091M4041H-Assignment 4

Algorithm Design and Analysis

Notice:

1. Please submit your answers in hard copy **AND** submit a digital version to UCAS website https://sep.ucas.ac.cn.
2. Hard copy should be submitted before 9 am. Dec. 21. and digital version should be submitted before 11pm. Dec. 21.
3. Please finish problems 7, 8 on OJ and choose at least two problems from Problem 1-6.
4. Integer Linear Programming is different from the classic Linear Programming that some extra constraints such as

$$x_i \text{ is an integer}, \text{ for all } i = 1, 2, \cdots, n$$

or

$$x_i \in \{0, 1\}, \text{ for all } i = 1, 2, \cdots, n$$

are added.

5. When you give the formulation of an LP or ILP, you should explain all mathematical symbols you are using if not appearing in the problem, and interpret the constraints if necessary.

1 Linear-inequality feasibility 可行解判断问题 / aux !

Given a set of m linear inequalities on n variables $x_1, x_2, \cdots, x_n$, the **linear-inequality feasibility problem** asks if there is a setting of the variables that simultaneously satisfies each of the inequalities.

Show that if we have an algorithm for linear programming, we can use it to solve the linear-inequality feasibility problem. The number of variables and constraints that you use in the linear-programming problem should be polynomial in n and m.

2 Interval Scheduling Problem

A teaching building has m classrooms in total, and n courses are trying to use them. Each course i ($i = 1, 2, \cdots, n$) only uses one classroom during time interval $[S_i, F_i)$ ($F_i > S_i > 0$). Considering any two courses can not be carried on in a same classroom at any time, you have to select as many courses as possible and arrange them without any time collision. For simplicity, suppose $2n$ elements in the set $\{S_1, F_1, \cdots, S_n, F_n\}$ are all different.

x_i indicate whether class is selected. / $\max \sum_{i=1}^{n} x_i$

constraint: avoid collision

classroom 1: $a_{11}x_1 + a_{12}x_2 + \cdots a_{1n}x_n$

同一时间的课不超过m?

$$2n \begin{pmatrix} \overbrace{a_{11} \quad \cdots \quad a_{1n}}^{n \text{ courses}} \\ a_{21} \quad \quad \quad \\ a_{31} \quad \quad \quad \\ \vdots \quad \quad \quad \vdots \\ a_{2n1} \quad \cdots \quad a_{2n\cdot n} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \le b = \begin{pmatrix} m \\ \vdots \\ m \end{pmatrix}$$

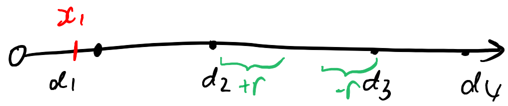
Prob 3. Target = $\min\left(\max_{1\le i\le n-1}\{x_{i+1}-x_i\}\right)$ let $t \ge x_{i+1}-x_i$

Constraints $-r \le x_1 - d_1 \le r$

$-r \le x_2 - d_2 \le r$

Please use ILP to solve this problem, then construct an instance and use GLPK or Gurobi or other similar tools to solve it.

3 Gas Station Placement



Let's consider a long, quiet country road with towns scattered very sparsely along it. Sinopec, largest oil refiner in China, wants to place gas stations along the road. Each gas station is assigned to a nearby town, and the distance between any two gas stations being as small as possible. Suppose there are n towns with distances from one endpoint of the road being $d_1, d_2, \cdots, d_n$. n gas stations are to be placed along the road, one station for one town. Besides, each station is at most r far away from its correspond town. $d_1, \cdots, d_n$ and r have been given and satisfied $d_1 < d_2 < \cdots < d_n$, $0 < r < d_1$ and $d_i + r < d_{i+1} - r$ for all i. The objective is to find the optimal placement such that the maximal distance between two successive gas stations is minimized.

Please formulate this problem as an LP.

4 Stable Matching Problem

n men $(m_1, m_2, \cdots, m_n)$ and n women $(w_1, w_2, \cdots, w_n)$, where each person has ranked all members of the opposite gender, have to make pairs. You need to give a stable matching of the men and women such that there is no unstable pair. (A matching is *unstable* if: there is an element A of the first matched set which prefers some given element B of the second matched set over the element to which a is already matched, and B also prefers A over the element to which B is already matched.) Please choose one of the two following known conditions, formulate the problem as an ILP, construct an instance and use GLPK or Gurobi or other similar tools to solve it.

1. You have known that for every two possible pairs (man m_i and woman w_j, man m_k and woman w_l), whether they are stable or not. If they are stable, then $S_{i,j,k,l} = 1$; if not, $S_{i,j,k,l} = 0$. $(i, j, k, l \in \{1, 2, \cdots, n\})$

2. You have known that for every man m_i, whether m_i likes woman w_j more than w_k. If he does, then $p_{i,j,k} = 1$; if not, $p_{i,j,k} = 0$. Similarly, if woman w_i likes man m_j more than m_k, then $q_{i,j,k} = 1$, else $q_{i,j,k} = 0$. $(i, j, k \in \{1, 2, \cdots, n\})$

5 Duality

Please write the dual problem of the MultiCommodityFlow problem in *Lec8.pdf*, and give an explanation of the dual variables.

Please also construct an instance, and try to solve both primal and dual problem using GLPK or Gurobi or other similar tools.

6 Dual Simplex Algorithm

For the problem

minimize

$$-7x_1 + 7x_2 - 2x_3 - x_4 - 6x_5$$

subject to:

$$\begin{aligned} 3x_1 - x_2 + x_3 - 2x_4 &= -3 \\ 2x_1 + x_2 + x_4 + x_5 &= 4 \\ -x_1 + 3x_2 - 3x_4 + x_6 &= 12 \\ x_i &\geqslant 0, (i = 1, ..., 6) \end{aligned}$$

Implement dual simplex algorithm with your favorate language to solve this problem, and make comparison result with GLPK or Gurobi or other similar tools.

7 Airplane Landing Problem

Original Problem

$\max\left(\min_{1\le i\le 4} \{ (x_{i+1} - x_i) \} \right)$

$s_i \le x_i \le t_i$

Alt. Problem

$\Downarrow$ let $d = \min_{1\le i\le 4} \{ (x_{i+1} - x_i) \}$

$\max d$

$s_i \le x_i \le t_i$

$d \le x_{i+1} - x_i$

$x_i - x_{i+1} + d \le 0$

With human lives at stake, an air traffic controller has to schedule the airplanes that are landing at an airport in order to avoid airplane collision. Each airplane i has a time window $[s_i, t_i]$ during which it can safely land. You must compute the exact time of landing for each airplane that respects these time windows. Furthermore, the airplane landings should be stretched out as much as possible so that the minimum time gap between successive landings is as large as possible.

For example, if the time window of landing three airplanes are [10:00-11:00], [11:20-11:40], [12:00-12:20], and they land at 10:00, 11:20, 12:20 respectively, then the smallest gap is 60 minutes, which occurs between the last two airplanes.

Given n time windows, denoted as $[s_1, t_1], [s_2, t_2], \cdots, [s_n, t_n]$ satisfying $s_1 < t_1 < s_2 < t_2 < \cdots < s_n < t_n$, you are required to give the exact landing time of each airplane, in which the smallest gap between successive landings is maximized.

Please formulate this problem as an LP, construct an instance and use GLPK or Gurobi or other similar tools to solve it.

8 Volunteer Recruitment

Suppose you will recruit a group of volunteers for a coming event. It is estimated that this event will take N days to complete, and the $i(th)$ day needs at least A_i volunteers. The number of kinds of volunteers is M. The volunteers of $i(th)$ kind can volunteer from the S_i day to the F_i day and the recruit fee is C_i. In order to do his job well, you hope to recruit enough volunteers with least money. Please formulate this problem and solve it using simplex algorithm with your favorate language.