# Feature Matching using Supersymmetric Geometric Constraints

#### Abstract

Feature matching is a challenging problem lying at the heart of numerous computer graphics and computer vision applications. We present here the SuperMatching feature matching algorithm for finding correspondences between two sets of features. Super-Matching finds matches of feature points by considering triangle or higher-order polygons formed by such points, going beyond the pointwise and pairwise approaches typically used. It is formulated as a supersymmetric-tensor-based matching scheme, casting feature matching as a higher-order graph matching problem. The super-10 symmetric tensor represents an affinity metric which takes into ac-11 count geometric constraints between features. SuperMatching exploits a compact form of this affinity tensor, whilst also taking advantage of supersymmetry to devise an efficient sampling strategy 14 to create the affinity tensor. Matching is performed by computing a 15 rank-one approximation of the tensor directly using a higher-order 16 power solution method. Experiments on both synthetic and real 17 captured data show that our algorithm provides accurate and robust 18 feature matching even in complex cases.

Keywords: Feature matching; Geometric constraints; Supersymmetric tensor

## 1 Introduction

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Building correspondences between two sets of features of 2D/3D shapes is a fundamental problem in many geometric processing, computer graphics and computer vision tasks. It arises in applications such as feature extraction [Johnson and Hebert 1999; Lowe 2004; Sun et al. 2009; Toler-Franklin et al. 2010; Leutenegger et al. 2011], shape matching [Berg et al. 2005; Brown and Rusinkiewicz 2007; Tevs et al. 2009; Ovsjanikov et al. 2010; Tevs et al. 2011; Sahillioglu and Yemez 2011; Windheuser et al. 2011], registration of 3D shapes [Gelfand et al. 2005; Aiger et al. 2008; Li et al. 2008; Zeng et al. 2010; van Kaick et al. 2011; Chang and Zwicker 2011], automatic shape understanding [Lipman and Funkhouser 2009; Sun et al. 2010; Kim et al. 2011], shape retrieval from database [Bronstein et al. 2011], and reconstruction [Brown and Rusinkiewicz 2007; Chang and Zwicker 2011].

In principle and practise, the correspondence building process mainly proceed in three steps [Lowe 2004; Leutenegger et al. 2011]: salient feature detection, high-quality descriptor, and accurately matching. The former two problems have been been widely attracted considerable attention as their importance is easy perceivable. However, even with the ideal feature detector and descriptor that capturing the most important and distinctive information content enclosed in the detected salient regions, the correspondences still could not ideally built in the state-of-the-art algorithms [van Kaick et al. 2011]. The reason is that the real input data is so inperfect and complex especially with symmetric and repetitive regions. The researchers are beginning to be aware of that the limitation could be effectively alleviated by feature matching. Some original idea (e.g. RANSAC-like algorithms [Tevs et al. 2009; Tevs et al. 2011]) has been extended, or indirectly sorting to shape embedding strategies, such as Generalized Multidimensional Scaling [Bronstein et al. 2011], Heat Kernel Map [Ovsjanikov et al. 2010], 108 Möbius Transformation [Lipman and Funkhouser 2009; Kim et al.

2011]. But, these previous algorithms still did not treat feature matching as one independent problem, although they agreed with that matching is not tightly coupled with feature detection and description.

In the paper, we mainly focus on feature matching problem, which is complementary to existing correspondence algorithms. As we know, a matching itself may comprise one or more items of a given kind. We may match single point to single point (point-single), point pairs separated by a fixed distance to other point pairs (segment-double), triples of points forming a triangle to other triples of points (triangle-triple), quadruples of points forming a quadrangle to other quadruples of points (quadrangle-quadruple), and so on. As pointed by [Conte et al. 2004], when single features are matched, we must solve a linear assignment problem, if multiple features are matched at once, a quadratic or higher-order assignment problem results.

Linear assignment matches single features in one set with single features in the other set, and could be mainly treated as single points linkage. Matching two feature sets by considering similarities of *single* features from each set can easily fail in the presence of ambiguities such as repeated elements, textures or relatively uniform local appearance.

Quadratic and higher-order assignment matches groups of features in one set simultaneously with groups from the other set, and requires a greater consistency between the information being matched, making it more reliable. As well as the features themselves, other constraints such as consistency of the distances between the features being matched are also enforced, greatly improving the reliability of higher order matching.

As a particular example of *quadratic* assignment, Leordanu and Hebert [Leordeanu and Hebert 2005] consider pairs of feature descriptors, and distances between pairs of features from each set to reduce the number of incorrect correspondences. Such pairwise distance constraints are particularly helpful in cases when the features themselves have low discriminative ability. The idea has been widely adopted in the 3D shape matching algorithms [Tevs et al. 2009; Ovsjanikov et al. 2010; Tevs et al. 2011; Kim et al. 2011; Sahillioglu and Yemez 2011; Windheuser et al. 2011].

Higher-order assignment further generalizes the assignment problem to include yet more complex constraints between features. For example, third-order potential functions, proposed in [Duchenne et al. 2009; Zeng et al. 2010; Chertok and Keller 2010], quantify the affinity between two point triples by measuring the similarity of the angles formed by such two feature tuples (triangles). However, this angular similarity value only calculates the total differences in corresponding angles, and does not change with the order of the input assignments. By changing the affinity tensor to a *supersymmetric* tensor [Kofidis and Regalia 2002], the unaccurate limitation could be further solved by our algorithm.

Thus, by formulating the higher-order problem using a supersymmetric affinity tensor, we propose an accurate and efficient higher-order supersymmetric matching algorithm, named *SuperMatching*. SuperMatching, addressed in Section 4, works well when matching a moderate number of features using triples or higher polygons of features. The contributions of this paper include:

We show how to define a compact higher-order supersymmet-

ric affinity tensor to express geometric consistency constraints the between tuples of features.

- Based on the supersymmetry of the affinity tensor, we propose a higher-order power iteration method, which efficiently solves the matching problem.
- The affinity tensor is physically created by using a new efficient sampling strategy for feature tuples, which avoids sampling repetitive items, reduces the number of feature tuples to be sampled while ensuring the accuracy of the power iteration result.

Our experiments in Section ?? on both synthetic and real captured data sets show that SuperMatching result is accurate and robust, and has competitive computational cost compared to previous algorithms. More important, the matching is performed by incorporating variant 2D and 3D feature descriptors, which demonstrate that SuperMatching is independent to descriptors and could be generally applied in computer graphics and computer vision fields.

#### 2 Related work

According to the applied constraints, previous approaches to feature matching can be classified into those which match single points to single points, those which match pairs of points to pairs of points, and so on.

Matching single points to single points is a linear assignment problem which only considers an affinity measure between two graph nodes, one from each set being matched; this measure is typically the feature distance between the two feature points. In concrete terms, the linear assignment problem may be expressed as: find a mapping  $f: P_1 \rightarrow P_2$ , Affinity measures used in computer vision and computer graphics tasks rely heavily on descriptors computed using local information around each feature point, e.g. SIFT [Lowe 2004], spin images [Johnson and Hebert 1999], heat diffusion signatures [Sun et al. 2009], and BRISK [Leutenegger et al. 2011]. It is apparent that point-to-point matching is weak in that wrong correspondences maybe readily be established.

Matching pairs of points in one set to pairs of points in the other set leads to a quadratic assignment problem. The usual approach is now to take into account both similarity of the point features *and* the Euclidean/Geodesic distance between the points in a pair, assuming the objects are related by a rigid body transform [Leordeanu and Hebert 2005], or at least an isometry [Li et al. 2008; Tevs et al. 2009; Ovsjanikov et al. 2010; Tevs et al. 2011; Sahillioglu and Yemez 2011; Windheuser et al. 2011]. The quadratic assignment problem seeks to find a mapping which represents the optimal assignment. Unfortunately, this problem is NP-hard, unreasonable matching results could not be avoided.

Several higher-order approaches have also been proposed. Such higher-order methods can significantly improve matching accuracy, but higher-order assignment problem is again NP-hard, and various approximate methods have again been developed. Zass and Sashua [Zass and Shashua 2008] consider a probabilistic model of soft hypergraph matching. They reduce the higher-order problem to a first-order one by marginalizing the higher-order tensor to a one dimensional probability vector. Duchenne et al. [Duchenne et al. 2009] introduced a third-order tensor in place of an affinity matrix to represent affinities of feature triples, and higher-order power iteration was used to achieve the final matching. Chertok et al. [Chertok and Keller 2010] treat the tensor as a joint probability of assignments, marginalize the affinity tensor to a matrix, and find optimal soft assignments by eigendecomposition of the matrix. Wang et al. [Wang et al. 2010] also build a third-order affinity tensor, and

obtain a final matching by rank-one approximation of the tensor. Higher-order assignment problems typically require large amounts of memory and computational resources. By reducing the number of elements needed to represent the affinity measures, the above approaches can efficiently match large numbers (many hundreds or more) of features. However, these approaches sparsify the affinity information to some degree, leading to a reduction in matching accuracy. When matching two feature sets which do not accurately meet the assumption that the datasets are related by an isometry, the matching results may become unstable.

A related idea using higher constraints in 3D registration, 4-points congruent sets method (4PCS), was proposed by Aiger et al. [Aiger et al. 2008]. It is a fast alignment scheme for 3D point sets that uses widely separated points. However, the coplanar 4 points and rigid-transformation constraints are some strong, and limit its application in real data. We solve both rigid and isometric shape matching problem by one more mathematical and formulated tensor-based algorithm.

## 3 Overview

Assume we are given two sets of feature points  $P_1$  and  $P_2$ , with  $N_1$  and  $N_2$  points respectively. The matching between these two feature sets can be represented by an assignment matrix X of size  $N_1 \times N_2$  whose elements are 0 or 1. X(i,j) = 1 when  $i \in P_1$  matches  $j \in P_2$  and X(i,j) = 0 otherwise. X can be row-wise vectorized to give an assignment vector  $\mathbf{x} \in \{0,1\}^{N_1N_2}$ .

Solving the higher-order matching problem is equivalent to find the optimal assignment vector  $\boldsymbol{x}^* = < x_{i_1}, \cdots, x_{i_N} > \in \{0,1\}^{N_1 N_2}$ , satisfying

$$\boldsymbol{x}^* = \arg\max_{\boldsymbol{x}} \sum_{i_1, \dots, i_N} \mathcal{T}_N(i_1, \dots, i_N) x_{i_1} \dots x_{i_N}.$$
 (1)

Here,  $i_n$  stands for an assignment  $(i_n^{'},j_n^{'}), (i_n^{'} \in P_1,j_n^{'} \in P_2)$ , and the product  $x_{i_1} \cdots x_{i_N}$  is 1 only if all assignments  $\{i_n\}_{n=1}^N$  are equal to 1, which means the tuple of features  $(i_1^{'},\cdots,i_N^{'})$  from  $P_1$  is matched correspondingly to the tuple of features  $(j_1^{'},\cdots,j_N^{'})$  from  $P_2$ .  $\mathcal{T}_N(i_1,\cdots,i_N)$  defines the affinity of the set of assignments  $\{i_n\}_{n=1}^N$ ; it also can be seen as the affinity measure between the ordered feature tuples  $(i_1^{'},\cdots,i_N^{'})$  and  $(j_1^{'},\cdots,j_N^{'})$ .

In the paper, we consider the one-to-many correspondence problem. We assume that a point in  $P_1$  is matched to exactly one point in  $P_2$ , but that the reverse is not necessarily true. If do we want to treat both datasets in the same way, we can first match  $P_1$  to  $P_2$ , then match  $P_2$  to  $P_1$ , and then combine the matching results by taking their union or intersection. Uniqueness of matches for  $P_1$  means that the assignment matrix X satisfies  $\sum_i X(i,j) = 1$ .

From Equ.(1) we can see that there are four issues to be considered when using higher-order matching algorithms. How should we:

- organize and express the affinity measures  $\mathcal{T}_N$ ? (see Section 4.1)
- approximately solve the optimal higher-order assignment problem efficiently? (see Section 4.2)
- define the affinity measure between two feature tuples, or equivalently, the higher-order potential function  $\phi_N$ ? (see Section 4.3)
- determine an appropriate sampling strategy that physically build the affinity tensor and influence good matching accuracy? (see Section 4.4)

## 4 SuperMatching

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260 261 We now discuss the first and second issues mentioned above, which are independent of application; later we turn to definition of affinity measure, which is application dependent, and sampling strategy.

#### 4.1 Supersymmetric affinity tensor

A tensor generalises vectors and matrices to higher dimensions: a vector is a tensor of order one, and a matrix is a tensor of order two. A higher-order tensor can be expressed as a multi-dimensional array [Kolda and Bader 2009]. Here we consider a higher-order supersymmetric affinity tensor, which represents a real-valued higher-order affinity between feature tuples.

**Definition 1 (Supersymmetric Tensor)** A tensor is called supersymmetric if its entries are invariant under any permutation of its indices [Kofidis and Regalia 2002].

For example, a third-order supersymmetric tensor  $\mathcal{T}_3$ , satisfies the relationships:  $\mathcal{T}_3(i_1,i_2,i_3) = \mathcal{T}_3(i_1,i_3,i_2) = \mathcal{T}_3(i_2,i_1,i_3) = \mathcal{T}_3(i_2,i_3,i_1) = \mathcal{T}_3(i_3,i_1,i_2) = \mathcal{T}_3(i_3,i_2,i_1)$ .

**Definition 2 (Supersymmetric Affinity Tensor)** Given two feature sets  $P_1$  and  $P_2$ , with  $N_1$  and  $N_2$  features respectively, the supersymmetric affinity tensor is an  $N^{th}$  order  $I_1 \cdots \times I_N$ , nonnegative tensor  $\mathcal{T}_N$ , where  $I_1 = I_2 = \cdots = I_N = \{1, \cdots, N_1 N_2\}$ , for which there exists a set of indices  $\theta_N$ , and an  $N^{th}$  order potential function  $\phi_N$ , such that

$$\mathcal{T}_{N}(i_{1},\ldots,i_{N}) = \begin{cases} \phi_{N}(\Omega(i_{1},\ldots,i_{N})) &, \forall (i_{1},\ldots,i_{N}) \in \theta_{N} \\ 0 &, \forall (i_{1},\ldots,i_{N}) \notin \theta_{N} \end{cases}$$
(2)

where  $\Omega$  stands for an arbitrary permutation of the vector, and  $\theta_N$  satisfies  $\forall (i_1, i_2, \ldots, i_N) \in \theta_N, \forall i_p \in \{i_1, \ldots, i_N\}$  and  $\forall i_q \in \{i_1, \ldots, i_N\} - \{i_p\}$  meets the requirement that  $i_p \neq i_q$ .

A tensor element with  $(i_1, i_2, ..., i_N) \in \theta_N$  is called a potential element, while other elements are called non-potential element.

Using Definition 2, we now can greatly reduce the amount of storage needed, representing every potential element  $\mathcal{T}_N(i_1,i_2,\ldots,i_N)$  by the canonical entry  $\mathcal{T}_N(\operatorname{sort}(i_1,i_2,\ldots,i_N)), \ \forall (i_1,i_2,\ldots,i_N) \in \theta_N$ . Each stored value thus provides the value for N! entries. As nonpotential elements all have value zero, there is no need to store them. This greatly reduces the sampling needed for feature tuples when creating the affinity tensor, as discussed in Section 4.4. At the same time, it can be used to make the power iteration process more efficient: see Section 4.2.

## 4.2 Higher-order Power Iteration Solving

Using Definition 2, Equ.(1) can be expressed as:

$$\boldsymbol{x}^{*} = \arg \max_{\boldsymbol{x}} \sum_{i_{1}, i_{2}, \dots, i_{N}} \mathcal{T}_{N}(i_{1}, \dots, i_{N}) x_{i_{1}}, \dots, x_{i_{N}}$$

$$= \max \langle \mathcal{T}_{N}, \boldsymbol{x}^{*N} \rangle$$
(3)

where  $\star$  is called the Tucker product [Kofidis and Regalia 2002], and  $\boldsymbol{x} \in \{0,1\}^{N_1N_2}$ . Solving Equ.(3) is an NP-complete problem, so it is common to relax the constraints: the binary assignment vector  $\boldsymbol{x} \in \{0,1\}^{N_1N_2}$  is replaced by an assignment vector  $\boldsymbol{u}$  with elements taking real values in [0,1]. This changes the optimization problem to one of computing the rank-one approximation of the affinity tensor  $\mathcal{T}_N$  [Kofidis and Regalia 2002], i.e.

**Algorithm 1** Higher-order power iteration method for a supersymmetric affinity tensor (with  $C_1$  norm)

```
Input: Nth-order supersymmetric affinity tensor \mathcal{T}_n
Output: unit-norm vector u
          1: Initialize u_0, k = 1
        2: repeat
        3:
                                                 for all (i_1, i_2, \cdots, i_N) \in \theta_N do
        4:
                                                                    for all m \in (i_1, \cdots, i_N) do
                                                                                       \begin{aligned} v_m^{(k)} &= (N-1)! \phi_N(i_1, \cdots, i_N) 2 v_m^{(k-1)} \, v_{i_1}^{2(k-1)} \cdots \\ v_{m-1}^{2(k-1)} \, v_{m+1}^{2(k-1)} \cdots v_{i_N}^{2(k-1)} \end{aligned}
          5:
        6:
          7:
                                                                    for i=1:N_1 do
                                                                                       v^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2) = \hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2+1)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2+1)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2+1)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1).\|
          8:
       9:
 10:
                                                 end
 11:
                                                 k = k + 1;
  12: until convergence;
                             Note: v^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2) denotes the slice of v^{(k)} with indices
                             from (i-1)\cdot N_2+1 to i\cdot N_2.
```

finding a scalar  $\lambda$  and a unit norm vector  $\boldsymbol{u} \in \mathbb{R}^{N_1 N_2}$ , such that the tensor  $\hat{\mathcal{T}_N} = \lambda \boldsymbol{u} \star \boldsymbol{u} \star \cdots \star \boldsymbol{u} = \boldsymbol{u}^{\star N}$  minimizes the function  $f(\hat{\mathcal{T}_N}) = \|\mathcal{T}_N - \hat{\mathcal{T}_N}\|$ . The final matching result is found by replacing each element of  $\boldsymbol{u}$  by 0 or 1 according to whichever it is closer to.

The higher-order power method is commonly used to find the rank-one tensor approximation; a version for supersymmetric tensors (S-HOPM) is given in [Kofidis and Regalia 2002]. The S-HOPM algorithm converges under the assumption of convexity for the functional induced by the tensor [Kofidis and Regalia 2002], which is so robust that could be satisfied in practical application. S-HOPM is performed in two iterative steps: high-order power iteration of  $\boldsymbol{u}$ , normalization of  $\boldsymbol{u}$  under the Frobenius norm. In the recent, one more effective improvement [Duchenne et al. 2009] was proposed by using  $\mathcal{C}_1$  norm to replace the traditional  $\mathcal{C}_2$  norm.

We directly take the  $C_1$  norm [Duchenne et al. 2009] and further revise S-HOPM as following. For the first step (high-order power

iteration of  $\boldsymbol{u}$ ),  $\hat{\boldsymbol{u}}^{(k)} = \mathcal{I}^{\mathcal{T}_N}_{\star} \left(\boldsymbol{u}^{(k-1)}\right)^{\mathcal{T}_N}_{\star} (N-1)$ ,  $\mathcal{T}_N$  is a so-called  $\mathcal{T}_N$ -product,  $\mathcal{I}$  is the unit tensor [Kofidis and Regalia 2002]. Then, for  $\hat{\boldsymbol{u}}^{(k)}$  with Nth-order supersymmetric affinity tensor, it can be formulated as:

$$\begin{split} \hat{\boldsymbol{u}}^{(k)} &= \mathcal{I} \overset{\mathcal{T}_{N}}{\star} \left( \boldsymbol{u}^{(k-1)} \right)^{\overset{\mathcal{T}_{N}}{\star} (N-1)} \Leftrightarrow \\ \forall m \in (i_{1}, \cdots, i_{N}), v_{m}^{(k)} &= \\ &\sum_{i_{1}, \dots, i_{N}} \mathcal{T}_{N}(i_{1}, \cdots, i_{N}) 2 v_{m}^{(k-1)} v_{i_{1}}^{2(k-1)} \cdots v_{m-1}^{2(k-1)} v_{m+1}^{2(k-1)} \cdots v_{i_{N}}^{2(k-1)} &= \\ &(N-1)! \phi_{N}(i_{1}, \cdots, i_{N}) 2 v_{m}^{(k-1)} v_{i_{1}}^{2(k-1)} \cdots v_{m-1}^{2(k-1)} v_{m+1}^{2(k-1)} \cdots v_{i_{N}}^{2(k-1)} & \\ & (4) \end{split}$$

where  $\boldsymbol{u}^{(k)} = \boldsymbol{v}^{2(k)}$ ,  $\phi_N$  is corresponding potential function. The deduction relies on two principles. First, we take advantage of the supersymmetry. Secondly, many of the elements of the affinity tensor are zero non-potential elements: it is much more efficient to perform the power iteration by just considering the non-zero potential elements. So, our supersymmetric higher-order power iteration solving is addressed as Algorithm 1.

This version excludes each non-potential element from the iteration process, so is more efficient, and the complexity of the whole iteration process only depends on the number  $|\theta_N|$  of affinities. Step 5

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in Algorithm 1 includes all permutations of each potential element  $\mathcal{T}_n(i_1,i_2,\cdots,i_n)$ , which are expressed by a single potential function  $\phi_n(i_1,i_2,\cdots,i_n)$ . Consequently, This method reduces memory costs while keeping accuracy. Please noted that, although the affinity tensor [Duchenne et al. 2009] was claimed as supersymmetric. [Duchenne et al. 2009] did not make full use of supersymmetry when creating the supersymmetric affinity tensor, nor did they take advantage of supersymmetry when accelerating the power iteration process. Fortunately, the limitation due to unbalanced and redundant tensor elements in [Duchenne et al. 2009] could be overcome by Algorithm 1. The later experiment proves the expectation.

Many initialization schemes have been proposed for the power method [Kofidis and Regalia 2002]. We simply use positive random values to initialize  $u_0$ , which ensures the algorithm's convergency.

## 4.3 Higher-order Potentials

 Different higher-order potentials are appropriate for different applications. Here we give two general higher-order potentials. One is used for the 2D cases, while the other is defined for 3D matching. The potentials are based on a Gaussian kernel which guarantees the tensor elements are non-negative and invariant to any permutation of the input assignments.

In 2D cases, we first restate a well-known 2D third-order geometric-similarity invariant potential  $\phi_3$  [Duchenne et al. 2009; Chertok and Keller 2010] linking two feature tuples, each having three features. Similarity of triangles formed by three points corresponds to invariance under scaling, rotation and translation—interior angles do not change. Thus  $\phi_3$  can be defined in terms of differences of corresponding interior angles:

$$\phi_{3}(i,j,k) = \phi_{3}(\{i_{1},i_{2}\},\{j_{1},j_{2}\},\{k_{1},k_{2}\})$$

$$= \exp(-1/\varepsilon^{2} \sum_{(l,l')} \|\alpha_{l} - \alpha_{l'}\|^{2})$$
(5) 353

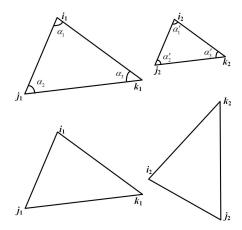
where  $\varepsilon>0$  is the is the kernel bandwidth, and  $\{\alpha_l\}_{l=1}^3$  and  $\{\alpha_l'\}_{l'=1'}^3$  are the angles triangles formed by feature triples  $(i_1,j_1,k_1)$  and  $(i_2,j_2,k_2)$ : see Fig. 1. Each element corresponds to one interior angle. We extend it to all general cases by using the internal angles formed by the triangle and higher polygons. It is easy to understand that he potential preserves invariance under similarity transformations in 2D field.

For 3D matching problem, we replace the internal angle by edge length, i.e., the geodesic distance between two neighboring points. Therefore, the feature matching is performed using the geometric constraints with invariance, which going beyond traditional the single (point) and pairwise features.

We will use these two high-order potentials to evaluate our algorithm. Fig. 1 illustrates the schematic diagram of third-order potential in 2D and 3D cases.

### 4.4 Sampling strategy

Algorithm 1 depends on all potential elements. Given two feature sets  $P_1$  and  $P_2$  with  $N_1$  and  $N_2$  features respectively, a potential element may be obtained by using two feature tuples sampled from each feature set separately. For Nth-order matching, a naive way to construct the potential elements is as follows: first find all feature tuples for  $P_1$  and  $P_2$ , as  $F_1$  and  $F_2$ ; then  $\forall (f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1) \in F_1$ , calculating the potentials for  $(f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1)$  with all feature tuples in  $F_2$ . This naive method is very expensive, which is why sampling is used. Choice of feature tuples to calculate the potentials directly determines the size  $|\theta_N|$  and influences the match-



**Figure 1:** Third-order potential diagram. The geometric constraints: internal angle invariance in 2D (up), and edge length invariance in 3D case (bottom).

ing accuracy. Different sampling strategies can be chosen for different applications, using random sampling for general feature correspondence problems where there is no guidance to provide a better sampling method.

Here we use a random sampling approach. In order to cover all features in  $P_1$  in  $F_1$ , we repeatedly take one feature as a required element, and then randomly choose  $t_1$  feature tuples containing thise required element. We repeat this process until all features in  $P_1$  have been chosen once as a required element. We do the same for  $P_2$ . We now have two feature-tuple sets for  $F_1$  and  $F_2$ , containing  $N_1t_1$  and  $N_2t_2$  feature tuples separately.

Now,  $\forall (f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1) \in F_1$ , we find the k most similar features in  $F_2$  to build k potential elements as  $\phi_i^k$ . Combining all the potential elements obtained, we form the desired potential element set  $\theta_N = \{\phi_i^k\}_{i=1}^{N_1t_1}$ , with the size  $|\theta_N| = N_1t_1k$ . The parameters  $t_1, t_2$  and k must be chosen according to the size of the feature sets. In practice for two point sets each with a hundred points, this approach works well when  $t_1 \approx t_2 \approx 50$  and  $k \approx 300$  for third-order or fourth-order matching. The sampling cost is  $O(n t k \log n)$ , where  $n = \max(N_1, N_2)$ ,  $t = \max(t_1, t_2)$ .

The most important part of the process is to use the supersymmetry of the affinity tensor. An  $N^{th}$ -order supersymmetric affinity tensor must satisfy:

$$\forall (i_1, i_2, \dots, i_N), (j_1, j_2, \dots, j_N) \in \theta_N, (i_1, i_2, \dots, i_N) \neq \Omega(j_1, j_2, \dots, j_N)$$
 (6)

where  $\Omega$  is an arbitrary permutation. Thus, we use a sampling constraint that the sets of feature tuples  $F_1$  and  $F_2$  obtained from the sampling process, should have no repetition, in the sense that

$$\forall (f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1), (f_{j_1}^1, f_{j_2}^1, \cdots, f_{j_N}^1) \in F_1, (f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1) \neq \Omega(f_{j_1}^1, f_{j_2}^1, \cdots, f_{j_N}^1)$$
 (7)

$$\forall (f_{i_1}^2, f_{i_2}^2, \cdots, f_{i_N}^2), (f_{j_1}^2, f_{j_2}^2, \cdots, f_{j_N}^2) \in F_2,$$

$$(f_{i_1}^2, f_{i_2}^2, \cdots, f_{i_N}^2) \neq \Omega(f_{j_1}^2, f_{j_2}^2, \cdots, f_{j_N}^2)$$
(8)

and  $\Omega$  is arbitrary permutation. This sampling constraint eliminates overlaps that may appear in the potential elements  $\theta_N$ .

When  $N_1$  is not equal to  $N_2$  (suppose  $N_1 < N_2$ ), we perform sampling as in [Duchenne et al. 2009], using random sampling for  $P_1$ 

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but full sampling for  $P_2$  (to find all  $N_2^N$  feature tuples). However we still use the sampling constraint for feature-tuple set  $F_1$ , so our strategy is different to that in [Duchenne et al. 2009]. The process of building the tensor elements is the same to the process stated above when  $N_1$  equals to  $N_2$ . In practice this works well when matching two feature sets with different numbers of features.

Earlier work [Duchenne et al. 2009; Zass and Shashua 2008] also adopted random sampling, but failed to impose any constraint on the sampling process, leading to the possibility that feature tuples may include repetition. For example, for third-order matching, it is 439 possible that a feature tuple  $(f_{i_1}^1, f_{i_2}^1, f_{i_3}^1)$  may be sampled from  $P_1$ and  $(f_{i_1}^2, f_{i_2}^2, f_{i_3}^2)$  from  $P_2$ , and also a feature tuple  $(f_{i_1}^1, f_{i_3}^1, f_{i_2}^1)$ sampled from  $P_1$  and  $(f_{i_1}^2, f_{i_3}^2, f_{i_2}^2)$  from  $P_2$ . That will create 442 two tensor elements  $\phi_3(s_{i_1}, s_{i_2}, s_{i_3})$  with index  $(s_{i_1}, s_{i_2}, s_{i_3})$  and 443  $\phi_3(s_{i_1}, s_{i_3}, s_{i_2})$  with index  $(s_{i_1}, s_{i_3}, s_{i_2})$ , which are the same. 444 However, we just need one tensor element to express the affinity measure on the assignment group  $(s_{i_1}, s_{i_2}, s_{i_3})$ . This problem not only makes the potential elements redundant but also affects the accuracy of the power iteration, because the numbers of each tensor element may not be equal and some elements may be used more than once during the iteration progress. Our method reduces the sampling cost, while also improving the accuracy of the power iteration.

## 5 Conclusion

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a novel matching algorithm named SuperMatching, which tackles the classic Computer Graphics and Computer Vision problem of matching various features for cases without any assumptions.

This paper has given an efficient higher-order matching algorithm based on the supersymmetric affinity tensor, based on an efficient power iteration method. Our main contributions are as follows. First, we use a higher-order supersymmetric affinity tensor with a compact form to express higher-order consistency constraints of features. Secondly, we derive an efficient higher-order power iteration method, which makes significant savings by taking advantage of supersymmetry. We also give an efficient sampling strategy for choosing feature tuples to create the affinity tensor. Finally, we give a fourth-order potential which possesses affine invariance. Our experiments on both synthetic and real image data sets show that our method has improved matching performance compared to state-of-the-art approaches.

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