

INTRODUCTION

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What is a data structure?

- An organization of information, usually in memory, for better algorithm efficiency.
- Or, a way to store and organize data in order to facilitate access and modifications.

$$2n^5 - 3n^2 + 11n + 27$$

0	1	2	3	4	5	6
27	11	-3	0	0	2	0

- 又可分為:
 - Linear data structure: 必須循序地存取 (如linked list, stack, queue)
 - Non-linear data structure: 可以不循序的存取 (如tree, graph)

What is an algorithm?

- An algorithm is any well-defined **computational procedure** that takes some value, or set of values, as **input** and produces some value, or set of values, as **output**.
- An algorithm is a tool for solving a well-specified **computational problem**.
- Computational problem → **input/output relationship**
- The algorithm describes a specific computational procedure for achieving that input/output relationship.

What is an algorithm?

- Example:
- Sorting problem:
- **Input:** A sequence of n numbers $\langle a_1, a_2, \dots, a_n \rangle$
- **Output:** A permutation (reordering) $\langle a'_1, a'_2, \dots, a'_n \rangle$ of the input sequence such that $a'_1 \leq a'_2 \leq \dots \leq a'_n$
- An **instance** of the sorting problem: $\langle 31, 41, 59, 26, 41, 58 \rangle$
- A sorting algorithm should return as output the sequence $\langle 26, 31, 41, 41, 58, 59 \rangle$.

What is an algorithm?

- All algorithm must satisfy the following criteria:
 - Input: 外部給的資訊(可以是零個或多個)
 - Output: 產生的結果(至少一個)
 - Definiteness: 每一個指令都是清楚而不模糊的
 - Finiteness: 所有的狀況下(所有的input), 演算法會在有限步驟之後結束
 - Effectiveness: 每一個指令都必須是簡單可以直接執行的 (必須可以執行)

Example

- Statement 1: **“Is $n=2$ the largest value of n for which there exist positive integers x , y , and z such that $x^n + y^n = z^n$ has a solution?”**

Definiteness

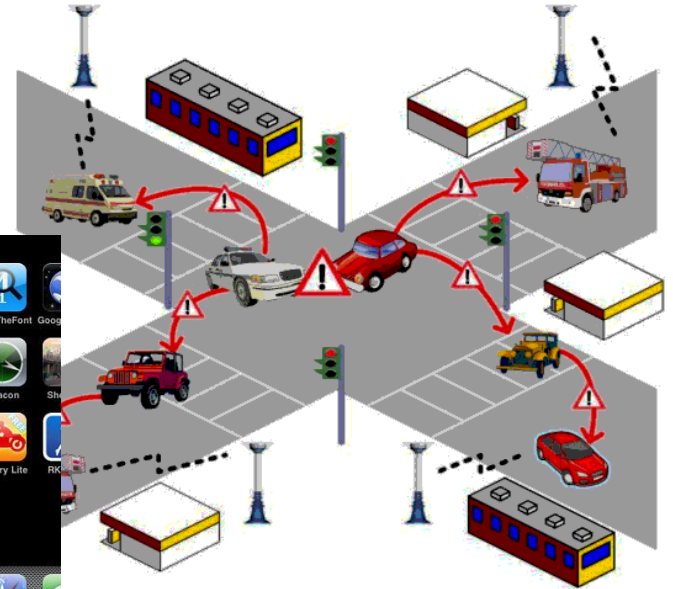
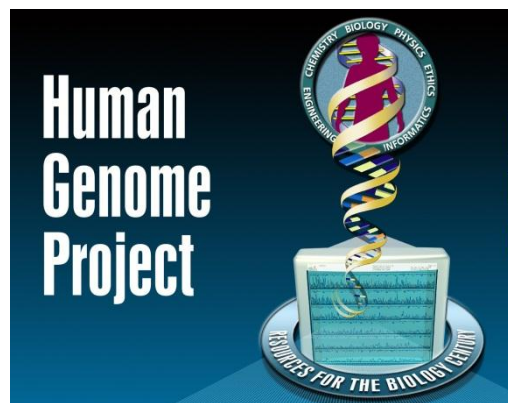
- Statement 2: **“Store 5 divided by zero into x and go to statement \neg .”**

Effectiveness

- Which criterion do they violate?
 - Input
 - Output
 - Definiteness
 - Finiteness
 - Effectiveness

Why are algorithms/data structure important?

- 它們被用在生活中的每個層面:



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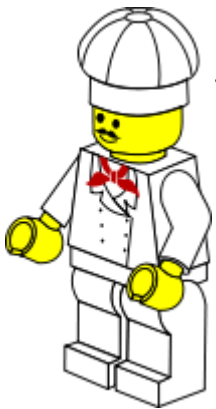


Why are algorithms/data structure important?

- Q: 如果電腦無限快/記憶體免錢，我們還需要研究資料結構與演算法嗎?
- A: Yes. 我們仍然需要確認我們想出來的解法會停止(不會無窮地執行下去)，而且**每次都產生正確**的答案。
- 在這個假想的世界中，任何**正確**的解法都適用，因此我們通常會選**最容易**實作的方法。
- 但是在真實的世界裡:
 - 電腦不是無限快(計算需要**時間**)
 - 記憶體不是免錢(儲存資料需要**空間**)
- 因此我們需要學習如何好好利用這些**資源**來解決問題!

How do we describe an algorithm?

- Human language (English, Chinese, ...)
- Programming language
- A mix of the above



1. 拿平底鍋
2. 拿沙拉油
 1. 我們有油嗎?
 1. 有的話, 倒一茶匙的沙拉油到鍋子裡
 2. 沒有的話, 我們想要買油嗎?
 1. 是的話, 就去全聯買一罐沙拉油
 2. 如果不想的話, 只好先不煮了.
3. 打開火爐, ...

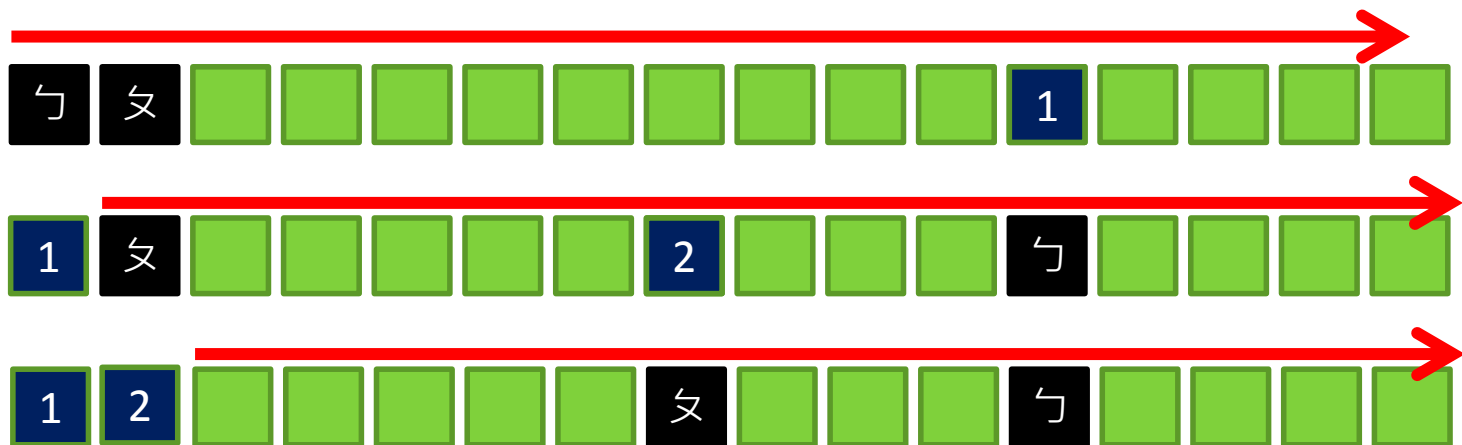
Example: Selection Sort

Sorting problem:

Input: A sequence of n numbers $\langle a_1, a_2, \dots, a_n \rangle$

Output: A permutation (reordering) $\langle a'_1, a'_2, \dots, a'_n \rangle$ of the input sequence such that $a'_1 \leq a'_2 \leq \dots \leq a'_n$

- Integers are stored in an array, list. The i -th integer is stored in $\text{list}[i]$, $0 < i < n$.
- Solution: From those integers that are currently unsorted, find the smallest and place it next in the sorted list.



Example: Selection Sort

- First attempt:

```
for (i=0; i<n; ++i) {
```

```
    Examine list[i] to list[n-1] and suppose that  
    the smallest integer is at list[min]; Task 1
```

```
        Interchange list[i] and list[min];
```

```
    }
```

Task 2

Task 2

Task 2

```
void swap(int *x, int *y) {  
    int temp = *x;  
    *x=*y;  
    *y=temp;  
}
```

Or

```
#define SWAP(x,y,t) ((t)=(x), (x)=(y), (y)=(t))
```

Task 1

Task 1

```
min=i;  
for (j=i; j<n; ++j)  
    if (list[j]<list[min])  
        min=j;
```

Final program

```
#include <stdio.h>
#include <math.h>

#define MAX-SIZE 101
#define SWAP(x,y,t) ((t) = (x), (x) = (y), (y) = (t))

void sort(int [],int); /*selection sort */
void main(void)
{
    int i,n;
    int list[MAX-SIZE];

    printf("Enter the number of numbers to generate: ");
    scanf (" %d", &n) ;
    if( n < 1 || n > MAX-SIZE) {
        fprintf(stderr, "Improper value of n\n");
        exit(EXIT_FAILURE);
    }

    for (i = 0; i < n; i++) { /*randomly generate numbers*/
        list[i] = rand() % 1000;
        printf("%d ",list[i]);
    }
```

```
    sort(list,n);  
    printf("\n Sorted array:\n ");  
    for (i = 0; i < n; i++) /* print out sorted numbers */  
        printf("%d ",list[i]);  
    printf("\n");  
}
```

```
void sort(int list[],int n)  
{  
    int i, j, min, temp;  
  
    for (i = 0; i < n-1; i++) {  
        min = i;  
        for (j = i+1; j < n; j++)  
            if (list[j] < list[min])  
                min = j;  
        SWAP(list[i],list[min],temp);  
    }  
}
```

How do we prove that it is correct?

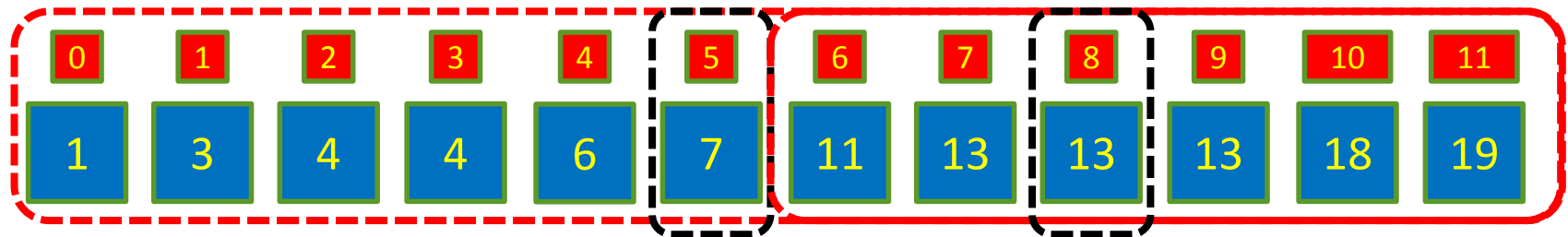
- **[Theorem]** Function $\text{sort}(\text{list}, n)$ correctly sorts a set of $n \geq 1$ integers. The result remains in $\text{list}[0], \dots, \text{list}[n-1]$ such that $\text{list}[0] \leq \text{list}[1] \leq \dots \leq \text{list}[n-1]$.
- **Proof:**
When the outer for loop completes its iteration for $i=q$, we have $\text{list}[q] \leq \text{list}[r], q < r < n$. Further, on subsequent iterations, $i > q$ and $\text{list}[0]$ through $\text{list}[q]$ are unchanged. Hence following the last iteration of the outer for loop (i.e., $i=n-2$), we have $\text{list}[0] \leq \text{list}[1] \leq \dots \leq \text{list}[n-1]$.

Example: Binary Search

- Input:
 - *searchnum*: the number to be found
 - *list*: sorted array, size n , and $list[0] \leq list[1] \leq \dots \leq list[n - 1]$
- Output:
 - -1 if *searchnum* is not found in *list*
 - the index of *searchnum* in *list*[] if *searchnum* is found

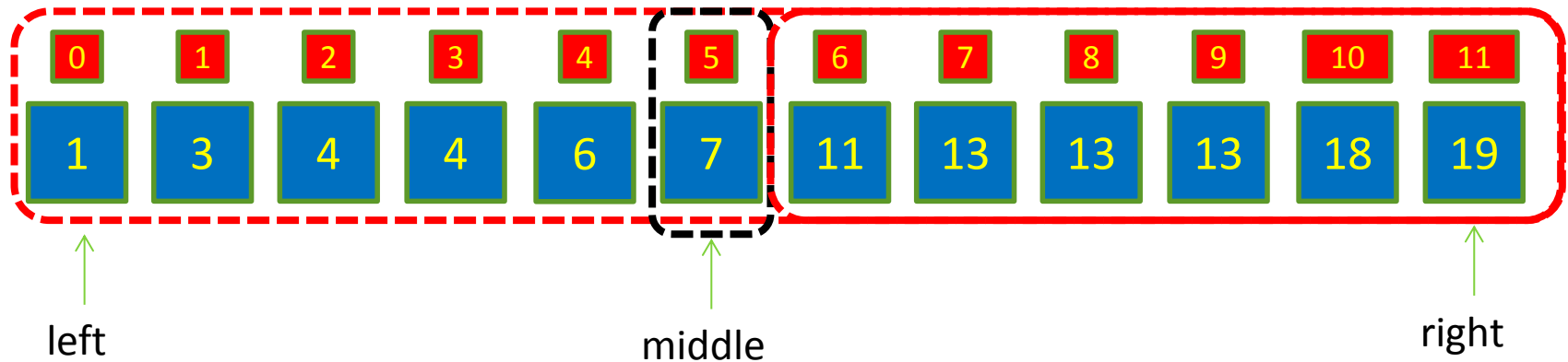
Example:

searchnum=13;



Example:

searchnum=13;

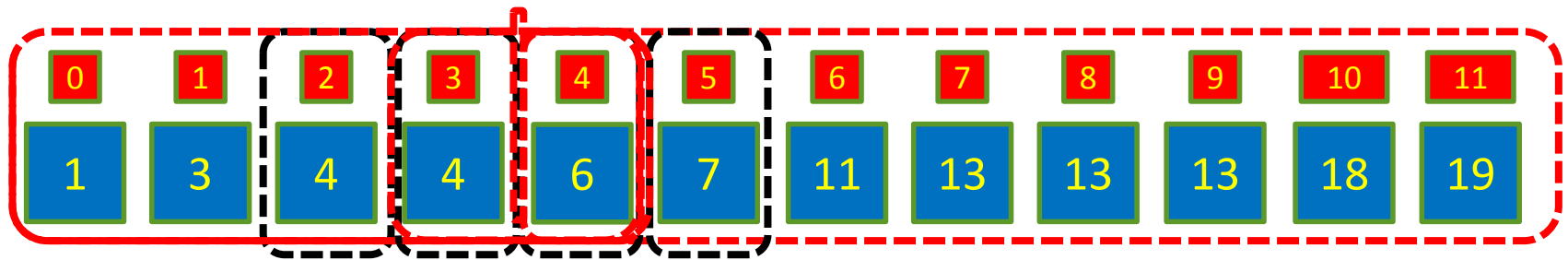


```
middle=(left+right)/2;
```

```
left=middle+1;
```

Example:

searchnum=5;



return -1;

```
int binsearch(int list[], int searchnum, int
left, int right) {
    int middle;
    while(left<=right) {
        middle=(left+right)/2;
        switch(COMPARE(list[middle], searchnum)) {
            case -1: left=middle+1; break;
            case 0: return middle;
            case 1: right=middle-1;
        }
    }
    return -1;
}
```

list: 存sort好數字的array

searchnum: 要找的數字

left, right: 正在找的範圍左邊和右邊邊界

What is a data Type?

- A data type is a collection of objects and a set of operations that act on those objects.
- 每種data type 有所占的記憶體大小及可表示的資料數值範圍
- Data types in C
 - char, int, float, long, double (unsigned, signed, ...)
 - Array

```
int iarray[16];
```
 - Structure (User-defined)

```
struct {  
    int a;  
    int b;  
    char str[16];  
    int * iptr;  
} blah;
```

Operations of Data Types

- Operations
 - +, -, *, /, %, ==
 - =, +=, -=
 - ? :
 - sizeof, - (negative)
 - gilgulu(int a, int b)

Data Type

- **Representation** of the objects of the data type
- Example: char
- `char blah='A';` ('A': ASCII code is 65(dec), or 0x41 (hex))

1 byte of memory:

01000001

Q: The maximum number which can be represented with a char variable?
A: 255.

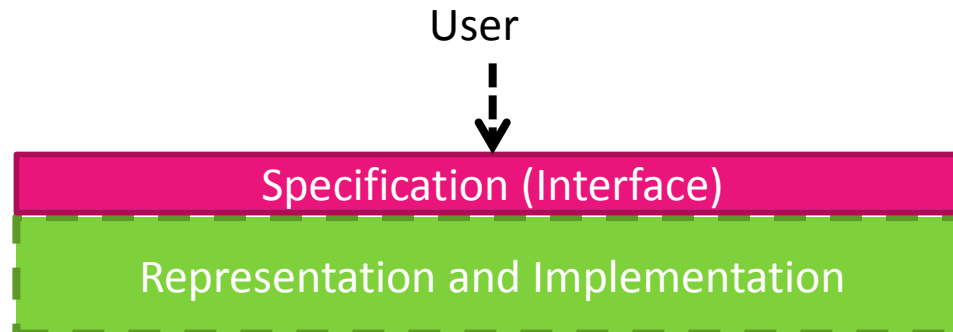
- How about char, int, long, float?

Data Type

- **Q:** 我們需要知道data type的representation嗎?
- **A:** 不一定.
 - 知道representation可能可以設計出更有效率的algorithm.
- 但是當data type的representation被修改以後，程式可能必須重新確認、修正、或**完全重寫**。 凹
- 移植到不同的平台上 (x86, ARM, embedded system, ...)
- 改變program或library的specification(規格) (ex. 16-bit int → 32-bit long)

Abstract Data Type

- “Abstract Data Type” (ADT):
- Separate the specifications from the representation and the implementation



Abstract Data Type

- Specifications:
 - Operations:
 - Name of the function and the description of what the function does
 - The type of the argument(s)
 - The type of the result(s) (return value)
 - Data (usually hidden)
- Function categories:
 - Creator/constructor
 - Transformers
 - Observer/reporter

Example

ADT NaturalNumber is

objects:

an ordered subrange of the integers starting at zero and ending at the maximum integer (INT-MAX) on the computer

functions:

for all $x, y \in \text{NaturalNumber}$; $\text{TRUE}, \text{FALSE} \in \text{Boolean}$ and where $+$, $-$, $<$, and $=$ are the usual integer operations

```

NaturalNumber Zero()           ::= 0
Boolean IsZero(x)              ::= if (x = 0) return TRUE
                                else return FALSE
Boolean Equal(x, y)            ::= if (x == y) return TRUE
                                else return FALSE
NaturalNumber Successor(x)     ::= if (x == INT-MAX) return x
                                else return x + 1
NaturalNumber Add(x, y)        ::= if ((x + y) <= INT-MAX)
                                return x + y
                                else return INT-MAX
NaturalNumber Subtract(x, y)   ::= if (x < y) return 0
                                else return x - y

```

end NaturalNumber

怎麼評估一個程式寫得好不好?

1. Does the program meet the original specifications of the task?
2. Does it work correctly?
3. Does the program **contain documentation** that shows how to use it and how it works?
4. Does the program effectively use **functions** to create **logical units**?
5. Is the program's code **readable**?



怎麼評估一個程式寫得好不好?

6. Does the program efficiently use primary and secondary storage?

Primary storage: memory?

Secondary storage: Hard drive, flash disk, etc.

7. Is the program's running time acceptable for the task?

Example: Network intrusion detection system



(1) 99.8% detection rate, 50 minutes to finish analysis of a minute of traffic



(2) 85% detection rate, 20 seconds to finish analysis of a minute of traffic

怎麼評估一個程式寫得好不好?

6. 程式是否有效地使用主要及次要的儲存?

Space complexity

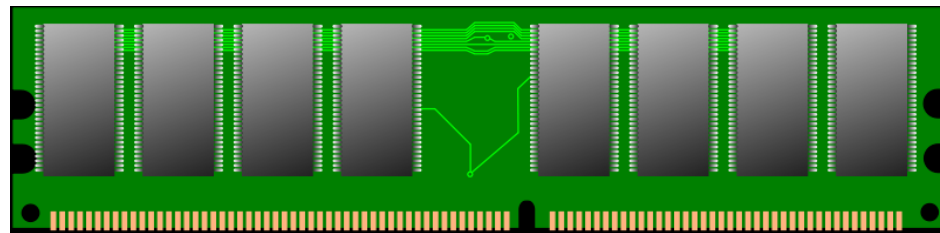
7. 程式的執行時間是否適合所需解決的工作內容?

Time complexity

空間及時間複雜度

- 程式的空間複雜度:
 - 程式執行完畢所需使用的所有空間(記憶體)
- 程式的時間複雜度:
 - 程式執行完畢所需使用的(執行)時間
- Goal: 找出執行時間/使用空間”如何”隨著input size變長(成長的有多快)
- 什麼是input size?
- 問題給的input的”元素數量”, 如:
 - Array大小
 - 多項式最高項的次方
 - 矩陣的長寬
 - 二進位數的位元數目

空間複雜度



- 程式所需空間:
 1. 固定的空間
 - 和input/output的大小及內容無關
 2. 變動的空間
 - 和待解問題P的某個input instance I(某一個input)有關
 - 跟recursive function會使用到的額外空間有關
- $S(P) = c + S_P(I)$

時間複雜度

- 一個程式 P 所需使用的時間:
 - Compile所需時間
 - 執行時間 (execution time or run time)
- Compile時間: 固定的. (例外?)
 - C (and other compiled programming languages)
→ One Compilation → Multiple Executions
- Run time: T_P
 - 和input instance的特性有關!



Reading materials for this lecture

- Cormen Chapter 1
- Horowitz Chapter 1.3-1.4