The Fairest Ransomware

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Table of Contents

Who am I?

2 Introduction

Definition and Notation





University of Tsukuba (Undergraduate)



- University of Tsukuba (Undergraduate)
- Authentication Platform Section, DCS Dept



- University of Tsukuba (Undergraduate)
- Authentication Platform Section, DCS Dept
- I'm interesting in Cryptography

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Ransomware

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Will ransomwares decrypt the data if the victim pays Bitcoin?

The Fairnest Ransomware

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If the victim pays some Bitcoins, their data will be decrypted under the probability on which they agreed.



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It's possible using cryptographic technique

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Symmetric Key Encryption (SKE) is a cryptographic scheme that uses the *same* key to encrypt and decrypt data like AES. An encryption function is denoted Enc, a decryption function is denoted Dec. The following equation holds for the symmetric key *k*.

$$x = \operatorname{Dec}_k \left(\operatorname{Enc}_k \left(x \right) \right)$$

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RSA Encryption is a cryptographic scheme that uses the *different* keys between encrypting and decrypting data. The key using encryption is called *public key* and The key using description is called *secret key*. The following holds for a public key (e, N) and the secret key d.

$$x = (x^e)^d = (x^d)^e \pmod{N}$$

Hash Function

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A hash function *H* has the following properties:

- Preimage Resistance A hash value h, it's difficult to find any message m such that h = H(m).
- Second Preimage Resistance An input m_1 , it's difficult to find different input m_2 such that $H(m_1) = H(m_2)$.
- Collision resistance It's difficult to find two different messages m_1 and m_2 such that $H(m_1) = H(m_2)$.

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Prover Alice has the secret key d for RSA chipher text c encrypted by public key (e, N). And she has chipher text $s := \operatorname{Enc}_k (c^d \mod N)$ and its symmetric key k.

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Bob want to verify as follows:

A preimage of the hash value H(k) is a symmetric key k that can decrypt SKE chipher text s.

without knowing the secret key d or symmetric key k.

We use *cut-and-choose protocol* where RSA chipher text c encrypted by public key (e, N) and its secret key d.

Thank you for your attention! Any question?