## 高斯消元

## **Gauss-Jordan Elimination**

Complexity:  $O(n^3)$  Code (浮点):

```
1
    namespace LA{
2
         int n;
         double a[N][N], b[N];
3
4
         void init(int nn){
5
             n = nn;
 6
             for(int i = 1; i <= n; i++){
                 b[i] = 0;
8
                 for(int j = 0; j <= n; j++) a[i][j] = 0;
9
10
         }
11
         bool Gauss(){
12
         // false: no solution or multiple solutions; true: a[][n+1] is the only solution
13
         /* a[1,1]x1 + a[1,2]x2 + ... + a[1,n]xn = b[1]
14
             a[2,1]x1 + a[2,2]x2 + ... + a[2,n]xn = b[2]
15
16
17
             a[n,1]x1 + a[n,2]x2 + ... + a[n,n]xn = b[3] */
18
             for(int i = 1; i <= n; i++) a[i][n+1] = b[i];
19
             for(int j = 1; j <= n; j++){
                 int r = j;
21
                 for(int i = j + 1; i <= n; i++)
                     if(a[i][j] > a[j][j])
23
24
                        r = i;
                 if(r != j) swap(a[r], a[j]);
25
                 if(a[j][j] == 0) return false;
26
                 for(int i = 1; i <= n; i++){
27
                     if(i == j) continue;
28
29
                     double div = a[i][j] / a[j][j];
                     for(int k = j; k \le n + 1; k++)
30
                         a[i][k] = div * a[j][k];
32
33
             for(int i = 1; i <= n; i++) a[i][n+1] /= a[i][i];
34
35
             return true;
36
37
         double det(){ // get determinant
38
             double res = 1;
39
             int flag = 1;
             for(int j = 1; j \le n; j++){
40
                 int r = j;
41
42
                 for(int i = j + 1; i \le n; i + +)
43
                     if(a[i][j] > a[j][j])
44
                         r = i;
45
                 if(r != j) swap(a[r], a[j]), flag = -flag;
46
                 if(a[j][j] == 0) return 0;
47
                 for(int i = 1; i <= n; i++){
48
                     if(i == j) continue;
49
                     double div = a[i][j] / a[j][j];
50
                     for(int k = j; k \le n; k++)
51
                         a[i][k] = div * a[j][k];
52
```

## Code (取模):

```
1
     namespace \ LA\{
2
         int n;
 3
         LL a[N][N], b[N];
 5
         void init(int nn){
 6
             n = nn;
             for(int i = 1; i <= n; i++){
 8
                 b[i] = 0;
 9
                 for(int j = 0; j \le n; j++) a[i][j] = 0;
10
11
12
         bool Gauss(){
13
         // false: no solution or multiple solutions; true: a[][n+1] is the only solution
14
         /* a[1,1]x1 + a[1,2]x2 + ... + a[1,n]xn = b[1]
15
             a[2,1]x1 + a[2,2]x2 + ... + a[2,n]xn = b[2]
16
17
             a[n,1]x1 + a[n,2]x2 + ... + a[n,n]xn = b[3] */
18
             for(int i = 1; i <= n; i++) a[i][n+1] = b[i];
19
20
             for(int j = 1; j \le n; j++){
                 int r = j;
21
                 for(int i = j + 1; i <= n; i++)
                     if(a[i][j] > a[j][j])
23
24
                         r = i;
                 if(r != j) swap(a[r], a[j]);
                 if(a[j][j] == 0) return false;
26
                 for(int i = 1; i <= n; i++){
27
                     if(i == j) continue;
28
29
                     LL div = a[i][j] * fpow(a[j][j], MOD-2) % MOD;
                      for(int k = j; k \le n + 1; k++){
                          a[i][k] \stackrel{-=}{-} div * a[j][k];
                          ((a[i][k] %= MOD) += MOD) %= MOD;
                     }
                 }
34
             for(int i = 1; i <= n; i++) (a[i][n+1] *= fpow(a[i][i], MOD-2)) %= MOD;
             return true;
38
         LL det(){ // get determinant
39
             LL res = 1;
40
41
             int flag = 1;
             for(int j = 1; j \le n; j++){
42
                 int r = j;
43
                 for(int i = j + 1; i <= n; i++)
44
                     if(a[i][j] > a[j][j])
45
                         r = i;
46
                 if(r != j) swap(a[r], a[j]), flag = -flag;
47
                 if(a[j][j] == 0) return 0;
48
                 for(int i = 1; i <= n; i++){
49
                     if(i == j) continue;
50
                     LL div = a[i][j] * fpow(a[j][j], MOD-2) % MOD;
51
                     for(int k = j; k \le n; k++){
52
                          a[i][k] = div * a[j][k] % MOD;
53
                          ((a[i][k] %= MOD) += MOD) %= MOD;
54
                     }
55
                 }
56
57
```

```
for(int i = 1; i <= n; i++) (res *= a[i][i]) %= MOD;
return flag > 0 ? res : MOD - res;
}
60  }
61 }
```