

Open Book and Notes

1. Prove that the language $L(G)$ is not regular where G is the following context-free grammar: $G = (\{S, A, B, C\}, \{a, b\}, \{S \rightarrow A|Bb, A \rightarrow C, B \rightarrow bbS, C \rightarrow a\}, S)$.

Note: You must first determine $L(G)$.

2. Eliminate all ϵ -productions in the following cfg G :

$$G = (\{S, A, B\}, \{a, b, c\}, \{S \rightarrow aA|BBBb, A \rightarrow cS|\epsilon, B \rightarrow Sb|\epsilon\}, S).$$

3. Construct a **reduced** dfa for the following extended regular expression over the alphabet $\{0, 1, 2\}$ (not $\{0, 1\}!$):

$$[(010)^* \cap \overline{0(10^*)^*}]^*$$

Note: You must first determine nfas for $(010)^*$ and $0(10^*)^*$, then handle the intersection and complementation, and then deal with the star. Finally reduce the resulting dfa. Consider de Morgan's laws!

4. Construct a Chomsky normal form grammar for $L(G)$ for the following cfg G :

$$G = (\{S, B\}, \{a, b, c, d\}, \{S \rightarrow aSBBBB|B|cd, B \rightarrow cSda|S|cba\}, S).$$

Note: You must first remove all unit productions.

5. Construct a Greibach normal form grammar for $L(G)$ for the following CNF G :

$$G = (\{S, A, B\}, \{a, b\}, \{S \rightarrow AA, A \rightarrow BSS|a, B \rightarrow SSA|b\}, S).$$

Note: First derive all the productions for S , A , and B . You may only indicate how the final result looks for whatever primed variables you obtain.

Points: 1: 15 2: 10 3: 30 4: 15 5: 30

Production:

$$S \rightarrow A|Bb$$

$$A \rightarrow \textcircled{B}$$

$$B \rightarrow bbS$$

$$C \rightarrow a$$

① determine $L(G)$

$$S \rightarrow A|Bb$$

$$S \rightarrow Bb \rightarrow bbSb \rightarrow bbBbb$$

$$\rightarrow A \rightarrow c$$

bbS

$$\downarrow$$

$$bbAb$$

$$bbab$$



$$bbbbSbb$$

$$\rightarrow \underline{bbbb}a\underline{bb} \rightarrow \underline{b^4}a\underline{b^4}$$

$$L(G) = \{ b^n a b^n \mid n \geq 0 \}$$

Assuming $L(G)$ is regular and belongs to dfa, and D will have n states total.

$$\text{If } x = b^n a b^n = K \cdot \Delta \quad K = b^n, \Delta = ab^n$$

$\because |K| = n \therefore$ pumping lemma applies. $K = k_1 k_2 k_3$

such that $|k_2| \geq 1 \therefore \tau(q_0, K) = \tau(q_0, k_1 (k_2)^i k_3)$

when $i = 0$

$|K_1 \cdot K_2| = n - |K_2| < n$, so $\tau(q_0, k_1 k_3)$ does Not belong to $L(G)$

$\therefore \tau(q_0, K) \in L(G)$, but $\tau(q_0, k_1 k_3) \notin L(G)$

This is the contradiction which makes the

$L(G)$ not regular

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$$2) \begin{aligned} S &\rightarrow aA | BBBb \\ A &\rightarrow cS | \epsilon \\ B &\rightarrow Sb | \epsilon \end{aligned}$$

$$A \rightarrow \epsilon$$

$$S \rightarrow aA | BBBb | a$$

$$A \rightarrow cS$$

$$B \rightarrow Sb$$

$$B \rightarrow \epsilon$$

$$S \rightarrow aA | BBBb | a | BBb | Bb | b$$

$$A \rightarrow cS$$

$$B \rightarrow Sb$$

$$S \rightarrow \epsilon \text{ (1)}$$

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$$S \rightarrow aA | BBBb | a | BBb | Bb | b$$

$$A \rightarrow cS | cX$$

$$B \rightarrow Sb | Bx \text{ (2)}$$

$$[(010)^* \cap 0(10^*)^*]^* \Rightarrow [(010)^* \cup 0(10^*)^*]^*$$

1 2 3 4 5 6 1 2 3 4 5 6

	0	1	2
0	1	-	-0
1	-	-	-1

	0	1	2
0	-2	-	0
2	-	-	-1

	0	1	2
0	3	-	-0
3	-	-	-1

	0	1	2
0	1	-	-0
1	-	2	-0
2	-	-	-1

	0	1	2
0	1	-	-0
1	-	2	-0
2	3	-	-0
3	-	-	-1

	0	1	2
0	4	-	-0
4	-	-	-1

	0	1	2
0	-5	-	0
5	-	-	-1

	0	1	2
0	6	-	-0
6	-	-	-1

	0	1	2
0	-5	-	0
5	6	-	-1
6	6	-	-1

	0	1	2
0	-5	-	1
5	6	5	-1
6	6	5	-1

	0	1	2
0	1	-	-1
1	-	2	-0
2	3	-	-0
3	1	-	-1

dFA

	0	1	2
0	1	-	-1
1	-	2	-0
2	-	-	-0
3	3	-	-0
4	1	-	-1

	0	1	2
0	4	-	-0
4	-	5	-1
5	6	5	-1
6	6	5	-1

	0	1	2
A	B	C	C
B	C	D	C
C	C	C	C
D	E	C	C
E	B	C	C

renamed

	0	1	2
A	B	C	C
B	C	D	C
C	C	C	C
D	E	C	C
E	B	C	C

	0	1	2
A	B,4	C	C
B	C	D	C
C	C	C	C
D	E	C	C
E	B	C	C
4	-	5	-1
5	6	5	-1
6	6	5	-1

	0	1	2
A	B,4	C	C
B,4	C	D,5	C
C	C	C	C
D,5	E,6	C,5	C
E,6	B,6	C,5	C
C,5	C,6	C,5	C
B,6	C,6	D,5	C
C,6	C,6	C,5	C

	0	1	2
1	2	3	3
2	3	4	3
3	3	3	3
4	5	6	3
5	7	6	3
6	8	6	3
7	8	4	3
8	8	6	3

complement

	0	1	2
1	2	3	3
2	3	4	3
3	3	3	3
4	5	6	3
5	7	6	3
6	8	6	3
7	8	4	3
8	8	6	3



	0	1	2
1	2	3	3
2	3	4	3
3	3	3	3
4	5	6	3
5	7	6	3
6	8	6	3
7	8	4	3
8	8	6	3

Reduce

1	2	3	4	5	6	7	8
---	---	---	---	---	---	---	---

1	2	3	4	5	6	7	8
---	---	---	---	---	---	---	---

(A)

(B)

	0	1	2
A	B	B	B
B	B	B	B

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*)

$$S \rightarrow aSB\bar{B}B|B|cd$$

$$B \rightarrow cSda|S|cba$$

$$B \rightarrow S$$

$$S \rightarrow aSB\bar{B}B|cd|aS\bar{B}B\bar{B}|aSS\bar{S}B|aSS\bar{S}S|aSB\bar{S}S$$

-1 missing / aSBBS|aSSBS

$$B \rightarrow cSda|cba$$

S → B - S

CNF

$$S \rightarrow \underbrace{X_a S \bar{B} B B}_4 | \underbrace{X_c X_d}_3 | \underbrace{X_a S \bar{S} B B}_3 | \underbrace{X_a S \bar{S} \bar{S} B}_3 | \underbrace{X_a S \bar{S} \bar{S} S}_3 | \underbrace{X_a S \bar{S} B \bar{S}}_3 | \underbrace{X_a S \bar{S} B S}_3$$

$$B \rightarrow X_c S X_d X_a | X_c X_b X_a$$

$$X_a \rightarrow a \quad X_b \rightarrow b \quad X_c \rightarrow c \quad X_d \rightarrow d$$

$$S \rightarrow X_a S_1 | X_c X_d | X_a \bar{S}_4 | X_a S_7 | X_a S_{10} | X_a S_{13}$$

$$S_1 \rightarrow S S_2 \quad S_9 \rightarrow S B$$

$$S_2 \rightarrow B S_3 \quad S_{10} \rightarrow S S_{11}$$

$$S_3 \rightarrow B B \quad S_{11} \rightarrow S S_{12}$$

$$S_4 \rightarrow S S_5 \quad S_{12} \rightarrow S S$$

$$S_5 \rightarrow S S_6 \quad S_{13} \rightarrow S S_{14}$$

$$S_6 \rightarrow B B \quad S_{14} \rightarrow B S_{15}$$

$$S_7 \rightarrow S S_8 \quad S_{15} \rightarrow S S$$

$$S_8 \rightarrow S S_9 \quad S_{16} \rightarrow S S_{17}$$

$$S_{17} \rightarrow B S_{18}$$

$$X_a S_{16} | X_a S_{19}$$

$$S_{18} \rightarrow B S$$

$$S_{19} \rightarrow S S_{20}$$

$$S_{20} \rightarrow S S_{21}$$

$$S_{21} \rightarrow B S$$

$$B \rightarrow X_c S \mid X_d X_a \mid X_c X_b X_a$$

$$B \rightarrow X_c B_1 \mid X_c B_3$$

$$B_1 \rightarrow S B_2$$

$$B_2 \rightarrow X_d X_a$$

$$B_3 \rightarrow X_b X_a$$

$$X_a \rightarrow a \quad X_b \rightarrow b \quad X_c \rightarrow c \quad X_d \rightarrow d$$

g



$$S \rightarrow AA$$

$$A \rightarrow BSS|a$$

$$B \rightarrow SSA|b$$

no useless
no ϵ -production
no unit production

$$i=1 \quad S \rightarrow AA \quad \text{no immediate left recursion}$$

$$i=2 \quad j=1 \quad A \rightarrow BSS|a \quad \text{no immediate left R.}$$

$$i=3 \quad j=1 \quad B \rightarrow \underline{A}A SA|b$$

$$j=2 \quad B \rightarrow \underline{BSS}ASA|a \quad ASA|b$$

immediate left R. found

$$B \rightarrow \underline{BSSASA} \mid \underbrace{aASA}_{B_1} \mid \underbrace{b}_{B_2}$$

$$B \rightarrow aASA|b|a \cancel{ASA} B' | bB'$$

$$B' \rightarrow SSASA | \cancel{SSASA} B'$$

$$S \rightarrow AA$$

$$A \rightarrow \underline{BSS}|a$$

$$B \rightarrow \underline{aASA}|b|a \underline{ASAB'}|bB'$$

$$B' \rightarrow SSASA | SSASAB'$$

$$A \rightarrow \underline{aASASS}|bSS|a \underline{ASAB'SS}|bB'SS|a$$

$$S \rightarrow \underline{aASASSA}|bSSA|a \cancel{ASAB'SS}A|AB'SSA|aA$$

$$B \rightarrow aASA|b|a \cancel{ASAB'}|bB'$$

$B' \longrightarrow$ back

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$B' \rightarrow a A s A s s A s A s A \mid b s s A s A s A \mid a A s A B' s s A s A s A$
 $\mid b B' s s A s A s A \mid a A s A s A \mid a A s A s s A s A s A B'$
 $b s s A s A s A B' \mid a A s A B' s s A s A s A B'$
 $b B' s s A s A s A B' \mid a A s A s A s A B'$