

Appendix B

Big O and Little o Notation

It is often useful to talk about the *rate* at which some function changes as its argument grows (or shrinks), without worrying too much about the detailed form. This is what the $O(\cdot)$ and $o(\cdot)$ notation lets us do.

A function $f(n)$ is “of constant order”, or “of order 1” when there exists some non-zero constant c such that

$$\frac{f(n)}{c} \rightarrow 1 \quad (\text{B.1})$$

as $n \rightarrow \infty$; equivalently, since c is a constant, $f(n) \rightarrow c$ as $n \rightarrow \infty$. It doesn't matter how big or how small c is, just so long as there is some such constant. We then write

$$f(n) = O(1) \quad (\text{B.2})$$

and say that “the proportionality constant c gets absorbed into the big O ”. For example, if $f(n) = 37$, then $f(n) = O(1)$. But if $g(n) = 37(1 - \frac{2}{n})$, then $g(n) = O(1)$ also.

The other orders are defined recursively. Saying

$$g(n) = O(f(n)) \quad (\text{B.3})$$

means

$$\frac{g(n)}{f(n)} = O(1) \quad (\text{B.4})$$

or

$$\frac{g(n)}{f(n)} \rightarrow c \quad (\text{B.5})$$

as $n \rightarrow \infty$ — that is to say, $g(n)$ is “of the same order” as $f(n)$, and they “grow at the same rate”, or “shrink at the same rate”. For example, a quadratic function $a_1 n^2 + a_2 n + a_3 = O(n^2)$, no matter what the coefficients are. On the other hand, $b_1 n^{-2} + b_2 n^{-1}$ is $O(n^{-1})$.

Big- O means “is of the same order as”. The corresponding little- o means “is ultimately smaller than”: $f(n) = o(1)$ means that $f(n)/c \rightarrow 0$ for any constant c . Recursively, $g(n) = o(f(n))$ means $g(n)/f(n) = o(1)$, or $g(n)/f(n) \rightarrow 0$. We also read $g(n) = o(f(n))$ as “ $g(n)$ is ultimately negligible compared to $f(n)$ ”.

There are some rules for arithmetic with big- O symbols:

- If $g(n) = O(f(n))$, then $cg(n) = O(f(n))$ for any constant c .
- If $g_1(n)$ and $g_2(n)$ are both $O(f(n))$, then so is $g_1(n) + g_2(n)$.
- If $g_1(n) = O(f(n))$ but $g_2(n) = o(f(n))$, then $g_1(n) + g_2(n) = O(f(n))$.
- If $g(n) = O(f(n))$, and $f(n) = o(h(n))$, then $g(n) = o(h(n))$.

These are not *all* of the rules, but they’re enough for most purposes.