

Exploration of Cryptography



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Class 02

Basic Notions of Cryptography



Cryptography

Protection of data by transformations that turn useful and comprehensible plaintext into scrambled and meaningless ciphertext under control of secret keys.



What Drives Cryptography?

Like every other field of engineering, cryptography responds to the requirements by using the available tools, driven by the imaginations of its practitioners.



Cryptography Guarantees Confidentiality

- Only authorized receivers who know secret keys can decrypt
- Eavesdropper may intercept message, but cannot understand it.

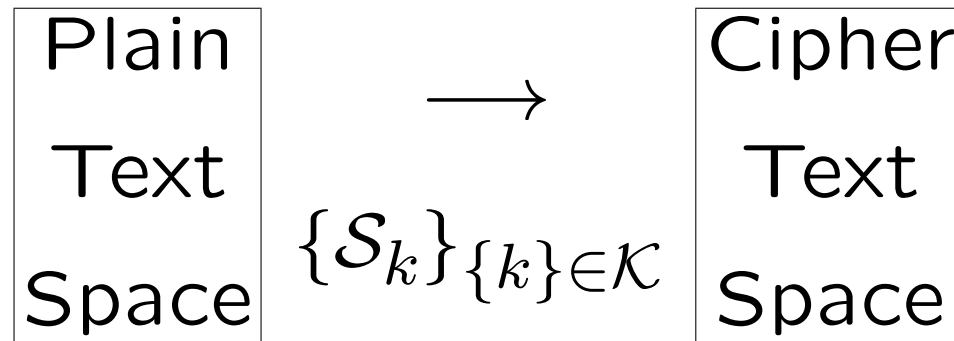


Cryptography Guarantees Authenticity

- Message sent by Intruder will decrypt to nonsense.
- Intruder may inject messages, but cannot 'get them accepted.'



Cryptographic System



Representation of a Cipher

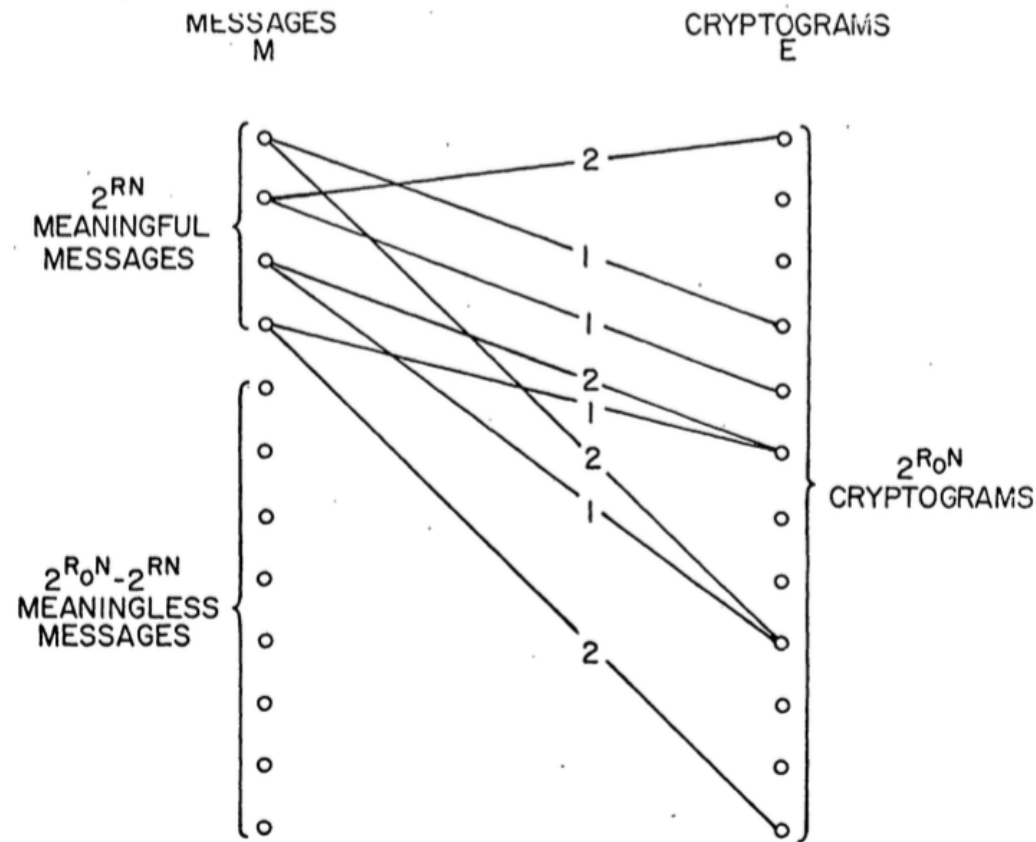


FIG. 1 REPRESENTATION OF A CIPHER



Representation of Encryption

Unkeyed map from plaintexts to plaintext representations.

Keyed map from plaintext representations to ciphertext representations.

Unkeyed map ciphertext representations to ciphertext.



Plaintext

Information In Usable
and Comprehensible Form

(Definition by intention.)



Examples of Plaintext

- Written human language
- Spoken Language
- Accounting Ledgers
- Seismic Data



Examples of Plaintext (Cont'd)

- Computer program
- Formal commands (EFT, C^2)
- Somebody else's ciphertext



Representations of Plaintext

- Messages drawn from a fixed set
- Fixed length strings over an alphabet
- Finite length strings ...



Characteristics of Plaintext

- Message (e.g., letter) frequencies
- Pair and triple frequencies
- Contact characteristics
- Gramatical restrictions



Ciphertext

Information in Useless
Scrambled Form

(Definition by intention.)



Characteristics of Ciphertext

Unpredictable (random) at some level.

Scrambled and Useless

- Unreadable
- ‘Destroyed’ by Alteration
- Looks Like Random Noise



Substitution

Permutation of the space of
plaintexts.



Full Substitution

Plaintext space, or message space, or alphabet: M of n elements.

$$k : M \mapsto M \quad \text{invertible}$$

Number of keys = $(\text{Card}(M))!$

where each $k \in S_{\text{Card}(M)}$.



Examples of Full Substitution

- Monoalphabetic substitution on Roman (or ASCII, etc.) alphabet: S-box.

$$k : \{A, \dots, Z\} \mapsto \{A, \dots, Z\}$$

(Induces function on Roman-alphabet text.) each $k \in S_{26}$.



Examples of Full Substitution (cont'd)

- Prearranged message code

sell short	TXDNL
go public	PULID
fire CEO	HADOS
file Chap. 11	AMBIT

...

...



Examples of Full Substitution (cont'd)

- Substitution on letters
- Substitution on words
- Substitution on phrases
- Substitution on Sentences



Maximal and Minimal Substitutions

Minimal Substitution: The Cyclic Shifts

Let: $a \rightarrow 0, b \rightarrow 1, \dots$ and let
 $k \in \{0, \dots, 25\}$ map $\ell \mapsto \ell + k$.

Full or Maximal Substitution

For any set $\{p_1, \dots, p_n\}$ of plaintexts
and any set of ciphertexts
 $\{c_1, \dots, c_n\}$ there is a key such that
 $k(p_i) = c_i$.



Transposition

Permute Elements of Message

w h a t i s i n i t f o r m e

h m i i t s i o t w r n e a f



Identification Friend or Foe

Mark XII IFF

Challenger

Responder

32-bit Random

32 bits in

Encrypt Locally

4 bits out

compare



General Systems and Specific Keys

- Cryptographic Systems
 - Standardized and Public
- Cryptographic Keys
 - Unique and Secret

Like Physical Locks and Keys



Cryptographic Systems

- Standardized
- Economy of Scale
- In Principle Public



Cryptographic Keys

- Very Large Supply
- Each Key Unique
- Must Be Kept Secret

All Secrecy Resides In Key



Primary and Secondary Keys

The primary key is the one changed most frequently, perhaps with every message.

The secondary key is changed less frequently, perhaps every day, week, or month.



Cryptanalysis

Produce Plaintext from Ciphertext (Analytic)

or

Produce Ciphertext from Plaintext (Synthetic)

(Without Prior Knowledge of the Key)



Analytic Cryptanalysis

Analytic Cryptanalysis

- Extract Information from Cryptograms
- Violates Confidentiality
- Better Known Problem



Synthetic Cryptanalysis

Synthetic Cryptanalysis

- Create False Messages
- Violates Authenticity
- Harder Problem



Analytic vs. Synthetic Cryptanalysis

Synthetic crypta can be much harder
than analytic.



Rolling Code Speech Scrambler

- Filter Speech into multiple bands.
- Rearrange the Bands.
- Change Rearrangement several times a second.



Food for Thought

Can synthetic ever be easier than analytic?



Applications of Cryptanalysis

- Entertainment
- Production of Intelligence
- Certification of Systems



Cryptanalytic Circumstances

Cryptanalyst is Always Presumed
to Know the General System

- Matter of Definition
- Sound Security Practice



Types of Attack

- Ciphertext Only
- Known Plaintext
- Chosen Plaintext



Ciphertext Only

- Statistics of Language
- Probable Words



Known Plaintext

- Common Words and Phrases
- Previously Secret Material
- Good Certification Assumption



Chosen Plaintext

- Best certification assumption
- Either plain or ciphertext
- Can be interactive



Oneway Functions

- Easy to Compute Forward
- Hard to Go Back



Easy to Raise Numbers to Powers

$$x \rightarrow x^7$$

Hard to Extract Roots

$$x \rightarrow \sqrt[7]{x}$$



Easy to Multiply

$$127 \times 997 = 126,619$$

Hard to Factor

126,619



Public Key Cryptography

- Keys Come In Inverse Pairs
- Given One — Can't Find the Other



Fundamental Property

- Public Key
 - In telephone directory
- Secret Key
 - In subscriber's safe
(or cryptoequipment)



Solves Both Problems

Key Distribution: Encrypt With Public Key

Digital Signature: Encrypt With Secret Key



Signed and Sealed Message



Information Theory

- Information resolves uncertainty
- Perfect secrecy
- Ideal secrecy



Finite Automaton

A finite automaton consists of a set of states S , a set of inputs I , a set of outputs O , and a distinguished starting state S_0 , together with two rules: a change of state map:

$$CS : I \times S \mapsto S$$

and an output map:



Finite Automaton (Cont'd)

$$Out : S \mapsto O \quad \text{Moore}$$

or

$$Out : S \times O \mapsto O \quad \text{Mealy}$$

