# SuperMatching: Feature Matching using Supersymmetric Geometric Constraints

Zhi-Quan Cheng, Yin Chen, Ralph R. Martin, Yukun Lai, Aiping Wang

Abstract—Feature matching is a challenging problem lying at the heart of numerous computer graphics and computer vision applications. We present here the *SuperMatching* algorithm for finding correspondences between two sets of features. It does so by considering triangles or higher-order tuples of points, going beyond the pointwise and pairwise approaches typically used. SuperMatching is formulated using a supersymmetric tensor representing an affinity metric which takes into account feature similarity and geometric constraints between features: feature matching is cast as a higher-order graph matching problem. SuperMatching takes advantage of supersymmetry to devise an efficient sampling strategy to estimate the affinity tensor, as well as to store the estimated tensor compactly. Matching is performed by an efficient higher-order power iteration approach which takes advantage of this compact representation. Experiments on both synthetic and real captured data show that SuperMatching provides more accurate feature matching than other state-of-the-art approaches for a wide range of 2D and 3D features, with competitive computational cost.

Index Terms—Feature matching, Geometric constraints, Supersymmetric tensor.

## 1 Introduction

Building correspondences between two sets of features belonging to a pair of 2D images or 3D shapes is a fundamental problem in many computer graphics, geometry processing, and computer vision tasks. It arises in applications such as registration of partial or entire 3D shapes [1], [2], [3], [4], [5], [6], [7], [8], shape retrieval from databases [9], shape matching [10], [11], [12], [13], [14], [15], [16], [17], shape reconstruction [11], [18], [19], [8], and automatic shape understanding [20], [21], [22], [23].

Determining correspondences is typically done in three steps [24], [25], [26], [27], [28]: (i) computing high-quality descriptors which serve to distinguish points from one another, (ii) choosing certain salient points with unusual feature descriptors, for matching, and (iii) determining the most suitable matching between the two sets of points. The former two problems have attracted considerable attention; their importance is clear. However, even supposing ideal feature descriptors and selectors that capture the most important and distinctive information about the neighborhood of each salient point, state-of-theart algorithms still find it challenging to determine the best matching [7]: real input data is noisy, and data may only be approximately in correspondence. The problem is further complicated by the presence

of symmetric and congruent regions. Various feature matching algorithms have been devised to be robust in the presence of such issues. RANSAC-like algorithms [13], [15] minimize the effects of outliers, while generalized multidimensional scaling [9] and heat kernel maps [14] consider the manifold in which the points are embedded. Möbius transformations [21], [23] also provide a powerful approach. However, these previous algorithms generally do not treat the matching step as an independent problem, even if matching is not tightly coupled with feature description and selection. This paper focuses on the feature matching problem as a problem in its own right.

Matching may be done pointwise (single points to single points), or using tuples of points: e.g. point pairs, separated by a fixed distance, to other point pairs, triples of points forming a triangle to other triples of points, and so on. As pointed out by [29], matching single features leads to a linear assignment problem, but if multiple features are matched simultaneously, a quadratic or higher-order assignment problem results. Matching two feature sets by considering similarities of single features from each set can easily fail in the presence of ambiguities such as repeated elements, or similar local appearance. Quadratic and higherorder assignment matches groups of features, enforcing other constraints such as consistency of distances between the points in each tuple being matched. Doing so helps to reject many false matches, greatly improving the matching output. Feature similarity and satisfaction of constraints may in general be expressed in terms of an affinity tensor relating pairs of point tuples.

As a particular example of *quadratic* assignment, Leordanu and Hebert [30] consider pairs of feature de-

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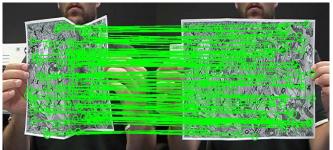


Fig. 1. Correspondences between datasets determined by SuperMatching, using feature points created simply by uniform sampling rigid scans on the left, and SIFT feature points on a deformable surface on the right. For clarity, only some representative matches are shown.

scriptors, and use *distances* between pairs of features from each set to reduce the number of incorrect correspondences. Such pairwise distance constraints are particularly helpful in cases when the features themselves have low discriminative ability. The idea has been widely adopted in 3D shape matching algorithms [13], [14], [15], [23], [16], [17].

Higher-order assignment includes yet more complex constraints between features. For example, third-order potential functions, used in [31], [32], [6], [33], [34], quantify the affinity between two point triples by measuring the similarity of the angles of the triangles formed by such triples. However, this angular similarity value only considers the *total* difference in corresponding angles, and does not change with reordering of elements in the tuple. When similarity is expressed in this way, the affinity tensor becomes a *supersymmetric* tensor [35].

Our *SuperMatching* algorithm also formulates higherorder matching problems using a supersymmetric affinity tensor. It can accurately match a moderate number of features (several hundreds) using triples or larger tuples of features. The contributions of this paper include:

- We show how to define a compact higher-order supersymmetric affinity tensor to express geometrically consistent constraints between feature tuples.
- Complete computation of the full affinity tensor is computationally infeasible. We efficiently estimate it using a sampling strategy which takes advantage of supersymmetry. This avoids sampling repetitive items, it allows the tensor to be stored compactly, and also improves the matching accuracy by avoiding imbalances in sampling.
- We make full use of the compactness of the affinity tensor to deduce a power iteration method which efficiently solves the matching problem.

Our experiments using both synthetic and real captured data sets show that SuperMatching is more accurate and robust than prior methods, yet has similar computational cost. Importantly, it is a general matching approach, independent of choice of 2D or 3D feature descriptors.

## 2 RELATED WORK

Previous approaches to feature matching can be classified into those which match single points to single points, those which match point pairs to point pairs, and so on.

Matching single points to single points is a linear assignment problem which only considers an affinity measure between two features, one from each set being matched. The affinity measure is typically defined as the feature distance between the feature vectors, which in turn are based on local information around each feature point, e.g. SIFT [25], spin images [24], heat diffusion signatures [26], and BRISK [28]. Point-to-point matching can give misleading results as wrong correspondences are readily established.

Matching point pairs in one set to point pairs in the other set leads to a quadratic assignment problem. Such methods now take into account both similarity of the point features, *and* either the Euclidean distance between the points in a pair, assuming the two sets of points are related by a rigid transformation [30], [36], or the geodesic distance, assuming isometry [4], [13], [14], [15], [16], [17]. Unfortunately, this quadratic assignment problem is NP-hard, and again, matches found are not always reliable.

Several higher-order approaches have also been proposed. While they can significantly improve matching accuracy, higher-order assignment is even more computationally demanding, so various approximate methods have been developed. [37] considered a probabilistic model of soft hypergraph matching. They reduce the higher-order problem to a first-order one by marginalizing the higher-order tensor to a one dimensional probability vector. [31], [32] introduced the use of a third-order tensor in place of an affinity matrix to represent affinities of feature triples, and

higher-order power iteration was used to achieve the final matching. [34] built a unified multiple higherorder affinity tensor, by extending the third-order tensor method [31], [32]. [33] treated the tensor as a joint probability of assignments, marginalized the affinity tensor to a matrix, and found optimal soft assignments by eigendecomposition of the matrix. Higherorder assignment problems typically require large amounts of memory and computational resources. By reducing the number of elements needed to represent the affinity measures, the above approaches can match moderate numbers (many hundreds) of features. However, these 2D approaches do not really take advantage of supersymmetry of the affinity tensor. SuperMatching does so, leading to an improvement in matching accuracy. 3D problems are even more challenging.

A related idea also using higher order constraints for 3D registration is the 4-points congruent sets method (4PCS) proposed in [3]. It is a fast alignment scheme for 3D point sets that uses widely separated points. However, the need to find coplanar 4-tuples of points and the assumption of rigid transformation limit its applicability. We solve both rigid and isometric shape matching problems with a single approach.

## 3 OVERVIEW

A tensor generalizes vectors and matrices to higher dimensions: a vector is a tensor of order one, and a matrix is a tensor of order two. A higher-order tensor can be expressed as a multi-dimensional array [38]. More formally, an  $N^{th}$ -order tensor is an element of the tensor product of N vector spaces, each with its own coordinate system.

Assume we are given two sets of feature points  $P_1$  and  $P_2$ , with  $N_1$  and  $N_2$  points respectively.  $i_n = (f_{i_n}^1, f_{i_n}^2)$  is a pair of points from  $P_1$  and  $P_2$ , respectively. Matching between these two feature sets can be represented by an assignment variable x which is a vector  $\in \{0,1\}^{N_1N_2}$ , with each element representing whether a pair  $i_n(f_{i_n}^1, f_{i_n}^2)$  is selected in the matching (if  $x_{i_n} = 1$ ) or not (if  $x_{i_n} = 0$ ). From the  $N^{th}$ -order tensor viewpoint, the matching problem is equivalent to finding the optimal assignment tensor  $x^* \in \{0,1\}^{N_1N_2}$ , satisfying [38]

$$\boldsymbol{x}^* = \arg\max_{\boldsymbol{x}} \sum_{i_1, \dots, i_N} \mathcal{T}_N(i_1, \dots, i_N) x_{i_1} \dots x_{i_N}.$$
 (1)

Here,  $i_n \in \{i_1, \cdots, i_N\}$  stands for an assignment in the  $n^{th}$  dimension of the N vector spaces. Let all feature tuples for  $P_1$  and  $P_2$  be  $F_1$  and  $F_2$ , then  $\forall (f^1_{i_1}, \cdots, f^1_{i_N}) \in F_1$ , there is a matching to corresponding feature tuples in  $F_2$ . For example, given a third-order tensor,  $i_n \in \{1,2,3\}$ , each index could be expressed as  $i_1 = (f^1_{i_1}, f^2_{i_1}), i_2 = (f^1_{i_2}, f^2_{i_2}), i_3 =$ 

 $(f_{i_3}^1,f_{i_3}^2)$ : pairs of potentially matched points. The product  $x_{i_1}\cdots x_{i_N}$  will be equal to 1 if the points  $(f_{i_1}^1,\cdots,f_{i_N}^1)$  are matched to the points  $(f_{i_1}^2,\cdots,f_{i_N}^2)$ , and otherwise 0.  $\mathcal{T}_N(i_1,\cdots,i_N)$  is the affinity of the set of assignments  $\{i_n\}_{n=1}^N$ , which is high if the features in tuple  $(f_{i_1}^1,\cdots,f_{i_N}^1)$  have similar descriptor values to the features in the tuple  $(f_{i_1}^2,\cdots,f_{i_N}^2)$ , and their distance constraints are similar. Note that the size of  $\mathcal{T}_N(i_1,\cdots,i_N)$  is  $(N_1N_2)^N$ . In this paper, the affinity measures expressing similarity of feature tuples are compactly represented and efficiently computed by using the supersymmetric tensor.

In the rest of the paper, we consider the one-to-many correspondence problem. We assume that each point in  $P_1$  is matched to exactly one point in  $P_2$ , but that the reverse is not necessarily true. If we *do* want to treat both datasets in the same way, we can first match  $P_1$  to  $P_2$ , then match  $P_2$  to  $P_1$ , and then combine the matching results by taking their union or intersection.

From Equ.(1) we can see that there are four issues to be considered when using higher-order matching algorithms. How should we:

- organize and express the affinity measures  $\mathcal{T}_N$  in a supersymmetric manner? (see Section 4.1)
- approximately solve the optimal higher-order assignment problem efficiently? (see Section 4.2)
- determine an appropriate sampling strategy to estimate the affinity tensor in a way which will give good matching accuracy (it is too large to compute fully)? (see Section 4.3)
- define a suitable affinity measure between two feature tuples? (see Section 4.4)

#### 4 SUPERMATCHING

We now discuss the first two issues mentioned above, which are independent of application; later we turn to sampling strategy and definition of affinity measure, which is application dependent.

#### 4.1 Supersymmetric Affinity Tensor

The supersymmetric higher-order affinity tensor is invariant under permutation of indices. The main motivation of using supersymmetry is to allow us to avoid redundant storage and computation.

Definition 1 (Supersymmetric Tensor): A tensor is supersymmetric if its entries are invariant under any permutation of its indices [35].

For example, a third-order supersymmetric tensor  $\mathcal{T}_3$ , satisfies the relationships:  $\mathcal{T}_3(i_1,i_2,i_3)=\mathcal{T}_3(i_1,i_3,i_2)=\mathcal{T}_3(i_2,i_1,i_3)=\mathcal{T}_3(i_2,i_3,i_1)=\mathcal{T}_3(i_3,i_1,i_2)=\mathcal{T}_3(i_3,i_2,i_1).$ 

Definition 2 (Supersymmetric Affinity Tensor): Given two feature sets  $P_1$  and  $P_2$ , with  $N_1$  and  $N_2$  features respectively, the supersymmetric affinity tensor is an  $N^{th}$  order  $I_1, \dots, I_N$ , nonnegative tensor  $\mathcal{T}_N$ , for which there exists a set of indices  $\theta_N$ , and an  $N^{th}$  order potential function  $\phi_N$ , such that

$$\mathcal{T}_{N}(i_{1},\ldots,i_{N}) = \begin{cases} \phi_{N}(\Omega(i_{1},\ldots,i_{N})) &, \ \forall (i_{1},\ldots,i_{N}) \in \theta_{N} \overset{3:}{4:} \\ 0 &, \ \forall (i_{1},\ldots,i_{N}) \notin \theta_{N} \overset{3:}{5:} \end{cases}$$
(2)

where  $\Omega$  stands for an arbitrary permutation of the vector.  $\theta_N$  satisfies  $i_m \neq i_n \ \forall (i_1,\ldots,i_N) \in \theta_N, \forall i_m \in \{i_1,\ldots,i_N\}$  and  $\forall i_n \in \{i_1,\ldots,i_N\} - \{i_m\}$ .

A tensor element with  $(i_1, \ldots, i_N) \in \theta_N$  is called a *potential element*, while other elements are called *non-potential elements*. A potential element represents one matching result out of all possible matching candidates. Potential elements are further detailed in Section 4.3.

Using Definition 2, we can reduce the amount of storage needed, by representing every potential element  $\mathcal{T}_N(i_1,\ldots,i_N)$  by its canonical entry  $\mathcal{T}_N(\operatorname{sort}(i_1,\ldots,i_N))$ ,  $\forall (i_1,\ldots,i_N) \in \theta_N$ . Each stored value thus provides the value for N! entries. Furthermore, as non-potential elements all have value zero, there is no need to store them. This greatly reduces both storage, and the amount of feature tuple sampling needed when estimating the affinity tensor, as discussed in Section 4.3. At the same time, it can be used to make the power iteration process more efficient: see Section 4.2.

## 4.2 Supersymmetric Higher-order Power Iteration

The higher-order tensor problem in Eq. (1) may be solved by tensor decomposition [38]; tensor decomposition originated in [39]. We utilize the rank-one higher-order power method [40] to approximately solve Eq. (1); as noted, an exact computation is infeasible. Eq. (1) can be expressed as:

$$\boldsymbol{x}^* = \underset{\boldsymbol{x}}{\operatorname{arg max}} \sum_{i_1, i_2, \dots, i_N} \mathcal{T}_N(i_1, \dots, i_N) x_{i_1} \dots x_{i_N}$$
$$= \underset{\boldsymbol{x}}{\operatorname{max}} \langle \mathcal{T}_N, \boldsymbol{x}^{\star N} \rangle$$
(3)

where  $\boldsymbol{x} \in \{0,1\}^N$ , and  $\star$  is called the Tucker product [35], which is a succession of mode-N products. To get an approximate solution, we relax the constraints: the binary assignment vector  $\boldsymbol{x} \in \{0,1\}^N$  is replaced by an assignment vector  $\boldsymbol{u}$  with elements taking real values in [0,1]. This changes the optimization problem to one of computing the rank-one approximation of the affinity tensor  $\mathcal{T}_N$  [35], i.e. finding a scalar  $\lambda$  and a unit norm vector  $\boldsymbol{u} \in \mathbb{R}^N$ , such that the tensor  $\hat{\mathcal{T}_N} = \lambda \boldsymbol{u} \star \boldsymbol{u} \star \dots \star \boldsymbol{u} = \boldsymbol{u}^{\star N}$  minimizes the Frobenius norm squared function  $f(\hat{\mathcal{T}_N}) = \|\mathcal{T}_r - \hat{\mathcal{T}_N}\|_F^2$ . The final

**Algorithm 1** Higher-order power iteration solution (with  $\ell^1$  norm) for the

supersymmetric affinity tensor

```
Input: N^{th}-order supersymmetric affinity tensor
Output: Unit \ell^1-norm vector \boldsymbol{u}
         1: Initialize v_0 to random values in [0,1], k=1
         2: repeat
                                            for all (i_1, \dots, i_N) \in \theta_N do
                                                              for all m \in (i_1, \cdots, i_N) do
                                                                            v_m^{(k)} = (N-1)! \phi_N(i_1, \cdots, i_N) 2v_m^{(k-1)} v_{i_1}^{2(k-1)} \cdots v_{m-1}^{2(k-1)} v_{m+1}^{2(k-1)} \cdots v_{i_N}^{2(k-1)}
         7:
                                                              for i = 1 : N_1 do
                                                                              v^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2) = \hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1).\|\hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1):i\cdot N_2)/\|\hat{v}^{(k)}(((i-1)\cdot N_2+1
         9:
  10:
                                             k = k + 1;
  12: until convergence;
  13: \boldsymbol{u}^{(k)} = \boldsymbol{v}^{2(k)}
                            Note: u^{(k)} = v^{2(k)}, and v^{(k)}(((i-1) \cdot N_2 + 1) : i \cdot N_2)
                            denotes the slice of v^{(k)} with indices from (i-1) \cdot N_2 + 1
                            to i \cdot N_2.
```

matching result is found by replacing each element of u by 0 or 1 according to whichever it is closer to.

The higher-order power method is commonly used for finding rank-one tensor approximations; a version for supersymmetric tensors (S-HOPM) is given in [35]. The S-HOPM algorithm converges under the assumption of convexity (or concavity) for the functional induced by the tensor [35], which is sufficiently robust for practical applications. S-HOPM is performed in two iterative steps: higher-order power iteration of u, followed by normalization of u under the Frobenius norm. A recent effective improvement [31], [32] uses the  $\ell^1$  norm to replace the traditional  $\ell^2$  norm.

We use the  $\ell^1$  norm, and further revise S-HOPM as follows. To perform higher-order power iteration of  $\boldsymbol{u}$ , we must compute  $\hat{\boldsymbol{u}}^{(k)} = \mathcal{I}^{\mathcal{T}_N}_{\star}(\boldsymbol{u}^{(k-1)})^{\mathcal{T}_N}_{\star}^{(N-1)}$ , where  $\boldsymbol{\dot{\tau}}^N$  is a so-called  $\mathcal{T}_N$ -product, and  $\mathcal{I}$  is the unit tensor [35]. For  $\hat{\boldsymbol{u}}^{(k)}$  belonging to an  $N^{th}$ -order supersymmetric affinity tensor, this can be formulated as follows:

$$\begin{split} \hat{\boldsymbol{u}}^{(k)} &= \mathcal{I} \overset{\mathcal{T}_{N}}{\star} (\boldsymbol{u}^{(k-1)})^{\overset{\mathcal{T}_{N}}{\star} (N-1)} \text{ implies that } \forall m \in (i_{1},...,i_{N}), \\ v_{m}^{(k)} &= \\ &\sum_{i_{1},...,i_{N}} \mathcal{T}_{N}(i_{1},...,i_{N}) 2v_{m}^{(k-1)} v_{i_{1}}^{2_{(k-1)}} ... v_{m-1}^{2_{(k-1)}} v_{m+1}^{2_{(k-1)}} ... v_{i_{N}}^{2_{(k-1)}} = \\ &(N-1)! \phi_{N}(i_{1},...,i_{N}) 2v_{m}^{(k-1)} v_{i_{1}}^{2_{(k-1)}} ... v_{m-1}^{2_{(k-1)}} v_{m+1}^{2_{(k-1)}} ... v_{i_{N}}^{2_{(k-1)}} \\ &(4) \end{split}$$

where  $u^{(k)} = v^{2_{(k)}}$ , and  $\phi_N$  is the potential function explained in Section 4.4. Eq. (4) is more compact than earlier expressions in the literature, as it handles all

symmetrically related potential elements as a single item using multiplication by (N-1)!.

Many initialization schemes have been proposed for the S-HOPM method [35]. We simply use positive random values between 0 and 1 to initialize  $u_0$ , which ensures convergence; proofs are detailed in [41], [35].

Our supersymmetric higher-order power iteration solution of Eq. (1) is performed by the SuperMatching algorithm—see Algorithm 1. Its efficiency relies on two principles. First, we take advantage of the supersymmetry to deduce u as in Eq. (4), using just a single canonical element for computation (see Step 5). Secondly, power iteration just consides the non-zero potential elements, and excludes each non-potential element from the iteration process. The complexity of the whole iteration process depends only on the number  $|\theta_N|$  of non-zero affinities. Consequently, this method reduces also memory costs while keeping accuracy.

Note that, although [31], [32] claimed to use a supersymmetric affinity tensor, their approach does not make full use of supersymmetry when creating the supersymmetric affinity tensor, nor does it take advantage of supersymmetry to accelerate the power iteration process. By doing so, we overcome limitations due to unbalanced and redundant tensor elements in [31], [32], as our experiments show later.

#### 4.3 Sampling Strategy

Algorithm 1 depends on the potential elements. We next discuss the issue of how to sample the feature tuples to build potential items, which determines the size  $|\theta_N|$  and influences matching accuracy (and speed).

Given the two feature sets  $P_1$  and  $P_2$ , a potential element may be obtained by using a feature tuple sampled from each feature set separately. For  $N^{th}$ -order matching, a naive way to construct the potential elements is as follows: first find all feature tuples for  $P_1$  and  $P_2$ , as  $F_1$  and  $F_2$ ; then  $\forall (f_{i_1}^1, \cdots, f_{i_N}^1) \in F_1$ , calculate the potentials for  $(f_{i_1}^1, \cdots, f_{i_N}^1)$  with all feature tuples in  $F_2$ . This is far too time-consuming, so sampling is used instead. We suggest random sampling for general feature matching problems, but this does not preclude more directed sampling if prior knowledge of the matching problems gives guidance.

Our sampling approach is to repeatedly randomly sample  $t_1$  feature tuples for each feature point from  $P_1$ , and fully sample  $P_2$ . For  $P_1$ , we take each feature in turn as a required element, and then randomly choose  $t_1$  feature tuples containing this required element. Thus, the number of feature tuples in  $F_1$  is  $N_1t_1$ , and  $N_2^N$  in  $F_2$ . Then, for each feature tuple in  $F_1$ , we find the k most similar feature tuples in  $F_2$  to build

k potential elements as  $\phi_i^k$ . Combining all the potential elements obtained, we form the desired potential element set  $\theta_N = \{\phi_i^k\}_{i=1}^{N_1t_1}$ , of size  $|\theta_N| = N_1t_1k$ . For  $P_1$ , the sampling cost is  $O(N_1t_1k\log N_2)$ , where the  $\log N_2$  arises from use of a KD-tree to search for the k most similar feature tuples in  $F_2$ . The parameters  $t_1$  and k must be chosen according to the size of the feature sets. In practice, for two feature sets each with hundreds points, we may take  $t_1 \approx 100$  and  $k \approx 300$  for third-order matching. Our experiments demonstrate that this sampling approach works well.

An important aspect of our sampling approach is to use the supersymmetry of the affinity tensor. Potential elements whose indices are permutations of each other have the same value, so should not be repeatedly sampled. Thus, we use a sampling constraint that the sets of feature tuples  $F_1$  obtained from the sampling process should have no repetition, in the sense that

$$\forall (f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1), (f_{j_1}^1, f_{j_2}^1, \cdots, f_{j_N}^1) \in F_1,$$

$$(f_{i_1}^1, f_{i_2}^1, \cdots, f_{i_N}^1) \neq \Omega(f_{j_1}^1, f_{j_2}^1, \cdots, f_{j_N}^1)$$

$$(5)$$

where  $\Omega$  is an arbitrary permutation.

Earlier work [37], [31], [34] adopted random sampling, but failed to impose any constraint on the sampling process to take into account supersymmetry, leading to the possibility that feature tuples may be sampled multiple times. For example, for third-order matching, it is possible that a feature tuple  $(f_{i_1}^1, f_{i_2}^1, f_{i_3}^1)$  may be sampled from  $P_1$  and  $(f_{i_1}^2, f_{i_2}^2, f_{i_3}^2)$  from  $P_2$ , and also a feature tuple  $(f_{i_1}^1, f_{i_3}^1, f_{i_2}^1)$  from  $P_1$  and  $(f_{i_1}^2, f_{i_3}^2, f_{i_2}^2)$ from  $P_2$ . That would create two tensor elements  $\phi_3(i_1, i_2, i_3)$  with index  $(i_1, i_2, i_3)$  and  $\phi_3(i_1, i_3, i_2)$  with index  $(i_1, i_3, i_2)$ , which are the same. However, we just need one tensor element to express the affinity on the assignment group  $(i_1, i_2, i_3)$  for any permutation of indices. This extra sampling is not only inefficient, but may also reduce the accuracy of the power iteration: one set of symmetrically related elements may be represented by a different number of samples than another set of symmetrically related elements, which unbalances the power iteration process, and can lead to inaccurate results. Our sampling method reduces the sampling cost, while also improving the accuracy of the power iteration.

#### 4.4 Higher-order Potentials

Different higher-order potentials are appropriate for different applications. Here we briefly give two simple examples of general higher-order potentials for 2D and 3D matching respectively; we use them later to evaluate our algorithm. The potentials are based on a Gaussian function which guarantees the tensor elements are non-negative and invariant under any permutation of the input assignments.

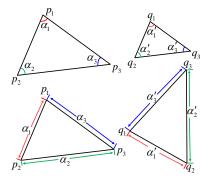


Fig. 2. Third-order potential. The geometric constraints are: internal angle invariance in 2D (above), and edge length invariance in 3D (below).

In 2D, we use a well-known third-order geometric-similarity invariant potential  $\phi_3$  [31], [32], [33] for linking point feature triples. Triangles formed by three points are *similar* under scaling, rotation and translation—interior angles are invariant. Thus  $\phi_3$  can be defined in terms of differences of corresponding interior angles:

$$\phi_{3}(i_{1}, i_{2}, i_{3}) = \phi_{3}(\{p_{1}, q_{1}\}, \{p_{2}, q_{2}\}, \{p_{3}, q_{3}\})$$

$$= \exp(-1/\varepsilon^{2} \sum_{(l, l')} \|\alpha_{l} - \alpha_{l'}\|^{2})$$
(6)

where  $\varepsilon>0$  is the kernel bandwidth, which is automatically determined as the average of the  $\ell^1$  norm of all differences, and  $\{\alpha_l\}_{l=1}^3$  and  $\{\alpha_l'\}_{l'=1'}^3$  are the angles formed by feature triples  $(p_1,p_2,p_3)$  and  $(q_1,q_2,q_3)$ : see Figure 2. Each point corresponds to one interior angle. We may extend it to higher order by using the internal angles formed by polygons with more than 3 sides.

For 3D matching problems, we may replace the internal angles by edge lengths, which for meshes are based on geodesic distance across the mesh in which the points are embedded. This corresponds to assuming an isometry transform relating the meshes. Geodesic distance may be computed by Dijkstra's algorithm [42]. See Figure 2.

#### 5 EXPERIMENTS

We have used synthetic data and real captured data to evaluate the SuperMatching algorithm. To demonstrate the SuperMatching algorithm's independence of feature descriptors, several descriptors were used. In some cases, we simply uniformly sampled feature points on the objects, and used a trivial feature descriptor of 1 for all points, meaning affinities are based simply on distances between feature points—this allows us to show our method is robust in the presence of ambiguous feature descriptors. In other cases, we still used uniformly distributed feature points, together with SIFT descriptors, which shows that



Fig. 3. Pairwise alignment of Coati and Dragon scans. From left to right: before alignment, our matching result, our alignment result, alignment result from [3].

feature points do not have to be carefully chosen. *Note that, we only use the feature descriptors to select the feature points, and do not employ them to build the matching.* Both the uniform and SIFT feature are not selective as they are widely used in the applications. The universal-applied descriptors proves that the SuperMatching is independent to the descriptors in some degree. The feature descriptors could be replaced by other types, especially proposed in recent years. We used third-order matching in our experiments, but it would be simple (but more costly) to use higher order.

# 5.1 3D Rigid Shapes Scans

Firstly, we use SuperMatching to build pairwise matchings between 3D rigid shape scans based on uniformly sampled feature points. Rigid transforms can be computed from each triple of compatible matching points. The transform which brings the most data points within a threshold distance of a point in the model is chosen as the optimal alignment transform [20]. As discussed in [2], such a voting scheme is guaranteed to find the optimal alignment between pairwise scans and is robust to the initial pose of the input scans.

To test the robustness of the SuperMatching, the noisy shapes used by the STAR rigid registration 4PCS algorithm [3] are all examined. The results demonstrate that the SuperMatching could produce the desirable registration for all of these shapes. For example, Figure 3 shows the initial state of input, the registration results by the SuperMatching and 4PCS [3]. Even from the Coati and Dragon examples, it is noticeable that the final alignment details from the SuperMatching is some better than the 4PCS.

Figure 4 shows some registration results for the Rooster model from [43]. The left column shows the original state, the two middle columns are our matching and alignment results, and the right column shows the corresponding result produced by 4PCS [3], which has clearly failed. As analyzed by our, the main failure reason of 4PCS was that it did not use the correlation among each 4 congruent points.

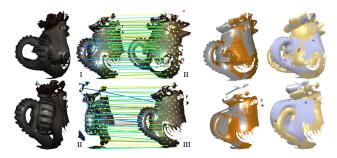


Fig. 4. Pairwise alignment of Rooster I-II and II-III scans. From left to right: before alignment, our matching result, our alignment result, alignment result from [3].

TABLE 1
Transformation matrix comparison

			d matrix		ground truth				
						-0.0043			
						0.9998			
II	0.6928	0.0206	0.70021	-50.05	0.7068	0.02067	0.7071	-54.1	
	0	0	0	1	0	0	0	1	
						-0.0220			
						0.9992			
Ш	1.0166	0.0362	0.0052	-176.6	0.9995	0.0322	0.0000	-186.0	
	0	0	0	1	0	0	0	1	

But the SuperMatching uses the global optimization scheme of all higher-order feature tuples at one time to compute the matching, as formulated by Eq. 1. The transformation matrix computed by SuperMatching is compared with the ground truth provided by [43] in Table 1, which demonstrates that our computed matrices for matching I-II and I-III are very close to the ground truth. Note that the matrices are used to transform and alignment from II to I and III to I transformations.

Next, we used SuperMatching to build a complete model from a set of scans from different viewpoints. For these multiple scans, third-order matching was first performed between each pair of consecutive scans. After doing so, the alignment was refined using the iterative closest point (ICP) algorithm [1]. Figure 5 illustrates the approach and shows the result.

#### 5.2 3D Depth Scans with Color Information

We next provide a further real-world, noisy, example of the use of SuperMatching. In this case, data with surface color information was captured using a Kinect camera [44]. Uniform samples and SIFT feature vectors were used as a basis for SuperMatching. This resulted in robust matches, as illustrated in Figure 6.

## 5.3 3D Articulated Shape Synthetic Data

A further application is registration of (approximately) articulated shapes. Such problems are com-

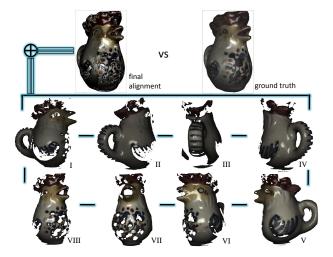


Fig. 5. Alignment of several Rooster scans from different viewpoints. Above: our final registered Rooster and the ground truth [43]. Below: 8 partial scans, the dark lines indicating the pairwise matches.

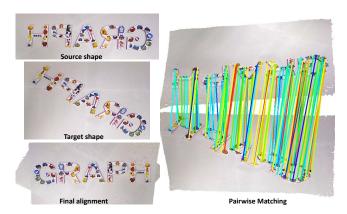


Fig. 6. 3D real depth scans with color information, captured using Kinect. Matching points are connected by colored lines.

mon in dynamic range scanning such as human motion capture. Given a sequence of range scans of a moving articulated subject, our method automatically registers all data to produce a complete 3D shape. Unlike many other methods, we do not need manual segmentation, markers, or a prior template. While the problem of non-rigid registration of deformable shapes is ill-posed and no algorithm is applicable to all scenarios, we believe that our approach pushes the limits of what can be achieved with minimal prior information. SuperMatching provides robust, accurate matching, even although the partial scans have holes and different poses.

Again uniformly sampled points were used. Registration of scans was performed by computing piecewise rigid transformations between matches. These transformations may be propagated from feature points to the entire set of points in each scan using nearest neighbor interpolation. Figure 7 shows a registration

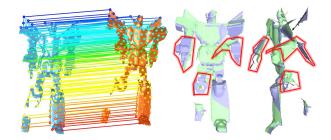


Fig. 7. Pairwise matching of an articulated Robot between two frames. Left: our result. Right: result produced by [5], from front and side views; red polygons indicate regions of large distortion.

example for an articulated model. On the left is our result, on the right is the result produced by the method [5].

## 5.4 Deformable Surfaces

Finally, we matched SIFT points on images of deforming surfaces<sup>1</sup> showing a cloth and a cushion. The surface of the cloth underwent relatively smooth deformation, while the surface of the cushion included sharp folds. This data comes with ground truth, which allows quantitative verification of the accuracy of the matches found. From each surface set we randomly chose two frames before and after a large deformation. We randomly chose 100 corresponding points on each surface, using the provided ground truth.

We used the above input data as a basis for comparison with the spectral algorithm [36] (a quadratic assignment algorithm), a third-order tensor algorithm [31], [32], and the hypergraph matching algorithm [37], using the authors' code in each case. All methods were executed in Matlab on a 2.3GHz Core2Duo with 2GB memory. To enable direct and fair comparison, [31], [32], [37] and SuperMatching used the same potential and the same tensor size.

In these tests, SuperMatching considered  $3 \times 10^6$  feature tuples, while the method of [31], [32] considered  $10 \times 10^6$  tuples and the method of [37] used  $4 \times 10^6$ . The difference mainly results from differences in sampling strategy; note that we have the lowest sampling cost. The average running time to match two feature sets each with 100 features was around 8s for SuperMatching, 13s for [31], 6.5s for [37], and 5s for [36]. SuperMatching takes less time than the third-order tensor algorithm in [31], [32] both because it uses fewer feature tuples and because of the more efficient supersymmetric higher-order power iteration solution.

Matching accuracy was assessed by the number of correctly matched points (known from the ground

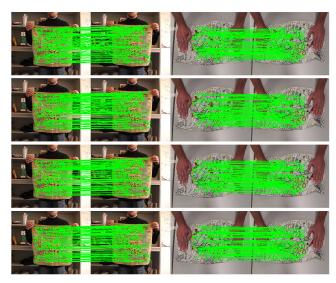


Fig. 8. Matching results. Left: cloth set, matching between frame 80 and 90, right: cushion set, matching between 144 and 156. Top to bottom, spectral method [36], hypergraph matching method [37], a third-order tensor method [31], [32], and SuperMatching.

TABLE 2
Accuracy of deformable surface matching.

Dataset	cloth				cushion				
Matching	F80-	F90-	F95-	F100-	F144-	F156-	F165-	F172-	Time
frames	F90	F95	F100	F105	F156	F165	F172	F188	(s)
Ours	83%	85%	84%	81%	66%	60%	69%	56%	8
[37]	73%	79%	70%	72%	44%	39%	54%	43%	6.5
[31], [32]	67%	77%	73%	65%	39%	31%	47%	42%	13
[36]	27%	29%	22%	27%	14%	5%	28%	7%	5

truth) divided by the total number of points that could be matched. The results are summarised in Table 2 and illustrated in Figure 8. Table 2 demonstrates that SuperMatching achieves a higher matching accuracy than previous algorithms. The worst matching result is produced by the spectral quadratic assignment algorithm [36], due to the lower discriminatory power of pairwise geometric constraints. Higher-order algorithms perform better due to the more complex geometric constraints. Nevertheless, SuperMatching also significantly outperforms the third-order algorithm [31], [32] and the hypergraph matching algorithm [37], as these do not tale proper advantage of supersymmetry.

## 6 CONCLUSIONS

This paper has presented the SuperMatching algorithm, which tackles the classic computer graphics and computer vision problem of feature matching, independently of feature description. It is an efficient higher-order matching algorithm which uses a compact form of the higher-order supersymmetric affinity tensor to express relatedness of features. Matching is

performed using an efficient power iteration method, which takes advantage of supersymmetry and avoids computing with zero elements. We also give an efficient sampling strategy for choosing feature tuples to create the affinity tensor. Experiments on both synthetic and real 2D and 3D data sets show that SuperMatching has greater accuracy than competing methods, whilst having competitive performance.

In future, we wish to further improve the performance of the SuperMatching algorithm. Random sampling could be executed in parallel. We also intend to apply it to more challenging deformable imperfect 3D data matching with real captured data. SuperMatching has applicability to many fields, as matching is a foundation for many computer graphics and computer vision applications.

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