

Part1

Question1

a)

Algorithm swapValueM(A, m)

Input A: array of n integers, m: an integer

Output: A

```
for i <- 0 to n-2 do
  if A[i] == m then
    A[i] = A[i] + A[i+1]
    A[i+1] = A[i] - A[i+1]
    A[i] = A[i] - A[i+1]
return A
```

b) Time complexity $O(N)$

c) Space complexity $O(1)$

Question2

Algorithm swapString(S)

Input: S, a string with n characters

Output: S_res, a string

n <- length of S

consonant <- string

repeat <- string

vowel <- string

for i <- 0 to n - 1 do

 if S[i] is consonant:

 repeated <- false

 for j <- i + 1 to n - 1 do

 if S[j] == S[i] do

 repeat += S[i]

 repeated = true

 break

 if repeated == true:

 repeat += S[i]

 else:

 consonant += S[i]

 else if S[i] is vowel:

 vowel += S[i]

return consonant+repeat+vowel

b) time complexity $O(N^2)$

c) space complexity $O(N)$

Question3

Algorithm tetradicNumber(Numbers):

Input: Numbers, an array of integers

Output: success1, success2, farthest1, farthest2

success1 <- integer, first successive number with largest difference

success2 <- integer, second successive number with largest difference

largest_difference <- integer

farthest1 <- integer, first number of two farthest numbers whose difference is ten

farthest2 <- integer, second number of two farthest numbers whose difference is ten

farthest <- integer

n <- length of array Numbers

for i <- 0 to n - 2:

if absolute value of Numbers[i] - Numbers[i+1] > largest_difference:

largest_difference = Numbers[i] - Numbers[i+1]

success1 = i

success2 = i+1

for i <- 0 to n-1:

for j <- i+1 to n-1:

if absolute value of Numbers[j] - Numbers[i] == 10 && j - i > farthest:

farthest = j - i

farthest1 = i

farthest2 = j

return [success1, success2, farthest1, farthest2]

b) This is the brute-force solution based on iterations to find out the two groups, first iteration finds the two successive number with the largest difference, and the second iteration finds the two farthest numbers with their difference being 10

iii) time complexity is the $O(N^2)$. The first loop takes the $O(N)$, the second loop takes the $O(N^2)$. Therefore the overall runtime is $O(N^2)$

iv) The stack growth is constant since this algorithm is not recursive and uses a constant number of variables, the stack usage is constant, and does not increase with respect to the input size.