

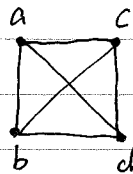
Problem set 1

1. a. distance (e, i) = 5

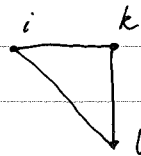
1. b. diameter = 3

1. c. average path length = $\frac{1}{\binom{4}{2}} [\text{distance}(a, b) + \text{distance}(a, c) + \text{distance}(a, d) + \text{distance}(b, c) + \text{distance}(b, d) + \text{distance}(c, d)]$
 $= \frac{1}{6} [1 + 1 + 1 + 2 + 2 + 2]$
 $= \frac{3}{2}$

2a. I find 2 cliques of network B:



clique 1



clique 2

2b. clustering coefficient of network B:

• common node:

• common node:

• common node:

There're 24 pairs of nodes that are linked to a common node, and only 15 of these are themselves linked, therefore, the clustering coefficient is $\frac{15}{24} = \frac{5}{8}$.

3. Since $S \cup (V \setminus S) = V$, nodes in V are either in S or $V \setminus S$, edges leaving the set S must have one endpoint in S and another endpoint in $V \setminus S$. Therefore, $|e(S)|$ the number of edges leaving the set S is the number of edges crossing the cut $C(S, V \setminus S)$, that is, $|C(S, V \setminus S)|$. For the same reason, $|e(V \setminus S)| = |C(S, V \setminus S)|$.

We can break down the definition as:

$$\min_{S \subseteq V} \frac{|C(S, V \setminus S)|}{\min\{|S|, |V \setminus S|\}} = \min \left\{ \min_{\substack{S \subseteq V \\ 1 \leq |S| \leq \frac{n}{2}}} \frac{|C(S, V \setminus S)|}{\min\{|S|, |V \setminus S|\}}, \min_{\substack{S \subseteq V \\ \frac{n}{2} \leq |S| \leq n}} \frac{|C(S, V \setminus S)|}{\min\{|S|, |V \setminus S|\}} \right\}$$

① For $1 \leq |S| \leq \frac{n}{2}$, $\min\{|S|, |V \setminus S|\} = |S|$

$$\Rightarrow \min_{\substack{S \subseteq V \\ 1 \leq |S| \leq \frac{n}{2}}} \frac{|C(S, V \setminus S)|}{\min\{|S|, |V \setminus S|\}} = \min_{\substack{S \subseteq V \\ 1 \leq |S| \leq \frac{n}{2}}} \frac{|e(S)|}{|S|}$$

② For $\frac{n}{2} \leq |S| \leq n$, $1 \leq |V \setminus S| \leq \frac{n}{2}$

$$\Rightarrow \min\{|S|, |V \setminus S|\} = |V \setminus S|$$

$$\Rightarrow \min_{\substack{S \subseteq V \\ \frac{n}{2} \leq |S| \leq n}} \frac{|C(S, V \setminus S)|}{\min\{|S|, |V \setminus S|\}}$$

$$= \min_{\substack{S \subseteq V \\ 1 \leq |V \setminus S| \leq \frac{n}{2}}} \frac{|e(V \setminus S)|}{|V \setminus S|}$$

Let $S' = V \setminus S$

$$= \min_{\substack{S' \subseteq V \\ 1 \leq |S'| \leq \frac{n}{2}}} \frac{|e(S')|}{|S'|}$$

From ① + ② $\Rightarrow \min_{S \subseteq V} \frac{|C(S, V \setminus S)|}{\min\{|S|, |V \setminus S|\}}$

$$= \min \left\{ \min_{\substack{S \subseteq V \\ 1 \leq |S| \leq \frac{n}{2}}} \frac{|e(S)|}{|S|}, \min_{\substack{S' \subseteq V \\ 1 \leq |S'| \leq \frac{n}{2}}} \frac{|e(S')|}{|S'|} \right\}$$

$$= \min_{\substack{S \subseteq V \\ 1 \leq |S| \leq \frac{n}{2}}} \frac{|e(S)|}{|S|}$$

4 a. There're $\binom{4}{2} = \frac{4!}{2!(4-2)!} = \frac{4 \times 3}{2} = 6$ possible links

(i) $P(\text{no edge}) = (1-p)^6$

(ii) $P(\text{one edge}) = \binom{6}{1} p(1-p)^5 = 6p(1-p)^5$

(iii) $P(\text{more than one edge})$
 $= 1 - P(\text{no edge}) - P(\text{one edge})$
 $= 1 - (1-p)^6 - 6p(1-p)^5$

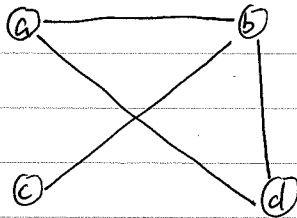
4 b. Let d_i denote the (random) degree of node i ($i=1, 2, 3, 4$), d_i can take any value in $\{0, 1, 2, 3\}$. The probability that $i \in V$ has exactly d links is:

$$P(d_i = d) = \binom{3}{d} p^d (1-p)^{3-d} = \binom{3}{d} \times 0.5^3$$

Therefore, the expected degree of each node in $G(4, 0.5)$ is:

$$\begin{aligned} E(d_i) &= 0 \times P(d_i=0) + 1 \times P(d_i=1) + 2 \times P(d_i=2) + 3 \times P(d_i=3) \\ &= 1 \times 3 \times 0.5^3 + 2 \times 3 \times 0.5^3 + 3 \times 1 \times 0.5^3 \\ &= 0.5^3 (3+6+3) \\ &= 1.5 \end{aligned}$$

4 c.



Pset1 Question 4. d

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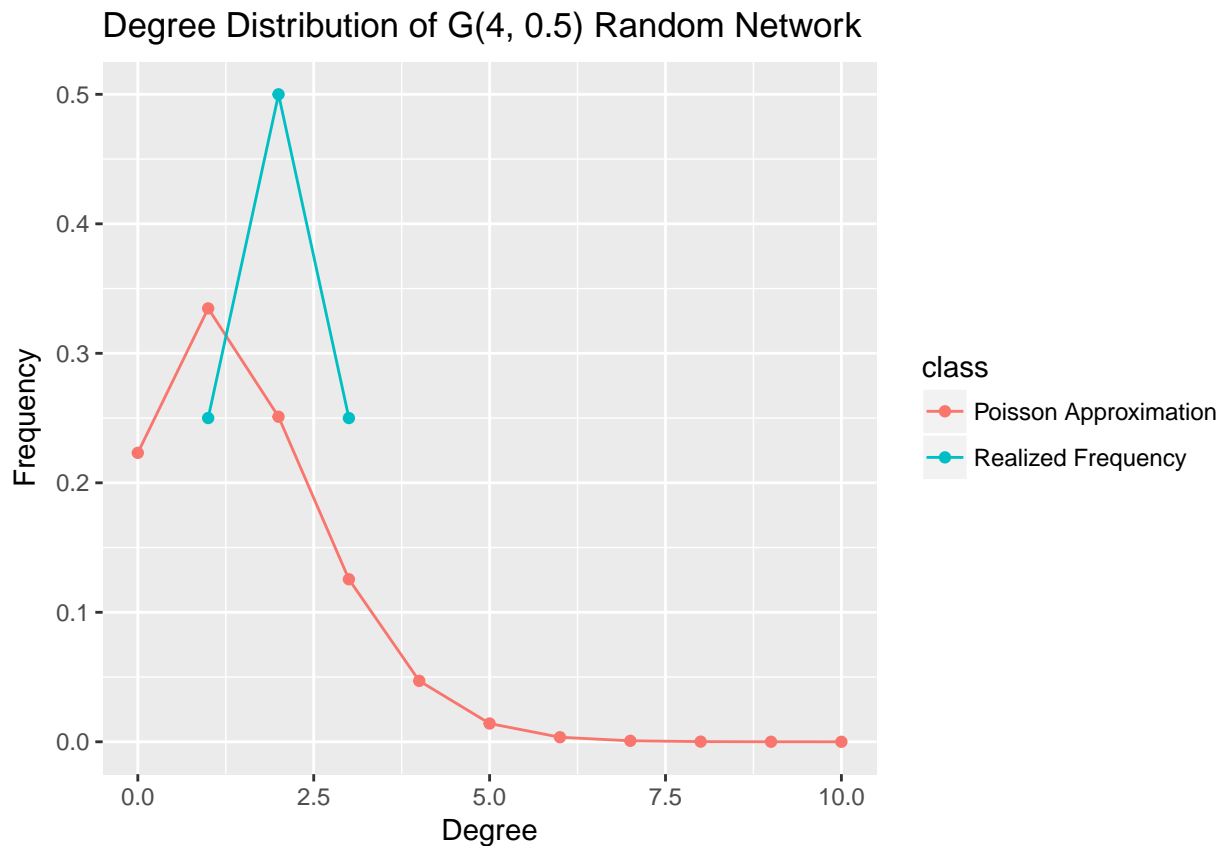
February 1, 2017

4 d

Frequency distribution of the randomly genrated network and the Poisson approximation for a probability of 0.5 on each link.

```
dat<-data.frame(degree=c(c(1,2,3),seq(0,10,1)),
                prob=c(c(0.25,0.5,0.25),dpois(seq(0,10,1),lambda = 1.5)),
                class=c(rep("Realized Frequency",3),rep("Poisson Approximation",11)))

library(ggplot2)
ggplot(dat,aes(x=degree, y=prob, color=class))+
  geom_point(aes())+geom_line()+
  xlab("Degree")+
  ylab("Frequency")+
  ggtitle("Degree Distribution of G(4, 0.5) Random Network")
```



Comment on the quality of the approximation: The Poisson distribution does not approximate the realized random network well, for such a small $n=4$ and large probability of link formation $p=0.5$.

5.a. Let $x \in [1, +\infty)$. $f(x) = (1 + \frac{1}{x})^x$

$$\Rightarrow \ln f(x) = x \ln(1 + \frac{1}{x})$$

$$\begin{aligned} \Rightarrow \frac{1}{f(x)} f'(x) &= \ln(1 + \frac{1}{x}) + x \cdot \frac{1}{1 + \frac{1}{x}} \cdot (-\frac{1}{x^2}) \\ &= \ln(1 + \frac{1}{x}) - \frac{1}{x+1} \quad (1) \end{aligned}$$

$$\text{Let } g(x) = \ln(1 + \frac{1}{x}) - \frac{1}{x+1}$$

$$g'(x) = \frac{1}{1 + \frac{1}{x}} \cdot (-\frac{1}{x^2}) + \frac{1}{(x+1)^2} = \frac{1}{(x+1)^2} - \frac{1}{x(x+1)} = -\frac{1}{x(x+1)^2}$$

for $x \geq 1$, we have $g'(x) < 0$, therefore, $g(x)$ is monotonically decreasing.

$$\text{And we also know } \lim_{x \rightarrow +\infty} g(x) = \lim_{x \rightarrow +\infty} \ln(1 + \frac{1}{x}) - \frac{1}{x+1} = 0$$

$$\text{Therefore, for } x \geq 1, g(x) > 0 \quad (2)$$

$$\text{From (1): } f'(x) = f(x) \cdot g(x)$$

$$> 0, \text{ for } x \geq 1$$

Therefore, $f(x)$ is monotonically increasing, for $x \geq 1$.

$$\Rightarrow f(x) \geq f(1), \text{ for } x \geq 1$$

$$\Rightarrow (1 + \frac{1}{x})^x \geq (1 + \frac{1}{1})^1 = 2, \text{ for } x \geq 1$$

$$\Rightarrow \text{Fix an integer } k \geq 1, (1 + \frac{1}{k})^k \geq 2$$

$$5.b. [(1 + \frac{1}{k})^k]^n \geq 2^n$$

$$\geq 2^{20} \quad (\text{for } n \geq 20)$$

$$= 2^{10+10} = 2^{10} \times 2^{10} = 1024 \times 1024$$

$$> 1000 \times 1000$$

$$= 10^3 \times 10^3$$

$$= 10^6$$

$$5.c. \text{ Return} = 1 \times (1 + p)^d = \left(1 + \frac{1}{k}\right)^{\frac{1}{p} \cdot pd}$$

Since $0 \leq p \leq 1$ (normally I would imagine nobody is willing to give $p > 1$)

$$\frac{1}{p} \geq 1, \text{ Let } k = \frac{1}{p} \geq 1$$

$$\Rightarrow \text{Return} = \left[\left(1 + \frac{1}{k}\right)^k\right]^{pd}$$

From (5b) we know that $pd \geq 20$ implies $\left[\left(1 + \frac{1}{k}\right)^k\right]^{pd} \geq 1 \text{ million}$

Therefore, $d \geq \frac{20}{p}$ implies return $\geq 1 \text{ million}$.

It would take you ceiling $\left(\frac{20}{p}\right)$ of days to become a millionaire.

For example, if daily compounded interest rate is 5%, then it only takes 40 days to become a millionaire (a little longer than a year!).