

## Extended Euclidean Algorithm (xgcd).

Fact: the gcd of  $a, b \in \mathbb{Z}$  is the smallest positive non-zero element of this set:

$$S = \{xa + yb \mid x, y \in \mathbb{Z}\}.$$

Sketch: say  $d = x^*a + y^*b$  is the smallest element of  $S$ . Note then that  $d$  divides all elements of  $S$ :

$$(xa + yb) = qd + r, \quad r \leq d$$

But  $\Rightarrow r = 0$ , as  $d$  was minimal:

$$xa + yb - qd = r$$

$$= a(x - qx^*) + b(y - qy^*)$$

$$(so, r \in S.)$$

$$\therefore d = \gcd(a, b). \quad \checkmark$$

[Goal for us: find  $x, y$  s.t.  
 $d = \gcd(a, b) = xa + yb$ .

Application: "modular inverses":

given  $a \in \{1, \dots, p-1\}$  for some large prime  $p$ ,

find  $e \in \{1, 2, \dots, p-1\}$  s.t.  $a \cdot x \not\equiv 0 \pmod{p} = 1 \quad x \equiv a^{-1}.$

To find  $x$ , we could just apply our goal to  $a, p$ .

$$xa + yp = 1 \Rightarrow xa = 1 - yp$$

$$xa \% p = (1 - yp) \% p = 1$$

## Details

inputs:  $a, b$

outputs:  $d, x, y$ .

outputs



```
int xgcd(int a, int b, int & x, int & y)
{
    if (b == 0) {
        x = 1;
        y = 0;
        return a; // d = a = 1*a + 0*b
    }
}
```

// Imagine  $xgcd$  works for all smaller  
// inputs (smaller values of  $b$ )

// How could the answer to  $xgcd(b, a \% b, -, -)$   
// help us?

int  $x', y'$ ;

int  $d = xgcd(b, a \% b, x', y')$ ;

//  $d = x' \cdot b + y' \cdot r$  ( $r = a \% b$ )

// how are  $x', y'$  useful???

// (Remember: we want  $x, y \in \mathbb{Z}$  s.t.  $d = xa + yb$ .)

Set  $a = qb + r$  ( $q = a/b$   $r = a \% b$ )

So,  $\underbrace{a - qb}_r = r$

$$d = x'b + y'r$$

$$= x'b + y'(a - qb)$$

$$= \underbrace{y'}_x a + \underbrace{(x' - y'q)}_y b$$

$$x = y';$$

$$y = x' - (a/b) \cdot y'$$

return d;

}