Kristina Poon August 13,2017

- 1 CDA, CDB, CDE
- @ PRIME: ABODE YES 3NF
- 3 No; A, BC, and ED are not superkeys

19.5) 
$$\bigcirc$$
 R<sub>1</sub>(ACBDE), A > B, C > D  
AC | BD | AC+ | a) INF  
ACBD | b) AB, CD, ACE  
R<sub>1</sub> AB | R<sub>2</sub> ACDE

$$\bigcirc$$
 R<sub>2</sub>(A,B,F) AC  $\rightarrow$  E, B  $\rightarrow$  F  
 $\bigcirc$  ACB  $\uparrow$  EF  $\bigcirc$  ACB  $\uparrow$  ACB  $\uparrow$  ACB  $\uparrow$  B) AB, BF

③ 
$$R_s(A,D)$$
  $D \rightarrow G$ ,  $G \rightarrow H$ 

D G HA  $D^+$  a) BCNF

DGH

Constitution of the second second	(S) AB>C, AB>D, C>A, D>B	
and a second second second	TABCD AT ABT / BC+ A ABCD BCAD	1 CD+ / AD+ /
	A ABCD BCAD	CDAB ADBC
	PRIME: ABCD	a) AB, AD, BC, CD
		b) 3NF
in the state of th	AC CBD	c) none
to the second		
	DB DC	
910	R(ABCD)	
	OB→C, D→A: BC and AD	
The second secon	BC EDA BOT TAC BD+	a) BD
	BDCA BDCA	b) unsatisfactory blc lossy
To a company when the same statutes	@ AB→C, C→A, C→D: ACD and BC	14 67 11 11
- institution region but the regions in the reco	BLACID B+ BA+ / BC+	a) AB, BC
	B BACD BCAD	b) unsatisfactory ble not
		dependency presoning
	3 A→BC, C→AD: ABC and AD	
a and a second second second	TACTED AT CADE	a) A, C
		b) not dependency preserving
	4 A→B, B→C,C→D: AB and ACD	
	ABCO A+ VAB+V	a) A
	ABCD ABCD	b) loss-less join decomposition
		but not dependency present
	(3) A>B, B>C, C>D: AB, AD, CD	
andrey be an arranged		a) A
discretization in the second		b) lossless but not dependency
		preserving
and the second second	The filtration of the state of	

16.2) R(X), W(X), R(Y), W(Y)

- O If T2 performs W(Y) before TI performs R(Y) and aborts, R(Y)'s value would be invalid and the abort will cascade.
- (a) There would be an exclusive lock on Typefore writing and the lock will hold until T2 commits or aborts. To cannot interfere by reading Y until T2 is done.

ensures safe interleaving
 easy to implement

16.3) (1) Ta: R(x), T2: R(Y), T2: W(X), T1: K(X)

□ TZ:R(X), TZ:R(Y), TI:R(X), TI:R(Y), TI:W(X)

Ta unrepeatable read on X

(3) T2:R(X), T2:R(Y), T1:R(X), T1:R(Y), T1:W(X), T2:W(X)

Ta overwrites uncommitted X 4)

(9) TI waits until T2 finishes
2) TR is already holding the lock; cannot be obtained by TI
3) TR is already holding the lock; cannot be obtained by TI