## McGill University ECSE 425 COMPUTER ORGANIZATION AND ARCHITECTURE Fall 2009 Midterm Examination

3:30 PM - 4:45 PM, October 16<sup>th</sup>, 2009

**Duration: 75 minutes** 

- There are 4 questions for a total of 100 points. There are 14 pages. Please check that you have all 14 pages.
- This is a closed-book exam. You can bring one single-sided sheet of hand-written notes. This sheet of notes must be entirely hand-written, no portions may be machine-produced or photocopied.
- Calculators are permitted, but no cell phones or laptops are allowed.
- Write your name and student number in the space below. Do the same on the top of each sheet of this exam.
- State any assumption.

Total:

• Write your answers in the space provided.

Name:		 
Student Number:		
Q1:	Q3:	 
Q2:	Q4:	 

Name:	ID:
Question 1. Short Answers (20 points). There are 10 sub questions (2 points each).	
For each question below, provide a short answer in	1-2 sentences.
(a) What factors contribute to the pipelining overh	nead?
(b) What is the impact of pipelining overhead?	
(c) Suppose that the CPU and the memory are two CPU has a MTTF of a, and the memory has a independently of the other. What is the overal	MTTF of b. Assume that each fails
· ·	·
(d) Why using ratio is a good way to compare cor	mputer performance?
(e) What is the difference between true depender	ncy and name dependency?
(2) Triat is the amerenee between true depender	io, and hame dependency:

Name:	ID:
For the next 3 sub	-questions, use the following code segment:
ADD.D S.D SUB.D	F0, F2, F4 F6, F0, F8 F6, 0(R1) F8, F10, F14 F6, F10, F8
Identify one of each o operands)	f the following hazards (name the instructions and the dependent
(f) RAW hazard:	
<b>(g)</b> WAR hazard:	
(h) WAW hazard:	
(i) In order to pres	serve data flow, what kinds of dependency need to be maintained?
(j) In precise exce are they handle	eption, are the exceptions handled as they occurred? If not, in what ordered?

Name: ID:
Question 2. Performance (20 points): There are two sub questions
(a) (10 points) Assume a four-stage pipeline where the branches are resolved at the end of the second cycle for unconditional branches and at the end of the third cycle for conditional branches. Suppose that 20% of all instructions are conditional branches (60% are taken) and 5% are unconditional branches or procedure calls. Ignore all other pipeline stalls.
1. What is the CPI for this computer if we don't use any branch prediction?
2. What is the CPI for this computer if the compiler predicts all branches as taken?

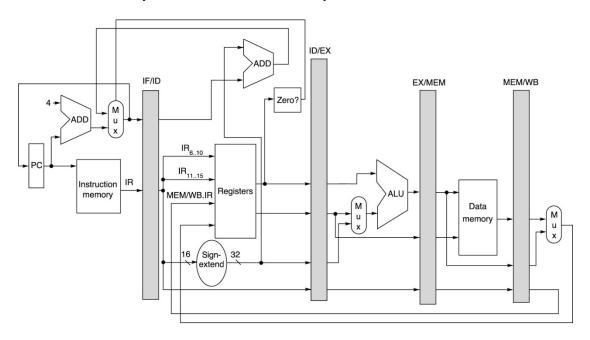
Na	me	ne: ID:	
(b)	sin wil ap	<b>10 points)</b> Your company has just bought a new dual single core one, and you are tasked with optimizing you will run two applications on this dual Pentium with une application needs 80% of the resources, and the othe hat 40% of the first application is parallelizable while	our software for this processor. You equal resource requirements. The first r only 20% of the resources. Assume
	1.	I. How much speedup would you achieve if you only	ran the first application?
	2.	2. How much overall system speedup would you obs	erve when running both applications?

## Question 3. Integer Pipelining (30 points). There are two parts to this question.

(a) (12 points) Consider the following code fragment which implements the integer code x[i] = 2x[i]+1;

Loop:	LD	R2,0(R1)
	DMUL	R2,R2,#2
	DADDI	R2,R2,#1
	SD	R2,0(R1)
	DADDI	R1,R1,#8
	SUB	R4,R3,R1
	BNEZ	R4,Loop

The code will run on the classic 5-stage (I, D, E, M, W) RISC pipeline shown in the figure below. Assume all memory accesses take 1 clock cycle.



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Fill in the chart on the next page to show the timing of this instruction sequence **without** any forwarding or bypassing hardware but assuming a register read and a register write in the same clock cycle. Assume that the branch is handled by flushing the pipeline and branches have a one-cycle delay.

Loop: LD R2,0(R1)

DMUL R2,R2,#2

DADDI R2,R2,#1

SD R2,0(R1)

DADDI R1,R1,#8

DADDI R1,R1,#8 SUB R4,R3,R1 BNEZ R4,Loop

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
LD	I	D	Ε	M	W														

Assume that the initial value of R3 is R1+72. How many clock cycle does the whole loop take to execute?

(b) (18 points) Consider the same code fragment:

Loop:	LD	R2,0(R1)
	DMUL	R2,R2,#2
	DADDI	R2,R2,#1
	SD	R2,0(R1)
	DADDI	R1,R1,#8
	SUB	R4,R3,R1
	BNEZ	R4,Loop

**Assume the branch is handled by a one-cycle delayed branch.** Show the scheduled code with delayed branch.

Assume now that the pipeline has full hazard detection and forwarding hardware (the forwarding paths are not shown). Assume a register read and a register write in the same clock cycle "forward" through the register file. Assume all memory accesses take 1 clock cycle.

Show the timing of the **scheduled code fragment with delayed branch** (derived above) for the RISC pipeline. Fill in the chart below to show the timing. Add arrows to indicate any forwarding. The first line is filled in for you.

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
LD	I	D	Ε	M	W														

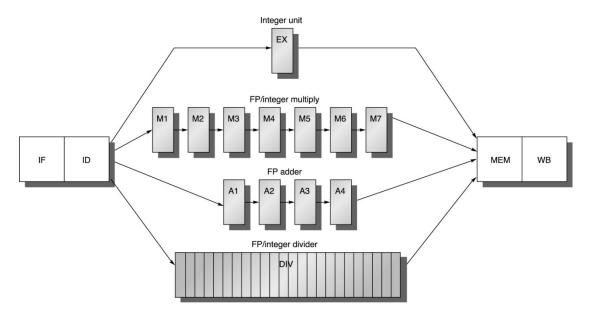
Assume the initial value of R3 as R1+72, how many clock cycles does this loop take to execute?

## Question 4. FP Pipelining and Loop Unrolling (30 points). There are 2 sub questions.

Consider the following loop which computes  $Y[i] = a \times X[i] + Y[i]$ , the key step in Gaussian elimination.

loop:	L.D	F0, 0(R1)
-	MULT.D	F0, F0, F2
	L.D	F4, 0(R2)
	ADD.D	F0, F0, F4
	S.D	F0, 0(R2)
	DADDUI	R1, R1, #-8
	DADDUI	R2, R2, #-8
	BNEZ	R1, loop

The code will run on the standard MIPS FP pipeline shown in the figure below.



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(a) (15 points) Assume that the pipeline has full hazard detection and forwarding hardware. Assume a register read and a write in the same clock cycle "forward" through the register file. Assume all memory accesses take 1 clock cycle. Assume that the branch is handled by flushing the pipeline. Fill in the chart on the next page to show the timing of 1 loop iteration.

ID:

How many clock cycles does a single iteration of this loop take?

Name:

## (b) (15 points) Here is the code fragment again

loop:	L.D	F0, 0(R1)
	MULT.D	F0, F0, F2
	L.D	F4, 0(R2)
	ADD.D	F0, F0, F4
	S.D	F0, 0(R2)
	DADDUI	R1, R1, #-8
	DADDUI	R2, R2, #-8
	BNEZ	R1, loop

Assume the pipeline latencies and initiation intervals as in Table 1. Unroll the loop as many times as necessary to schedule it without any delays. Assume the branch is now handled by a one-cycle delayed branch. Show the unrolled loop. How many clock cycles does it take per original loop iteration?

Functional unit	Latency	Initiation interval
Integer ALU	0	1
FP add	3	1
FP multiply	6	1
FP divide	24	25

Table 1. Latency and Initiation Interval of FP unit

Name:	ID:

Additional page for Q4b.

Name:	ID:

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