

## K-subsets

$$\mathcal{P}_k(S) = \{T \in \mathcal{P}(S) \mid |T| = k\}.$$

→ E.g. if  $S = \{1, 2, 3\}$ ,  $\mathcal{P}_2(S) = \{\{1, 2\}, \{1, 3\}, \{2, 3\}\}$ .

obvious way, since we already know how to compute  $\mathcal{P}(S)$ :

Just filter out the ones you want (size = k).

Say  $|S| = n$ . Then  $|\mathcal{P}_k(S)| = \binom{n}{k} \ll |\mathcal{P}(S)|$

for  $k$  close to 0  
or  $n$ .

$$\sum_{k=0}^n \binom{n}{k} = |\mathcal{P}(S)| = 2^n = (1+1)^n = \sum_{k=0}^n \binom{n}{k} 1^k 1^{n-k}$$

---

Algorithm outline. 2 size parameters:  $k, |S|$ .

Base case(s)?

if  $k = 0$ , answer =  $\{\{\}\}$ .

if  $k > |S|$ , answer =  $\{\}$ .

Now assume we have a magic box that sets correct answers  
for  $\mathcal{P}_{k'}(S')$  when  $k' < k$  or  $|S'| < |S|$ .

Example from above:

$$S = \{1, 2, 3\}, \quad \mathcal{P}_2(S) = \{\{1, 2\}, \{1, 3\}, \{2, 3\}\}.$$

Would  $\mathcal{P}_2(S \setminus \{3\})$  help? Sure!  $\mathcal{P}_k(S') \subseteq \mathcal{P}_k(S)$   
if  $S' \subseteq S$ .

all 2-subsets w/o 3.

What's missing? All 2-subsets that have 3...

Start w/  $P_{k-1}(S \setminus \{3\})$ , and then add 3 to each one.

$$P_2(S \setminus \{3\}) = \{\{1, 2\}\}$$

$$P_1(S \setminus \{3\}) + 3 = \{\{1\}, \{2\}\} + 3$$

$$= \{\{1, 3\}, \{2, 3\}\}.$$

$$P_k(S) \quad \checkmark$$

(Recall, for a set of sets  $\mathcal{P}$  and element  $x$ , we defined  $\mathcal{P} + x \triangleq \{T \cup \{x\} \mid T \in \mathcal{P}\}.$ )

---

Regarding last project: think of it as a multi-dimensional  
Cartesian Set