## 树上启发式合并

## DSU on Tree

Idea: 如果暴力做,我们会整体做一次 dfs 以遍历每个点,然后对遍历的每个点 dfs 它的子树,计算答案,把答案记录抹去,算下一个点。这样做是  $O(n^2)$  的。为优化,我们遍历到一个点时,先遍历它的轻儿子,计算轻儿子的答案并抹去记录,最后计算它的重儿子的答案且不抹去记录,如此在计算它本身的答案时,只需要暴力算出轻儿子的答案,合并上原本就存在的重儿子答案即可。

代码中的 dsu() 函数对应暴力算法的整体 dfs, getData() 函数对应暴力算法的"对遍历的每个点 dfs 计算子树答案"。

Complexity:  $O(n \lg n)$ 

Code:

```
int fa[N], sz[N], son[N];
 2
     void dfs(int x, int f){
         fa[x] = f, sz[x] = 1, son[x] = 0;
3
         for(int i = head[x]; i; i = edge[i].nxt){
4
             if(edge[i].to == f) continue;
             dfs(edge[i].to, x);
6
7
             sz[x] += sz[edge[i].to];
             if(!son[x] \parallel sz[edge[i].to] > sz[son[x]])
8
9
                 son[x] = edge[i].to;
10
         }
     }
11
12
     LL ans[N], mx, sum, cnt[N]; // GLOBAL variants to store the answer
13
14
     int mark; // mark the heavy son which needs to be ignored
15
     void getData(int x, int val){ // get data with brute-force
16
17
         cnt[c[x]] += val;
         if(mx < cnt[c[x]]) mx = cnt[c[x]], sum = c[x];
18
19
         else if(mx == cnt[c[x]])
                                      sum += c[x];
20
21
         for(int i = head[x]; i; i = edge[i].nxt){
             if(edge[i].to == fa[x]) continue;
22
             if(edge[i].to == mark) continue; // ignore the marked subtree
2.3
24
             getData(edge[i].to, val);
25
         }
26
27
     void dsu(int x, bool opt){ // opt == true: answer needs to be erased
2.8
         for(int i = head[x]; i; i = edge[i].nxt){
             if(edge[i].to == fa[x] || edge[i].to == son[x]) continue;
29
             dsu(edge[i].to, true); // solve for light sons first
30
31
32
         if(son[x]) dsu(son[x], false), mark = <math>son[x]; // solve for heavy son
33
         // now the global variants have already stored heavy son's answer
34
         getData(x, 1);
35
         mark = 0;
36
         // now the global variants store the answer for vertex \boldsymbol{\boldsymbol{x}}
37
38
         ans[x] = sum;
39
         if(opt){ // erase the answer
40
41
             getData(x, -1);
42
             mx = 0, sum = 0;
43
44
     }
45
46
     int main(){
         scanf("%d", &n);
47
48
         for(int i = 1; i <= n; i++)
             scanf("%d", &c[i]);
49
50
         for(int i = 1; i < n; i++){
             int u, v; scanf("%d%d", &u, &v);
51
52
             addEdge(u, v), addEdge(v, u);
53
         dfs(1, 0);
54
55
         dsu(1, true);
         for(int i = 1; i <= n; i++)
56
57
             printf("%lld ", ans[i]);
58
         return 0:
```