5.11.1

Given the address stream shown, and the initial TLB and page table states provided above, show the final state of the system. Also list for each reference if it is a hit in the TLB, a hit in the page table, or a page fault.

Final TLB and Page Table:

Addr	Addr in Binary	Page number	Page	Final status
			number(base10)	
4669	0001001000111101	0001	1	Page Fault
2227	0000100010110011	0000	0	Page Hit
13916	0011011001011100	0011	3	TLB Hit
34587	1000011100011011	1000	8	Page Fault
48870	1011111011100110	1011	11	Page Hit
12608	0011000101000000	0011	3	TLB Hit
49225	110000001001001	1100	12	Page Hit

Valid Bit	Tag	Physical page number
1	12	15
1	8	14
1	3	6
1	11	12

Valid	Physical Page	
1	5	
1	13	
0	Disk	
1	6	
1	9	
1	11	
0	Disk	
1	4	
1	14	
0	Disk	
1	3	
1	12	
1	15	
	·	

5.11.2 Repeat 5.11.1, but this time use 16 KiB pages instead of 4 KiB pages. What would be some of the advantages of having a larger page size? What are some of the disadvantages?

Addr	Addr in Binary	Page number	Page	Final status
			number(base10)	
4669	0001001000111101	0000	0	Page Fault
2227	0000100010110011	0000	0	TLB Hit
13916	0011011001011100	0000	0	TLB Hit
34587	1000011100011011	0010	2	Page Fault
48870	1011111011100110	0010	2	TLB Hit
12608	0011000101000000	0000	0	TLB Hit
49225	110000001001001	0011	3	TLB Hit

Final TLB and Page table:

Valid Bit	Tag	Physical page number
1	2	13
1	7	4
1	3	6
1	0	5

The advantage of a larger page size could bring the cache associative mapping be increased and be more efficient. And the virtual address size is decreased. Therefore, the mapping of the block offset is also reduced.

However, the disk searching latencies could be increased to be high, which could slow down the process.

Valid	Physical Page
1	5
0	Disk
0	Disk
1	6
1	9
1	11
0	Disk
1	4
0	Disk
0	Disk
1	3
1	12
1	13

5.11.3 Show the final contents of the TLB if it is 2-way set associative. Also show the contents of the TLB if it is direct mapped. Discuss the importance of having a TLB to high performance. How would virtual memory accesses be handled if there were no TLB.

Addr	Addr in Binary	Page number	Page	Final status
			number(base10)	
4669	0001001000111101	0001	1	Page Fault
2227	0000100010110011	0000	0	Page Hit
13916	0011011001011100	0011	3	TLB Hit
34587	1000011100011011	1000	8	Page Fault
48870	1011111011100110	1011	11	Page Hit
12608	0011000101000000	0011	3	TLB Hit
49225	110000001001001	1100	12	Page Hit

Final TLB for 4669:

Valid Bit	Tag	Physical
		page number
		number
1	11	12
1	0	13
1	3	6
0	4	9

Final TLB for 2227:

Valid Bit	Tag	Physical page number
1	0	5
1	0	13
1	3	6
0	4	9

Final TLB for 13916:

Valid Bit	Tag	Physical
		page number
		number
1	0	5
1	0	13
1	3	6
1	0	6

Final TLB for 34587:

Valid Bit	Tag	Physical
		page number
		number
1	2	14
1	0	13
1	3	6
1	0	6

Final TLB for 48870:

Valid Bit	Tag	Physical
		page number
1	2	14
1	0	13
1	3	6
1	2	12

Final TLB for 12608:

Valid Bit	Tag	Physical
		page number
		number
1	2	14
1	0	13
1	3	6
1	0	6

Final TLB for 12608:

Valid Bit	Tag	Physical
		page number
1	3	15
1	0	13
1	3	6
1	0	6

The high performance for TLB could reduce the time for physical memory lookup mapping. If there is no TLB, the virtual memory would have to lookup and map each memory location once a time, which would lead to a inefficient performance.

5.11.4

Given the parameters shown above, calculate the total page table size for a system running 5 applications that utilize half of the memory available.

Number of page table entries = $(2^32)/(2^10 * 2^3) = 2^19$ Page Table Size = $2^19 * 2^2 = 2^18$ Sine it only utilize half of the memory available: Utilized size = 1^18 Therefore, Total page size used = 1^18