

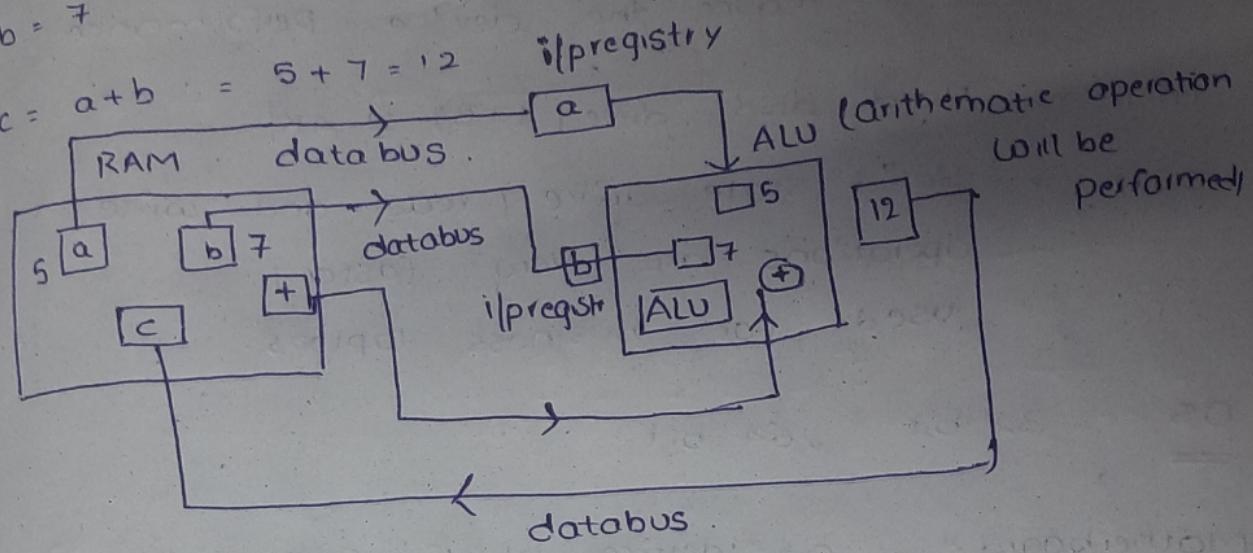
23/09/2023

Highlevel \rightarrow machine level
(0's & 1's)

$$a = 5$$

$$b = 7$$

$$c = a+b = 5+7=12$$



Desktop :- Specifications

$$1K = 10^3$$

$$1 \times 10^9 HZ$$

$$2 \times 10^9 HZ$$

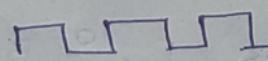
$$1M = 10^6$$

$$1 \times 10^9 CC/SEC$$

$$2 \times 10^9 CC/SEC$$

$$1G = 10^9$$

1 clock cycle = 1 instruction.



A can send 10^9 Instructions/sec and b can send 2×10^9 Instructions/sec

RAM \rightarrow Size/speed (we send data via buses we check how fast it is responding)

Speed

Cores (multicores helps run tasks parallelly)

Cache (L1, L2, L3 (storage and the quick retrieval of frequently used data))

OS

Motherboard

GPU : (VRAM)

Context switching :-

OS
P1
P2
P3

→ If we want P4 (a new program) to run we have to use the context switching but this might reduce the device performance.

HDD → Hard disk drive (1TB) \cong 4K/-

SSD → Solid state drive (read and write operations are very faster). Used for business laptops.

OS :- 32 bit, 64 bit.

Motherboard :- a device which connects the hard ware components with wires and make the hard ware software interaction.

Bit to Byte Conversion :-

$$1 \text{ Byte} = 8 \text{ Bit}$$

$$1 \text{ KB} = 1024 \text{ Bytes} = 2^{10} \text{ Bytes}$$

$$1 \text{ MB} = 1024 \text{ KB} = 2^{20} \text{ Bytes} = 2^{10} \text{ KB}$$

$$1 \text{ GB} = 1024 \text{ MB} = 2^{10} \text{ MB} = 2^{20} \text{ KB} = 2^{30} \text{ Bytes}$$

4 GB RAM uses

8GB / 4GB which has more performance for 32 bit Sys

$$\text{Sol: } 4 \text{ GB} = 4 \times 2^{30} = 2^2 \times 2^{30} = 2^{32}$$

for 2 Bits we have 4 Combinations

2 Bits : 0 0

0 1

1 0

1 1

For 3 bits we have 8 combinations.

3 Bits: 000

001

010

011

100

101

110

111

If we derive the formula:-

For n bit system it has 2^n combinations.

N Bitsystem can generate 2^n Bytes.

By Reverse Engineering:

$$4\text{GB} = 2^{32} \text{Bytes}$$

32 bit system generates 2^{32} Bytes.

so 4GB is compatible with 32 bit system.

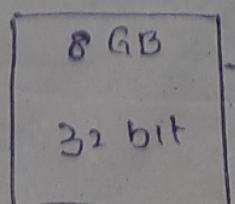
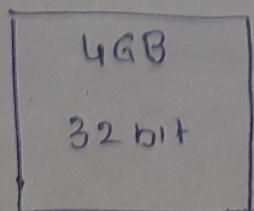
$$8\text{GB} = 2^{\underline{30}}$$

$= 2^{33}$ (33 bit system is not available).



The Space will be there to operate on more than 32 bits. but in the question it was asked for only 32 bit system.

So both are feasible. and the performance is same.



both function in a same way.

Convert 24 into Binary :-

$$\begin{array}{r} 24 \\ 2 \Big| 12 - 0 \\ 2 \Big| 6 - 0 \quad \uparrow \\ 2 \Big| 3 - 0 \\ 2 \Big| 1 - 1 \end{array}$$

$$\begin{array}{r} 00011000 \\ + 65 + 3 + 1 \\ \hline 0 \end{array}$$

Method : 2

$$24 \rightarrow \begin{array}{r} 0 \ 0 \ 0 \ 1 \ 1 \ 0 \ 0 \ 0 \\ - - - - - - - \\ 2^7 \ 2^6 \ 2^5 \ 2^4 \ 2^3 \ 2^2 \ 2^1 \ 2^0 \end{array}$$

$2^5 > 24$, $2^4 < 24$ (fill 1 at 2^4 th place)

$$24 - 2^4 = 24 - 16$$

$$= 8.$$

$2^3 = 8$, so fill 0 at 2^3 position

$$8 - 8$$

$$= 0.$$

$$19 \rightarrow \begin{array}{r} 0 \ 0 \ 0 \ 1 \ 0 \ 0 \ 1 \ 1 \\ - - - - - - - \\ 2^7 \ 2^6 \ 2^5 \ 2^4 \ 2^3 \ 2^2 \ 2^1 \ 2^0 \end{array}$$

$$2^5 < 19$$

$$2^4 > 19$$

$$19 - 16 = 3. \quad \textcircled{1}$$

$$2^3 > 3$$

$$2^2 > 3$$

$$2 < 3$$

$$3 - 2 = 1 \quad \textcircled{2}$$

$$2^0 = 1$$

$$1 - 1 = 0 \quad \textcircled{3}$$

Signed Magnitude representation :-

The MSB bit will represent the sign.

- (negative) \rightarrow 1 in MSB bit

+ (positive) \rightarrow 0 in MSB bit.

$$\begin{array}{r} 24 \rightarrow 00011000 \\ -19 \rightarrow 10010011 \\ \hline \end{array}$$

$1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1$ $\rightarrow 128$

$2^7 \quad 2^6 \quad 2^5 \quad 2^4 \quad 2^3 \quad 2^2 \quad 2^1 \quad 2^0$

$1 + 2 + 8 + 32 = 43$

represents \ominus sign

so result will be -43 .

but $24 - 19 \neq -43 \rightarrow$ ambiguity ①

$$\begin{array}{r} 0 \rightarrow 00000000 \\ -0 \rightarrow 10000000 \\ \hline \end{array}$$

" -0 " does not exists. so will create an ambiguity. \rightarrow ambiguity 2

To overcome this issue we are going for the 2's Compliment.

$$2^s \text{ Compliment} = 1^s \text{ Compliment} + 1$$

$$19 \rightarrow 00010011$$

$$1^s \text{ Compliment} \rightarrow 11101100$$

of 19

$$\begin{array}{r} 1 \rightarrow 00000001 \\ -19 \rightarrow 11101101 \\ \hline \end{array}$$

Mathematical operation on $24 - 19$

$$= 24 \rightarrow \begin{array}{r} 00011000 \\ -19 \rightarrow \begin{array}{r} 11101101 \\ \hline 00000101 \end{array} \end{array}$$

discarded

$$00000101 \rightarrow 5$$

$$24 - 19 = 5$$

Understanding the -19 Binary representation
and find the value of MSB Bit

$$-19 = \begin{array}{r} 11101101 \\ \downarrow \end{array}$$

$$-19 = x + 2^6 + 2^5 + 2^3 + 2^2 + 1$$

$$-19 = x + 109$$

$$x = -(109 + 19)$$

$$x = -128 = -2^7$$

$$\cancel{x = -127}$$

Formula derivation:-

for n bits the MSB represents -2^{N-1}

mathematical operation on $24 + 19$

$$24 \rightarrow 00011000$$

(-24) 's compliment of 24 $\rightarrow \begin{array}{r} 11110100 \\ \square \end{array}$

$$19 \rightarrow 00010011$$

$$\begin{array}{r} 11110100 \\ 00010011 \\ \hline 11111010 \end{array}$$

$2^6 2^5 2^4 2^3 2^2 2^1 2^0$

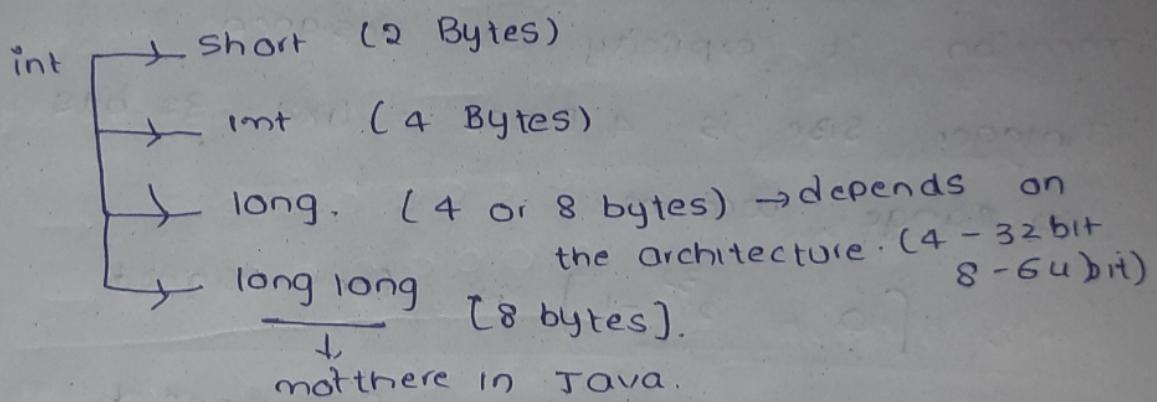
$$-128 + 64 + 32 + 16 + 8 + 2 =$$

1st Complement of 24

$$\begin{array}{r}
 \text{Comp} \oplus 2^4 = 0 \cdot 0 0 0 0 0 0 1 \\
 + 1 = \hline
 - 24 = 1 \cdot 1 1 0 1 0 0 0 \\
 + 19 = \hline
 \end{array}$$

$$\begin{array}{r}
 -128 + 64 + 32 + 8 + 4 + 2 + 1 \\
 - 8 + 3 \\
 = -5 \\
 - 24 + 19 = -5
 \end{array}$$

Datatypes:-



float → 4 bytes.

char → 1 byte / 2 bytes (Java).

Bool → 1 byte.

double → 8 bytes

16 bytes.

Deriving the ranges for the datatypes:-

	3bit	0 0 0 = 0	0 0 0 → 0	3bit Signed ranges
unsigned	0 0 1	= 1	0 0 1 → 1	from
ranges	0 1 0	= 2	0 1 0 → 2	[-7 3]
from	0 1 1	= 3	0 1 1 → 3	
	1 0 0	= 4	1 0 0 → -4	
[0 7]	1 0 1	= 5	1 0 1 → -3	
	1 1 0	= 6	1 1 0 → -2	
	1 1 1	= 7	1 1 1 → -1	

BITS	unsigned		Signed	
	Min	Max	Min	Max
2 Bits	0	3	-2	3
3 Bits	0	7	-4	7
4 Bits	0	15	-8	127
8 Bits	0	255	-128	$2^7 - 1$
N Bits	0	$2^n - 1$	-2^{n-1}	$+2^{n-1} - 1$

By default int is signed.
 If we want only +ve integers we have to mention it explicitly as unsigned.

Integer size is 4 Bytes = $4 \times 8 = 32$ bits.

so Range becomes:

$$[0 \quad 2^{32} - 1] \rightarrow \text{for unsigned}$$

$$[-2^{31} \quad 2^{31} - 1] \rightarrow \text{for signed.}$$

long long Range:-

8 bytes: $8 \times 8 = 64$ bits.

$$\text{ULL} \rightarrow [0 \quad 2^{64} - 1]$$

$$\text{LL} \rightarrow [-2^{63} \quad 2^{63} - 1]$$

Converting the powers of 2 to powers of 10 for easy understanding of range:-

$$2^{10} \rightarrow 10^3$$

$$2^{20} \rightarrow 10^6$$

$$2^{30} \rightarrow 10^9$$

$$2^{40} \rightarrow 10^{12}$$

$$2^{50} \rightarrow 10^{15}$$

$$2^{60} \rightarrow 10^{18}$$

int range = $[0 \quad 2^{32} - 1]$

for unsigned

 $= [0 \quad 2^{32} - 1]$
 $\approx [0 \quad 10^9]$

signed int range: $[-2^{31} \quad 2^{31} - 1]$

 $= [-2^{30} \quad 2^{30} - 1]$
 $= [-2 \times 10^9 \quad 2 \times 10^9 - 1]$
 $\approx [-10^9 \quad 10^9]$

long long range :-

ULL : $[0 \quad 2^{64} - 1]$

 $= [0 \quad 2^4 \times 2^{60} - 1]$
 $= [0 \quad 2^4 \times 10^{18} - 1]$
 $\approx [0 \quad 10^{18}]$

ll : $[-2^{63} \quad 2^{63} - 1]$

 $= [-2^3 \times 2^{60} \quad 2^3 \times 2^{60} - 1]$
 $\approx [-2^3 \times 10^{18} \quad 2^3 \times 10^{18} - 1]$
 $\approx [-10^{18} \quad 10^{18}]$

float :- 4 bytes :

Since it has decimal part & fraction part the range will be definitely less than int.

The range depends on the precision.

double : The range will be definitely less than long

but it depends on the precision.

Points to be considered when solving a

problem :-

Problem Statement

I/p format

O/p format

constraints

Sample I/o

Explanation.

Importance of constraints with example :-

sum of elements in an array with the
below constraints :-

$$1 \leq N \leq 10^3$$

$$-10^7 \leq arr[i] \leq 10^7$$

If all the 10^7 values has the value of 10^3 then the sum of all elements will be $10^7 \times 10^3 = 10^{10}$

so the range of varies from -10^{10} to 10^{10} result.

which is not in integer range and we have to select the datatype appropriately.