可持久化并查集

Persistent Disjoint Set Union

Idea: 所谓并查集,在实现上无非就是两个数组 fa[] 和 sz[](按秩合并),只需要将这两个数组可持久化即可。把 y 合并到 x 上时,我们需要在某个历史版本上更改 fa[findfa(y)]为 findfa(x),并将 sz[findfa(x)]加上 sz[findfa(y)]。而所谓可持久化数组,仍旧是可持久化线段树来实现。

Complexity: $O(n \lg^2 n)$

Code:

```
#include<algorithm>
    #include<iostream>
    #include<cstdio>
 3
 4
    using namespace std;
 5
 6
 7
    const int N = 200005;
    int n, m;
 8
 9
10
    struct segTree{
        int lson, rson;
11
        int fa, sz;
12
    tr[N * 30];
13
    int cnt, root[N];
14
    void build(int id, int l, int r){
15
        if(l == r){
16
            tr[id].fa = l, tr[id].sz = 1;
17
18
            return;
19
        }
        tr[id].lson = ++cnt, tr[id].rson = ++cnt;
20
        int mid = (l + r) \gg 1;
21
22
        build(tr[id].lson, l, mid);
        build(tr[id].rson, mid+1, r);
23
24
    int queryID(int id, int l, int r, int pos){ // query the ID of
25
    pos'th leaf in tr[id]
        if(l == r) return id;
26
        int mid = (l + r) \gg 1;
27
```

```
28
        if(pos <= mid) return queryID(tr[id].lson, l, mid, pos);</pre>
29
                return queryID(tr[id].rson, mid+1, r, pos);
30
    }
    void modify(int cur, int pre, int l, int r, int pos, int fa){
31
32
        tr[cur] = tr[pre];
        if(l == r){ tr[cur].fa = fa; return; }
33
        int mid = (l + r) \gg 1;
34
        if(pos <= mid){</pre>
35
            tr[cur].lson = ++cnt;
36
            modify(tr[cur].lson, tr[pre].lson, l, mid, pos, fa);
37
38
        }
39
        else{
            tr[cur].rson = ++cnt;
40
            modify(tr[cur].rson, tr[pre].rson, mid+1, r, pos, fa);
41
        }
42
43
    void add(int id, int l, int r, int pos, int val){
44
        if(l == r){ tr[id].sz += val; return; }
45
        int mid = (l + r) \gg 1;
46
        if(pos <= mid) add(tr[id].lson, l, mid, pos, val);</pre>
47
                add(tr[id].rson, mid+1, r, pos, val);
        else
48
49
    }
50
    int findfa(int cur, int x){
51
        int xid = queryID(cur, 1, n, x); // x's id in tr[cur]
52
        return tr[xid].fa == x ? x : findfa(cur, tr[xid].fa);
53
54
    void unionn(int cur, int pre, int x, int y){
55
56
        tr[cur] = tr[pre]; // ATT!
        x = findfa(pre, x), y = findfa(pre, y);
57
        if(x == y) return;
58
        int xid = queryID(pre, 1, n, x), yid = queryID(pre, 1, n, y);
59
        if(tr[xid].sz < tr[yid].sz) swap(x, y), swap(xid, yid);</pre>
60
        modify(cur, pre, 1, n, y, x);
61
        add(cur, 1, n, x, tr[yid].sz);
62
    }
63
64
65
    int main(){
        scanf("%d%d", &n, &m);
66
        build(root[0] = 0, 1, n);
67
        for(int i = 1; i <= m; i++){
68
69
            int opt, a, b;
            scanf("%d", &opt);
70
            if(opt == 1){
71
72
                scanf("%d%d", &a, &b);
                 root[i] = ++cnt;
73
                unionn(root[i], root[i-1], a, b);
74
```

```
75
           else if(opt == 2){
76
77
              scanf("%d", &a);
               root[i] = root[a];
78
           }
79
           else{
80
               scanf("%d%d", &a, &b);
81
               root[i] = root[i-1];
82
               puts(findfa(root[i], a) == findfa(root[i], b) ? "1" :
83
   "0");
84
          }
85
      }
return 0;
87 }
```