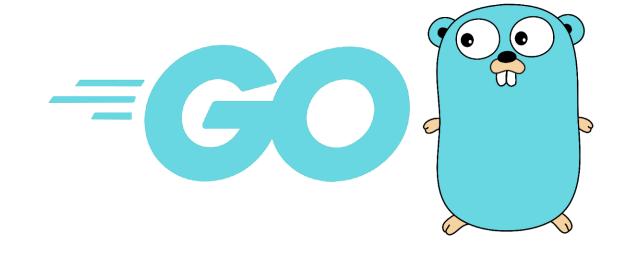
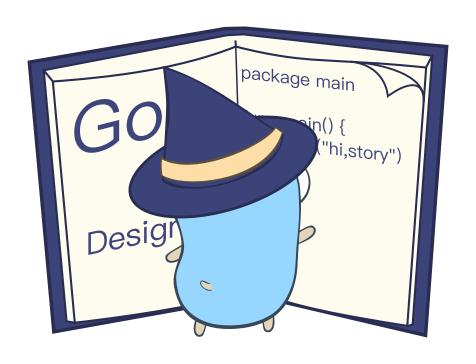
Go设计与实现

Let's Go!



Go的历史 ->



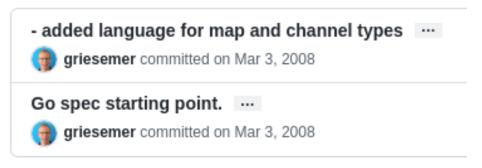
Go创造者



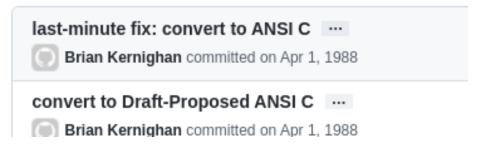
Robert Griesemer	Rob Pike	Ken Thompson	
V8/Go	UTF-8/Plan9/Go	Unix/B/Plan9/UTF-8/Go	

Go历程

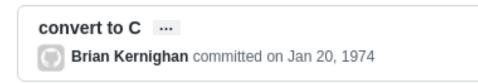
- 2008.03.03 Griesemer 起草
 Go is an attempt at a new systems programming language.
- 2008.03.07 The Go Programming Language. Language Specification
- 2008.03.28 完成Go语言文法和语法的制定,随后着手开发编译器
- 2008.06.05 提交了第一份编译器相关代码(C)
- 2008.09 实习生Russ Cox 加入
- weekly.2009-11-06 ~ weekly.2011-03-07.1(r56第一个稳定版本)
- 2012-03-27 Go1.0发布;基本每年两个版本;最新版本Go1.17.5



-o- Commits on Apr 1, 1988



- Commits on Jan 20, 1974



-o- Commits on Jul 19, 1972

```
hello, world ...

Brian Kernighan committed on Jul 19, 1972
```

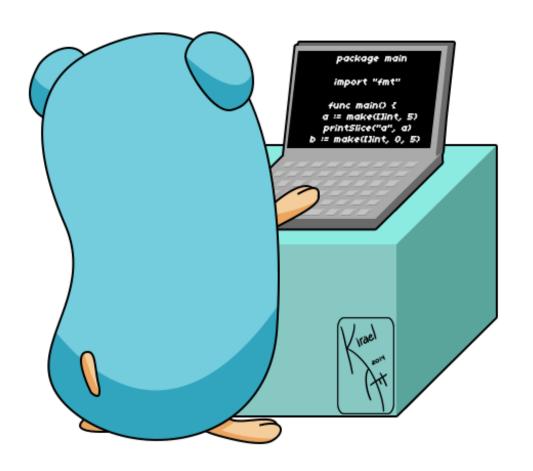
Hello World

• Brain Kernighan "版权所有"

```
main( ) {
          extrn a, b, c;
          putchar(a); putchar(b); putchar(c); putchar('!*n');
}
a 'hell';
b 'o, w';
c 'orld';
```

```
main() {
    printfC("hello, world");
}

/*
调试一段代码的难度是编写它们的两倍,因此如果你的代码写的尽可能巧妙,
按照定义而言,你可能没有能力来调试它了— Kernighan Law
*/
```



Go编译原理

Guiding principles

Go is a new systems programming language intended as an alternative to C++ at Google. Its main purpose is to provide a productive and efficient programming environment for compiled programs such as servers and distributed systems.

The design is motivated by the following guidelines:

- very fast compilation (1MLOC/s stretch goal); instantaneous incremental compilat
- procedural
- strongly typed
- concise syntax avoiding repetition
- few, orthogonal, and general concepts
- support for threading and interprocess communication
- garbage collection
- container library written in Go
- reasonably efficient (C ballpark)

The language should be strong enough that the compiler and run time can be written in itself.

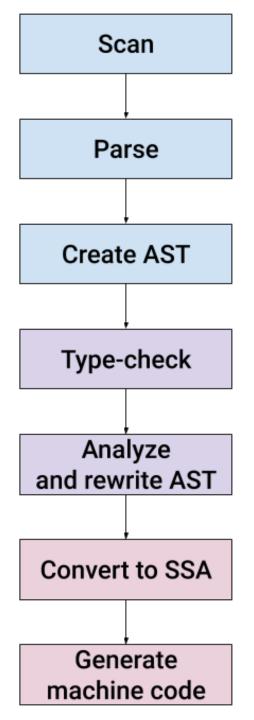
Go设计原则

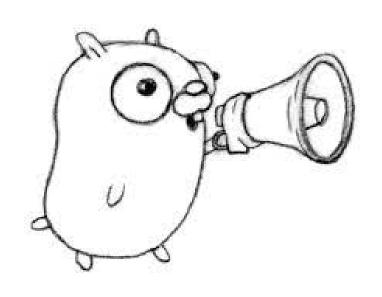
- C++一样的系统编程语言,用于 服务端或分布式系统
- 构建要十分快,瞬时增量编译
- 面向过程(继承 Pascal/C的传统)
- 强类型; 简洁的语法避免重复
- 很少,正交,通俗的概念
- 支持线程和进程间通信
- GC
- 用 Go 编写的容器库(没有模板)
- 相当高效,相当于C

2021-12-30 自举

编译流程

- 1. 词法器扫描
- 2. 语法器解析
- 3. 创建AST
- 4. 类型检查
- 5. 分析重写AST
- 6. 转换成Static Single-Assignment
- 7. 生成机器码





语法是什么?

语法是一门编程语言自身的描述!

语法 = 文法 + 语义

• 文法: 描述语言的句法规则

• 语义:解释文法的含义,限定文

法推导的一个子集合法

$$\int_a^b f(x) = \sum (fx) dx$$

上下文无关文法

- 描述编程语言的语言
- BNF(Backus-Naur Form-巴科斯范式) 或者 EBNF
 - 。 数学语言
 - 。 递归的思想
 - 。精确,无歧义
- 文法规则
 - 产生式由左部是一个非终结符,右部是由非终结符和终结符组成的一个符号串
 - 分隔符:=;操作符 |,(),[],{};终结符、非 终结符

EBNF

Extended Backus-Naur Form

Version of Jul 26, 2021

Table of Contents

Introduction Index expressions
Notation Slice expressions
Source code representation
Characters Calls

Letters and digits Passing arguments to ... parameters

Operators

Return statements

Lexical elements

Comments Arithmetic operators Tokens Comparison operators Semicolons Logical operators Identifiers Address operators Receive operator Keywords Operators and punctuation Conversions Integer literals Constant expressions Floating-point literals Order of evaluation

Imaginary literals Statements

Rune literals Terminating statements String literals Empty statements Constants Labeled statements Variables Expression statements Send statements Types IncDec statements Method sets Boolean types Assignments If statements Numeric types Switch statements String types Array types For statements Go statements Slice types Select statements Struct types

Pointer types Return statements
Function types Break statements
Interface types Continue statements
Map types Goto statements
Channel types Fallthrough statements
Type identity Built-in functions

Assignability

Representability Length and capacity

Blocks

Pointer types

Declarations and scope
Label scopes
Blank identifier
Predeclared identifiers

Making slices, maps and channels
Appending to and copying slices
Deletion of map elements
Manipulating complex numbers

Close

Allocation

Exported identifiers Handling panics
Uniqueness of identifiers Bootstrapping
Constant declarations Packages

 Iota
 Source file organization

 Type declarations
 Package clause

 Variable declarations
 Import declarations

 Short variable declarations
 An example package

Function declarations Program initialization and execution

Method declarations The zero value Expressions Package initialization Operands Program execution

Qualified identifiers Errors

Composite literals Run-time panics
Function literals System considerations
Primary expressions Package unsafe

Selectors Size and alignment guarantees

Method expressions Method values

Go的语法

```
Production = production_name "=" [ Expression ] "." .

Expression = Alternative { "|" Alternative } .

Alternative = Term { Term } .

Term = production_name | token [ "..." token ] | Group | Option | Repetition .

Group = "(" Expression ")" .

Option = "[" Expression "]" .

Repetition = "{" Expression "}" .
```

```
ForStmt = "for" [ Condition | ForClause | RangeClause ] Block .
Condition = Expression .
```

语法量

- Go1.17 / 88 / free
- C11/180/\$60
- C++17 / 1500 / \$116

字符集

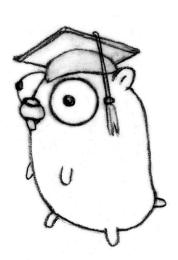
uft-8编码的Unicode字符集,变长1~4个字节

中文编程如何?

```
package main
     import "fmt"
     const(
            真 = true
            假 = false
            嗯 = true
            是的 = true
     func 打印(输出的信息 ...interface{}) {
            fmt.Println(输出的信息)
     func main() {
            你好:= "Golang"
            打印(你好);
            吃饱了吗:= 真
            if(吃饱了吗 == 嗯) {
                   打印("写点代码吧!")
            } else {
                   打印("吃饱了再来。")
            return
2021-12}30
```

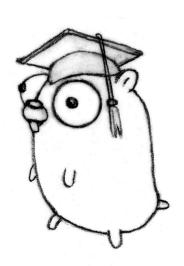
词法元素

- 注释
 - // 或者 /**/
 - ∘ rune 和 string中不有注释
- 分号
 - an identifier
 - an integer, floating-point, imaginary, rune, or string literal
 - one of the keywords break, continue, fallthrough, or return
 - one of the operators and punctuation ++, --,),], or }
- token
 - identifiers, keywords, operators and punctuation and literals.
 - Integer, Floating-point, Imaginary, Rune, String



语法分析

- 自上而下 LL
 - 。 从语法树根开始,使用产生式将非终结符号向左展开
 - 。 实现简单
- 自下而上 LR, LR(n)
 - 从语法树的叶子开始,匹配产生式从右往左规约到非终结符号,最终到达树根
 - 向前预读 n 个token,帮助产生式冲突时识别正确的文法规则
 - 借助规约栈,规约表,实现复杂



Go语法结构1

```
//语法树根
    SourceFile
                     = PackageClause ";" { ImportDecl ";" } { TopLevelDecl ";" } .
    //包条款
    PackageClause
                   = "package" PackageName .
    PackageName
                   = identifier .
    //导入包
    ImportDecl
                    = "import" ( ImportSpec | "(" { ImportSpec ";" } ")" ) .
    ImportSpec = [ "." | PackageName ] ImportPath .
    ImportPath
                     = string_lit .
    //顶层声明
    TopLevelDecl = Declaration | FunctionDecl | MethodDecl .
    //声明
    Declaration = ConstDecl | TypeDecl | VarDecl .
    FunctionDecl = "func" FunctionName Signature [ FunctionBody ] .
    FunctionName = identifier .
    FunctionBody = Block .
    MethodDecl = "func" Receiver MethodName Signature [ FunctionBody ] .
    Receiver = Parameters .
2021-12-30
```

16

Go语法结构2

```
//语句块
Block = "{" StatementList "}" .
StatementList = { Statement ";" } .
Statement =
       Declaration | LabeledStmt | SimpleStmt |
        GoStmt | ReturnStmt | BreakStmt | ContinueStmt | GotoStmt |
        FallthroughStmt | Block | IfStmt | SwitchStmt | SelectStmt | ForStmt |
        DeferStmt .
// 简单语句
SimpleStmt = EmptyStmt | ExpressionStmt | SendStmt |
             IncDecStmt | Assignment | ShortVarDecl .
//表达式语句
ExpressionStmt = Expression .
//表达式
Expression = UnaryExpr | Expression binary_op Expression
```

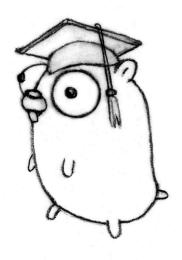
Block 产生式的左边还有一个Block,因此嵌套自然诞生;Expression同理

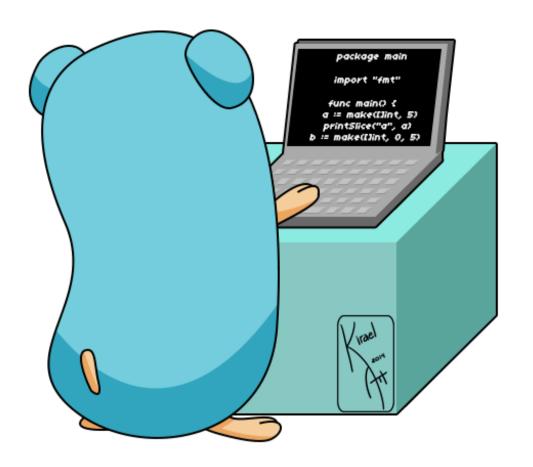
Go编译器源码

```
func (p *parser) parseStmt() (s ast.Stmt) {
           switch p.tok {
               case token.CONST, token.TYPE, token.VAR:
                       s = &ast.DeclStmt{p.parseDecl()}
           case token.GO:
                       s = p.parseGoStmt()
               case token.DEFER:
                       s = p.parseDeferStmt()
           //...
}
       func (p *parser) parseGoStmt() ast.Stmt {
               pos := p.expect(token.G0)
               call := p.parseCallExpr()
               p.expectSemi()
               if call == nil {
                       return &ast.BadStmt{pos, pos + 2} // len("go")
               return &ast.GoStmt{pos, call}
       func (p *parser) parseDeferStmt() ast.Stmt {
               pos := p.expect(token.DEFER)
               call := p.parseCallExpr()
               p.expectSemi()
               if call == nil {
                       return &ast.BadStmt{pos, pos + 5} // len("defer")
               return &ast.DeferStmt{pos, call}
2021-12+30
```

Go语法结构

- 语法树推导,非AST
- 每条产生式都有对应的语义说明
- LL-从树根到叶子,方便阅读,从整体到局部
- LR-从叶子到树根,不太方便阅读,从细节到整体
- 入门适合从LL开始,遇到具体语法问题反过来使用LR分析





Go数据结构

- array, slice
- string
- map
- channel

数组

单一类型的多个元素序列。

```
ArrayType = "[" ArrayLength "]" ElementType .
ArrayLength = Expression .
ElementType = Type .

Type = ArrayType //其余省略
Expression->"..."
```

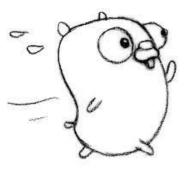
语义:

- 长度是数组的一部分,长度不同,类型不同
- 长度必须是可计算为非负int类型常量的表达式,... 编译器自动推导
- 可以使用[0,len-1]索引访问
- len()返回其长度
 - ▲可以嵌套成多维粉组

数组结构

```
a := [3]int{1,2,3}
b := a // 拷贝
func fun(p [3]int) {}
fun(a) // 拷贝
```

- 使用静态类型检查和运行时runtime.panicIndex 防止越界
- 数组的赋值, 函数参数的传递都是值拷贝
- 运行时不可改变长度
- 指针和切片优化



22

切片

切片是底层数组的连续段的描述符。提供对该数组中连续元素序列的访问。

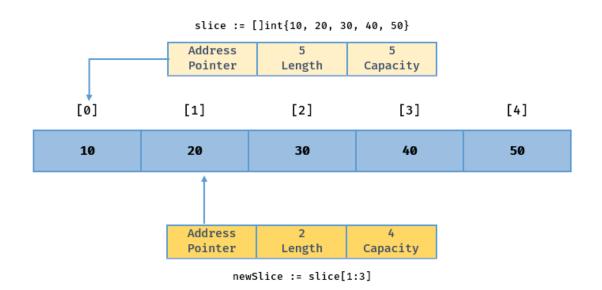
```
SliceType = "[" "]" ElementType .
make([]T, length, capacity)
```

- 与数组不同,它可能会在执行过程中发生变化
- len()返回长度,不可能是负值
- 一旦初始化就关联一个底层数组,并和关联数组的其他切片共享存储
- make 创建的切片总是分配一个新的、隐藏的数组,返回的切片值引用该数组
- 对于多维数组,低维数组长度总相同;但是多维切片的低维切片可能动态变化

23

切片结构

- 未初始化的切片=nil
- cap = len + extent
- 切片的Cap超过数组的大小,会 创建新的底层数组,并复制数据
- len>cap 时,触发扩容



切片扩容

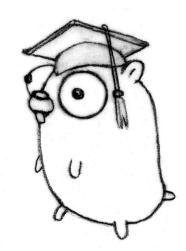
```
slice = append(slice, elems...)
```

扩容策略

- 期望容量大于当前容量2倍,使用期望容量
- 当前切片长度<1024,容量翻倍
- 当前切片长度>=1024,每次增加新容量值的25%,直到容量大于期望容量
- 依据元素类型,进行内存对齐

字符串

- 字符串类型表示字符串值的集合。
- 值是一个字节序列(非字符,可能是空),长度
- 不可变的,一旦创建不可再更改。
- 字符串常量,则长度是编译时常量。
- 字节可以通过索引访问。
- 取一个元素的地址非法: &s[i]无效。



字符串结构

```
// a separate, correctly typed pointer to the underlying data.
type StringHeader struct {
    Data uintptr
    Len int
}
```

- StringHeader 不可变,原因是string要作为map的key
- 字符串的追加元素通过创建一个新的StringHeader和数据复制完成

map

由唯一键索引的另一种元素类型的无序集合

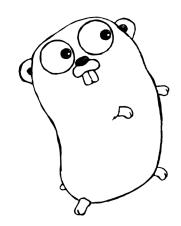
```
MapType = "map" "[" KeyType "]" ElementType .
KeyType = Type .
```

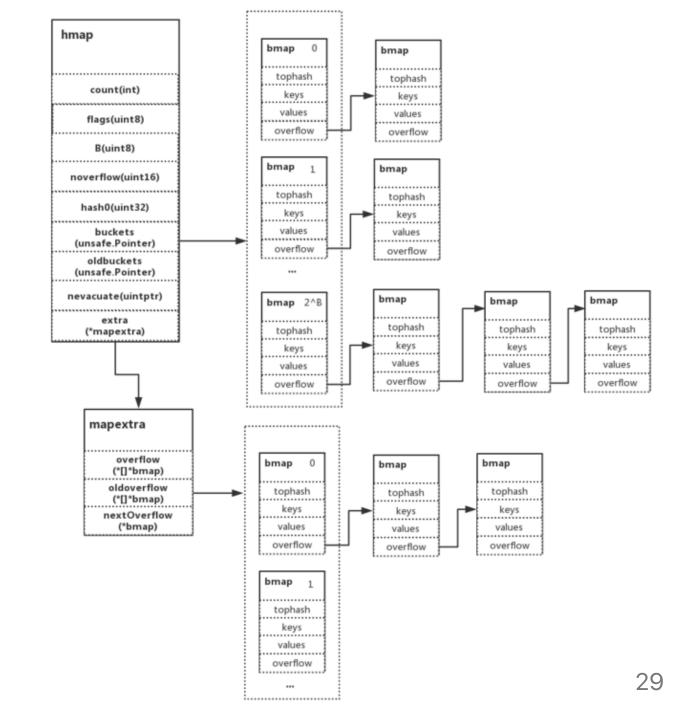
- keyType的!=和==必须是明确定义;因此,function, map, slice不能作为主键
- keyType是interface类型,动态键值的比较运算必须已定义。否则,panic
- 未初始化map为nil,不能添加元素,make创建空map
- [key]索引或赋值/ len() / delete

28

map结构

• 链地址法哈希表





Channel

通道为并发执行函数提供了一种机制,通过发送和接收指定元素类型的值来进行通信。

```
ChannelType = ( "chan" | "chan" "<-" | "<-" "chan" ) ElementType .
chan T // 声明一个双向通道
chan<- T // 声明一个只能用于发送的通道
<-chan T // 声明一个只能用于接收的通道
chan<- chan int  // same as chan<- (chan int)</pre>
chan<- <-chan int // same as chan<- (<-chan int)</pre>
<-chan <-chan int // same as <-chan (<-chan int)
//发送语句
SendStmt = Channel "<-" Expression .</pre>
Channel = Expression .
//接受操作
Expression = UnaryExpr | Expression binary_op Expression .
UnaryExpr = PrimaryExpr | unary_op UnaryExpr .
unary_op = "<-" //其余省略
```

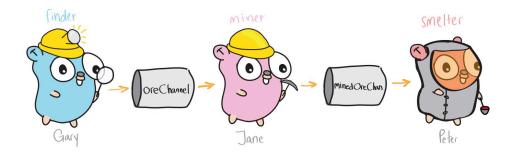
Channel

- make(T, cap) / close
- cap = 0/缺省
 - 。 无缓冲通道,只有当发送方和接收方都准备好时,通信才能成功,否则阻塞
- cap > 0
 - 有缓冲通道,若缓冲区未满(发送)或不为空(接收),则通信成功而不会阻塞
- 通道可以被任意个goroutine执行发送语句,接收操作,len, cap而不用执行同步操作
- FIFO

31

Channel结构

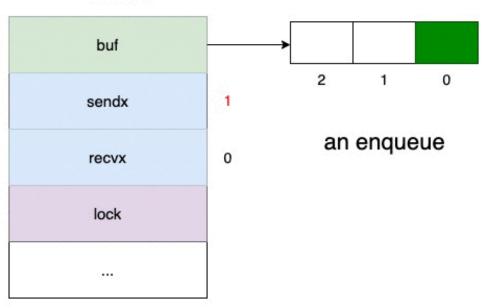
```
// 管道
     type hchan struct {
     qcount
           uint
                           // 元素的数量
      dataqsiz uint
                           // 缓冲区的大小
      buf
             unsafe.Pointer // 循环队列
      elemsize uint16 // 每个元素的大小
      closed uint32 // 管道状态
      elemtype *_type // 元素类型
      sendx
             uint // 发送的索引
             uint
      recvx
      recvq waitq // 等待接收的队列
      sendq waitq // 等待发送的队列
      lock mutex //互斥锁
     // 等待队列
     type waitq struct {
      first *sudog
      last *sudog
     type sudog struct {
      g *g
      next *sudog
      prev *sudog
2021-12}30
```



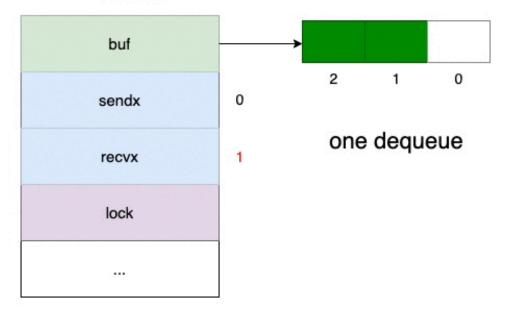
发送

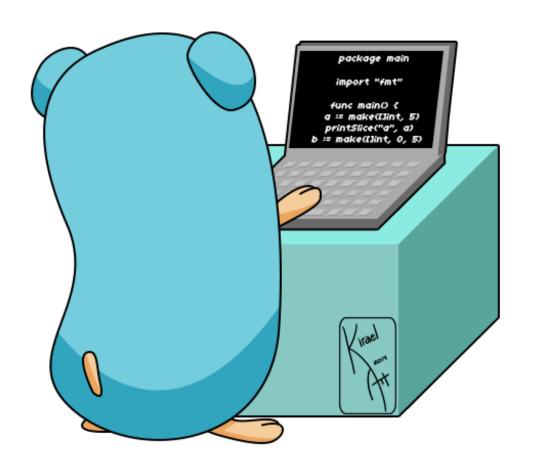
接收

hchan



hchan





Go的设计

- defer
- goroutine
 - 。 并发
 - 。 调度
- GC

defer

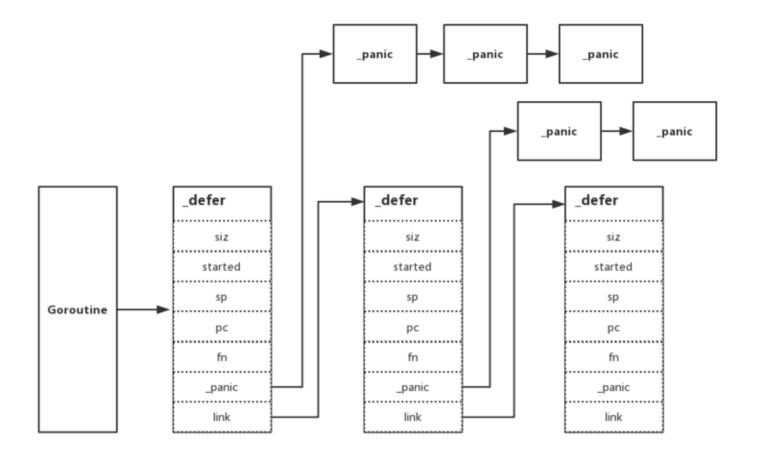
中文:延期;——延迟执行函数。

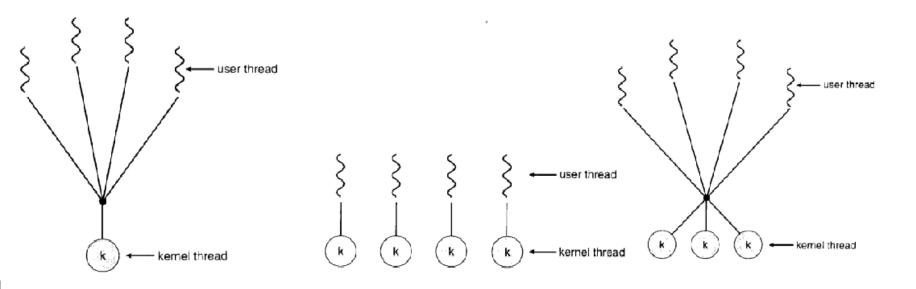
DeferStmt = "defer" Expression . //Expression必须是函数或方法调用

- 函数body结束,执行return语句,或者panic时执行
- 遇到 defer 时,执行正常解析计算,不立即执行函数调用
 - 函数返回时才能够发现 defer 函数会不会 panic
 - 如果一个 defer 在retrun语句之后,那么它不会被解析
 - 函数体内的多个 defer,按照其在代码中的先后顺序,FILO栈原则
 - defer 可以嵌套和递归,此递归不存在递归返回。
- 函数入参压栈,出参丢弃;defer函数可以修改函数的返回值

defer 实现

- deferproc/deferreturn
- panic只执行G的defer





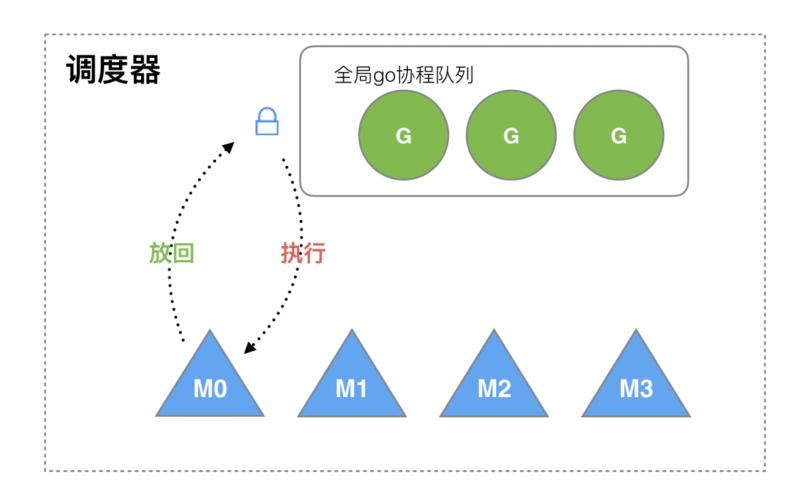
线程模型

- N:1-用户
- 1:1-内核
- N:M-两级

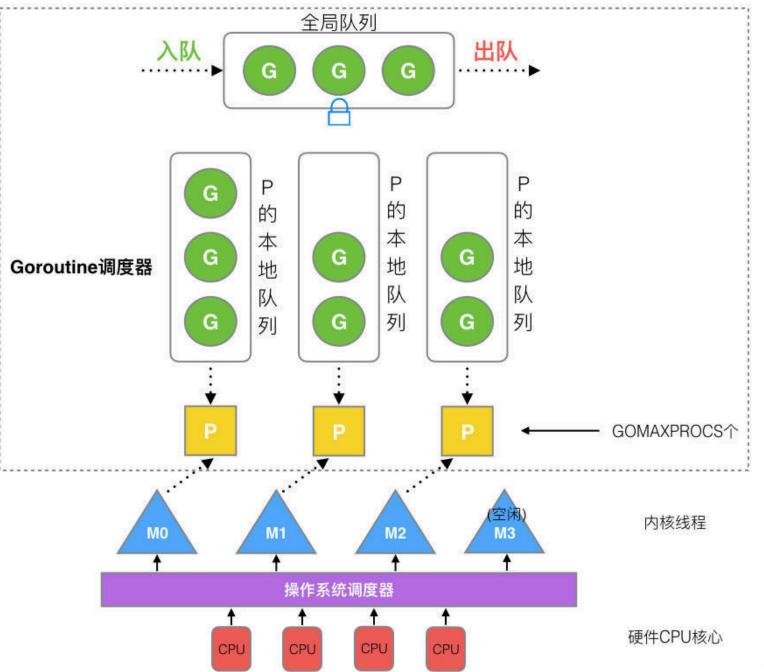
内存 + CPU

G-M模型

- 竞争
- 局部化



G-M-P模型



同步与通信

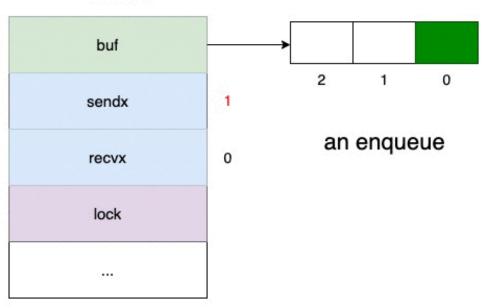
Do not communicate by sharing memory; instead, share memory by communicating. 不要通过共享内存来通信,而应通过通信来共享内存。

- Channel
- WaitGroup/Mutex/RWMutex/Cond/Once

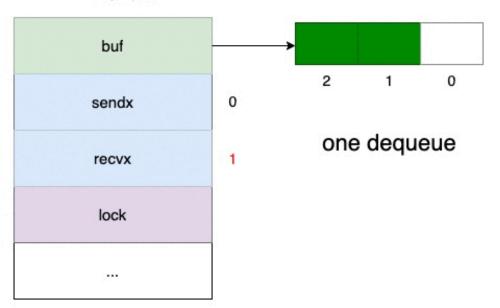
发送

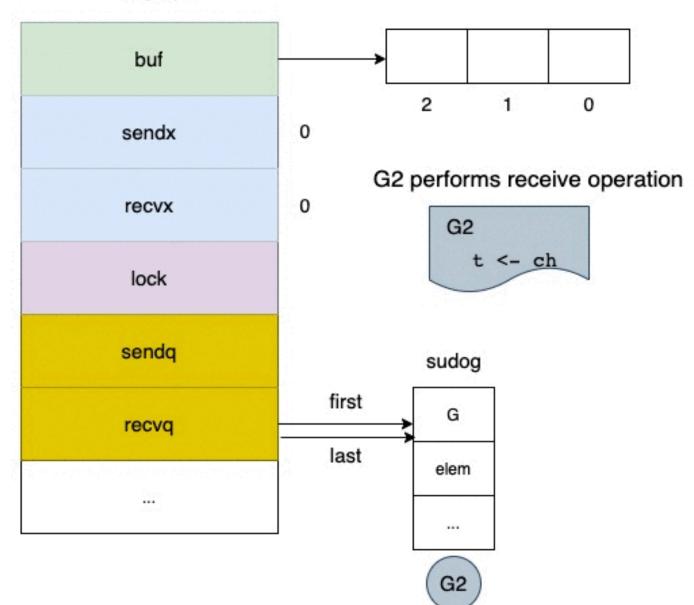
接收

hchan



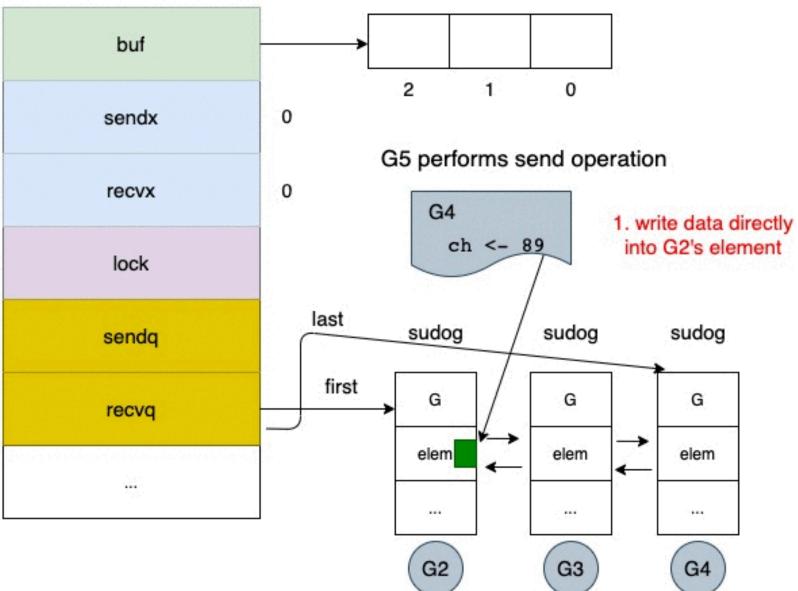
hchan



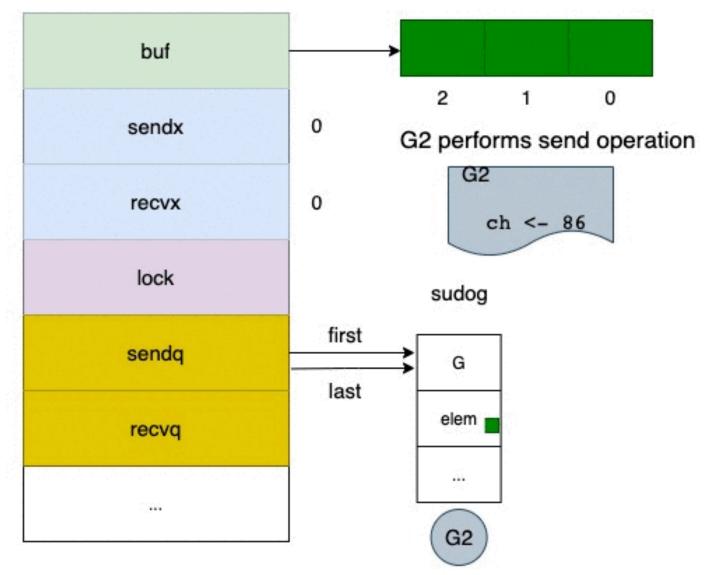


接收者阻塞

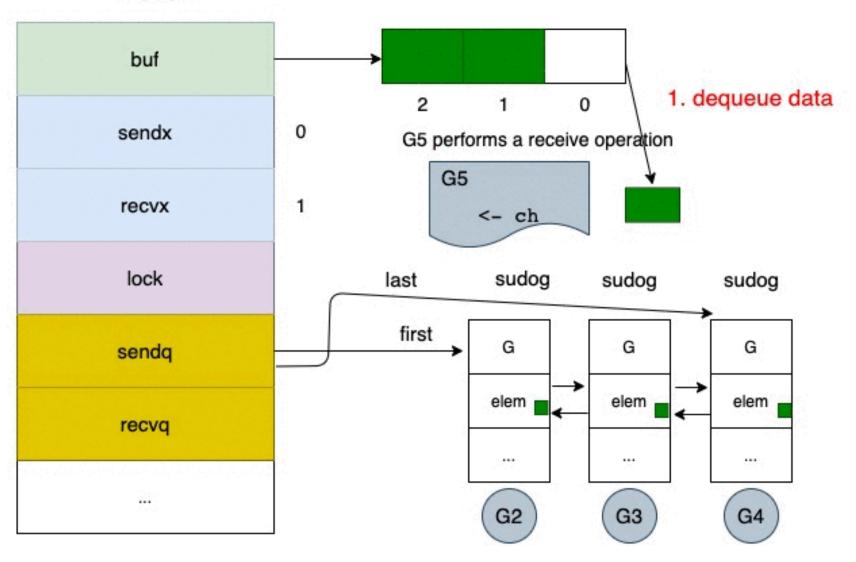
42



接收者唤醒



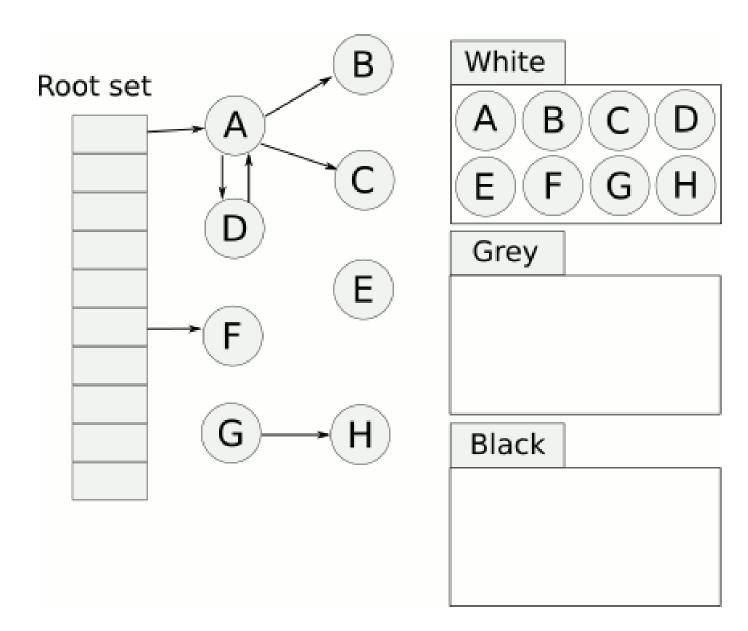
发送者阻塞

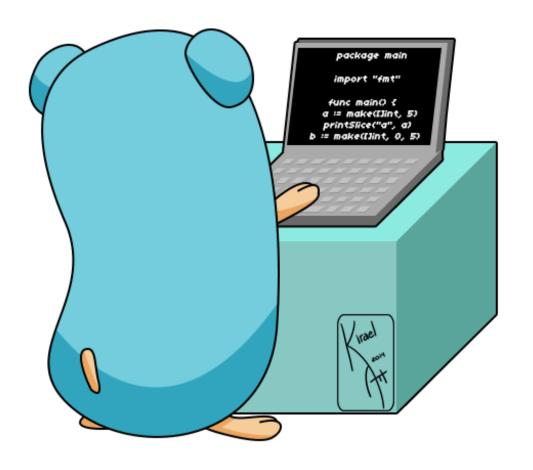


发送者唤醒

GC

- 所有对象开始都是白色
- 从 root 开始找到所有可 达对象,标记为灰色, 放入待处理队列
- 遍历灰色对象队列,将 其引用对象标记为灰色 放入待处理队列,自身 标记为黑色。
- 处理完灰色对象队列, 执行清扫工作。





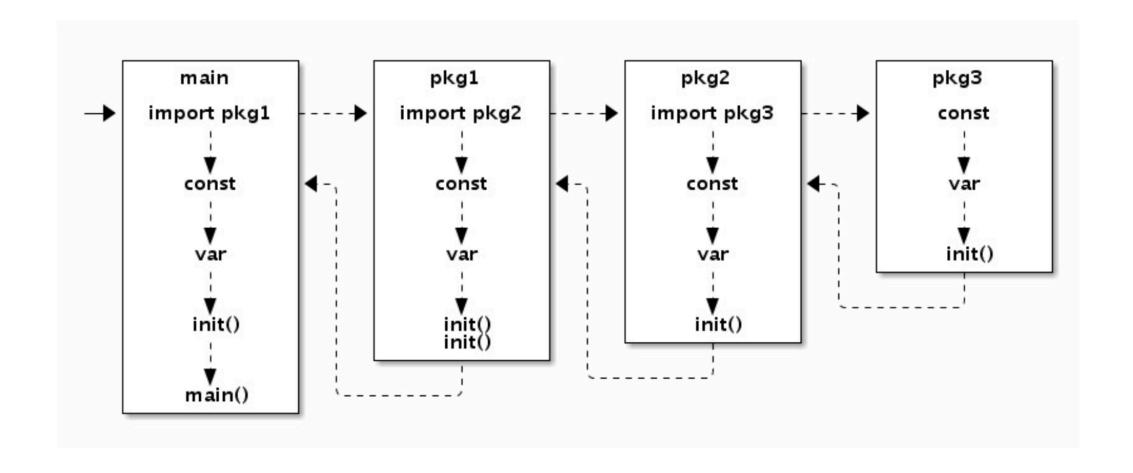
Go的生态

- package
- CGO

包

- 一个包只在一个目录下,一般包名和目录名相同
- import 包
 - 包路径 import "lib/math"
 - 重命名 import ms "lib/math"
 - 本地化 import . "lib/math"
- GOPATH/GOROOT | vendor | module

包的初始化



站在C/C++的肩上

- 包名后紧接着是/**/包含的完成C代码,然后 import "C"
- C++是不可以直接使用的,需要用C代码包装一层
- 强类型需要适配

```
#include <stdlib.h>

CPPObject NewCPPObject()
{
   return new CPPObject();
}

void Config(void* trans_api, char *user, char *passwd, char *ip, int port)
{
   ((CPPObject *)trans_api)->Config(user, passwd, ip, port);
}
```

```
package cppboject
/*
#cgo LDFLAGS: -L../lib -lcppboject
#cgo CFLAGS: -I ../include
#include <stdlib.h>
*/
import "C"
import (
 "fmt"
 "unsafe"
func (ta *AccTransApi) Init(user, passwd, addr string, port, timeout int32) error {
 obj = C.NewCPPObject()
  cUser := C.CString(user)
 defer C.free(unsafe.Pointer(cUser))
  cPasswd := C.CString(passwd)
 defer C.free(unsafe.Pointer(cPasswd))
  cAddr := C.CString(addr)
 defer C.free(unsafe.Pointer(cAddr))
 C.Config(obj, cUser, cPasswd, cAddr, C.int(port))
  return nil
```



Thanks