

2.19 Assume the following register contents:

`$t0 = 0xAAAAAAAA, $t1 = 0x12345678`

First, we transform `$t0` and `$t1` from Hex to Binary

`$t0 = 10101010101010101010101010101010`

`$t1 = 00010010001101000101011001111000`

2.19.1 [5] <\$2.6> For the register values shown above, what is the value of `$t2` for the following sequence of instructions?

```
sll $t2, $t0, 44
or  $t2, $t2, $t1
```

We translate instruction to high level language

`reg $t2 = reg $t0 << 44 bits`

`reg $t2 = reg $t2 | reg $t1`

The `sll` instruction is wrong because the largest shift amount is 31 in R-type instruction. We use remainder “13” to solve this question.

`$t2 = 0101 0101 0101 0101 0100 0000 0000 0000`

`$t1 = 0001 0010 0011 0100 0101 0110 0111 1000`

`$t2 = 0101 0111 0111 0101 0101 0110 0111 1000`

Some of students still use 44 to solve this question, the answer is

`$t2 = 0000 0000 0000 0000 0000 0000 0000 0000`

`$t1 = 0001 0010 0011 0100 0101 0110 0111 1000`

`$t2 = 0001 0010 0011 0100 0101 0110 0111 1000`

2.19.2 [5] <\$2.6> For the register values shown above, what is the value of `$t2` for the following sequence of instructions?

```
sll  $t2, $t0, 4
andi $t2, $t2, -1
```

We translate instruction to high level language

`reg $t2 = reg $t0 << 4 bits`

`reg $t2 = reg $t2 & (-1)`

`$t2 = 1010 1010 1010 1010 1010 1010 1010 0000`

`$t2 = 1010 1010 1010 1010 1010 1010 1010 0000`

`(-1) = 1111 1111 1111 1111 1111 1111 1111 1111`

`$t2 = 1010 1010 1010 1010 1010 1010 1010 0000 = 0xAAAAAA0`

2.23 [5] <\$2.7> Assume \$t0 holds the value 0x00101000. What is the value of \$t2 after the following instructions?

```

        slt  $t2, $0,  $t0
        bne  $t2, $0,  ELSE
        j    DONE
ELSE:    addi $t2, $t2, 2
DONE:

```

We translate instruction to high level language

\$t0 = 1000000001000000000000 = 1052672

If \$0 < reg \$t0 (0<1052672) then reg \$t2 = 1

If reg \$t0 != \$0 then go to ELSE (reg \$t2 = reg \$t2 + 2)

else go to DONE

Step 1 : \$t2 = 1

Step 2 : go to ELSE

Step 3 : \$t2 = \$t2 + 2 = 3

2.27 [5] <\$2.7> Translate the following C code to MIPS assembly code. Use a minimum number of instructions. Assume that the values of a, b, i, and j are in registers \$s0, \$s1, \$t0, and \$t1, respectively. Also, assume that register \$s2 holds the base address of the array D.

```

    for(i=0; i<a; i++)
        for(j=0; j<b; j++)
            D[4*j] = i + j;

    add $t0, $t0, $zero
    add $t1, $t1, $zero
LOOP1: slt $t2, $t0, $s0
       beq $t2, $zero, EXIT
       add $t1, $zero, $zero
       addi $t0, $t0, 1
LOOP2: slt $t3, $t1, $s1
       beq $t3, $zero, LOOP1
       add $t4, $t0, $t1
       muli $t5, $t1, 4
       sll $t5, $t5, 2
       add $t5, $t5, $s2
       sw $t4, $t5($s2)
       addi $t0, $t1, 1
       j LOOP2

```

2.34 Translate function `f` into MIPS assembly language. If you need to use registers `$t0` through `$t7`, use the lower-numbered registers first. Assume the function declaration for `func` is “`int f(int a, int b);`”. The code for function `f` is as follows:

```
int f(int a, int b, int c, int d){  
    return func(func(a,b),c+d);  
}
```

```
f:  addi $sp, $sp, 8  
    sw $ra, 4($sp)  
    sw $s0, 0($sp)  
    move $s0, $a2  
    jal func  
    move $a0, $v0  
    move $a1, $s0  
    jal func  
    lw $ra, 4($sp)  
    lw $s0, 0($sp)  
    addi $sp, $sp, 8  
    jr $ra
```