

# COSC 3340/6309

## Examination 2

Friday, June 22, 2012, 4 – 6 pm

Open Book and Notes

93

15/15

1. Prove that the language  $L(G)$  is not regular where  $G$  is the following context-free grammar:  $G = (\{S, A, B, C\}, \{a, b\}, \{S \rightarrow A|bB, A \rightarrow C, B \rightarrow Sb, C \rightarrow a\}, S)$ .

Note: You must first determine  $L(G)$ .

2. Eliminate all  $\epsilon$ -productions in the following cfg  $G$ :

10  $G = (\{S, A, B\}, \{a, b, c\}, \{S \rightarrow aA|BBBb, A \rightarrow cS|\epsilon, B \rightarrow Sb|\epsilon\}, S)$ .

30/30

3. Construct a **reduced** dfa for the following extended regular expression over the alphabet  $\{0,1,2\}$  (not  $\{0,1\}^*$ ):

$$[(101)^* \cap \overline{1(01^*)^*}]^*$$

Note: You must first determine nfas for  $(101)^*$  and  $1(01^*)^*$  over  $\{01,2\}$ , then handle the intersection and complementation, and then deal with the star. Finally reduce the resulting dfa. Consider de Morgan's laws!

4. Construct a Chomsky normal form grammar for  $L(G)$  for the following cfg  $G$ :

10  $G = (\{S, B\}, \{a, b, c, d\}, \{S \rightarrow BBBbS|B|bcd, B \rightarrow cSda|S|cba\}, S)$ .

Note: You must first remove all unit productions.

5. Construct a Greibach normal form grammar for  $L(G)$  for the following CNF  $G$ :

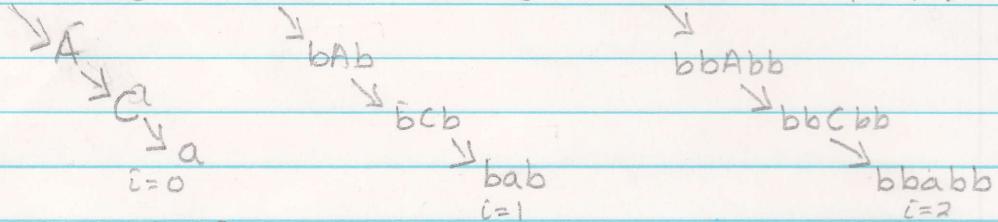
28  $G = (\{S, A, B\}, \{a, b\}, \{S \rightarrow AA, A \rightarrow BaA|a, B \rightarrow SSA|b\}, S)$ .

Note: First derive all the productions for  $S$ ,  $A$ , and  $B$ . You may only indicate how the final result looks for whatever primed variables you obtain.

Points: 1: 15 2: 10 3: 30 4: 15 5: 30

(V)  $S \rightarrow A|bB \quad A \rightarrow C \quad B \rightarrow Sb \quad C \rightarrow a$

$S \rightarrow bB \rightarrow bsb \rightarrow bbBb \rightarrow bbSbb \rightarrow \dots$



$$L(G) = \{ b^i a b^i \mid i \geq 0 \}$$

To prove  $L$  is not regular, assume  $L$  is regular

$\exists$  DFA  $D$  accepting  $L$

Let  $N$  be the number of states in  $D$

Consider the case  $X = b^i a b^i$  with  $|X| = N$

$X = w \cdot v$ ,  $w = b^N$  and  $v = a b^N$ ,  $|w| = N$

Pumping-Lemma:

$w = w_1 \cdot w_2 \cdot w_3$  s.t.  $|w_2| \geq 1$  and

$\tau(q_0, w) = \tau(q_0, w_1 (w_2)^s \cdot w_3) \quad \forall s \geq 0$

where  $w \cdot v \in L \cdot w_1 (w_2)^s w_3 \in L \quad \forall s \geq 0$

If  $s = 0$ ,  $w_1 (w_2)^0 w_3 = w_1 w_3$

$|w_1 (w_2)^0 w_3| = N - |w_2| < N, \quad |w_2| \geq 1$

$\Rightarrow b^{N-|w_2|} a b^N \notin L$  (not accepting the automata)

Also,  $\tau(q_0, v) = \tau(q_0, w_1 v)$

$= \tau(q_0, w, v)$

$= \tau(q_0, w, w_3, v)$

$\in L(D)$ , which is accepting the automata.

But  $b^{N-|w_2|} a b^N \notin L$

There is a contradiction  $\Rightarrow L(G)$  is not regular

(2)  $S \rightarrow aA \mid BBBa$

$A \rightarrow cS \mid \epsilon$

$B \rightarrow Sb \mid \epsilon$

Eliminate  $\epsilon$ -productions:

When  $B \rightarrow \epsilon$ :  $S \rightarrow aA \mid BBBa \mid BBa \mid Ba \mid a$   
 $A \rightarrow cS \mid \epsilon$   
 $B \rightarrow Sb$

When  $A \rightarrow \epsilon$ :  $S \rightarrow aA \mid a \mid BBBa \mid BBa \mid Ba \mid \times$   
 $A \rightarrow cS$   
 $B \rightarrow Sb$

Therefore

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$S \rightarrow aA \mid a \mid BBBa \mid BBa \mid Ba$   
 $A \rightarrow cS$   
 $B \rightarrow Sb$

$$(3) [(101)^* \cap \overline{(01^*)^*}]^* \equiv \left[ \frac{123}{(101)^*} \cup \frac{456}{(01^*)^*} \right]^* \text{ over } \{0, 1, 2\}$$

$$\begin{array}{c} (101)^* \\ \rightarrow 0 \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline - & 1 & 0 \\ \hline 2 & - & 0 \\ \hline 1 & - & 0 \\ \hline 3 & - & 0 \\ \hline \end{array} \end{array}$$

$$\begin{array}{c} 123 \quad 456 \\ \rightarrow 0 \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline - & 4 & 0 \\ \hline 4 & - & 1 \\ \hline \end{array} \quad \rightarrow 0 \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 5 & - & 0 \\ \hline 5 & - & 1 \\ \hline \end{array} \quad \rightarrow 0 \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 6 & - & 0 \\ \hline 6 & - & 1 \\ \hline \end{array} \end{array}$$

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$$\begin{array}{c} (101)^* \\ \rightarrow 0 \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline - & 1 & 1 \\ \hline 2 & - & 0 \\ \hline 1 & - & 0 \\ \hline 3 & - & 0 \\ \hline \end{array} \end{array}$$

$(101)^*$  dfa ↓

$$\begin{array}{c} \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 7 & 1 & 7 \\ \hline 7 & 7 & 7 \\ \hline 1 & 0 & 1 \\ \hline 2 & 0 & 1 \\ \hline 2 & 7 & 7 \\ \hline 3 & 7 & 7 \\ \hline 3 & 7 & 1 \\ \hline \end{array} \end{array}$$

Let 7 be null

$$\begin{array}{c} \text{rename} \downarrow \\ \rightarrow A \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline B & C & B \\ \hline B & B & B \\ \hline C & D & B \\ \hline D & B & E \\ \hline E & B & C \\ \hline \end{array} \end{array}$$

$(101)^* \cup 1(01^*)^*$

$$\begin{array}{c} \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 7 & 1,4 & 7 \\ \hline 1 & 2 & 7 \\ \hline 2 & 7 & 3 \\ \hline 3 & 7 & 1 \\ \hline 4 & - & - \\ \hline 5 & 5 & - \\ \hline 6 & 5 & 6 \\ \hline 7 & 7 & 7 \\ \hline \end{array} \end{array}$$

dfa

$$\begin{array}{c} \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 7 & 1,4 & 7 \\ \hline 7 & 7 & 7 \\ \hline 1,4 & 2,5 & 7 \\ \hline 2,5 & 5,7 & 3,6 \\ \hline 5,7 & 5,7 & 6,7 \\ \hline 3,6 & 5,7 & 1,6 \\ \hline 6,7 & 5,7 & 6,7 \\ \hline 1,6 & 2,5 & 6,7 \\ \hline \end{array} \end{array}$$

$$\begin{array}{c} \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 7 & 1,4 & 7 \\ \hline 7 & 7 & 7 \\ \hline 1,4 & 2,5 & 7 \\ \hline 2,5 & 5,7 & 3,6 \\ \hline 5,7 & 5,7 & 6,7 \\ \hline 3,6 & 5,7 & 1,6 \\ \hline 6,7 & 5,7 & 6,7 \\ \hline 1,6 & 2,5 & 6,7 \\ \hline \end{array} \end{array}$$

$$\begin{array}{c} \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline A & B & C \\ \hline B & B & B \\ \hline C & D & B \\ \hline D & E & F \\ \hline E & E & G \\ \hline F & E & H \\ \hline G & E & G \\ \hline H & D & G \\ \hline \end{array} \end{array}$$

$$\begin{array}{c} A \\ \begin{array}{|c|c|} \hline A & BCDEFGH \\ \hline A & BCDEFGH \\ \hline \end{array} \\ \begin{array}{c} 1 \\ 2 \end{array} \end{array}$$

$$\begin{array}{c} \begin{array}{|c|c|c|} \hline 0 & 1 & 2 \\ \hline 1 & 2 & 2 \\ \hline 2 & 2 & 2 \\ \hline \end{array} \end{array}$$

EG

(4)  $S \rightarrow BBBB \alpha S \mid B \mid bcd$   
 $B \rightarrow cSda \mid S \mid cba$

There does not exist an  $\epsilon$ -production  
 nor useless symbol in this case

Eliminate Unit production =

$S \rightarrow B \mid S \rightarrow BBBB \alpha S \mid BBBB \alpha B \mid bcd$   
 $B \rightarrow cSda \mid cBda \mid cba \mid S \mid \times$

21 missing prod<sup>o</sup>!  
 -5

$B \rightarrow S$

$S \rightarrow BBBB \alpha S \mid BBBB \alpha S \mid BSBBA \alpha S \mid BBSB \alpha S \mid BBBS \alpha S \mid BSSBB \alpha S \mid$   
 $SBSSB \alpha S \mid SBBS \alpha S \mid SSSB \alpha S \mid SSBSS \alpha S \mid SSSS \alpha S \mid bcd$

$B \rightarrow cSda \mid cBda \mid cba$

Convert to CNF

$S \rightarrow BBBB \alpha S \mid SBBB \alpha S \mid BSBBA \alpha S \mid BBSB \alpha S \mid BBBS \alpha S \mid SSBBA \alpha S \mid$   
 $SBSSB \alpha S \mid SBBS \alpha S \mid SSSB \alpha S \mid SSBSS \alpha S \mid SSSS \alpha S \mid X_b X_c X_d$

$B \rightarrow X_c S X_d X_a \mid X_c B X_d X_a \mid X_c X_b X_a$

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$S \rightarrow BS_1 \mid SS_5 \mid BS_9 \mid BS_{13} \mid BS_{17} \mid SS_{21} \mid SS_{25} \mid SS_{29} \mid SS_{33} \mid SS_{37} \mid SS_{41} \mid X_b S_{45}$

$S_1 \rightarrow BS_2 \quad S_5 \rightarrow BS_6 \quad S_9 \rightarrow SS_{10} \quad S_{13} \rightarrow BS_{14} \quad S_{17} \rightarrow BS_{18} \quad S_{21} \rightarrow SS_{22}$   
 $S_2 \rightarrow BS_3 \quad S_6 \rightarrow BS_7 \quad S_{10} \rightarrow BS_{11} \quad S_{14} \rightarrow SS_{15} \quad S_{18} \rightarrow BS_{19} \quad S_{22} \rightarrow BS_{23}$   
 $S_3 \rightarrow BS_4 \quad S_7 \rightarrow BS_8 \quad S_{11} \rightarrow BS_{12} \quad S_{15} \rightarrow BS_{16} \quad S_{19} \rightarrow SS_{20} \quad S_{23} \rightarrow BS_{24}$   
 $S_4 \rightarrow X_a S \quad S_8 \rightarrow X_a S \quad S_{12} \rightarrow X_a S \quad S_{16} \rightarrow X_a S \quad S_{20} \rightarrow X_a S \quad S_{24} \rightarrow X_a S$

$S_{25} \rightarrow BS_{26} \quad S_{29} \rightarrow BS_{29} \quad S_{33} \rightarrow SS_{34} \quad S_{37} \rightarrow SS_{38} \quad S_{41} \rightarrow SS_{42} \quad S_{45} \rightarrow X_c X_d$   
 $S_{26} \rightarrow SS_{27} \quad S_{30} \rightarrow BS_{30} \quad S_{34} \rightarrow SS_{35} \quad S_{38} \rightarrow BS_{39} \quad S_{42} \rightarrow SS_{43}$   
 $S_{27} \rightarrow BS_{28} \quad S_{31} \rightarrow SS_{32} \quad S_{35} \rightarrow BS_{36} \quad S_{39} \rightarrow SS_{40} \quad S_{43} \rightarrow SS_{44}$   
 $S_{28} \rightarrow X_a S \quad S_{32} \rightarrow X_a S \quad S_{36} \rightarrow X_a S \quad S_{40} \rightarrow X_a S \quad S_{44} \rightarrow X_a S$

$B \rightarrow X_c B_1 \mid X_c B_3 \mid X_c B_5$   
 $B_1 \rightarrow S B_2 \quad B_3 \rightarrow BB_4 \quad B_5 \rightarrow X_b X_a$   
 $B_2 \rightarrow X_d X_a \quad B_4 \rightarrow X_d X_a$

$X_a \rightarrow a$

$X_b \rightarrow b$

$X_c \rightarrow c$

$X_d \rightarrow d$

It's in CNF so no useless symbols,  $\epsilon$ , and Unit production

(5)

$S \rightarrow AA \quad \checkmark \quad i=1 \quad S \rightarrow AA$

$A \rightarrow BaA/a \quad \checkmark \quad i=2 \quad j=1 \Rightarrow A \rightarrow BaA/a$

$B \rightarrow \underline{SSA}/b \quad i=3 \Rightarrow j=1 =$

$\hat{B} = B \quad j=1 \Rightarrow B \rightarrow \underline{AASA}/b$

$S \rightarrow B$

$i=3 \quad j=2 \quad B \rightarrow aASA/BaAASA/b$

$A \rightarrow B \quad B_1 \quad a \quad B_2$

Eliminate immediate left recursion

$B \rightarrow b/bB'/aASA/aASAB'$

$B' \rightarrow aAASA/aAASAB'$

GNF

$A \rightarrow b\textcircled{a}A/bB'aA/aASAaA/aASAB'aA/a$

$S -$

$S \rightarrow b\textcircled{a}AA/bB\textcircled{a}AA/aASAaAA/aASAB'aAA/aA$

