

# Stream Join in Flink: from Discrete to Continuous

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## About Me



2017.02 - First Pull Request to Apache Flink



2018.05 – Flink Committer (mainly on SQL/Table API)



2018.06 – Ph.D. in CS from Shandong University (supervised by Prof. Xiaohui Yu)



2018.10 – Postdoc at York University



## Outline

#### 1. The Stream Join APIs

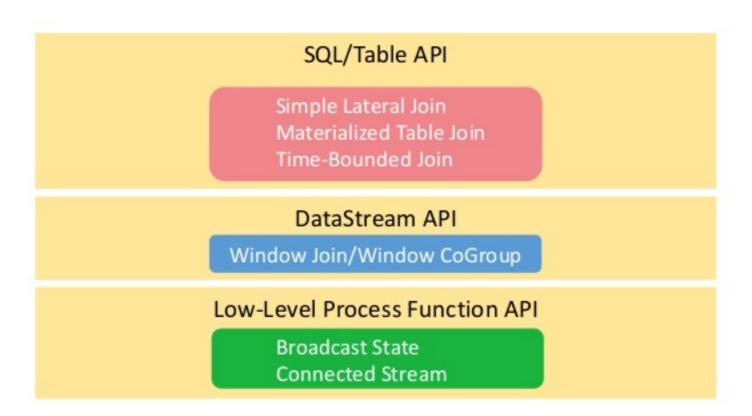
- Window Join in DataStream API
- Broadcast State in Low-Level API
- Joins in SQL/Table API

#### 2. Details about the Time-Bounded Join

- Why Time-Bounded
- How to Perform
- · What's More
- Future Work

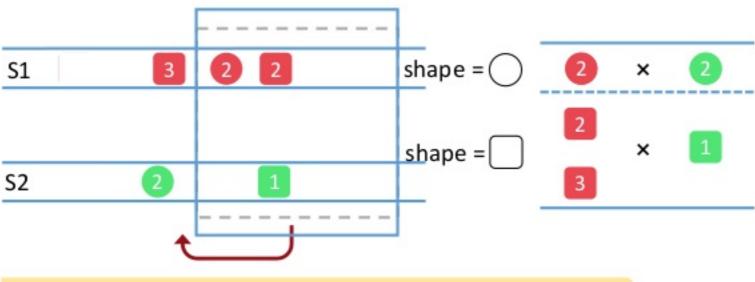


## Overview of the Stream Join APIs





## Window Join/Window CoGroup



```
S1.join(S2)
    .where(<shape1>).equalTo(<shape2>)
    .window(<window>)
    .apply {new JoinFunction()/FlatJoinFunction()};

S1.coGroup(S2)
    .where(<shape1>).equalTo(<shape2>)
    .window(<window>)
    .apply {new CoGroupFunction()};
```

#### JoinFunction

```
OUT join(
IN1 first,
IN2 second);
```

#### 22

#### <u>FlatJoinFunction</u>

```
void join(
    IN1 first,
    IN2 second,
    Collector<OUT> out);
```

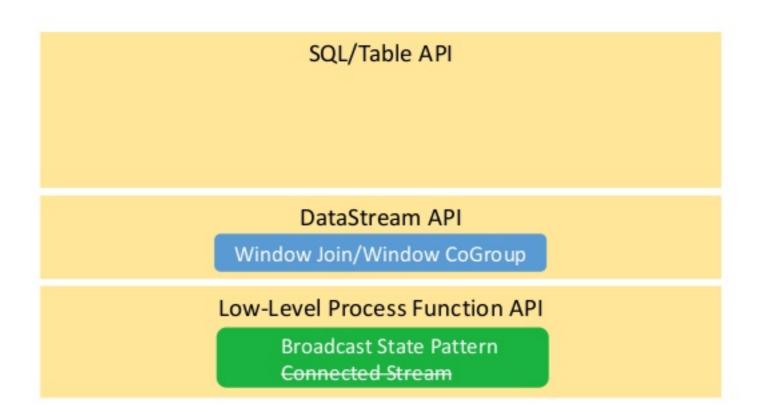
#### CoGroupFunction

```
void coGroup(
    Iterable<IN1> first,
    Iterable<IN2> second,
    Collector<OUT> out);
```

1

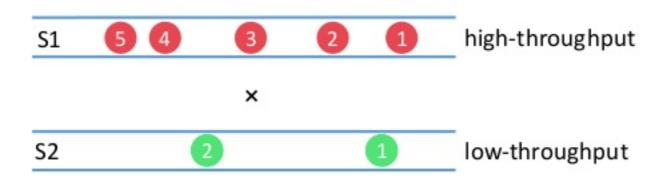


## Overview of the Stream Join APIs





## **Broadcast State**

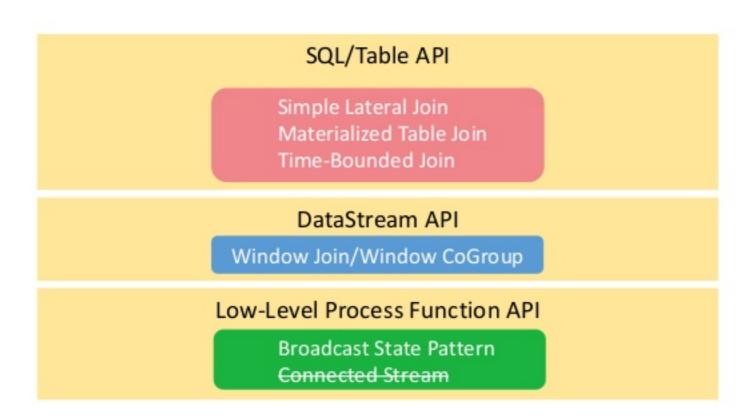


```
MapStateDescriptor state = new MapStateDescriptor();
BroadcastStream broadcast = S2.broadcast(state);
S1.keyBy(<key selector>);
    .connect(broadcast)
    .process(new KeyedBroadcastProcessFunction());
```

```
KeyedBroadcastProcessFunction
   IN1 value,
   ReadOnlyContext ctx,
   Collector<OUT> out)
   Context ctx,
   Collector<OUT> out)
processElement(...)
processBroadcastElement(...)
processElement(...)
                                 31
processBroadcastElement(...)
```

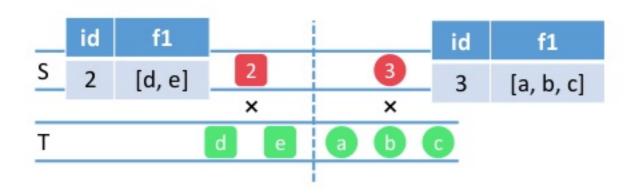


## Overview of the Stream Join APIs





# Simple Lateral Join

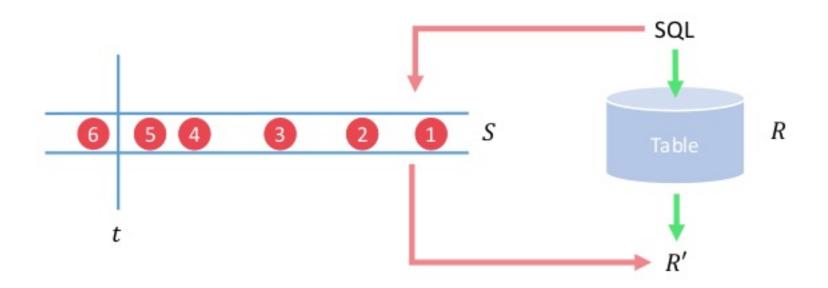


```
SELECT id, f2
FROM S, LATERAL TABLE(udtf(S.f1)) AS T(f2)

SELECT id, f2
FROM S, UNNEST(S1.f1) AS T(f2)
```



# SQL on Streams





## Stream to Table Conversion

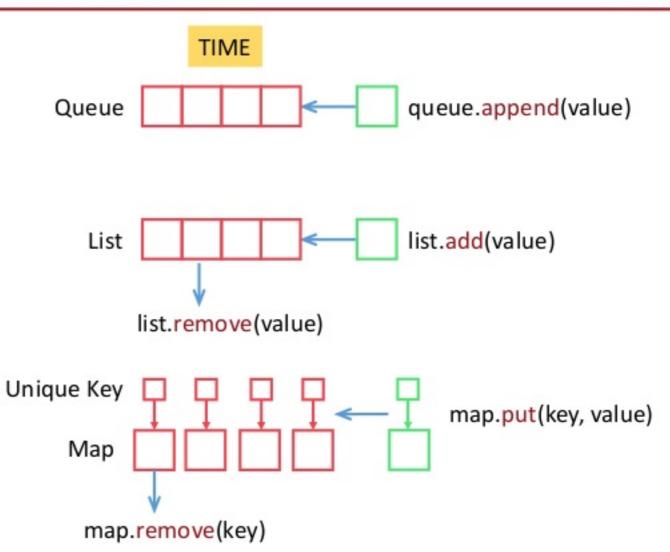
Append Stream

Time-Bounded Join

Retract Stream

Materialized Stream Join

Upsert Stream





#### Materialized Table Join

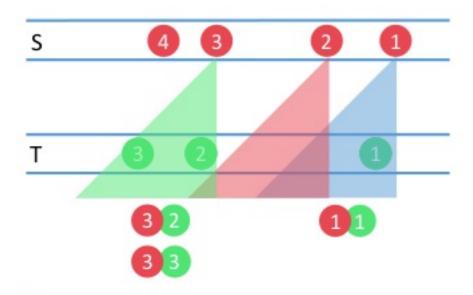
```
SELECT S.id, T.name
FROM S, T
WHERE S.id = T.id
```

```
State TTL Config
streamQueryConfig
.withIdleStateRetentionTime(Time.hours(1), Time.hours(2));
```

- 1. You know the table size will never exceed the state capacity (have duplications or rows will be retracted in time).
- The rows will be naturally expired after a period time.
- You don't really care about the accuracy of the join result (best-effort approach).



## Time-Bounded Join



```
SELECT S.id, T.time
FROM S, T
WHERE S.id = T.id AND
T.time BETWEEN S.time - INTERVAL '4' HOUR AND S.time
```



# Summary of the Stream Join APIs

Name	Description	Outer Support	Non- Equi	Conti- nuous	API Level	Launched Version
Window Join/ Window CoGroup	Cut two streams into chunks and join data in the same chunk.	-/LRF	N	N	DataStream	1.0
Connected Stream	Connect two streams and jointly process them with a CoProcessFunction.				ProcessFunction	1.2
Simple Lateral Join	Join a stream with an unnested array or a user-defined table function.	L	Υ	Υ	SQL/Table	1.2
Time-Bounded Join	Join two append only streams with a predicate that bounds the time on both sides.	LRF	N	Υ	SQL/Table	1.4
Retract Stream Join	Join two retract streams (with idle state retention time config).	LRF	N	Υ	SQL/Table	1.5
Broadcast State Pattern	Join a high-throughput stream with a low-throughput stream.	LRF	Υ	Υ	ProcessFunction	1.5



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# Append Stream Join

#### Time-Ordered Table

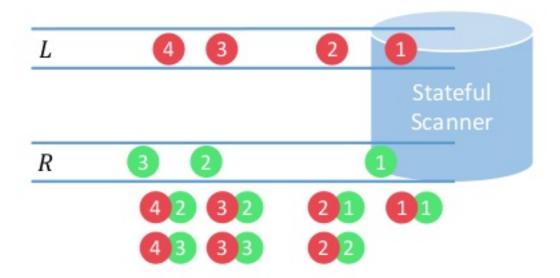
Append Stream = Sequential Scan

time	id	
3	Α	
8	S	

Join two time-ordered table with only one sequential scan on both sides.



## Make the Join Feasible

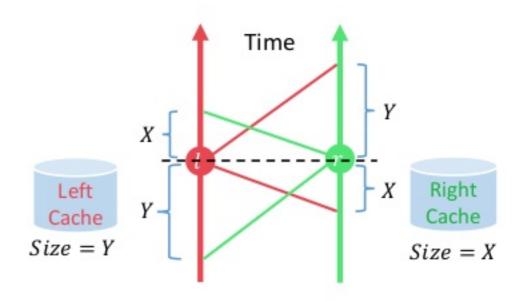


A sliding window that slides smoothly.



## The Predicates for Time-Bounded Join

```
SELECT L.id, R.time
FROM L, R
WHERE L.id = R.id AND
R.time BETWEEN L.time - X AND L.time + Y
L.time BETWEEN R.time - Y AND R.time + X
```



- 1. An equi-predicate to make the join operator scalable.
- 2. A time-bound predicate to keep the qualified records near.



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## The Connected Stream

```
KeyedCoProcessFunction
void processElement1(
    IN1 value,
   Context ctx,
   Collector<OUT> out)
void processElement2(
    IN2 value,
   Context ctx,
   Collector<OUT> out)
void onTimer(
    long timestamp,
   OnTimerContext ctx,
    Collector<OUT> out)
```

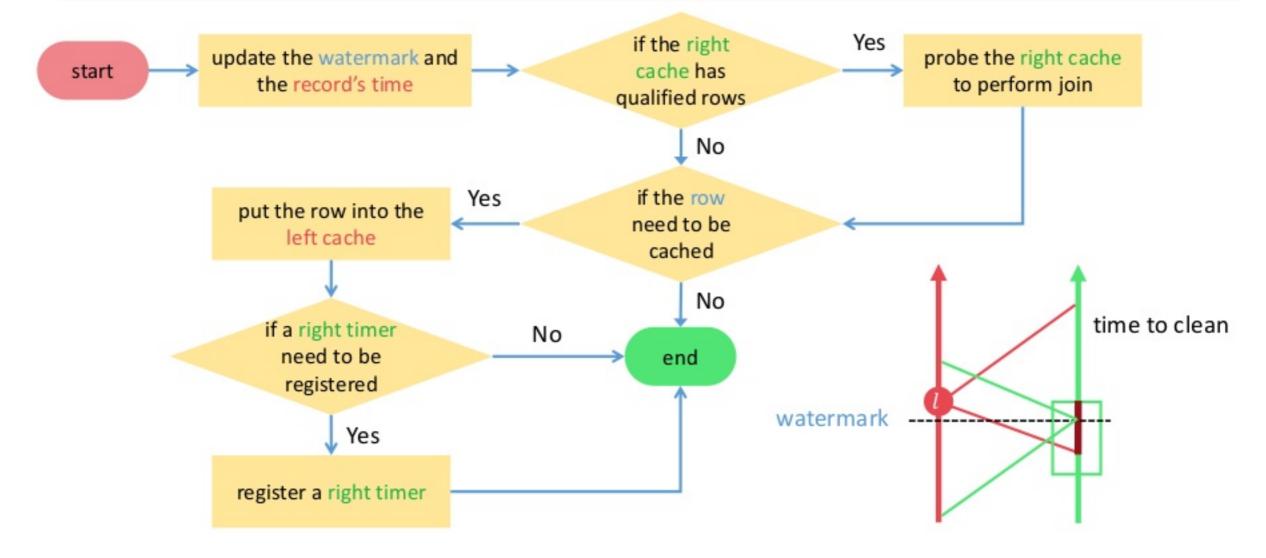
```
.connect(rightStream)
                 .keyBy(<key selector for the equi-predicate>)
                 .process(new KeyedCoProcessFunction())
record's time
                    watermark
                      cache
                      (state)
                 Map[Long, List[Row]]
right timer
                 time -> records of that time
```

leftStream

left timer

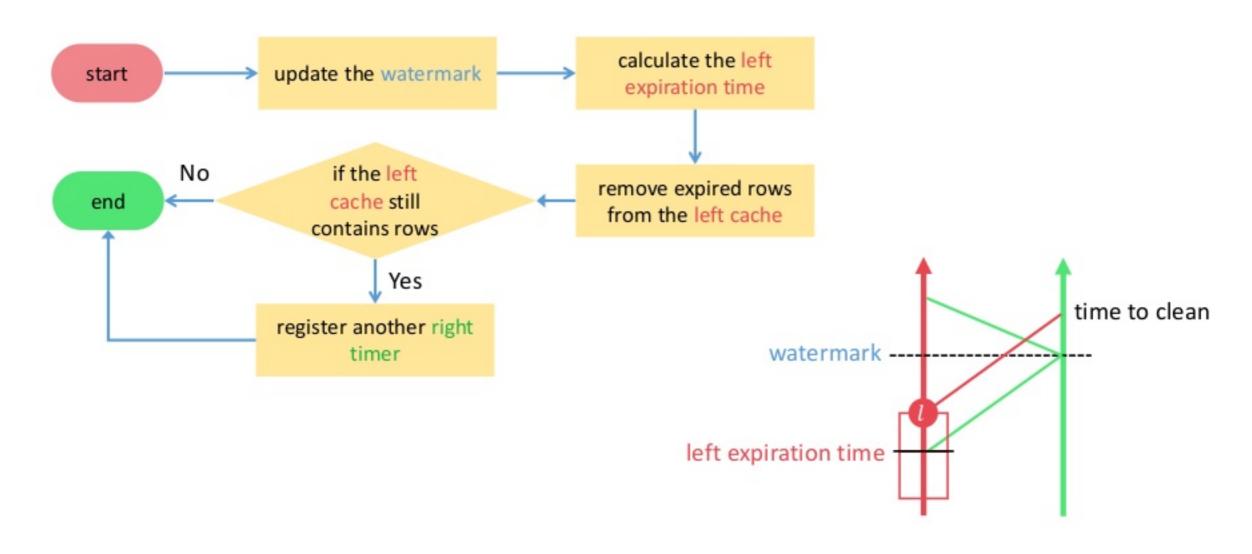


# The Logic for processElement()





# The Logic for onTimer()





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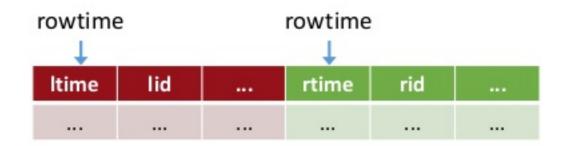
#### 3. Future Work



### About the Join Result



1. The result for a time-bounded join is still an append stream.

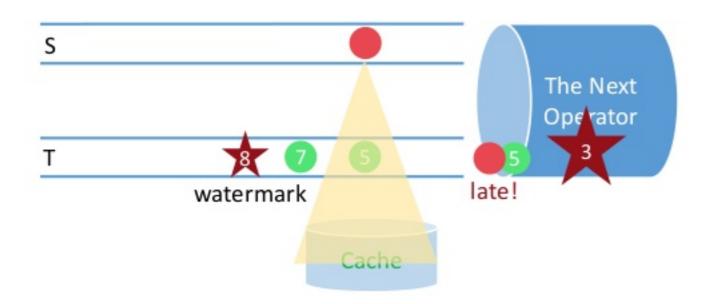


2. When converting the join result to a DataStream, at most one rowtime field can be reserved.

```
SELECT ltime, CAST(rtime AS TIMESTAMP), ... FROM ...
```



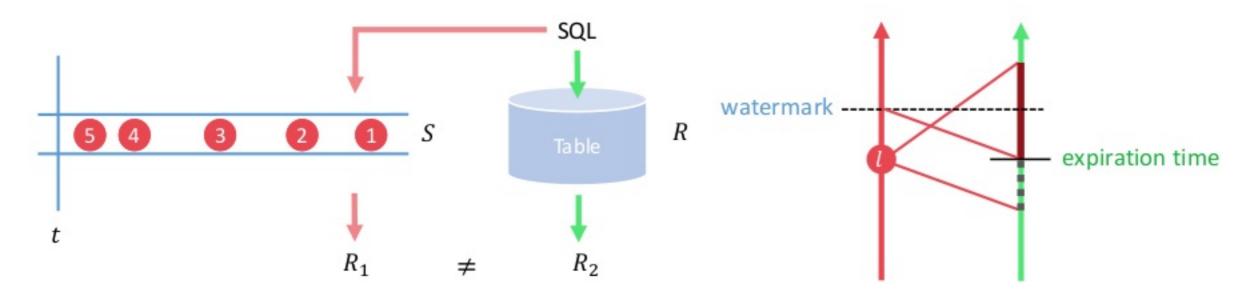
# Holding Back Watermarks



Subtract a fixed value from each watermark.



# The Integrity Problem



- A. Drop l as if it never comes (or use side output).
- B. Keep l to produce as many results as possible.

 $R_1$  could be affected by lateness, thus may not be identical with  $R_2$ .



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#### 3. Future Work



## P1. More Kinds of Stream Join

An efficient policy to clean the cache.

- 1. Time Versioned Join (WIP)
- 2. One-to-One Join (query hints?)



## P2. Better State

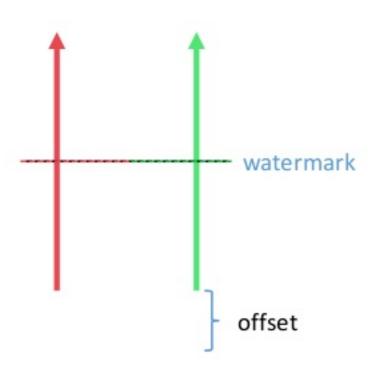
1. Sorted MapState

```
mapState.getByRange(start, end);
mapState.removeByRange(start, end);
```

2. RocksDB Backend for Operator State

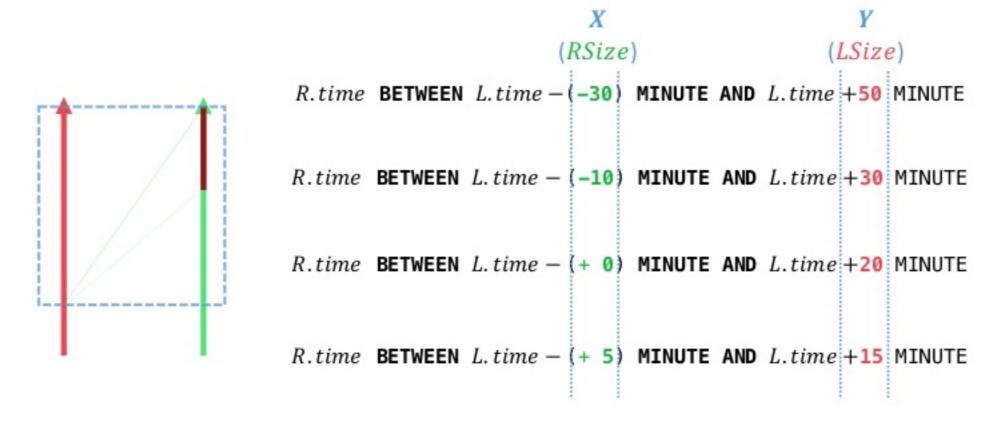


# P3. Separated Watermarks





# Cache Size Optimization



The size of the two caches can be optimized to be the minimum value, X + Y.



# P4. Query Optimization



# Q&A





https://www.youtube.com/watch?v=yDUfz3OFcOo



https://baike.baidu.com/item/%E5%8E%8B%E9%9D%A2%E6%9C%BA/3477203