

# PEP in Bregman Method

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## 1 Choice of Bregman Divergence

Finished:

- (last time) Derived Bound for Bregman Proximal Point
- Complete proof for BPG

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In progress:

- Bregman Halpern iteration for Shannon entropy (numerically unbounded)
- general step size  $\lambda_i$  for BPP's PEP (no time)

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## 1 Choice of Bregman Divergence

Classical kernel functions:

- $h(x) = \frac{1}{2}\|x\|^2, \nabla h(x) = x$
- $h(x) = \sum_{i=1}^d x^{(i)} \log x^{(i)}, \nabla h(x) = (1 + \log x^{(i)})_i \ (x \geq 0)$
- $h(x) = \sum_{i=1}^d -\log x^{(i)}, \nabla h(x) = (-\frac{1}{x^{(i)}})_i \ (x \geq 0)$
- $h(x) = \frac{1}{4}\|x\|^4 + \frac{1}{2}\|x\|^2, \nabla h(x) = \|x\|^2 x + x$

Last time we discussed the Shannon entropy:

- $h(x) = \sum_{i=1}^d x^{(i)} \log x^{(i)}, \nabla h(x) = (1 + \log x^{(i)})_i \ (x \geq 0)$

$$\nabla h(x) \leq x$$

But this cannot be used in PEP since  $\nabla h(x)$  and  $x$  are encoded in the Gram matrix. ( $G^{s,x}, s_i = \nabla h(x_i)$ )

Problem:

$$\min_{x \in \mathbb{R}^d} \phi(x) \quad (\text{P})$$

Method: Bregman Halpern's iteration

$$x_{k+1} = \nabla h^*[\lambda_k \nabla h(x_0) + (1 - \lambda_k) \nabla h(R_{\partial\phi}^h(x_k))] \quad (\text{BH})$$

Bregman resolvent  $R_T^h = (\nabla h + \alpha T)^{-1} \circ \nabla h$



# Formulation

(relaxed) PEP:

$$\max \quad \min D_h(x_{i+1}, x_i) \quad (\text{sdp-PEP})$$

$$\text{s.t.} \quad \begin{cases} h_{i-1} - h_i + G_{i-1,i}^{s,x} - G_{i-1,i-1}^{s,x} \leq 0 \\ h_{i+1} - h_i + G_{i+1,i}^{s,x} - G_{i+1,i+1}^{s,x} \leq 0 \\ h_i - h_* + G_{i,*}^{s,x} - G_{i-1,i-1}^{s,x} \leq 0 \\ \phi_{i+1} - \phi_i + G_{i+1,i}^{w,x} - G_{i+1,i+1}^{w,x} \leq 0 \\ \phi_i - \phi_* + G_{i,*}^{w,x} - G_{i,i}^{w,x} \leq 0 \\ \langle s_k + \alpha w_k - \lambda_{k-1}(s_0 + \alpha w_0) - (1 - \lambda_{k-1})s_{k-1}, x_j \rangle = 0 \end{cases}$$

(assume  $x_* = \phi_* = h_* = 0$ )

$$P = (P_x, P_s, P_w) = (x_0, \dots, x_N, s_*, s_0, \dots, s_N, w_*, w_0, \dots, w_N), \\ G = P^T P, \quad G^{s,x} = P_s^T P_x, \quad G^{w,x} = P_w^T P_x.$$

We turn to find constraints on  $\langle \nabla h(x), x \rangle$ :

- $h(x) = \frac{1}{2}\|x\|^2$ ,  $\langle \nabla h(x), x \rangle = 2h(x) = \|x\|^2$
- $h(x) = \sum_{i=1}^d x^{(i)} \log x^{(i)}$ ,  $\langle \nabla h(x), x \rangle = \sum_{i=1}^d x^{(i)} + h(x)$  ( $x \geq 0$ )
- $h(x) = \sum_{i=1}^d -\log x^{(i)}$ ,  $\langle \nabla h(x), x \rangle = -d \geq -\sum_{i=1}^d x^{(i)} - h(x)$  ( $x \geq 0$ )
- $h(x) = \frac{1}{4}\|x\|^4 + \frac{1}{2}\|x\|^2$ ,  $\langle \nabla h(x), x \rangle = 4h(x) - \|x\|^2$

# Bregman Halpern's PEP

For Euclidean case, we still not recover the result given by [Lieder, 2021].  
The SDP construction in [Lieder, 2021] is different for it considers the general mapping  $T$ .  
(Weird)

For the entropy case, we have not found a way to put in the constraints  $x > 0$  and relative to  $\sum_{i=1}^d x^{(i)}$ , even we already turn to  $\langle \nabla h(x), x \rangle$ .  
(Turn SDP to QCQP?)

Have tried to use slack variables  $s_k + t_k = x_k$ , but  $t_k$  still appears in inner product.

Have tried to write  $x_k^{(i)} = (z_k^{(i)})^2$  to guarantee nonnegative, but inner product is not expressible.

In Bregman's case, every convexity condition requires inner products with  $x_i$ .

But the algorithm iteration can only be expressed in terms of  $s_i = \nabla h(x_i)$ , so there's no simple way to eliminate  $x_i$  from the inner product (unlike in GM Euclidean case in Drori's and Kim's)

So it seems impossible to add constraints on  $x_i$  alone. It is always occupied.



Lieder, F. (2021).

On the convergence rate of the halpern-iteration.

*Optimization Letters*, pages 1–14.