

Q²⁷ Given an array, write an algorithm to reverse the same array.

Note: You should not create a new array.

Ex \rightarrow [5, 9, 1, 8, 2, 3]

ans \rightarrow [3, 2, 8, 1, 9, 5]

B.f [5, 9, 1, 8, 2, 3]

index \rightarrow 0 1 2 3 4 5

A.f [3, 2, 8, 1, 9, 5] n=6

$i=0 \rightarrow 6-0-1 \rightarrow 5$

$i=1 \rightarrow 6-1-1 \rightarrow 4$

$i=2 \rightarrow 6-2-1 \rightarrow 3$

$i=3 \rightarrow 6-3-1 \rightarrow 2$

Before Reversing \rightarrow element \rightarrow index $\rightarrow i$



After Reversing \rightarrow element \rightarrow index \rightarrow $n-i-1$

We need to somehow move the elements from their original index i to $n-i-1$

0 1 2 3 4 5
 [3, 2, 8, 1, 9, 5]

↑ ↑

0 1 2 3 4 5
 [3, 2, 8, 1, 9, 5]

B.f →

element	index
5	0
3	5
9	1
2	4

Af

element	index
5	5
3	0
9	4
2	1

Both the terminal elements are swapping their

spots.

swap(arr, i, j)

0 1 2 3 4 5 ←
[3, 2, 8, 1, 9, 5]
↑ ↑
j i

```
while (i <= j) {  
    swap(arr, i, j);  
    i += 1;  
    j -= 1;  
}
```

We are tracking
stock price

Amazon

2nd 4th
↓ ↓
[7, 1, 5, 3, 6, 4] → prices
↑ ↑ ↑ ↑
0th 1st 3rd 5th

prices[i] → what is the cost of the
amazon stock on the ith day.

Buy the stock on any one day & sell it later

0 1 2 3 4 5
[7, 1, 5, 3, 6, 4] //

if we
buy
now

sell here

$$(-1 + 6) \rightarrow \underline{\underline{5}}$$

Brute Force

$\Rightarrow [7, 1, 5, 3, 6, 4]$
0 1 2 3 4 5

① buy $\rightarrow 7$
② buy $\rightarrow 1$
③ buy $\rightarrow 5$

\hookrightarrow if we just try to consider all possible

cases-

$$\text{profit} = \$4/5$$

candidate

[7, 1, 5, 3, 6, 4]
0 1 2 3 4 5

$$\text{curr-min} = 7$$

last best price to buy

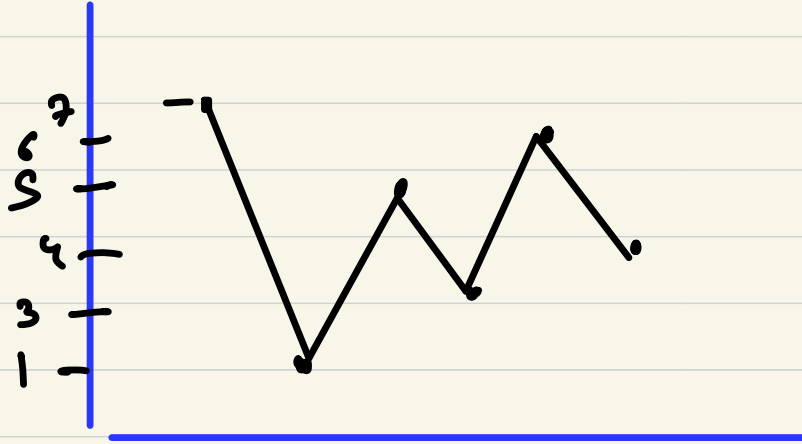
should I sell

→ we should buy as low as possible

→ we should sell as high as possible, after buy.

∴ Start tracking min values.

[7, 1, 5, 3, 6, 4]
0 1 2 3 4 5



Intuition

↓
Small value as
6th to buy.

Profit = ~~7~~

^{0 1 2 3 4 5 6}
[7, 6, 4, 3, 1, 8, 6]

↑
i -

(curr-min = ~~7~~)
↘ 4

Best time to

Buy -

if (price[i] > curr-min) {
 profit = Math.max(profit, price[i] - curr-min);

}
if (price[i] < curr-min) {
 curr-min = price[i];

}

on any i^{th} day, I ask a question, Should I
sell today ??

for selling we need min stock price before the
 i^{th} day.

nums1 [1, 2, 3, 4, 5,]

increasing

$m = 5$

nums2 [2, 3, 8, 9]

increasing

$n = 4$

C = [1, 2, 2, 3, 3, 4, 5, 8, 9]

don't return C

increasing order

nums1 [1, 3, 4, 8, 9] $m=5$

Indices: 0, 1, 2, 3, 4

nums2 [0, 3, 5]

Indices: 0, 1, 2

```

if (nums1[i] < nums2[j])
{
    result[k] = nums1[i];
    k++; i++;
} else {
    result[k] = nums2[j];
    k++; j++;
}

```

inc ↑
 result → [0, 1, 3, 3, 4, 5, 8, 9, ...] $m+n \rightarrow 8$

Indices: 0, 1, 2, 3, 4, 5, 6, 7, 8, ...

↑ k

after we complete the algorithm how many elements will be inside result.

On the 0th index, Smallest among both the arrays will be Present.

nums1 [a, b, c, d] → inc

nums2 [d, e, f] → inc

result → [a,
 0]

```
while (i < m && j < n) {
```

this condition
will be false
when nums2
is exhausted.

```
    if (nums1[i] < nums2[j])  
        { result[k] = nums1[i];  
          k++; i++;
```

```
    } else {
```

```
        result[k] = nums2[j]  
        k++; j++;  
    }
```

```
}
```

this will be true if &
only if we have
elements in nums1

```
while (i < m) {
```

```
    result[k] = nums1[i]  
    i++; k++
```

```
}
```

```
while (j < n) {
```

```
    result[k] = nums2[j]  
    k++; j++
```

```
}
```

if nums2 still
has elements

