Revising the simple measures of assortativity in complex networks

Xiao-Ke Xu,[1,2,][*] Jie Zhang,[2] Junfeng Sun,[3] and Michael Small[2]

[1]*School of Communication and Electronics Engineering, Qingdao Technological University, Qingdao 266520, China*
[2]*Department of Electronic and Information Engineering, Hong Kong Polytechnic University, Hung Hom, Kowloon, Hong Kong*
[3]*Med-X Research Institute, Shanghai Jiao Tong University, 1954 Hua Shan Road, Shanghai 200030, China*

We find that traditional statistics for measuring degree mixing are strongly affected by superrich nodes. To counteract and measure the effect of superrich nodes, we propose a paradigm to quantify the mixing pattern of a real network in which different mixing patterns may appear among low-degree nodes and among high-degree nodes. This paradigm and the simple revised measure uncover the true complex degree mixing patterns of complex networks with superrich nodes. The alternate method indicates that some networks show a false disassortative mixing induced by superrich nodes and have no tendency to be genuinely disassortative. Our results also show that the previously observed fragility of scale-free networks is actually greatly exacerbated by the presence of even a very small number of superrich nodes.

I. INTRODUCTION

With increasing evidence of the ubiquity of scale-free networks, attention has recently shifted to the particular structure of experimentally observed networks [1–6]. One key observation to arise has been that some networks display a propensity for high-degree nodes to connect to other high-degree nodes (*assortativity*) [7,8]. Conversely in certain other types of networks the reverse is true: high-degree nodes connect to low-degree nodes (*disassortativity*). In particular, numerical evidence from experimental data has shown that many technological (i.e., communication networks), biological (e.g., protein and neural networks), and certain social (online communities [9,10]) networks are found to exhibit a negative assortativity coefficient and are therefore claimed to be examples of disassortative mixing [7,11].

A widely accepted way to determine the mixing pattern of complex networks is to calculate the Pearson correlation coefficient of the degrees at both ends of the edges [3,12,13]. Although Newman's assortativity coefficient is not a perfect measure of assortativity, it is a very simple measure to approximately assess the degree mixing pattern of a network and has been employed broadly. Another natural approach to quantify mixing patterns is to calculate the correlations between two nodes connected by an edge [13]. There are two ways to characterize the node degree correlations. One is the conditional probability $P(k'|k)$ that an arbitrary edge connects a node with degree k to another node with degree k' [14]. The other is the joint degree distribution $P(k,k')$ that measures whether nodes with a given degree k prefer (or avoid) to connect to nodes with degree k' [1,2,15]. Although they can capture the mixing patterns by the two-dimensional degree correlation plots [16–18], it is difficult for the two methods to quantitatively evaluate the assortativities of complex networks.

While the existing measures for assortativity are simple, and the mathematical definition of what is meant by assorta-

tivity and disassortativity is clear and consistent, this is not equivalent to the general understanding of these phenomena. In particular, disassortativity can easily and often does arise in situations when the degree of neighbors matches as closely as possible. This is completely at odd with what is commonly understood, and this needs to be addressed. The current usage of degree correlation measures means that researchers often conclude that node degree between neighbors in physical and social systems is mismatched while in fact the opposite is true.

In this paper we show that in a finite-size network superrich nodes cause a network to show an observed disassortative mixing in most cases, while they sometimes let a strong disassortative network appear less disassortative property. And superrich nodes often limit the assortativity coefficient to be in a narrow range and mask the genuine mixing pattern of complex networks. The superrich nodes, which exist in many real networks, refer to the nodes whose degrees are far larger than most other nodes. The assortativity coefficient in certain experimentally measured networks is therefore due to a small fraction of superrich nodes and is not caused by a genuinely assortative (or disassortative) mixing. The mechanism underlying degree mixing has not been properly understood and many of these common indices misrepresent the true assortativity of a network. We find that the assortativity coefficient will give a "false" result of mixing patterns for the networks with superrich nodes.

To overcome the effect of these aberrant and uncommonly connected nodes, a new and more robust measure of degree-degree mixing is needed. We choose a modification of Newman's assortativity coefficient so that our measure will be simple and also as similar as possibly to what is already being employed. We deliberately choose to change the existing measure as little as possible, so that the measure can be best understood. We are not suggesting that the existing tools be abandoned and only that they be computed twice: once as is done now and again after removing the contribution of superrich nodes.

The superrich nodes can, and often will, have an extremely great effect on the network structure. Our results also show that the previously observed fragility of scale-free net-

works is actually greatly exacerbated by the presence of even a very small number of superrich nodes. An attack targeting the *superrich* nodes (rather than just the *rich*) of a network can very quickly fragment, or at least stretch, a scale-free network.

II. SUPERRICH NODES AFFECT DEGREE MIXING PATTERNS

A. Superrich nodes in scale-free network

The superrich nodes refer to the nodes whose degrees are far larger than most other nodes. Our finding in this study can be applied to any degree distribution network and is not limited for power-law and exponential degree distribution networks, so there is no need to judge whether the degree distribution obeys a power law strictly. While it is necessary to supply a simple and operative definition of superrich nodes in scale-free networks, for the power-law degree distribution broadly exists in real networks [19]. In terms of networks with approximately power-law (or exponential) degree distribution, superrich nodes are defined as the nodes whose degrees are larger than the natural cutoff value. It should be noticed that the nodes with degree that is predicted by the power law only are rich nodes.

The natural cutoff degree k_c is an important concept in finite-size scale-free networks [20]. It is defined as the value of the degree k_c above which one expects to find at most one node [21], that is,

$$N\int_{k_c}^{\infty} P(k)dk \sim 1, \tag{1}$$

where N is the number of nodes. For a scale-free network, this expression provides a dependence of the natural cutoff with N and the slope γ as

$$k_c(N) \sim N^{1/(\gamma-1)}. \tag{2}$$

Here, we obtain the slope γ by fitting the real data excluding the potential superrich nodes.

B. Adding superrich nodes to BA model

First we utilize the BA model [19] to demonstrate the significant effect of superrich nodes. We generate a small size of BA network (nodes $n=200$ and average degree $\langle k \rangle =6$) and then the maximal assortative mixing (MAM) and maximal disassortative mixing (MDM) networks are generated from the original BA network using the rewiring method [22]. These three networks have the same degree distributions but with different mixing patterns. The MAM network reveals a beautiful assortative structure in which nodes with similar degrees connect to one another as shown in Fig. 1(a). In contrast, the MDM network shows that high-degree nodes link to low-degree ones in Fig. 1(b). However, the addition of only one superrich node ($k=100$ and random linking) severely affects the entire topology of the network, as shown in Figs. 1(c) and 1(d). Our results in Fig. 1 also show that the clustering coefficient [23] and the average path length are extremely changed by the only one superrich node.
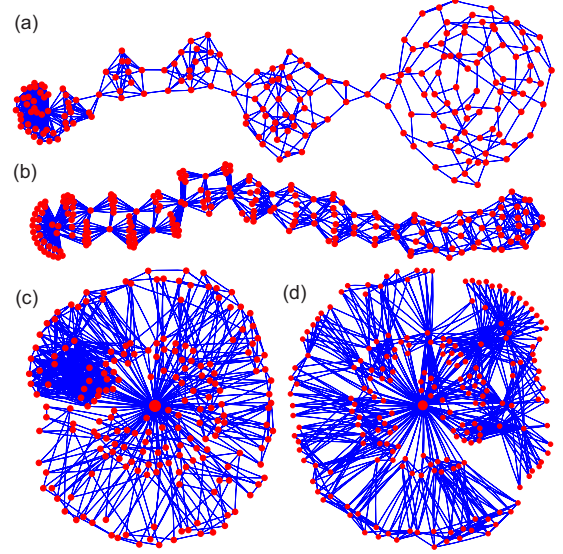

FIG. 1. (Color online) Networks without and with a superrich node show greatly different topological structure, average path length l, and clustering coefficient c: (a) maximal assortative mixing network, $l=9.73$ and $c=0.23$; (b) maximal disassortative mixing network, $l=8.19$ and $c=0.02$; (c) maximal assortative mixing network with one superrich node ($k=100$) added randomly, $l=2.88$ and $c=0.30$; and (d) maximal disassortative mixing network with one superrich node ($k=100$) added randomly, $l=2.75$ and $c=0.10$.

Now we demonstrate how superrich nodes bias the mixing pattern of a network. The assortativity coefficient r [7] is given by

$$r = \frac{M^{-1}\sum_i j_i k_i - \left[M^{-1}\sum_i \frac{1}{2}(j_i+k_i)\right]^2}{M^{-1}\sum_i \frac{1}{2}(j_i^2+k_i^2) - \left[M^{-1}\sum_i \frac{1}{2}(j_i+k_i)\right]^2}, \tag{3}$$

where j_i and k_i are the degrees of the two end points of the ith edge and M is the total number of edges in the network. If $r>0$, the network is claimed to be assortative mixing; while, if $r<0$, the network is called disassortative mixing.

A large-scale BA network ($n=5000$ and $\langle k \rangle =6$) is generated, and the MAM and MDM networks are obtained by the rewiring method [22]. The r values for these networks are listed in Table I. The above three networks exhibit the reasonable mixing coefficients: assortative, neutral, and disassortative. Then we add five superrich nodes ($k=1000$) to each of the three networks. For the MAM network, superrich nodes are connected to high-degree nodes, while for the MDM network, superrich nodes are linked to low-degree nodes; for the original BA network, random linking is adopted. As can be seen in Table I, all r are negative for these networks, which runs against our intuition. Take the MAM network, for example; the strong positive r (0.320) is replaced with a negative value (-0.136), although the superrich nodes are attached to the high-degree nodes (assortative adding). Moreover, the fluctuation of r becomes narrow for the networks with superrich nodes ($[-0.189,0.320]\Rightarrow[-0.169,-0.136]$). This result shows that superrich nodes can-

TABLE I. Assortativity coefficients r of different mixing patterns for BA networks without and with superrich nodes.

Network	r
Maximal disassortative mixing	-0.189 ± 0.029
Original BA network	-0.047 ± 0.004
Maximal assortative mixing	0.320 ± 0.088
Maximal disassortative mixing with five superrich nodes (disassortative adding)	-0.169 ± 0.002
Original BA network with five superrich nodes (random adding)	-0.163 ± 0.002
Maximal assortative mixing with five superrich nodes (assortative adding)	-0.136 ± 0.002

not only make a strong assortative network to show a negative assortativity coefficient, but can also make a strong disassortative network to appear less disassortative.

C. Superrich nodes in experimental networks

Many complex networks exhibit a scale-free degree distribution [19,24], such as the two real networks in Figs. 2(a) and 2(b). And the slope of power law is obtained by fitting the real data excluding the potential superrich nodes. For the collaboration network in computational geometry [25], its degree fits the power-law distribution perfectly, while for the language network [26], from Darwin's "The Origin of Species," superrich nodes exist in it because the maximum degree ($k_{max}=2568$) is far larger than the natural cutoff ($k_c=105$). The results of r for four mixing patterns of the two
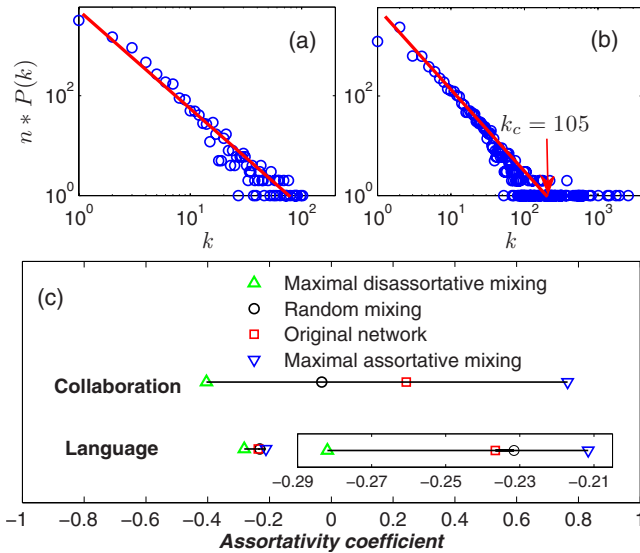

FIG. 2. (Color online) Degree distributions of two real networks and their assortativity coefficients for different mixing patterns: (a) author collaboration network in computational geometry (nodes $n=7343$, edges $m=11\,898$, and maximum degree $k_{max}=102$) [25]; (b) language network, which is a word adjacency network of texts from Darwin's The Origin of Species ($n=7724$, $m=46\,281$, and $k_{max}=2568$) [26]; and (c) assortativity coefficients of different mixing patterns for the above two networks.
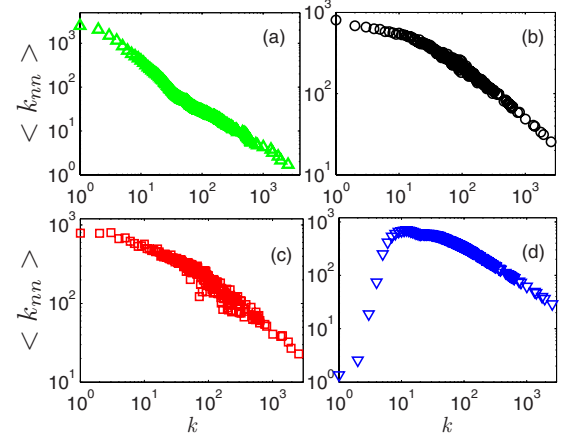

FIG. 3. (Color online) Average degree $\langle K_{nn} \rangle$ of the nearest neighbors of a node depending on its degree k for different mixing patterns of the language network: (a) maximal disassortative mixing, (b) random mixing, (c) original network, and (d) maximal assortative mixing.

networks are shown in Fig. 2(c). The original collaboration network is assortative like most social networks [7,8], while for all mixing patterns of the language network, r is strong negative and is confined to a narrow range, which is similar to the BA network with superrich nodes. According to the results of the theoretical model (BA network) and the language network, we conclude that r cannot determine mixing patterns of networks with superrich nodes accurately, and superrich nodes lead r to be strong negative and within a very narrow range.

One method to characterize the node degree correlations is the conditional probability $P(k'|k)$ that an arbitrary edge connects a node with degree k to another node with degree k' [14]. It is difficult for this method to quantitatively evaluate mixing patterns of complex networks with superrich nodes because the finite network size and the small amount of superrich nodes will lead to an unstable result [13,22]. This problem can be partially solved by calculating the average degree of the nearest neighbors of nodes with a given degree k [17,18], which is given by

$$\langle K_{nn} \rangle = \sum_{k'} k' P(k'|k). \tag{4}$$

Here, an increasing $\langle K_{nn} \rangle$ with k indicates that nodes with high degree tend to connect to nodes with high degree, and the network is classified as assortative, whereas a decreasing $\langle K_{nn} \rangle$ with k indicates that nodes with high degree tend to connect to nodes with low degree, and the network is disassortative.

The average degree $\langle K_{nn} \rangle$ for different mixing patterns of the language network is shown in Fig. 3. For the MDM, original, and random mixing networks, $\langle K_{nn} \rangle$ all decrease with k, which shows a disassortative property. For the MAM network, $\langle K_{nn} \rangle$ first increases and then decreases, and thus it is difficult to tell whether the network is assortative or disassortative. These results indicate that neither r nor $\langle K_{nn} \rangle$ can characterize the intrinsic mixing patterns of complex net-
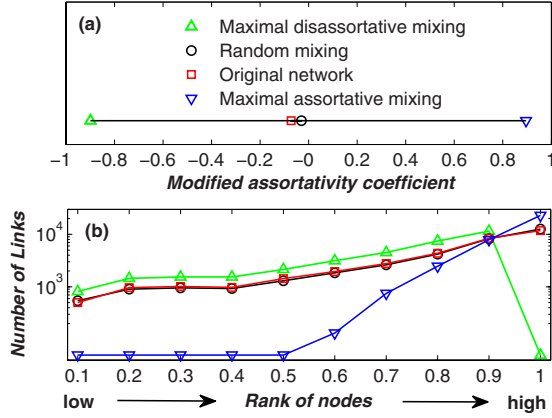
FIG. 4. (Color online) Results of two alternative measures for different mixing patterns of the language network: (a) modified assortativity coefficient to quantitatively measure mixing patterns among the nodes whose degrees are below natural cutoff degree and (b) qualitative measure for the links between superrich nodes and other nodes (including links among superrich nodes).

works with superrich nodes, for the degree of a superrich node is far larger than those of other nodes.

III. ALTERNATIVE PARADIGM OF MEASURING MIXING PATTERNS

A. Alternative paradigm

Since superrich nodes are widely observed in real networks, a new statistic is needed to appropriately classify whether a network with superrich nodes is genuinely assortative or not. In this paper, we manage to determine the mixing patterns of complex networks with two measures. The first one is a modified definition of assortativity coefficient, which is given by

$$r_c = \frac{M_c^{-1}\sum_i j_i k_i - \left[M_c^{-1}\sum_i \frac{1}{2}(j_i + k_i)\right]^2}{M_c^{-1}\sum_i \frac{1}{2}(j_i^2 + k_i^2) - \left[M_c^{-1}\sum_i \frac{1}{2}(j_i + k_i)\right]^2}, \quad (5)$$

where M_c, which is different to the original M, is the total number of edges among the nodes whose degrees are lower than k_c. In accordance with r, positive r_c indicates that a network is assortative, and negative r_c represents a network is disassortative. When calculating r_c in Eq. (5), we exclude superrich nodes whose degrees are greater than k_c. In contrast with the failure of r, r_c can determine the genuine mixing patterns of complex networks with superrich nodes and distinguish the four mixing patterns very effectively as demonstrated in Fig. 4(a). Especially, r_c of the original language network is neutral, which is different to the strong negative r that we have observed in Fig. 2(c).

Although our statistic r_c can effectively measure the mixing patterns among the nodes whose degree is below the natural cutoff value, it is not an optimal solution to neglect a very small number but the most important superrich nodes. Moreover, a real network may exhibit assortative mixing

among low-degree nodes and disassortative mixing among superrich nodes, and it is not accurate to assert whether the network is assortative or disassortative. It is necessary to propose a paradigm to measure the mixing pattern of low- and high-degree nodes, respectively, especially for the network with superrich nodes, because the assortativity coefficient of a very small number of superrich nodes will mask the mixing pattern of larger number of low-degree nodes.

We have used the modified assortativity coefficient to quantitatively measure the mixing pattern of low-degree nodes, so the second measure is proposed to quantify the mixing pattern between low-degree nodes and superrich nodes (including links among superrich nodes). Because the degrees of superrich nodes are far larger than most other nodes and superrich nodes have to link to many low-degree nodes, it is meaningless to say that the mixing pattern of superrich nodes is always disassortative. To find the genuine mixing pattern of superrich nodes, here we test that superrich nodes tend to link to high or low *rank* nodes. If superrich nodes tend to link to high rank nodes, we consider that they are assortative mixing. Conversely, if superrich nodes tend to link to low rank nodes, we believe that they show a disassortative property.

First nodes are ranked in an increasing order according to the degrees. If two nodes have the same degrees, we rank their orders randomly. The nodes are divided into ten bins according to their ranks. Then the number of links connecting superrich nodes to the nodes in each bin is calculated. We cannot calculate the assortativity coefficient of the orders after ordering all nodes because of a false result given by the large number of links of superrich nodes. For different mixing patterns, the distributions of linking numbers are shown in Fig. 4(b). The MAM network shows assortative for superrich nodes tend to connect to *high rank* nodes. On the contrary, the MDM network shows disassortative for superrich nodes tend to link to *low rank* nodes. The original and the random mixing networks are both neutral mixing for superrich nodes have no obvious tendency to connect with high or low rank nodes.

Figures 4(a) and 4(b) show the coincident results for different mixing patterns, which indicates that the two measures are both effective. Furthermore, our two measures compose an effective paradigm to quantify a real network in which different mixing patterns may appear among low-degree nodes and among high-degree nodes. They quantify the effect of superrich nodes on the mixing pattern of complex networks. The alternative paradigm can capture the rich nature of the mixing properties in a real network and distinguish whether the disassortativity of a network is derived from superrich nodes.

B. Comparison with correlation profile

The joint degree distribution $P(k,k')$ is another way to characterize the node degree correlations, which can measure whether nodes with a given degree k prefer (or avoid) to connect to nodes with degree k'. The two-dimensional degree correlation profile [1,15] is a suitable way to use the joint degree distribution to capture the mixing patterns. We
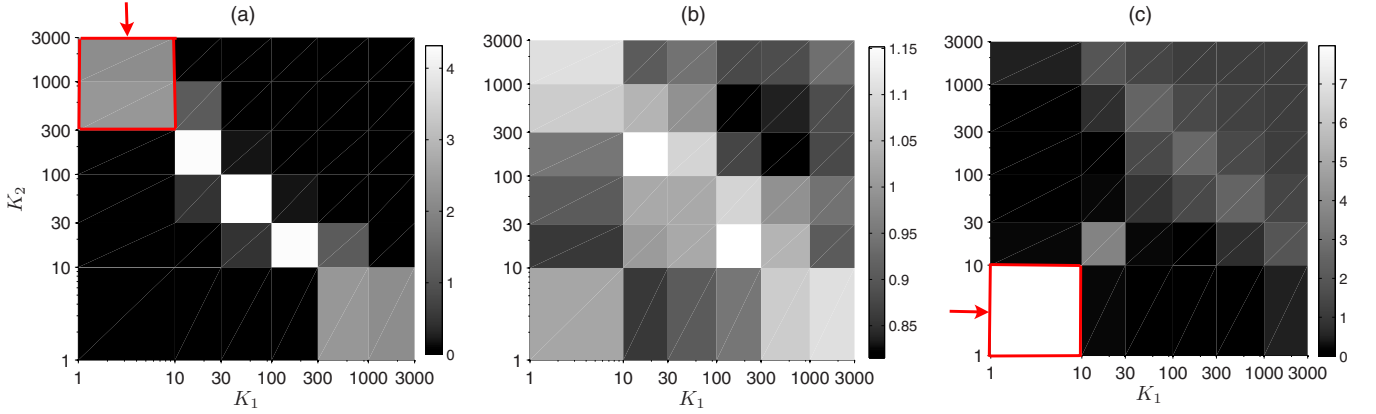
FIG. 5. (Color online) Correlation profile for different mixing patterns of the language network: (a) maximal disassortative mixing, (b) original network, and (c) maximal assortative mixing. These figures show the ratio $R(k_1,k_2)=N(k_1,k_2)/N_r(k_1,k_2)$, where $N(k_1,k_2)$ is the total number of edges in the links of the language network with degrees k_1 and k_2, while $N_r(k_1,k_2)$ is the same value in a randomized version of this network.

visualize the correlation profile for different mixing patterns of the language network by plotting the ratio $R(k_1,k_2)$ $=N(k_1,k_2)/N_r(k_1,k_2)$ in Fig. 5, where $N(k_1,k_2)$ is the total number of edges in the links of the language network with degrees k_1 and k_2, while $N_r(k_1,k_2)$ is the same value in a randomized version of this network.

Observing the results in Fig. 5, we find the correlation profile effectively captures the degree correlations of neighboring nodes. The left-upper corner of Fig. 5(a) shows that high-degree nodes tend to link to low-degree nodes (disassortative) in the MDM network. Conversely, the MAM network obviously shows that low-degree nodes prefer to connect to low-degree nodes (assortative) in the left-lower corner of Fig. 5(c). All the values of $R(k_1,k_2)$ for the original language network are in a very narrow range [0.82,1.15] in Fig. 5(b), which means this network is near neutral mixing. These results are consistent with those in Figs. 4(a) and 4(b), which shows that the two-dimensional degree correlation profile is an effective way to capture the mixing patterns and it is more robust than original r and $\langle K_{nn} \rangle$ to capture the mixing patterns of complex networks with superrich nodes.

Comparing the correlation profile with our paradigm, we can find that the advantage of two-dimensional correlation profile is describing the detailed information of degree correlations. But the correlation profile is only a qualitative method and is difficult to quantitatively evaluate the assortativities of complex networks. The assortativity coefficient is a very simple and the most widely accepted quantitative way to determine the degree mixing pattern. We choose a modification of Newman's assortativity coefficient r_c as the measure for real networks, so that our measure will be simple and also as similar as possible to what is already being employed. And this measure avoids replacing the existing simple and useful measures with something entirely different.

IV. RESULTS OF EXPERIMENTAL NETWORKS

The results of r and r_c for nine real networks are listed in Table II. Like most social networks, networks (a) and (b)

have positive r, for the nodes tend to connect to one another belonging to the same community [8], and there are no superrich nodes. But r of Wealink.com is negative, which is consistent with other online communities [9,10]. And k_c of Wealink.com is 150, so a person who has more than 150 links is a *superstar*. Although the number of superstars in Wealink.com is only 166 (0.07%), surprisingly r_c becomes more negative ($-0.08 \Rightarrow -0.42$) after excluding these superstars. The above result indicates that a very small number of superrich nodes mask the genuine degree mixing information of the online community. And superrich nodes that do not always make a strong assortative network to show a negative assortativity coefficient. On the contrary, sometimes they will make a strong disassortative network to appear as a less disassortative property, just like in this case.

Usually it is believed that technological and biological networks tend to be disassortative [7,11], such as networks (e), (f), and (i). We find that superrich nodes lead r of these networks to be negative. After filtering superrich nodes, (e) and (i) show neutral mixing properties, while r_c of (f) turns more negative. The above results show that r_c can reveal the intrinsic mixing patterns masked by superrich nodes. Since previous works do not take the effect of superrich nodes into consideration, many networks exhibit "false" negative r and decreasing $\langle K_{nn} \rangle$. The advantage of r_c over the previous measures is that it is not affected by superrich nodes. The modified definition of assortativity coefficient r_c is more suitable to characterize the mixing pattern of any degree distribution network, especially for a network with superrich nodes.

One possible explanation for the disassortative mixing in these nonsocial networks is the structural constraint: two nodes have no more than one edge connecting them [15]. But the structural constraint cannot explain why not all nonsocial networks are disassortative. Actually, further research has found that the prohibitively multiedged mechanism cannot generate the same correlation as the real Internet, and only part of degree correlations can be obtained in this way [16]. And it is difficult to distinguish which part of disassortativity is derived from the structural constraint. As far as we know, the intrinsic mechanism why these nonsocial networks ex-

TABLE II. Properties of real undirected networks: number of nodes n, edges m, natural cutoff degree k_c, maximal degree k_{max}, original assortativity coefficient r, and modified assortativity coefficient r_c. Note that, for the networks without superrich nodes, $k_c = k_{max}$ and $r_c = r$, and thus the values of k_c and r_c are not listed. Social networks: (a) network of electronic mail interchanges between members in a university [27]; (b) collaboration network of scientists who work on the condensed matter [28]; (c) online social network of Wealink.com [29]. Technological networks: (d) network of Roget's thesaurus [25]; (e) network of articles by and citing Lederberg from 1945 to 2002 [25]; (f) Internet snapshot at the level of autonomous systems [28]. Biological networks: (g) network of metabolic pathways for *E. Coli* [30]; (h) network of protein folding [30]; (i) neural network of *C. Elegans* [23].

Network	n	m	k_c	k_{max}	r	r_c
(a) e-mail	1134	10902		71	0.08	
(b) arXiv.org	40421	175693		278	0.19	
(c) Wealink.com	223624	273395	150	1657	−0.08	−0.42
(d) Roget's thesaurus	1022	5103		28	0.18	
(e) Citation network	8843	41609	105	1103	−0.10	−0.02
(f) Internet	22963	48436	110	2390	−0.20	−0.29
(g) Metabolic pathway	896	964		18	0.15	
(h) Protein folding	1287	33813		319	0.17	
(i) Neural net	307	2359	38	134	−0.16	0.03

hibit disassortative properties is not entirely clear [3,4].

The effect of superrich nodes on mixing patterns can explain why some technological and biological networks tend to have a negative r. Superrich nodes exist in these networks more commonly than in social networks, so many technological and biological networks exhibit a negative r induced by superrich nodes and are therefore claimed to be examples of disassortative mixing. Furthermore, we find that the technological and biological networks (without superrich nodes) have no tendency to be disassortative, such as networks (d), (g), and (h), for a positive r is commonly found in these networks. Our results indicate the conjecture that the disassortativity of degree is the normal state of a network in [16] may not be right. The normal state of degree mixing pattern in nonsocial networks is more like the neural mixing based on our paradigm.

V. CONCLUSION AND DISCUSSION

We show that in a finite-size network a very small number of superrich nodes will bias traditional measures of mixing patterns. We not only report that the traditional measures are not perfect measures but also propose a revised measure, which is better than the original one and reveals the real pattern of assortativity. First we propose a straightforward modification to the existing assortativity coefficient to both measure and counteract the effect of superrich nodes. Then, we develop an effective paradigm to quantify a real network in which different mixing patterns may appear among low-degree nodes and among high-degree nodes. Our method can capture the rich nature of the mixing properties in a real network. We can distinguish that the disassortativity of a network derives from the structural constraint or other reasons such as social and engineering factors.

Our results also indicate that the "robust yet fragile" nature of real networks (e.g., Internet) does not depend on the power-law degree distribution only, and superrich nodes have the same effect. For the random failure of the networks, a very small fraction of nodes with a very large degree will make any degree distribution network more robust than random graphs. On the other hand, superrich nodes are the real Achilles' heel of the Internet. The attack targeting these richest nodes causes the Internet to collapse faster than the [Erdhós and Rényi (ER)] graph, even the scale-free model without superrich nodes [Barabási and Albert (BA) network] as shown in Fig. 3 in [31]. Moreover, the cascade failure of one superrich node can lead more than 20% nodes of the Internet to be disconnected [32].

We demonstrate that superrich nodes critically change the way in which complex networks behave, and we have revised the false disassortative mixing in some nonsocial networks induced by superrich nodes. Revealing the intrinsic mixing patterns of complex networks masked by superrich nodes is crucial to study epidemic spreading, percolation on complex networks, and error and attack tolerance of real physical systems. Superrich nodes are one of the principal factors determining many aspects of the behavior of the overall network, although they are minority in number. We suggest that greater attention should be paid to the richest nodes when analyzing finite-size network data.

[1] S. Maslov and K. Sneppen, Science **296**, 910 (2002).

[2] L. K. Gallos, C.-M. Song, and H. A. Makse, Phys. Rev. Lett. **100**, 248701 (2008).

[3] M. E. J. Newman, SIAM Rev. **45**, 167 (2003).

[4] M. E. J. Newman, Phys. Today **61**(11), 33 (2008).

[5] L. Li, D. Alderson, J. C. Doyle, and W. Willinger, Internet Math. **2**, 431 (2006).

[6] P. Holme and J. Zhao, Phys. Rev. E **75**, 046111 (2007).

[7] M. E. J. Newman, Phys. Rev. Lett. **89**, 208701 (2002).

[8] M. E. J. Newman and J. Park, Phys. Rev. E **68**, 036122 (2003).

[9] P. Holme, C. R. Edling, and F. Liljeros, Soc. Networks **26**, 155 (2004).

[10] Y. Y. Ahn *et al.*, *Proceedings of the WWW2007 Conference* (IW3C2 and University of Calgary, Alberta, Canada, 2007), p. 835.

[11] M. E. J. Newman, Phys. Rev. E **67**, 026126 (2003).

[12] S. Boccaletti *et al.*, Phys. Rep. **424**, 175 (2006).

[13] L. da F. Costa *et al.*, Adv. Phys. **56**, 167 (2007).

[14] M. Boguñá and R. Pastor-Satorras, Phys. Rev. E **66**, 047104 (2002).

[15] S. Maslov, K. Sneppen, and A. Zaliznyak, Physica A **333**, 529 (2004).

[16] J. Park and M. E. J. Newman, Phys. Rev. E **68**, 026112 (2003).

[17] R. Pastor-Satorras, A. Vázquez, and A. Vespignani, Phys. Rev. Lett. **87**, 258701 (2001).

[18] A. Vázquez, M. Boguñá, Y. Moreno, R. Pastor-Satorras, and A. Vespignani, Phys. Rev. E **67**, 046111 (2003).

[19] A.-L. Barabási and R. Albert, Science **286**, 509 (1999).

[20] M. Boguñá, R. Pastor-Satorras, and A. Vespignani, Eur. Phys. J. B **38**, 205 (2004).

[21] S. N. Dorogovtsev and J. F. F. Mendes, Adv. Phys. **51**, 1079 (2002).

[22] S. Zhou and R. J. Mondragón, New J. Phys. **9**, 173 (2007).

[23] D. J. Watts and S. H. Strogatz, Nature (London) **393**, 440 (1998).

[24] R. Albert and A.-L. Barabási, Rev. Mod. Phys. **74**, 47 (2002).

[25] http://vlado.fmf.uni-lj.si/pub/networks/data/

[26] R. Milo *et al.*, Science **303**, 1538 (2004).

[27] http://deim.urv.cat/aarenas/data/welcome.htm

[28] http://www-personal.umich.edu/mejn/netdata/

[29] http://www.wealink.com/

[30] http://pil.phys.uniroma1.it/gcalda/cosinsite/

[31] R. Albert, H. Jeong, and A.-L. Barabási, Nature (London) **406**, 378 (2000).

[32] A. E. Motter and Y.-C. Lai, Phys. Rev. E **66**, 065102(R) (2002).