

DFS (Depth First Search)

Derinlik Öncelikli Arama

DFS (G, s)

mark(s);

$L = \{s\};$

while $L \neq \emptyset$

$u = \text{last}(L);$

→ there exists

if $\exists (u, v)$ such that v is unmarked

choose (u, v) with v of smallest index;

→ E

mark(v);

$L = L \cup \{v\};$ // push(L, v)

else

$L = L \setminus \{u\};$ // pop(L)

DFS de L isin stack (Last in First out)
kullanilir.

L (stack)

Marked

2
 2, 1 ^{→top}
 2, 1, 4
 2, 1, 4, 3 ^{→top}
 2, 1, 4, 3, 7
 2, 1, 4, 3
 2, 1, 4, 3, 8
 2, 1, 4, 3
 2, 1, 4
 2, 1, 4, 6
 2, 1, 4, 6, 5
 2, 1, 4, 6
 2, 1, 4
 2, 1
 2
 \emptyset

$s = 2$

1

4

3

7

—

8

—

—

6

5

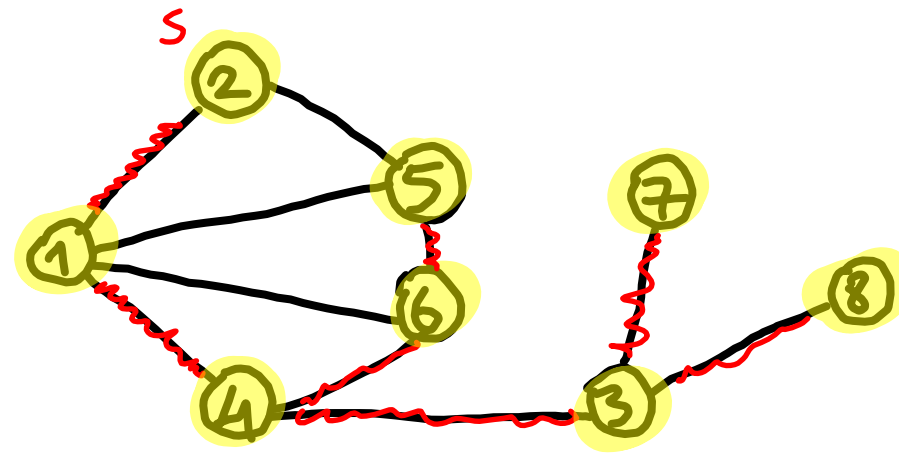
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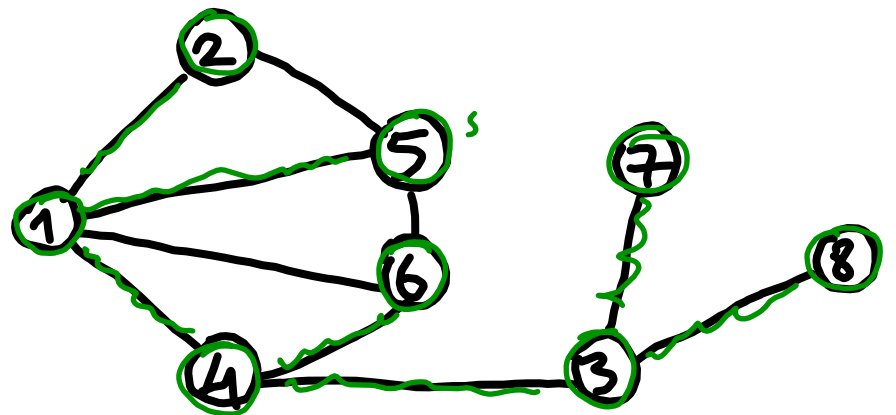
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while $L \neq \emptyset$ ^{→top}
 $u = \text{last}(L);$
 if $\exists (u, v)$ such that v is unmarked
 choose (u, v) with v of smallest index
 mark $(v);$
 $L = L \cup \{v\};$ // push (L, v)
 else
 $L = L \setminus \{u\};$ // pop (L)

Çalışma Zamanı:

$$T(n, |E|) = \Theta(E + V)$$



BFS (Breadth First Search)

Genişlik Öncelikli Arama

$BFS(G, s)$

mark(s);

$L = \{s\};$ // L kuyruktur.

while $L \neq \emptyset$

$u = \text{first}(L);$

 if $\exists (u, v)$ such that v is unmarked

 choose (u, v) with v of smallest index;

 mark(v);

$L = L \cup \{v\};$ // enqueue(L, v)

 else

$L = L \setminus \{u\};$ // dequeue(L)

L için kuyruk veri yapısı (First in first out)
kullanılır.

L (kuyruk)

Marked

front → 2 rear →

2, 1

2, 1, 5

1, 5

1, 5, 4

1, 5, 4, 6

5, 4, 6

4, 6

4, 6, 3

6, 3

3

3, 7

3, 7, 8

7, 8

8

∅

s = 2

1

5

—

4

6

—

—

3

—

—

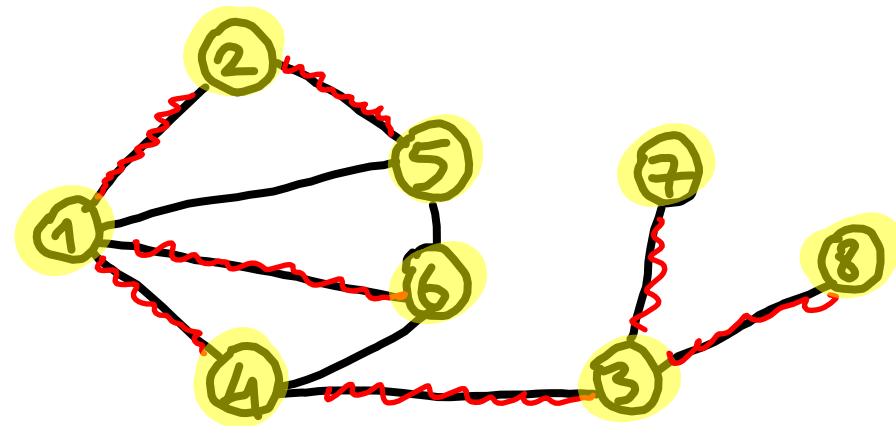
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8

—

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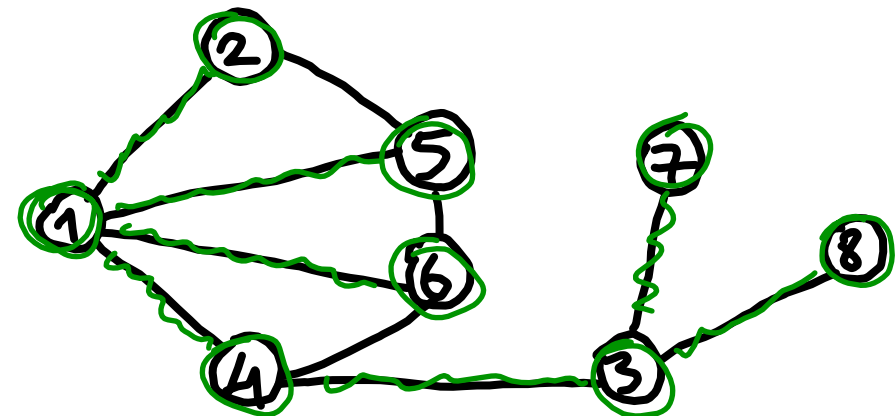
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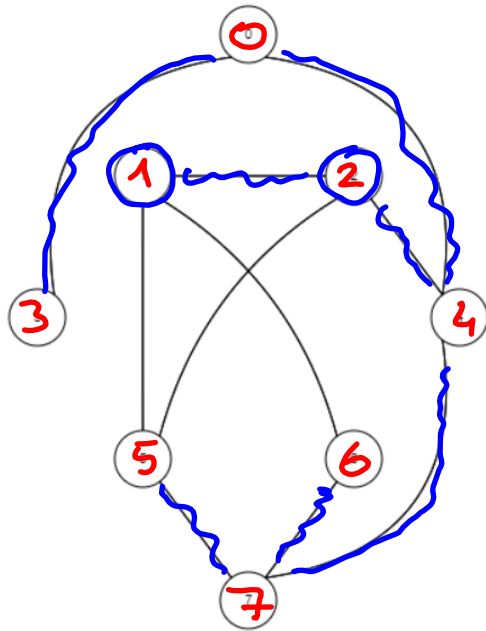
while $L \neq \emptyset$
 $u = \text{first}(L);$
 if $\exists (u, v)$ such that v is unmarked
 choose (u, v) with v of smallest index;
 mark $(v);$
 $L = L \cup \{v\};$ // enqueue (L, v)
 else
 $L = L \setminus \{u\};$ // dequeue (L)

Görüşme Zamanı:

$$T(|V|, |E|) = \Theta(E + V)$$



DFS: stack



BFS: queue

