## 作业八

1 假设以邻接表表示法作为图的存储结构,设计图的深度优先遍历递归算法。

```
template⟨class T⟩
inline void ALGraph<T>::DFS(int v, vector<bool>& visited)
    cout << adjlist[v].data << " ";</pre>
    visited[v] = true;
    for (EdgeNode* p = adjlist[v].firstedge; p; p = p->nextedge)
        if (!visited[p->adjvex])
            DFS(p->adjvex, visited);
template < class T>
inline void ALGraph<T>::DFS()
   vector<bool> visited(vexnum, false);
    for (int i = 0; i < vexnum; i++)
        if (!visited[i])
            DFS(i, visited);
    2. 试基于图的广度优先搜索策略编写一种算法,判别以邻接表方式存储的有向图中是否
存在由顶点 v_i 到顶点 v_i 的路径 (i \neq j)。
template⟨class T⟩
inline bool ALGraph<T>::ifExistPath BFS(T a, T b)
    int i = GetSubscriptOfVex(a), j = GetSubscriptOfVex(b);
    int e, k;
    vector<bool> visited(vexnum, false);
    queue⟨int⟩ q;
    q. push(i);
    while (!q.empty())
    {
        e = q. front();
            q. pop();
        visited[e] = true;
        for (EdgeNode* p = adjlist[e].firstedge; p; p = p->nextedge)
            k = p->adjvex;//当前指向顶点的位置
            if (k == j)
                return true;
            else if (!visited[k])
                q. push(k);
        }
```

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return false;
    3. 试修改 Prim 算法, 使之能在邻接表存储结构上实现求有向图的最小生成森林(森林
的存储结构为孩子兄弟链表)。
template \langle class T \rangle
inline void ALGraph<T>::Prim(int v, vector<bool>& visited, vector<Couple<T>>& mintree)
    vector<Edge<T>> miniedges(vexnum);//建立该点邻接点的数组,记录到达所需权值
    for (int i = 0; i < vexnum; i++)
        miniedges[i].adjvex = adjlist[v].data;
        miniedges[i].lowcost = GetWeight(v, i);
    miniedges[v].lowcost = 0;
    for (int i = 1; i < vexnum; i++)
        int k = 0;
        int min = INT_MAX;
        for (int j = 0; j < vexnum; j++)
            if (miniedges[j].lowcost < min && miniedges[j].lowcost != 0)//注意!=0
            {
                min = miniedges[j].lowcost;
                k = j;
        Couple<T> tmp;
        tmp.parent = miniedges[k].adjvex;
        tmp.child = adjlist[k].data;
        visited[GetSubscriptOfVex(tmp.parent)] = true;
        visited[GetSubscriptOfVex(tmp.child)] = true;
        mintree.push_back(tmp);
        miniedges[k].lowcost = 0;//此处排除已进入的点
        for (int j = 0; j < vexnum; j++)
            int costkj = GetWeight(k, j);
            if (costkj < miniedges[j].lowcost && miniedges[j].lowcost != 0 && costkj !=
0)//注意!=0
            {
                miniedges[j].adjvex = adjlist[k].data;
                miniedges[j].lowcost = costkj;
}
```

}

```
template<class T>
inline void ALGraph<T>::MinSpanForest Prim()
    vector<bool> visited(vexnum, false);
    vector<vector<Couple<T>>> mintrees(vexnum);
    vector<Couple<T>> minforest;
    for (int i = 0; i < vexnum; i++)
        if (!visited[i])
            Prim(i, visited, mintrees[i]);
    for(int i=0;i<vexnum;i++)</pre>
        if (!mintrees[i].empty())
            Couple<T> tmp;
            tmp.parent = 'R';
            tmp. child = adjlist[i]. data;
            minforest.push_back(tmp);
            int len = mintrees[i].size();
            for (int j = 0; j < len; j++)
                 minforest.push_back(mintrees[i][j]);
    for (int i = 0; i < minforest.size(); i++)</pre>
        cout << minforest[i].parent << "," << minforest[i].child << endl;</pre>
    //此处应使用森林生成算法(二元组), 待补
}
    4. 采用邻接表存储结构,编写一个判别无向图中任意两个给定的顶点之间是否存在一条
长度为k的简单路径的算法。
template⟨class T⟩
inline bool ALGraph<T>::ExistPath_DFS_Lenth(int i, int j, int k, int &n, vector<bool>
&visited)
    EdgeNode* p;
    visited[i] = true;
    if (i == j && n == k)
        return true;
    n++;
    p = adjlist[i].firstedge;
    while (p)
    {
        if (!visited[p->adjvex])
            if (ExistPath_DFS_Lenth(p->adjvex, j, k, n, visited))
                 return true;
        p = p-nextedge;
    visited[i] = false;
```

```
n--;
    return false;
template ⟨class T⟩
inline bool ALGraph<T>::ExistPath_Lenth(int i, int j, int k)
    vector<bool> visited(vexnum, false);
    int n = 0;
    return ExistPath_DFS_Lenth(i, j, k, n, visited);
    5. 找出邻接矩阵表示的图中的一个回路
template⟨class T⟩
inline void MGraph<T>::FindLoop_DFS(int vi, int v, vector<bool>& visited, vector<int>&
path)
{
    if (v != vi)
        visited[v] = true;
    path. push back(v);
    for (int i = 0; i < vexnum; i++)
         if (edges[v][i] ==1 \&\& i == vi)
         {
             path. push back(i);
             for (int i = 0; i < path. size(); i++)</pre>
                  cout << path[i] << " ";</pre>
             cout << endl;</pre>
         }
         if (edges[v][i] ==1 && !visited[i])
             FindLoop_DFS(vi, i, visited, path);
    path. pop_back();
}
template⟨class T⟩
inline void MGraph<T>::FindLoop(int vi)
{
    vector(int) path;
    bool exist = false;
    vector<bool> visited(vexnum, false);
    FindLoop_DFS(vi, vi, visited, path);
    if (path.empty())
         cout << vi << " 顶点无回路" << endl;
}
```