# Machine-Level Programming IV: Data

# **Today**

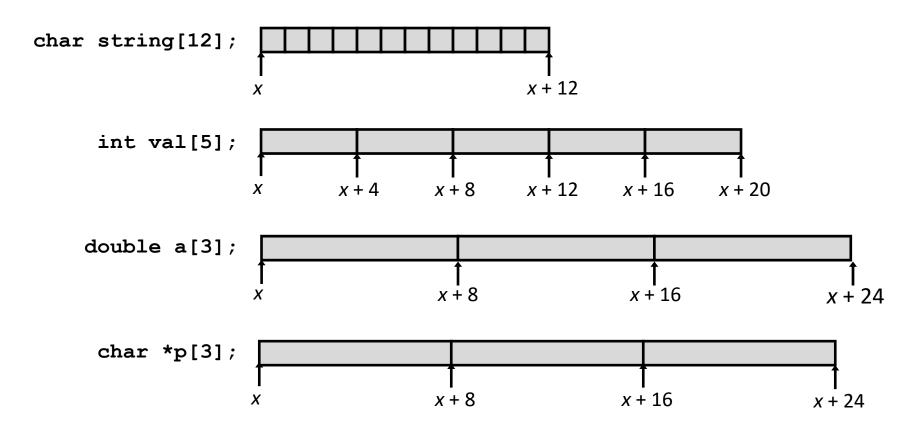
- Arrays
  - One-dimensional
  - Multi-dimensional (nested)
  - Multi-level
- Structures
  - Allocation
  - Access
  - Alignment
- Ploating Point

### **Array Allocation**

#### Basic Principle

 $T \mathbf{A}[L];$ 

- Array of data type T and length L
- Contiguously allocated region of L \* sizeof (T) bytes in memory

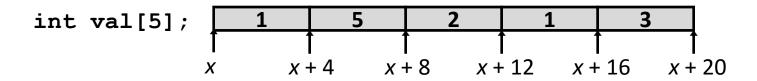


#### **Array Access**

#### **&** Basic Principle

```
T A[L];
```

- Array of data type T and length L
- **⊘**Identifier **A** can be used as a pointer to array element 0: Type *T\**

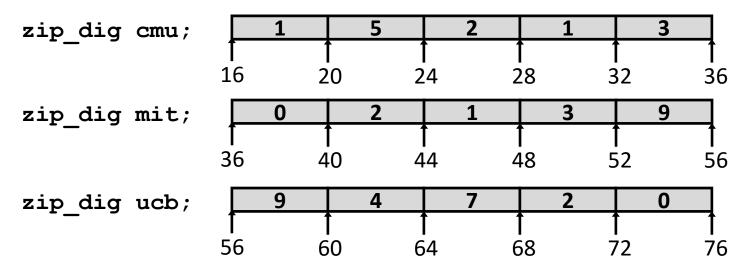


Reference	Type	Value
val[4]	int	3
val	int *	X
val+1	int *	x + 4
&val[2]	int *	x + 8
<b>v</b> al[5]	int	??
*(val+1)	int	5
val + i	int *	x + 4i

### **Array Example**

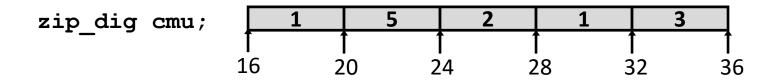
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration "zip dig cmu" equivalent to "int cmu[5]"
- **Example arrays were allocated in successive 20 byte blocks** 
  - Not guaranteed to happen in general

### **Array Accessing Example**



```
int get_digit
  (zip_dig z, int digit)
{
  return z[digit];
}
```

#### **IA32**

```
# %rdi = z
# %rsi = digit
movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register %rdi contains starting address of array
- Register %rsi contains array index
- Desired digit at %rdi + 4\*%rsi
- Use memory reference (%rdi,%rsi,4)

### **Array Loop Example**

```
void zincr(zip_dig z) {
   size_t i;
   for (i = 0; i < ZLEN; i++)
      z[i]++;
}</pre>
```

```
# %rdi = z
 movl $0, %eax
                        \# i = 0
                         # goto middle
 jmp .L3
.L4:
                         # loop:
 addl $1, (%rdi,%rax,4) # z[i]++
 addq $1, %rax
                         # i++
.L3:
                         # middle
 cmpq $4, %rax
                        # i:4
                         # if <=, goto loop</pre>
 jbe .L4
 rep; ret
```

# **Multidimensional (Nested) Arrays**

#### Declaration

 $T \mathbf{A}[R][C];$ 

- 2D array of data type T
- R rows, C columns
- Type *T* element requires *K* bytes

#### Array Size

**?** R \* C \* K bytes

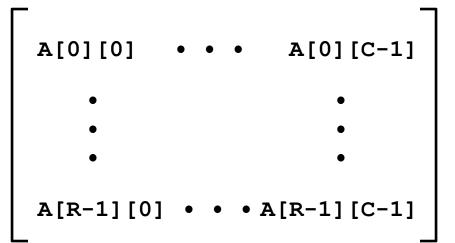
#### Arrangement

Row-Major Ordering

#### int A[R][C];

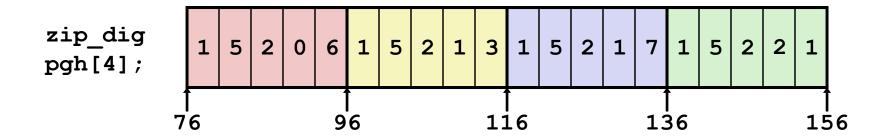
A [0] [0]	• • •	A [0] [C-1]	A [1] [0]	• • •	A [1] [C-1]	•	•	•	A [R-1] [0]	• • •	A [R-1] [C-1]	
-----------------	-------	-------------------	-----------------	-------	-------------------	---	---	---	-------------------	-------	---------------------	--

4\*R\*C Bytes



### **Nested Array Example**

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
   {{1, 5, 2, 0, 6},
    {1, 5, 2, 1, 3},
    {1, 5, 2, 1, 7},
    {1, 5, 2, 2, 1 }};
```

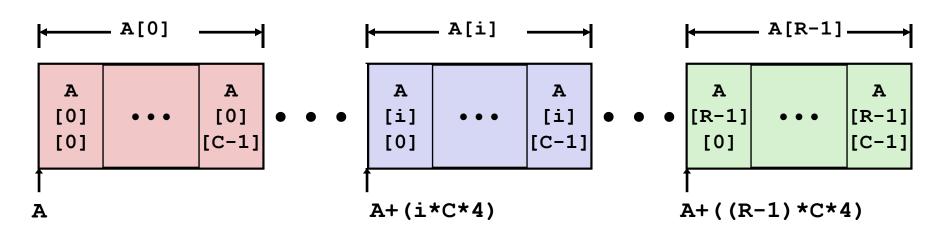


- "zip\_dig pgh[4]" equivalent to "int pgh[4][5]"
  - Variable **pgh**: array of 4 elements, allocated contiguously
  - Each element is an array of 5 int's, allocated contiguously
- "Row-Major" ordering of all elements in memory

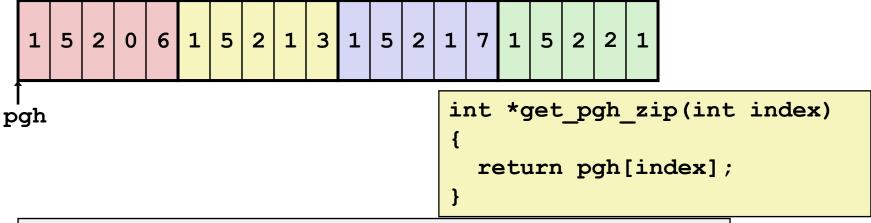
### **Nested Array Row Access**

- Row Vectors
  - **A**[i] is array of *C* elements
  - **Each** element of type *T* requires *K* bytes
  - $\diamond$  Starting address **A** + i \* (C \* K)

#### int A[R][C];



### **Nested Array Row Access Code**



```
# %rdi = index
leaq (%rdi,%rdi,4),%rax # 5 * index
leaq pgh(,%rax,4),%rax # pgh + (20 * index)
```

#### Row Vector

- **pgh[index]** is array of 5 **int**'s
- Starting address pgh+20\*index

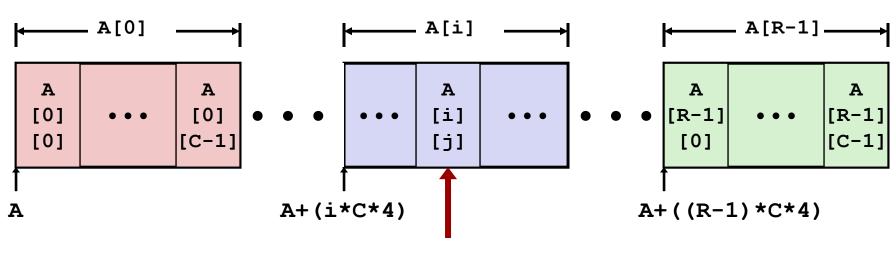
#### Machine Code

- Computes and returns address
- Compute as pgh + 4\*(index+4\*index)

### **Nested Array Element Access**

- Array Elements
  - $\mathbf{A}[\mathbf{i}][\mathbf{j}]$  is element of type T, which requires K bytes
  - **Address A** + i \* (C \* K) + j \* K = A + (i \* C + j) \* K





$$A+(i*C*4)+(j*4)$$

### **Nested Array Element Access Code**

```
1 5 2 0 6 1 5 2 1 3 1 5 2 1 7 1 5 2 2 1

pgh

int get_pgh_digit
    (int index, int dig)
{
    return pgh[index][dig];
}
```

```
leaq (%rdi,%rdi,4), %rax # 5*index
addl %rax, %rsi # 5*index+dig
movl pgh(,%rsi,4), %eax # M[pgh + 4*(5*index+dig)]
```

#### Array Elements

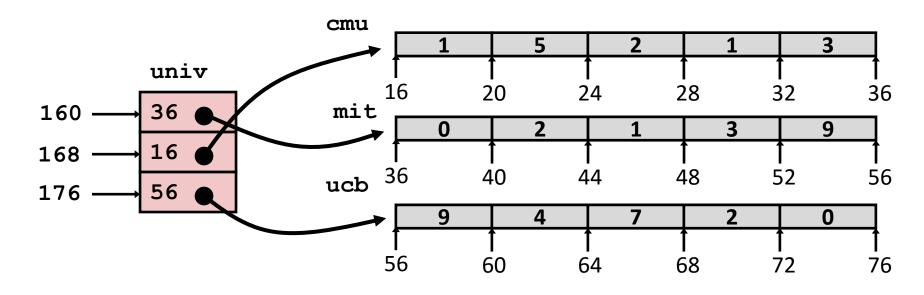
- pgh[index][dig] is int
- Address: pgh + 20\*index + 4\*dig
  - e = pgh + 4\*(5\*index + dig)

### **Multi-Level Array Example**

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

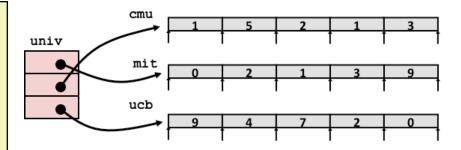
```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable univ denotes array of 3 elements
- Each element is a pointer
  - 8 bytes
- Each pointer points to array of int's



### **Element Access in Multi-Level Array**

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```



```
salq $2, %rsi # 4*digit
addq univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl (%rsi), %eax # return *p
ret
```

#### Computation

- Element access Mem [Mem [univ+8\*index]+4\*digit]
- Must do two memory reads
  - First get pointer to row array
  - Then access element within array

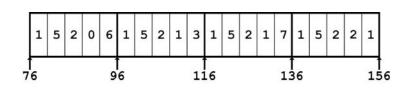
### **Array Element Accesses**

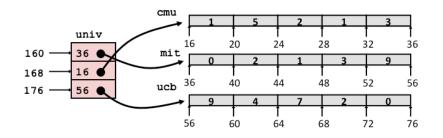
#### **Nested array**

```
int get_pgh_digit
  (size_t index, size_t digit)
{
  return pgh[index][digit];
}
```

#### Multi-level array

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```





Accesses looks similar in C, but address computations very different:

Mem[pgh+20\*index+4\*digit] Mem[Mem[univ+8\*index]+4\*digit]

# N X N Matrix Code

- Fixed dimensions
  - Know value of N at compile time
- Variable dimensions, explicit indexing
  - Traditional way to implement dynamic arrays
- Variable dimensions, implicit indexing
  - Now supported by gcc

#### 16 X 16 Matrix Access

#### Array Elements

- Address **A** + i \* (C \* K) + j \* K
- C = 16, K = 4

```
/* Get element a[i][j] */
int fix_ele(fix_matrix a, size_t i, size_t j) {
  return a[i][j];
}
```

```
# a in %rdi, i in %rsi, j in %rdx
salq $6, %rsi  # 64*i
addq %rsi, %rdi  # a + 64*i
movl (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]
ret
```

#### n X n Matrix Access

#### Array Elements

- Address **A** + i \* (C \* K) + j \* K
- C = n, K = 4
- Must perform integer multiplication

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
   return a[i][j];
}
```

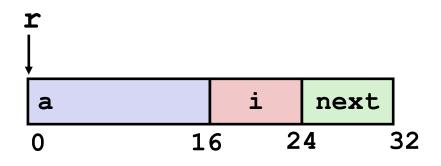
```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq %rdx, %rdi  # n*i
leaq (%rsi,%rdi,4), %rax # a + 4*n*i
movl (%rax,%rcx,4), %eax # a + 4*n*i + 4*j
ret
```

# **Today**

- Arrays
  - One-dimensional
  - Multi-dimensional (nested)
  - Multi-level
- Structures
  - Allocation
  - Access
  - Alignment
- Ploating Point

#### **Structure Representation**

```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```



- Structure represented as block of memory
  - **Big enough to hold all of the fields**
- Fields ordered according to declaration
  - Even if another ordering could yield a more compact representation
- Compiler determines overall size + positions of fields
  - Machine-level program has no understanding of the structures in the source code

### **Generating Pointer to Structure Member**

```
struct rec {
    int a[4];
    size_t i;
    struct rec *next;
};
```

```
r r+4*idx
a i next
0 16 24 32
```

#### Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as r + 4\*idx

```
int *get_ap
  (struct rec *r, size_t idx)
{
   return &r->a[idx];
}
```

```
# r in %rdi, idx in %rsi
leaq (%rdi,%rsi,4), %rax
ret
```

# **Following Linked List**

C Code

```
void set_val
  (struct rec *r, int val)
{
  while (r) {
    int i = r->i;
    r->a[i] = val;
    r = r->next;
  }
}
```

```
struct rec {
   int a[4];
   int i;
   struct rec *next;
};
```

i

24

next

32

#### ) [ 16 Element i

Register	Value
%rdi	r
%rsi	val

```
.L11:
                            # loop:
                            \# i = M[r+16]
 movslq 16(%rdi), %rax
         %esi, (%rdi,%rax,4) # M[r+4*i] = val
 movl
         24(%rdi), %rdi
                            # r = M[r+24]
 movq
                            #
 testq %rdi, %rdi
                                Test r
                            #
         .L11
                                if !=0 goto loop
 jne
```

r

a

### **Structures & Alignment**

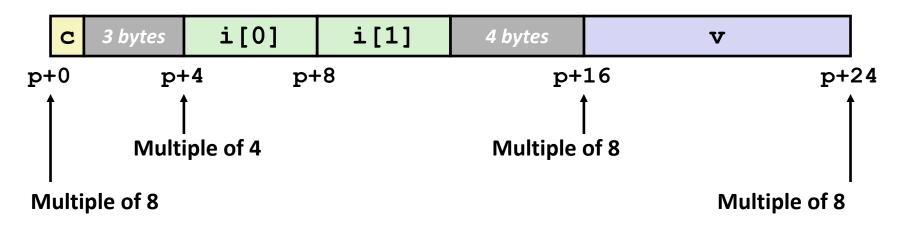
#### **Unaligned Data**

```
c i[0] i[1] v
p p+1 p+5 p+9 p+17
```

```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

#### Aligned Data

- Primitive data type requires **K** bytes
- Address must be multiple of K



### **Alignment Principles**

#### Aligned Data

- Primitive data type requires **K** bytes
- Address must be multiple of K
- Required on some machines; advised on x86-64

#### Motivation for Aligning Data

- Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
  - Inefficient to load or store datum that spans quad word boundaries
  - Virtual memory trickier when datum spans 2 pages

#### Compiler

Inserts gaps in structure to ensure correct alignment of fields

# **Specific Cases of Alignment (x86-64)**

- 1 byte: char, ...
  - no restrictions on address
- 2 bytes: short, ...
  - lowest 1 bit of address must be 02
- 4 bytes: int, float, ...
  - lowest 2 bits of address must be 00<sub>2</sub>
- 8 bytes: double, long, char \*, ...
  - lowest 3 bits of address must be 000<sub>2</sub>
- **16 bytes: long double** (GCC on Linux)
  - lowest 4 bits of address must be 0000<sub>2</sub>

# Satisfying Alignment with Structures

#### Within structure:

Must satisfy each element's alignment requirement

#### Overall structure placement

- Each structure has alignment requirement K
  - **& K** = Largest alignment of any element
- Initial address & structure length must be multiples of K

#### **2** Example:

K = 8, due to double element

```
        c
        3 bytes
        i [0]
        i [1]
        4 bytes
        v

        p+0
        p+4
        p+8
        p+16
        p+24

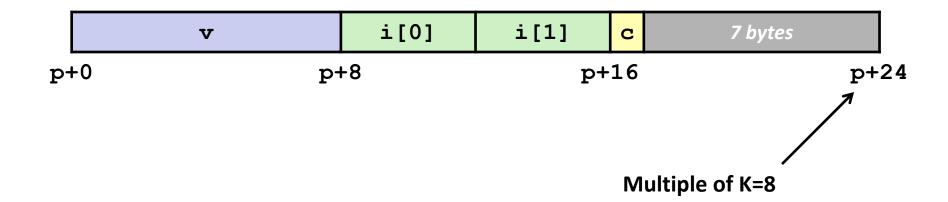
        Multiple of 4
        Multiple of 8
        Multiple of 8

Multiple of 8
```

### **Meeting Overall Alignment Requirement**

- For largest alignment requirement K
- **Overall structure must be multiple of K**

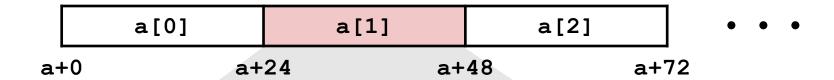
```
struct S2 {
  double v;
  int i[2];
  char c;
} *p;
```

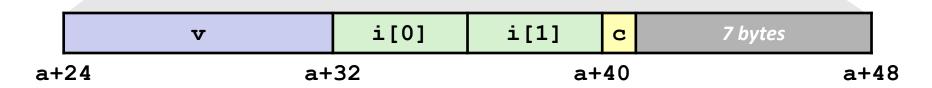


### **Arrays of Structures**

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

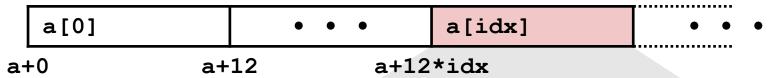
```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```





### **Accessing Array Elements**

- Compute array offset 12\*idx
  - **sizeof(S3)**, including alignment spacers
- Element j is at offset 8 within structure
- Assembler gives offset a+8
  - Resolved during linking



```
i 2 bytes v j 2 bytes a+12*idx+8
```

```
short get_j(int idx)
{
   return a[idx].j;
}
```

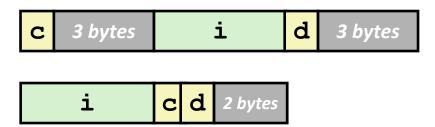
```
# %rdi = idx
leaq (%rdi,%rdi,2),%rax # 3*idx
movzwl a+8(,%rax,4),%eax
```

# **Saving Space**

Put large data types first

```
struct S4 {
  char c;
  int i;
  char d;
} *p;
struct S5 {
  int i;
  char c;
  char d;
} *p;
```

Effect (K=4)



# **Today**

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  - One-dimensional
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- Floating Point

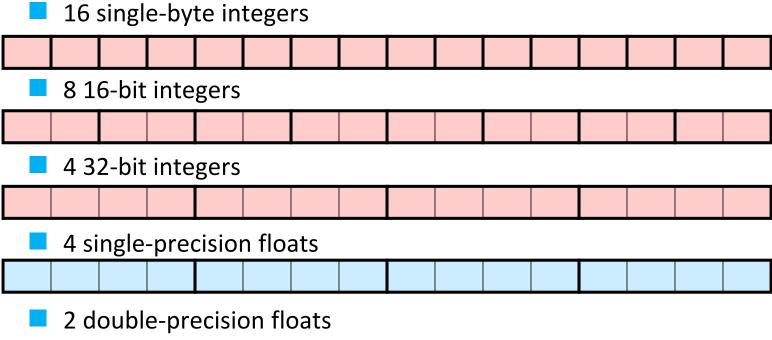
### **Background**

- History
  - **2** x87 FP
    - Legacy, very ugly
  - SSE FP
    - Supported by Shark machines
    - Special case use of vector instructions
  - AVX FP
    - Newest version
    - Similar to SSE
    - Documented in book

### **Programming with SSE3**

#### **XMM Registers**

16 total, each 16 bytes





- 1 single-precision float
- 1 double-precision float

### Scalar & SIMD Operations

Scalar Operations: Single Precision addss %xmm0,%xmm1 %xmm0 %xmm1 SIMD Operations: Single Precision addps %xmm0,%xmm1 %xmm0 %xmm1 Scalar Operations: Double Precision addsd %xmm0,%xmm1 %xmm0 %xmm1 35

#### **FP Basics**

- Arguments passed in %xmm0, %xmm1, ...
- Result returned in %xmm0
- All XMM registers caller-saved

```
float fadd(float x, float y)
{
    return x + y;
}
```

```
double dadd(double x, double y)
{
    return x + y;
}
```

```
# x in %xmm0, y in %xmm1
addss %xmm1, %xmm0
ret
```

```
# x in %xmm0, y in %xmm1
addsd %xmm1, %xmm0
ret
```

### **FP Memory Referencing**

- Integer (and pointer) arguments passed in regular registers
- **Problem** Problem Prob
- Different mov instructions to move between XMM registers, and between memory and XMM registers

```
double dincr(double *p, double v)
{
    double x = *p;
    *p = x + v;
    return x;
}
```

```
# p in %rdi, v in %xmm0
movapd %xmm0, %xmm1  # Copy v
movsd (%rdi), %xmm0  # x = *p
addsd %xmm0, %xmm1  # t = x + v
movsd %xmm1, (%rdi) # *p = t
ret
```

### Other Aspects of FP Code

- Lots of instructions
  - Different operations, different formats, ...
- Floating-point comparisons
  - Instructions ucomiss and ucomisd
  - Set condition codes CF, ZF, and PF
- **Using constant values** 
  - Set XMM0 register to 0 with instruction xorpd %xmm0, %xmm0
  - Others loaded from memory

### **Summary**

#### Arrays

- Elements packed into contiguous region of memory
- Use index arithmetic to locate individual elements

#### Structures

- Elements packed into single region of memory
- Access using offsets determined by compiler
- Possible require internal and external padding to ensure alignment

#### Combinations

Can nest structure and array code arbitrarily

#### Floating Point

Data held and operated on in XMM registers

Decl		An		*A <i>n</i>			
	Cmp	Bad	Size	Cmp	Bad	Size	
int A1[3]							
int *A2							

Cmp: Compiles (Y/N)

Bad: Possible bad pointer reference (Y/N)

Size: Value returned by sizeof

Decl		An		*A <i>n</i>			
	Cmp	Bad	Size	Cmp	Bad	Size	
int A1[3]	Y	N	12	Y	N	4	
int *A2	Y	N	8	Y	Y	4	



Cmp: Compiles (Y/N)

Bad: Possible bad pointer reference (Y/N)

Size: Value returned by sizeof

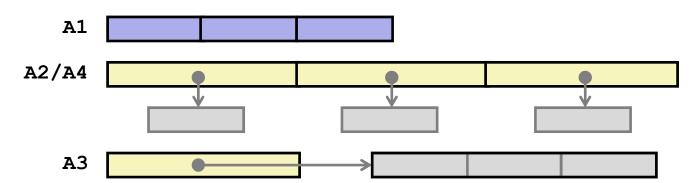
Decl	An				*A <i>n</i>			**A <i>n</i>		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size	
int A1[3]										
int *A2[3]										
int (*A3)[3]										
int (*A4[3])										

Cmp: Compiles (Y/N)

Bad: Possible bad pointer reference (Y/N)

Size: Value returned by sizeof

Decl	An				*A <i>n</i>			**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size	
int A1[3]	Y	N	12	Y	N	4	N	-	-	
int *A2[3]	Y	N	24	Y	N	8	Y	Y	4	
int (*A3)[3]	Y	N	8	Y	Y	12	Y	Y	4	
int (*A4[3])	Y	N	24	Y	N	8	Y	Y	4	

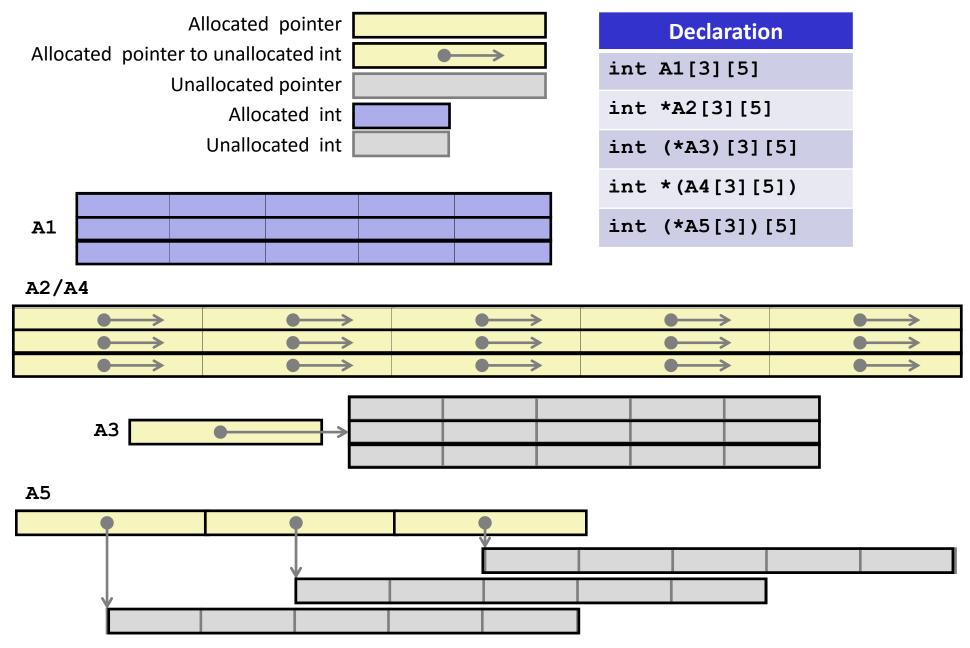


	Allocated pointer
	Unallocated pointer
	Allocated int
43	Unallocated int

Decl	An				*An		**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3][5]									
int *A2[3][5]									
int (*A3)[3][5]									
int *(A4[3][5])									
int (*A5[3])[5]									

- Cmp: Compiles (Y/N)
- Bad: Possible bad pointer reference (Y/N)
- Size: Value returned by sizeof

Decl	***An					
	Cmp Bad Siz					
int A1[3][5]		_				
int *A2[3][5]						
int (*A3)[3][5]						
int *(A4[3][5])						
int (*A5[3])[5]						



Decl	An				*An		**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3][5]	Y	N	60	Y	N	20	Y	N	4
int *A2[3][5]	Y	N	120	Y	N	40	Y	N	8
int (*A3)[3][5]	Y	N	8	Y	Y	60	Y	Y	20
int *(A4[3][5])	Y	N	120	Y	N	40	Y	N	8
int (*A5[3])[5]	Y	N	24	Y	N	8	Y	Y	20

- Cmp: Compiles (Y/N)
- Bad: Possible bad pointer reference (Y/N)
- Size: Value returned by sizeof

Decl	***An				
	Cmp	Size			
int A1[3][5]	N	-	-		
int *A2[3][5]	Y	Y	4		
int (*A3)[3][5]	Y	Y	4		
int *(A4[3][5])	Y	Y	4		
int (*A5[3])[5]	Y	Y	4		