

Principal of Mathematical Induction (PMI)

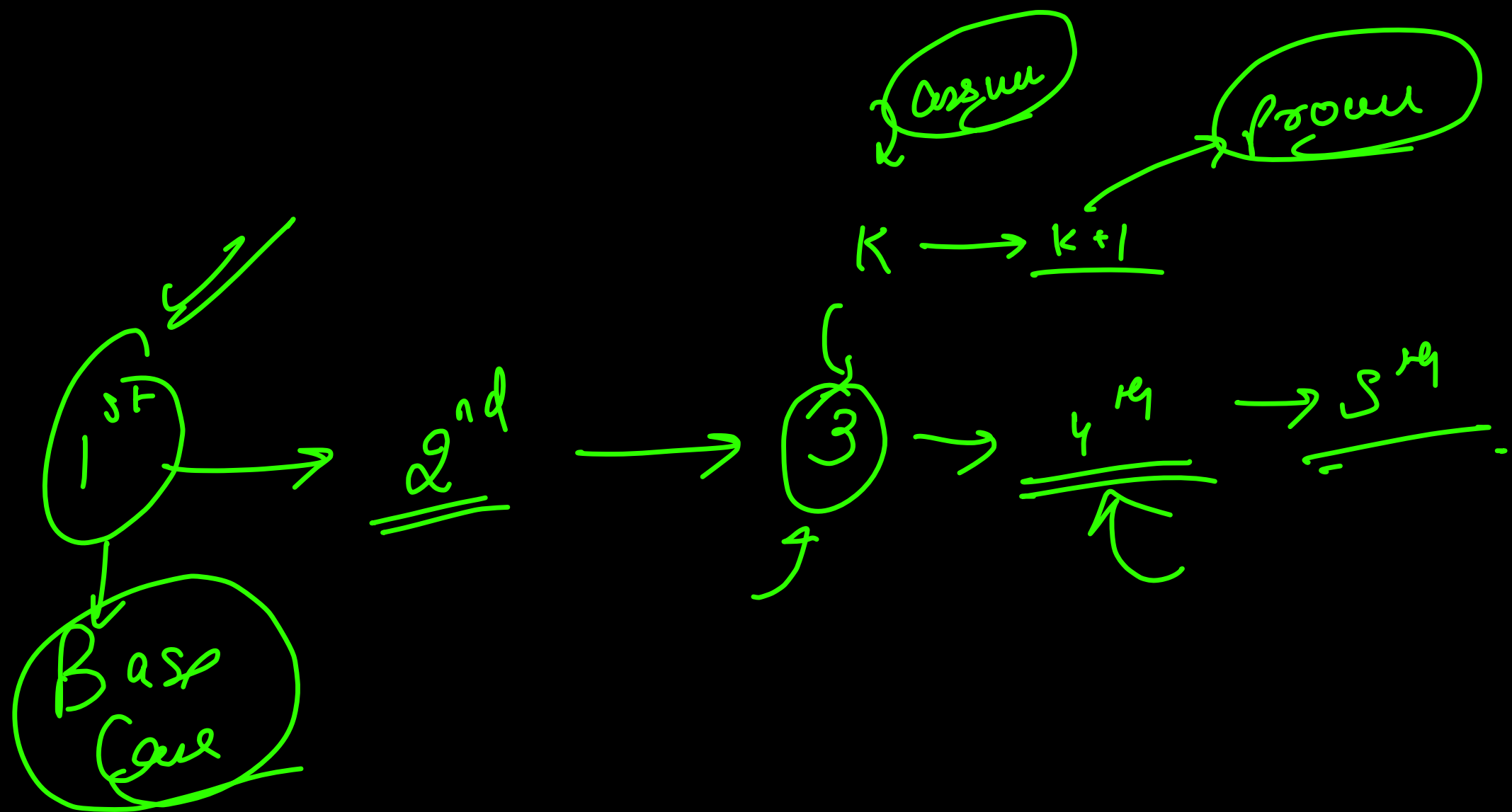
Q₁ Prove that sum of first n natural numbers,
is equal to $(n \times (n+1))/2$ $\rightarrow \frac{(k+1)(k+2)}{2}$

1) Base Case \rightarrow It is the smallest input value for which
we already know the ans.

$\hookrightarrow \underline{n=1}$

2) Assumption \rightarrow let's assume formulae work correctly for
 $n=k$

3) Self work → Using the fact that formula works for $n=k$, we will try to prove that formula works for $n=k+1$ also



assumption \rightarrow for $n=k$
i.e. Sum of first k natural no. is equal to $\frac{k \times (k+1)}{2}$

Selfwork \rightarrow Prove for $n=k+1$.

Sum of first $k+1$ natural no ??

$$1 + 2 + 3 + 4 + \dots + (k-1) + (k) + (k+1)$$

\hookrightarrow Sum of first k natural no.

$$= \frac{k \times (k+1)}{2} + (k+1) \Rightarrow (k+1) \left[\frac{k}{2} + 1 \right] \rightarrow \frac{(k+1)(k+2)}{2}$$

H.P

$$\underbrace{1 + 2 + 3 + 7}_{\downarrow} 6$$

Recursion

A child couldn't sleep, so her mother told a story about a little frog,
who couldn't sleep, so the frog's mother told a story about a little bear,
who couldn't sleep, so the bear's mother told a story about a little weasel
...who fell asleep.
...and the little bear fell asleep;
...and the little frog fell asleep;
...and the child fell asleep.

Recursion

→ It is a programming + math concept.

→ What is Recursion ??

Recursion is a technique using which we solve bigger problems by calculating ans of smaller subproblems. We generally denote the bigger problem as a funcⁿ, & some arguments, & then call the same funcⁿ with diff. arguments denoting smaller subproblems. So we get the ans of smaller subproblem & build the ans for bigger problem.

Recursion is funcⁿ calling itself.

We try to solve bigger problems using ans of smaller

subproblems

Ex factorial

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

$$7! = 7 \times 6 \times 5 \times 4 \times 3 \times 2 \times 1$$

$$1! = 1$$

$$n! \rightarrow n \times (n-1) \times (n-2) \times (n-3) \dots \dots \dots 2 \times 1$$

f_n Given a value n , calculate $n!$ recursively

Let's say f is a function which takes a value n
as an argument & can calculate $n!$

Base Case
↑

if ($n=1$) → 1

$f(n) = n \times f(n-1)$
↓
Calc → $n!$
↓
implement

$n \times$
↓
self-call

$f(n-1)$
↓
assume
 $f(n-1)$ works
correctly

$$\begin{aligned}
 5! &= 5 \times 4 \times 3 \times 2 \times 1 \\
 4! &= 4 \times 3 \times 2 \times 1 \\
 3! &= 3 \times 2 \times 1
 \end{aligned}$$

$$(k+1)! \quad ??$$

Base Case \rightarrow for $n=1$ we already know that $f(n)$ will be 1.

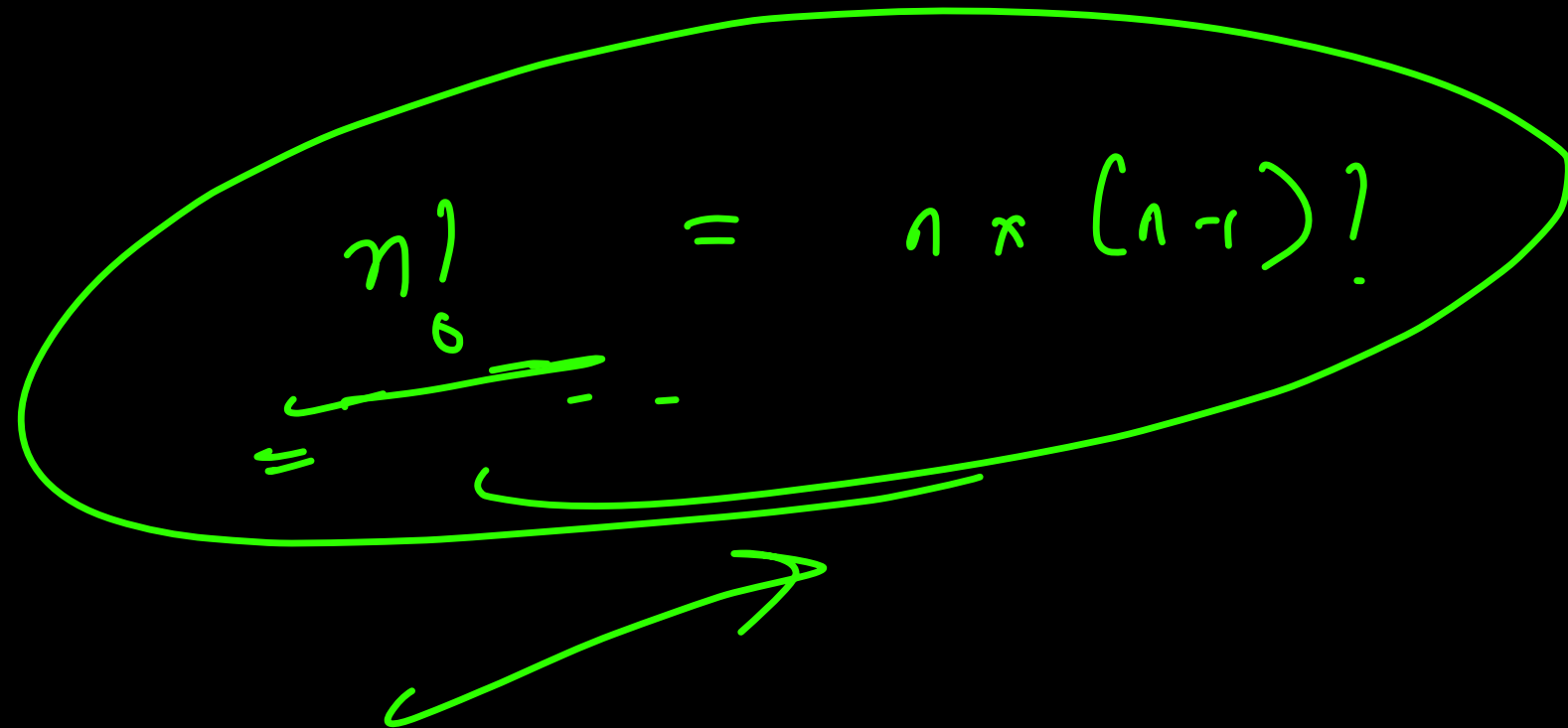
Assumption assume funcⁿ works correctly for some value k . $f(k) \rightarrow \checkmark \checkmark$

Selfwork $(n = k+1) \rightarrow \underline{(k+1) \times f(k)}$

$$(k+1)! \Rightarrow (k+1) \times \underbrace{(k) \times (k-1) \times \dots \times 3 \times 2 \times 1}_{k!}$$

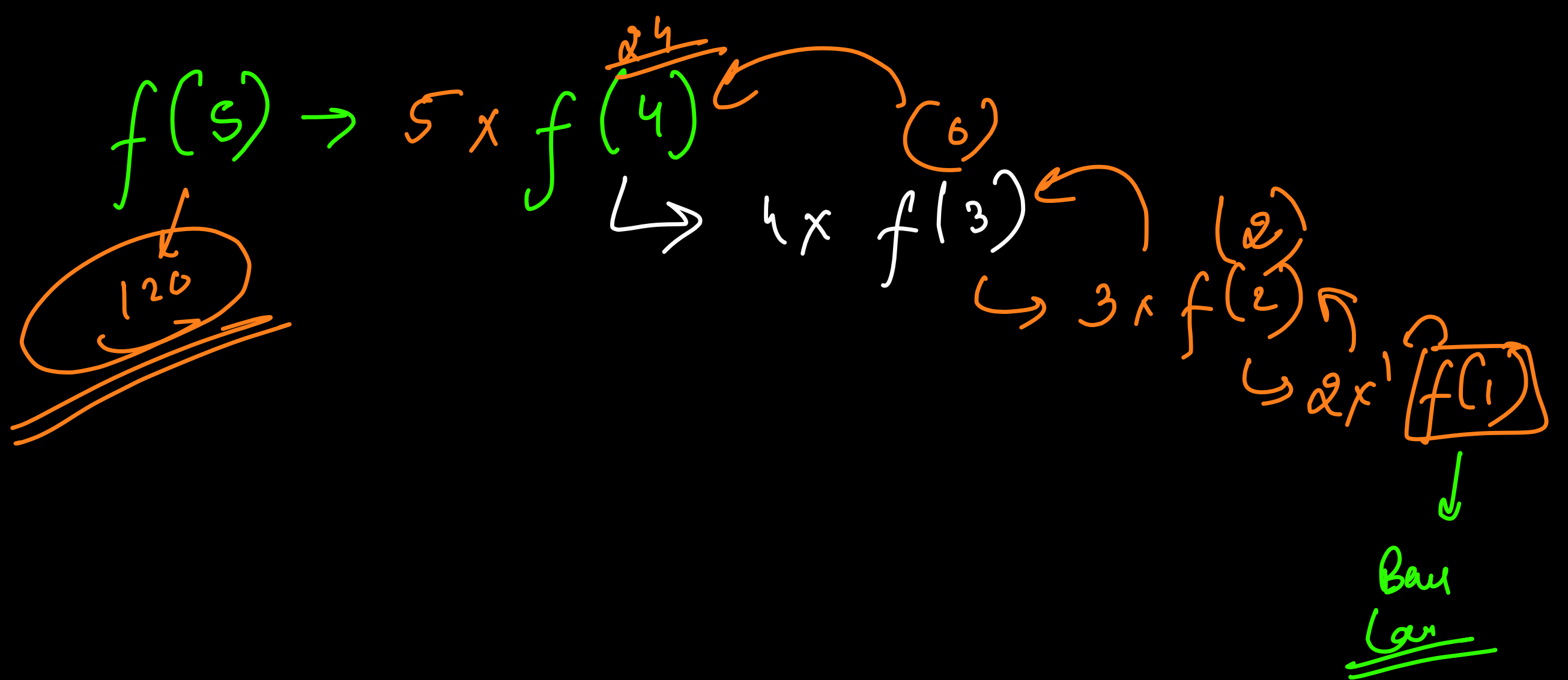
$k!$ \rightarrow we already know

$$(k+1)! \Rightarrow (k+1) \times k!$$

$$\eta_0 = n \pi (n-1)!$$


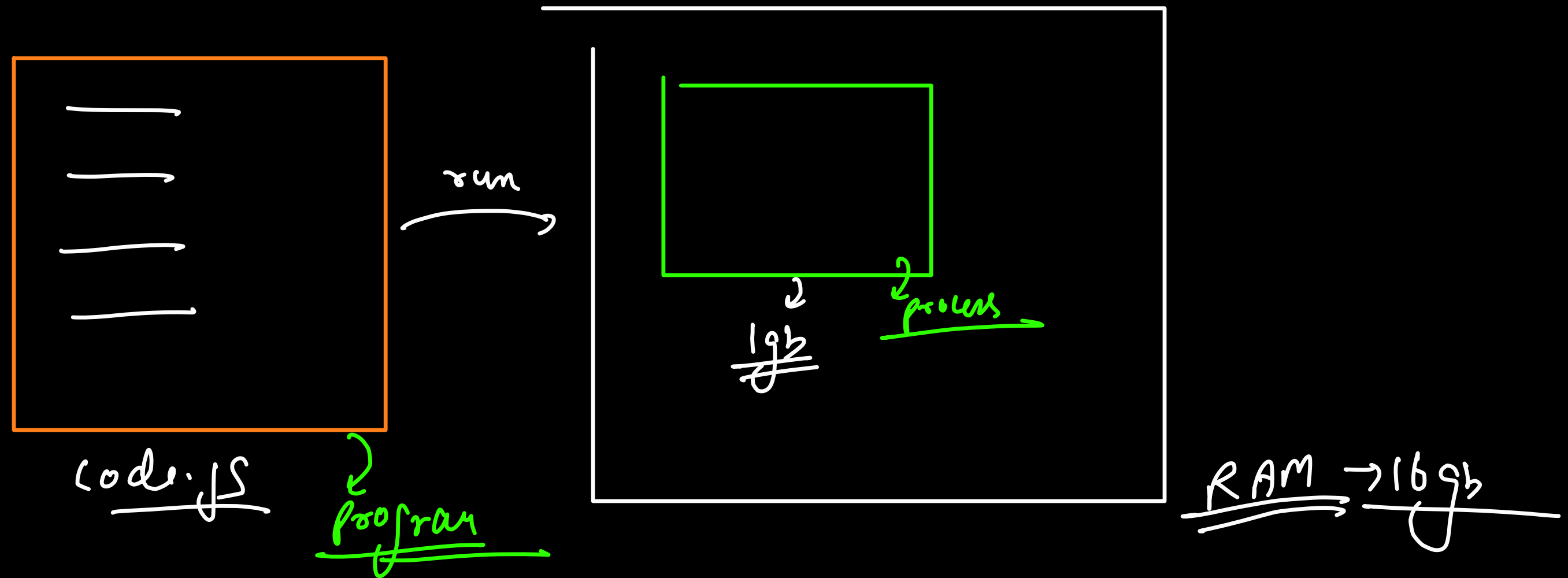
$$f(n) = n \times f(n-1)$$

$n!$

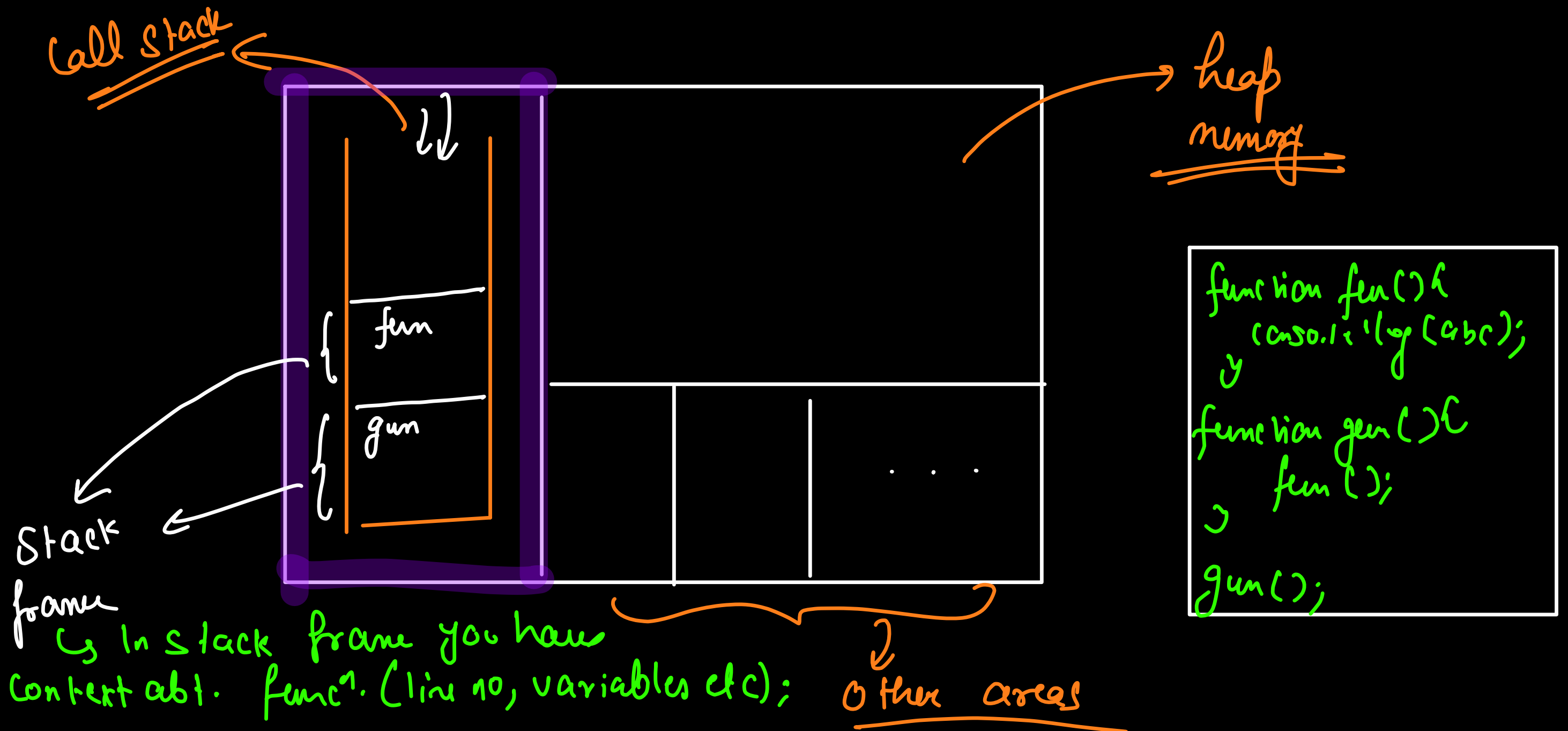


$f(n)$
 \downarrow
n-1

The magic comes with func.

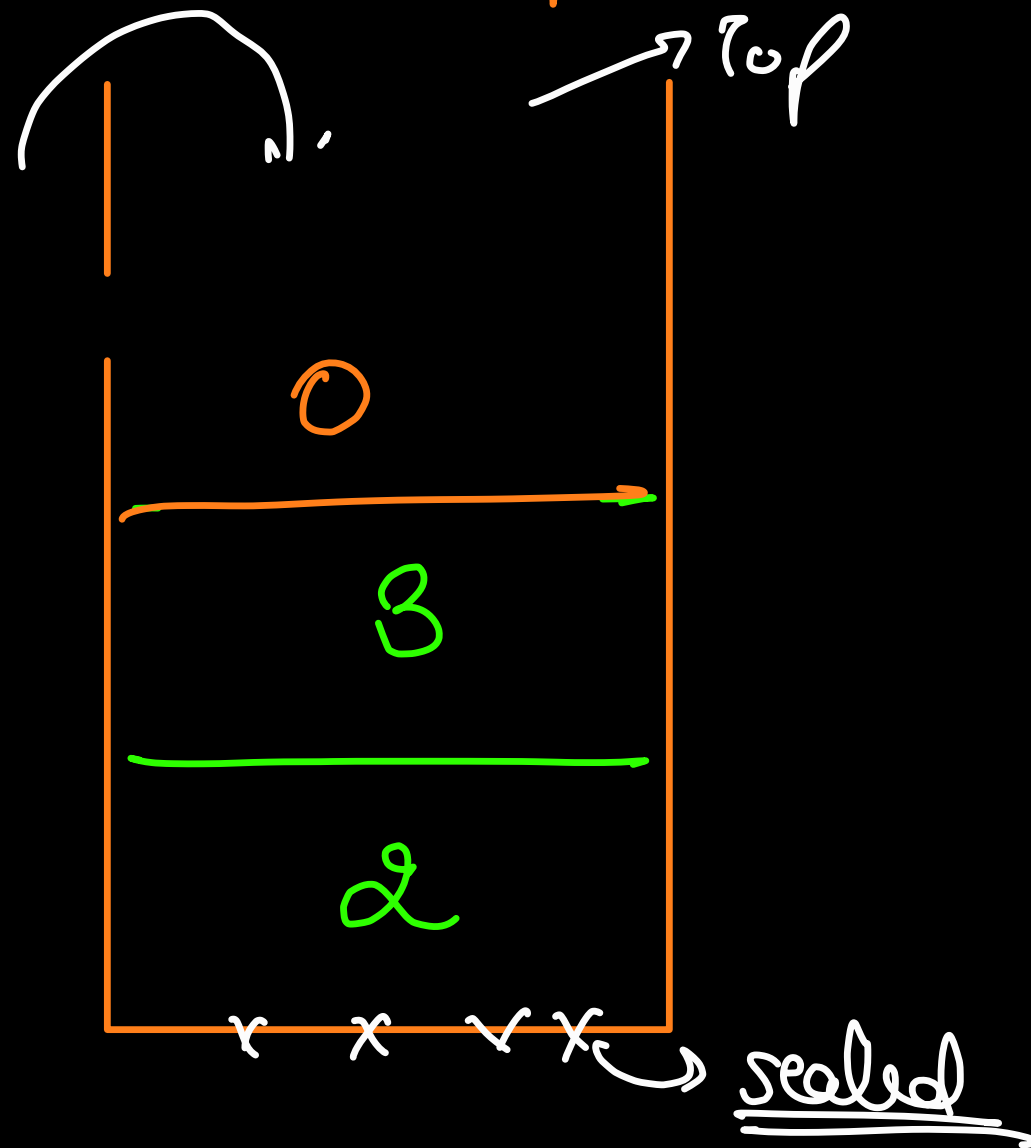


Program in a running state is process.



whenever we call a function from anywhere in the code, it adds a new entry in the call stack called as stack frame

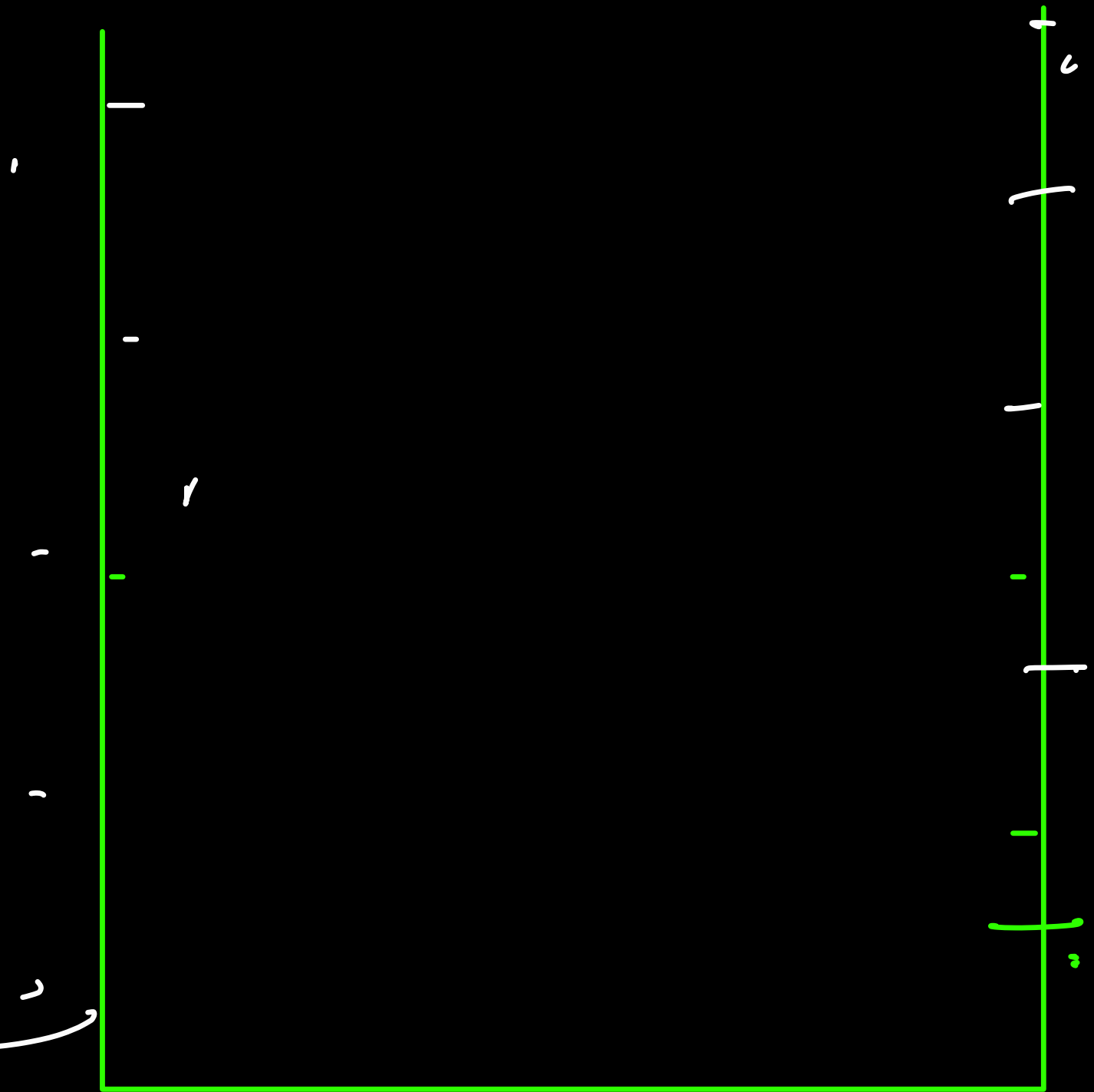
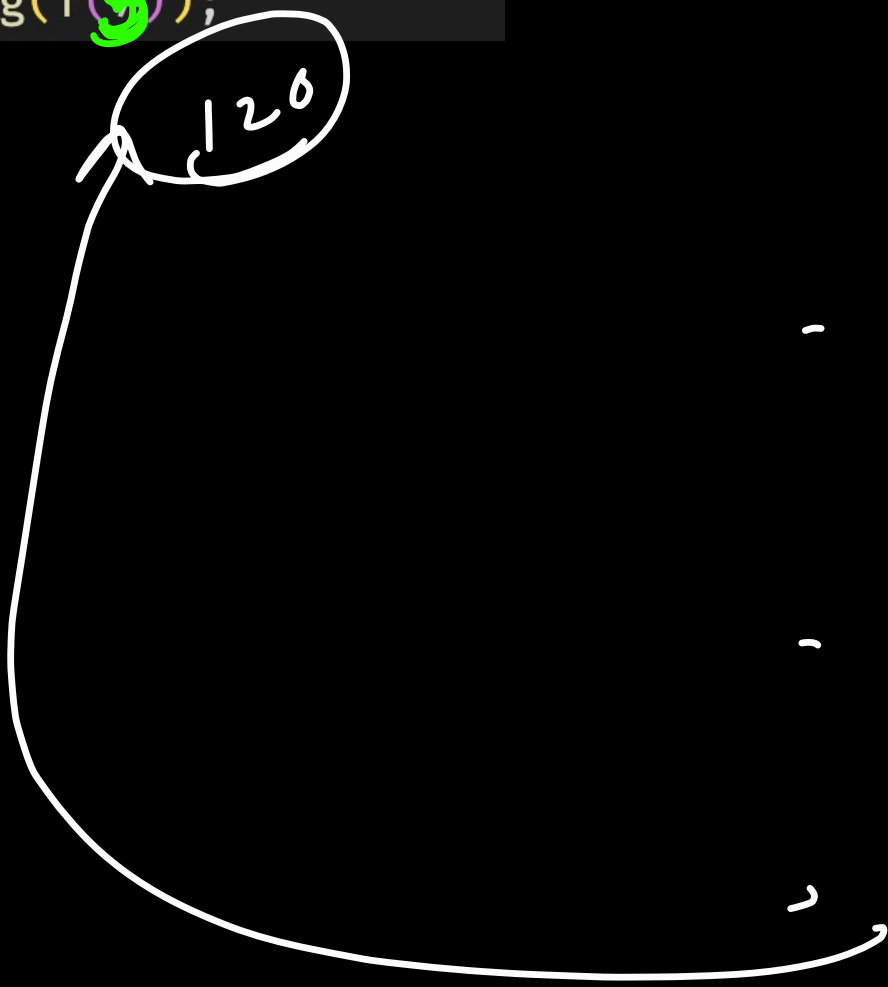
Stacks → linear data structure (mental model to store data in different fashion) in which we can add / remove / get data from the top only.



```
1 function f(n) {  
2     // base case  
3     if(n = 1) {  
4         return 1;  
5     }  
6     return n * f(n-1);  
7 }  
8  
9 console.log(f(5));
```

Sx 24

120



Q Given a value n (+ve integer), calculate the n^{th} fibonacci, recursively.

Ex $\rightarrow n = 5$

ans \rightarrow 5

starter

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...
↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓
0th 1st 2nd 3rd 4th 5th 6th 7th 8th 9th

Ex $\rightarrow n = 6$

ans \rightarrow 8

n

Let's say we have a function f , that takes an argument n , & calculates n^{th} fib.

$$f(n) = f(n-1) + f(n-2)$$

$f(n)$ is the n^{th} fib. (indicated by a double arrow from $f(n)$ to n^{th} fib)

$f(n-1)$ is the $(n-1)^{\text{th}}$ fib. (indicated by a double arrow from $f(n-1)$ to $(n-1)^{\text{th}}$ fib)

$f(n-2)$ is the $(n-2)^{\text{th}}$ fib. (indicated by a double arrow from $f(n-2)$ to $(n-2)^{\text{th}}$ fib)

Assume we get this value correctly from f .

Self work

if ($n == 0$ / $n == 1$)
return n ;

Base Case

if ($n == 1$) return 1;
if ($n == 0$) return 0;

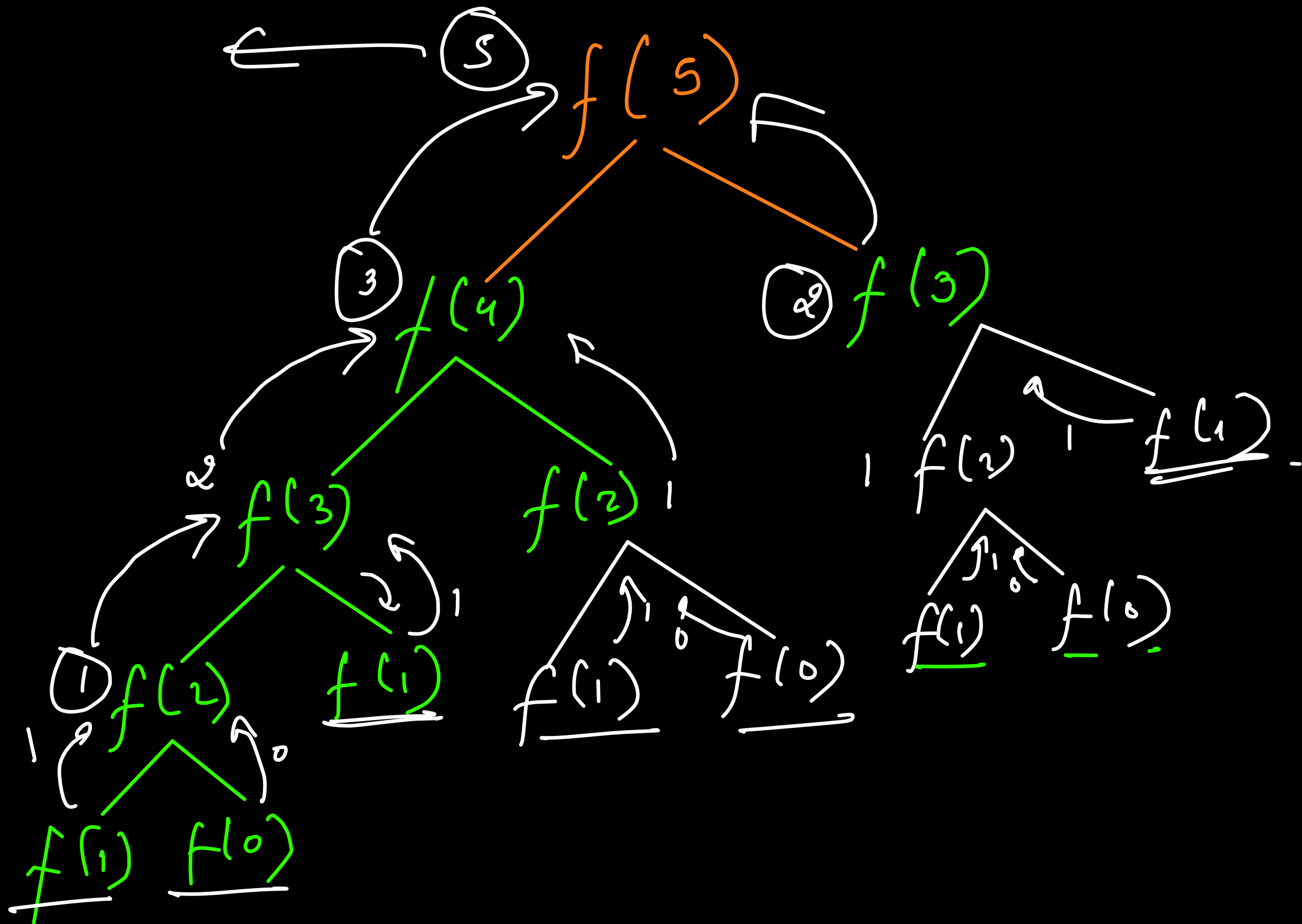
> if ($n == 0$ || $n == 1$)
return n;

assumption, \rightarrow let's assume that the function
works correctly for $f(n-1)$ and $f(n-2)$

$f(n-1) \rightarrow (n-1)^{\text{th}}$ fib

$f(n-2) \rightarrow (n-2)^{\text{th}}$ fib

Self work \rightarrow add $f(n-1)$ & $f(n-2)$



```

1 function f(n) {
2   if(n = 0 || n = 1) return n;
3   return f(n-1) + f(n-2);
4 }
5
6 console.log(f(4));

```

