

Quick Navigation

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Premium

Solution

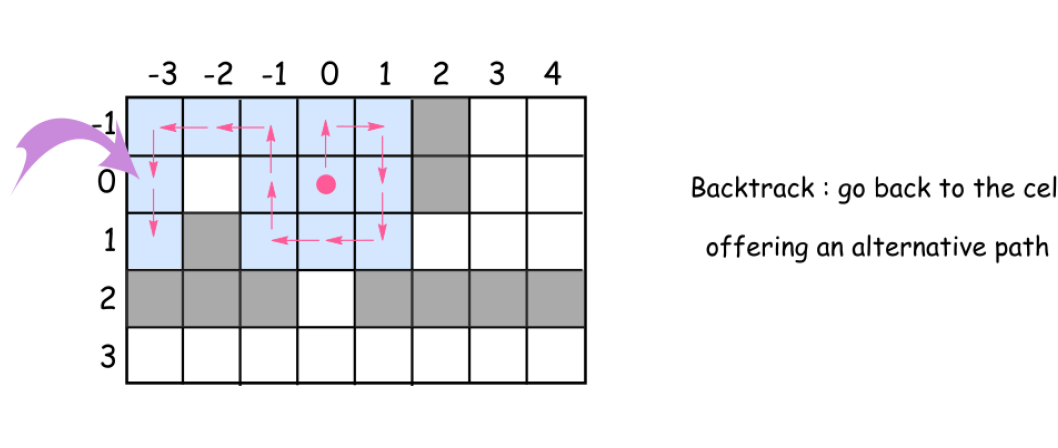
Approach 1: Spiral Backtracking

Concepts to use

Let's use here two programming concepts.

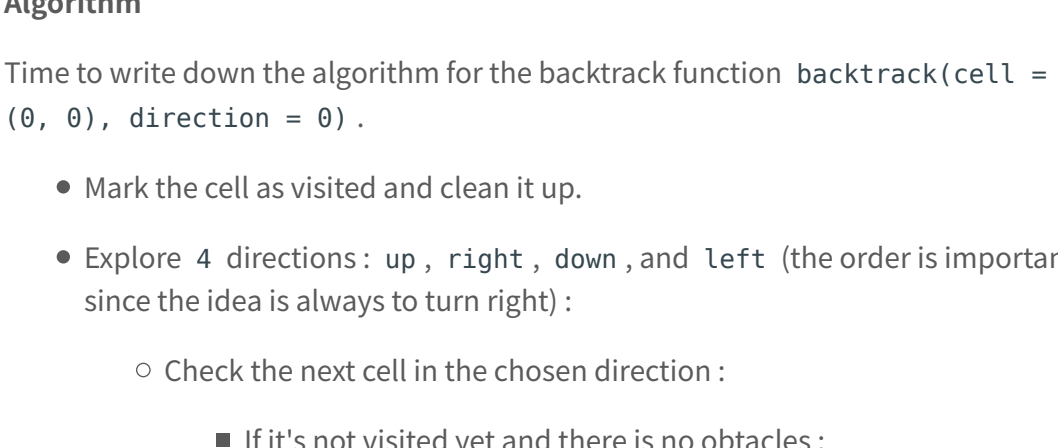
The first one is called *constrained programming*.

That basically means to put restrictions after each robot move. Robot moves, and the cell is marked as `visited`. That propagates *constraints* and helps to reduce the number of combinations to consider.



The second one called *backtracking*.

Let's imagine that after several moves the robot is surrounded by the visited cells. But several steps before there was a cell which proposed an alternative path to go. That path wasn't used and hence the room is not yet cleaned up. What to do? To *backtrack*. That means to come back to that cell, and to explore the alternative path.



Intuition

This solution is based on the same idea as maze solving algorithm called [right-hand rule](#). Go forward, cleaning and marking all the cells on the way as visited. At the obstacle *turn right*, again go forward, *etc*. Always *turn right* at the obstacles and then go forward. Consider already visited cells as virtual obstacles.

What do do if after the right turn there is an obstacle just in front ?

Turn *right* again.

How to explore the alternative paths from the cell ?

Go back to that cell and then *turn right* from your last explored direction.

When to stop ?

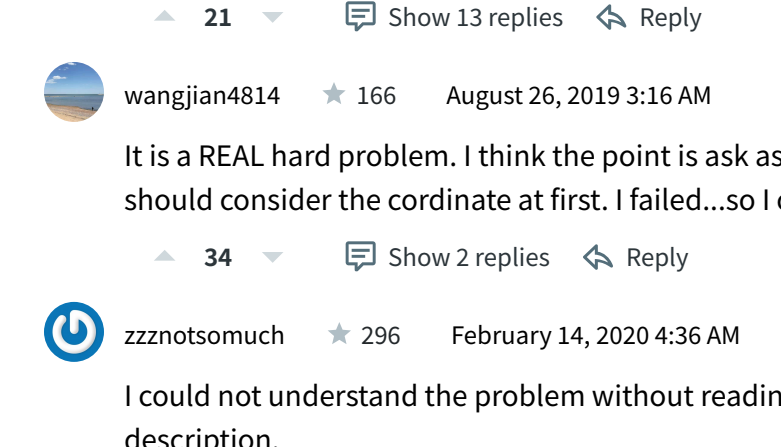
Stop when you explored all possible paths, *i.e.* all 4 directions (up, right, down, and left) for each visited cell.

Algorithm

Time to write down the algorithm for the backtrack function `backtrack(cell = (0, 0), direction = 0)`.

- Mark the cell as visited and clean it up.
- Explore 4 directions : up , right , down , and left (the order is important since the idea is always to turn right) :
 - Check the next cell in the chosen direction :
 - If it's not visited yet and there is no obstacles :
 - Move forward.
 - Explore next cells `backtrack(new_cell, new_direction)`.
 - Backtrack, *i.e.* go back to the previous cell.
 - Turn right because now there is an obstacle (or a virtual obstacle) just in front.

Implementation



Time to write down the algorithm for the backtrack function `backtrack(cell = (0, 0), direction = 0)`.

- Mark the cell as visited and clean it up.
- Explore 4 directions : up , right , down , and left (the order is important since the idea is always to turn right) :
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 - Backtrack, *i.e.* go back to the previous cell.
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Complexity Analysis

- Time complexity: $O(N \cdot M)$, where N is a number of cells in the room and M is a number of obstacles.
 - We visit each non-obstacle cell once and only once.
 - At each visit, we will check 4 directions around the cell. Therefore, the total number of operations would be $4 \cdot (N \cdot M)$.

- Space complexity: $O(N \cdot M)$, where N is a number of cells in the room and M is a number of obstacles.
 - We employed a hashtable to keep track of whether a non-obstacle cell is visited or not.

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calvinchankf 4279 April 4, 2019 10:27 AM

i am confused, why we need `new_d = (d + 1) % 4` ?

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wangjian4814 166 August 26, 2019 3:16 AM

It is a REAL hard problem. I think the point is ask as to figure out a small project should consider the coordinate at first. I failed...so I copy the answer.

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zznotsomuch 296 February 14, 2020 4:36 AM

I could not understand the problem without reading the solution. Please improve description.

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ka199 503 July 3, 2019 2:29 AM

i love this problem

28 Show 1 reply Reply

youjahan 47 June 24, 2019 12:01 AM

Each node is visited only once why the time complexity $O(4^N \cdot M)$

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hardfault 462 January 7, 2020 2:47 PM

I think this question is not hard to implement. But hard to understand xD

7 Reply

XeQR_DeV 9 March 2, 2020 9:25 PM

The solution starts the cleaning process from (0,0) if my understanding is correct supposed to be starting from a given position?

Can someone please help me out here.

6 Show 1 reply Reply

abilityfun 35 April 7, 2019 11:10 AM

@andvary, you are wrong on the time complexity. It is $M \cdot N$. The branching factor visited set. Hence, you will not visit the same square twice and complexity is $O(N \cdot M)$.

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coderonin 12 September 3, 2020 6:56 PM

There were a few things that were ambiguous in this question which I didn't know the solution:

- Where does the robot start? The question states that it can start anywhere coordinates passed as parameters to the `cleanRoom()` function. I see that as (0,0) but it is not obvious I am supposed to do that from reading the question. I would know if (0,0) is even a valid cell to start instead of an obstacle.
- Speaking of being blind, it is not very clear how blind my robot should be keeping a "visited" mapping is considered cheating, since it will exist only by looking at the solution can I know for sure that it is allowed.

3 Reply

jpakhar77 21 May 11, 2019 2:20 PM

It's called concepts, not conceptions, IMO.

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