## 可持久化并查集

## **Persistent Disjoint Set Union**

Idea: 所谓并查集,在实现上无非就是两个数组 fa[] 和 sz[] (按秩合并),只需要将这两个数组可持久化即可。把 y 合并到 x 上时,我们需要在某个历史版本上更改 fa[findfa(y)] 为 findfa(x),并将 sz[findfa(x)] 加上 sz[findfa(y)]。而所谓可持久化数组,仍旧是可持久化线段树来实现。

**Complexity**:  $O(n \lg^2 n)$ 

Code:

```
1
    #include<algorithm>
    #include<iostream>
    #include<cstdio>
5
    using namespace std;
 6
    const int N = 200005;
    int n, m;
8
    struct segTree{
10
        int lson, rson;
11
         int fa, sz;
12
    }tr[N * 30];
13
    int cnt, root[N];
14
    void build(int id, int l, int r){
15
         if(l == r){
16
17
             tr[id].fa = l, tr[id].sz = 1;
18
             return;
19
         tr[id].lson = ++cnt, tr[id].rson = ++cnt;
21
         int mid = (l + r) \gg 1;
22
         build(tr[id].lson, l, mid);
23
         build(tr[id].rson, mid+1, r);
24
    int queryID(int id, int l, int r, int pos){ // query the ID of pos'th leaf in tr[id]
25
         if(l == r) return id;
26
         int mid = (l + r) \gg 1;
27
         if(pos <= mid) return queryID(tr[id].lson, l, mid, pos);</pre>
28
29
               return queryID(tr[id].rson, mid+1, r, pos);
30
31
     void modify(int cur, int pre, int l, int r, int pos, int fa){
32
         tr[cur] = tr[pre];
33
         if(l == r){ tr[cur].fa = fa; return; }
34
         int mid = (l + r) \gg 1;
35
         if(pos <= mid){</pre>
36
             tr[cur].lson = ++cnt;
37
             modify(tr[cur].lson, tr[pre].lson, l, mid, pos, fa);
38
39
         else{
40
             tr[cur].rson = ++cnt;
41
             modify(tr[cur].rson, tr[pre].rson, mid+1, r, pos, fa);
42
43
44
     void add(int id, int l, int r, int pos, int val){
45
         if(l == r){ tr[id].sz += val; return; }
46
         int mid = (l + r) \gg 1;
47
         if(pos <= mid) add(tr[id].lson, l, mid, pos, val);</pre>
48
         else add(tr[id].rson, mid+1, r, pos, val);
```

```
49
50
51
     int findfa(int cur, int x){
52
          int xid = queryID(cur, 1, n, x); // x's id in tr[cur]
          return tr[xid].fa == x ? x : findfa(cur, tr[xid].fa);
53
54
55
     void unionn(int cur, int pre, int x, int y){
          tr[cur] = tr[pre]; // ATT!
56
57
         x = findfa(pre, x), y = findfa(pre, y);
58
          if(x == y) return;
59
          int xid = queryID(pre, 1, n, x), yid = queryID(pre, 1, n, y);
          if(tr[xid].sz < tr[yid].sz) \ swap(x, \ y), \ swap(xid, \ yid);\\
60
61
          modify(cur, pre, 1, n, y, x);
         add(cur, 1, n, x, tr[yid].sz);
62
63
     }
64
65
     int main(){
66
          scanf("%d%d", &n, &m);
67
          build(root[0] = 0, 1, n);
68
          for(int i = 1; i <= m; i++){
69
              int opt, a, b;
70
              scanf("%d", &opt);
71
              if(opt == 1){
72
                  scanf("%d%d", &a, &b);
73
                  root[i] = ++cnt;
74
                  unionn(root[i], root[i-1], a, b);
75
76
              else if(opt == 2){
77
                  scanf("%d", &a);
78
                  root[i] = root[a];
79
80
              else{
                  scanf("%d%d", &a, &b);
81
82
                  root[i] = root[i-1];
                  puts(findfa(root[i], a) == findfa(root[i], b) ? "1" : "0");
83
84
85
86
          return 0;
87
```