

Binary Search

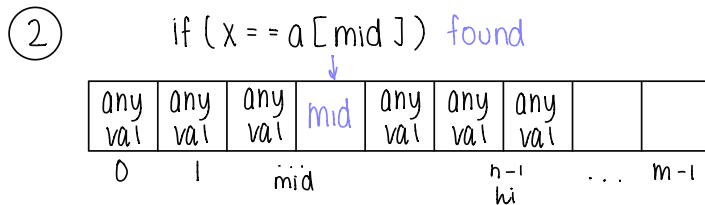
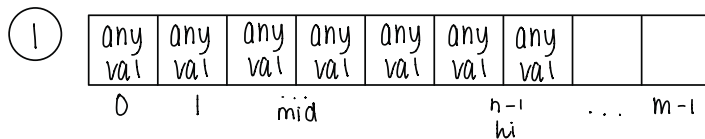
Idea: halve the space over which we search

- at each halving, eliminate $\frac{1}{2}$ of remaining possibilities
- max # of steps required is $O(\log n)$ ← halving is base 2

Precondition: Must be presorted

Demonstration:

3 indices \rightarrow lo, mid, hi



if ($x < a[mid]$)
search below mid

if ($x > a[mid]$)
search above mid

Code analysis:

```
bool binarySearch(int a[], int n, int x)
{
    int lo, mid, hi;

    lo = 0; // index vals
    hi = n-1;
    while (lo <= hi)
    {
        mid = (lo+hi)/2;

        if (a[mid] == x) // found
            return true;

        if (a[mid] < x) // search in lower half
            lo = mid+1;
        else // search in upper half
            hi = mid-1;
    }
    return false;
}
```