * (bits[3]-1) === 0;
}
```

Demo

```
thekev@zktest $ yarn test --grep "Num2FourBits"
yarn run v1.22.17
$ mocha -r ts-node/register 'test/**/*.ts' --grep Num2FourBits

Num2FourBits

Testing x: 12, expect output 0,0,1,1 (58ms)

1 passing (129ms)

Done in 2.35s.
```

Demo: Running our Solver-based test

```
thekev@zktest $ yarn test --grep "Num2FourBits"
yarn run v1.22.17
$ mocha -r ts-node/register 'test/**/*.ts' --grep Num2FourBits
 Num2FourBits
    1) Testing x: 12, expect output 0,0,1,1
  0 passing (2s)
  1 failing
  1) Num2FourBits
       Testing x: 12, expect output 0,0,1,1:
     AssertionError: Error: solutions might Not be unique
      at Object.checkConstraints (node_modules/circom_r1cs_tester/index.js:23:9)
      at Context.<anonymous> (test/demo.test.ts:30:34)
error Command failed with exit code 1.
info Visit https://yarnpkg.com/en/docs/cli/run for documentation about this command.
```

Limitations

 Grobner basis solver can only handle ~10 variables in reasonable amount of time

Future work

- Compilers for new proving systems
- Finding a good intermediate representation (IR)