可持久化并查集

Persistent Disjoint Set Union

Idea: 所谓并查集,在实现上无非就是两个数组 fa[] 和 sz[](按秩合并),只需要将这两个数组可持久化即可。把 y 合并到 x 上时,我们需要在某个历史版本上更改 fa[findfa(y)] 为 findfa(x),并将 sz[findfa(x)] 加上 sz[findfa(y)]。而所谓可持久化数组,仍旧是可持久化线段树来实现。

Complexity: $O(n \lg^2 n)$

Code:

```
#include<algorithm>
     #include<iostream>
     #include<cstdio>
3
     using namespace std;
6
     const int N = 200005;
8
     int n, m;
9
10
     struct segTree{
11
         int lson, rson;
12
         int fa, sz;
     }tr[N * 30];
13
14
     int cnt, root[N];
     void build(int id, int l, int r){
15
16
         if(l == r){
17
             tr[id].fa = l, tr[id].sz = 1;
18
             return;
19
2.0
         tr[id].lson = ++cnt, tr[id].rson = ++cnt;
21
         int mid = (l + r) \gg 1;
         build(tr[id].lson, l, mid);
22
23
         build(tr[id].rson, mid+1, r);
24
     int queryID(int id, int l, int r, int pos){ // query the ID of pos'th leaf in tr[id]
25
26
         if(l == r)
                      return id;
         int mid = (l + r) >> 1;
27
28
         if(pos <= mid) return queryID(tr[id].lson, l, mid, pos);</pre>
                return queryID(tr[id].rson, mid+1, r, pos);
29
3.0
31
     void modify(int cur, int pre, int l, int r, int pos, int fa){
         tr[cur] = tr[pre];
32
         if(l == r){ tr[cur].fa = fa; return; }
33
34
         int mid = (l + r) >> 1;
35
         if(pos <= mid){</pre>
36
             tr[cur].lson = ++cnt:
37
             modify(tr[cur].lson, tr[pre].lson, l, mid, pos, fa);
38
39
         else{
40
             tr[cur].rson = ++cnt;
41
             modify(tr[cur].rson, tr[pre].rson, mid+1, r, pos, fa);
42
43
     void add(int id, int l, int r, int pos, int val){
44
45
         if(l == r){ tr[id].sz += val; return; }
         int mid = (l + r) \gg 1;
46
47
         if(pos <= mid) add(tr[id].lson, l, mid, pos, val);</pre>
               add(tr[id].rson, mid+1, r, pos, val);
48
49
     }
50
     int findfa(int cur, int x){
51
52
         int xid = queryID(cur, 1, n, x); // x's id in tr[cur]
         return tr[xid].fa == x ? x : findfa(cur, tr[xid].fa);
53
54
55
     void unionn(int cur, int pre, int x, int y){
         tr[cur] = tr[pre]; // ATT!
56
57
         x = findfa(pre, x), y = findfa(pre, y);
         if(x == y) return;
58
         int xid = queryID(pre, 1, n, x), yid = queryID(pre, 1, n, y);
59
         if(tr[xid].sz < tr[yid].sz) \ swap(x, \ y), \ swap(xid, \ yid);\\
60
```

```
modify(cur, pre, 1, n, y, x);
61
62
         add(cur, 1, n, x, tr[yid].sz);
63
    }
64
65
     int main(){
         scanf("%d%d", &n, &m);
66
         build(root[0] = 0, 1, n);
67
         for(int i = 1; i <= m; i++){
68
            int opt, a, b;
69
70
             scanf("%d", &opt);
71
             if(opt == 1){
                 scanf("%d%d", &a, &b);
72
                 root[i] = ++cnt;
73
74
                 unionn(root[i], root[i-1], a, b);
75
76
             else if(opt == 2){
                 scanf("%d", &a);
77
78
                 root[i] = root[a];
79
             }
80
             else{
81
                 scanf("%d%d", &a, &b);
                 root[i] = root[i-1];
82
83
                 puts(findfa(root[i], a) == findfa(root[i], b) ? "1" : "0");
            }
84
85
         return 0;
86
87
```