- 1. We extend the orthogonal transforms in QR decomposition algorithms to complex matrices.
  - (a) Find the complex Householder matrix. Given a nonzero  $v \in \mathbb{C}^n$ , find a matrix  $H \in \mathbb{C}^{n \times n}$  such that  $HH^* = H^*H = I$  and Hv is a multiple of  $e_1$ , where the superscript \* stands for conjugate transpose.
  - (b) Find the complex Givens rotation matrix. More precisely, given  $\boldsymbol{x} \in \mathbb{C}^2$ , find a matrix  $\boldsymbol{G} \in \mathbb{C}^{2 \times 2}$  such that  $\boldsymbol{G}\boldsymbol{G}^* = \boldsymbol{G}^*\boldsymbol{G} = \boldsymbol{I}$  and  $\boldsymbol{G}\boldsymbol{x} = \begin{bmatrix} \times \\ 0 \end{bmatrix}$ , where the superscript \* stands for conjugate transpose and  $\times$  is a complex number.

(A) For 
$$V = (V_1, V_2, ..., V_n)^T \in \mathbb{C}^n$$
, let  $Hu = -e^{-i\theta}(I - 2uu^*)$ .

here u is unit of  $V + e^{i\theta} \|V\|_2 e_1$ ,  $e^{i\theta} = \frac{V_1}{\|V_1\|}$ .

Then

①  $Hu Hu^* = (I - 2uu^*)(I - 2uu^*)$ 

=  $I - 4uu^* + 4u(u^*u)u^*$  (by u is unit)

=  $I - 4uu^* + 4uu^*$ 

=  $I$ 

and also  $Hu^* + 4u = I$ 

hence 
$$\text{Huv} = -e^{-i\theta} \text{U} + e^{-i\theta} \text{U} + \text{IIVI}_2 e_1 = \text{IIV}_2 e_1$$
  
(b) Let  $G = \begin{pmatrix} c & s \\ -\bar{s} & \bar{c} \end{pmatrix}_1$  then  $G$  satisfies  $G = G^* = G^* G$ .
$$G = \begin{pmatrix} c & s \\ -\bar{s} & \bar{c} \end{pmatrix} \begin{pmatrix} \bar{c} & -s \\ \bar{s} & c \end{pmatrix} = \begin{pmatrix} c\bar{c} + s\bar{s} & 0 \\ 0 & s\bar{s} + c\bar{c} \end{pmatrix} = I.$$

We have |c|2 + 15|2 =1.

For 
$$X=(X_1,X_2)^T \in \mathbb{C}^2$$
,  $GX=\begin{pmatrix} cX_1+cX_2\\ -\overline{c}X_1+\overline{c}X_2 \end{pmatrix} = \begin{pmatrix} X\\ 0 \end{pmatrix}$ ,

then 
$$-\bar{s} x_1 + \bar{c} x_2 = 0$$
. We choose  $c = \frac{\bar{x}_1}{||x||_2}$ ,  $s = \frac{\bar{x}_2}{||x||_2}$  and  $G = \begin{pmatrix} \bar{x}_1/_{(|x||_2)} & \bar{x}_2/_{(|x||_2)} \\ - & \bar{x}_1/_{(|x||_2)} & \bar{x}_1/_{(|x||_2)} \end{pmatrix}$ .

- 2. Consider the least squares problem  $\min_{\boldsymbol{x} \in \mathbb{R}^n} \|\boldsymbol{A}\boldsymbol{x} \boldsymbol{b}\|_2^2$ , where  $\boldsymbol{A} \in \mathbb{R}^{m \times n}$  with m > n is of full column rank.
  - (a) Prove: x is the solution of the least squares problem if and only if there exists a vector  $y \in \mathbb{R}^m$  such that

$$\begin{bmatrix} \boldsymbol{I} & \boldsymbol{A} \\ \boldsymbol{A}^T & \boldsymbol{0} \end{bmatrix} \begin{bmatrix} \boldsymbol{y} \\ \boldsymbol{x} \end{bmatrix} = \begin{bmatrix} \boldsymbol{b} \\ \boldsymbol{0} \end{bmatrix}$$
 (1)

This provides another way to solve the least squares problem.

- (b) Prove that the coefficient matrix in (1) is non-singular and symmetric indefinite. (A matrix is indefinite if the associated quadratic form can be either positive or negative.)
- (c) Prove that, if A is properly scaled such that the minimum eigenvalue of  $A^TA$  is 1, then the condition number of the coefficient matrix in (1) is upper bounded by  $C \cdot \sqrt{\text{Cond}(A^TA)}$  for some universal constant C > 0, where  $\text{Cond}(\cdot)$  is the condition number. Therefore, (1) is better to solve than the normal equation, though it is indefinite.

(a) Since 
$$A \in \mathbb{R}^{m \times n}$$
 and  $m > n$ , Mull  $(A^T)$  is nontrivial.

$$\Leftrightarrow \exists y \in \mathbb{R}^m$$
 s.t.  $b-Ax=y$  and  $y \in Ran(A)^{\perp} = Null(A^{\top})$ 

$$\Leftrightarrow$$
  $\begin{pmatrix} I & A \\ A^{\mathsf{T}} & O \end{pmatrix} \begin{pmatrix} \mathsf{y} \\ \mathsf{x} \end{pmatrix} = \begin{pmatrix} \mathsf{b} \\ \mathsf{o} \end{pmatrix}$ 

(b) 
$$\bigcirc$$
 Since  $\begin{pmatrix} I & O \\ -A^T & I \end{pmatrix} \begin{pmatrix} I & A \\ A^T & O \end{pmatrix} \begin{pmatrix} I & -A \\ O & I \end{pmatrix} = \begin{pmatrix} I & o \\ o & -A^T A \end{pmatrix}$ ,

$$det \begin{pmatrix} I & A \\ A^T & O \end{pmatrix} = det (-A^T A) \neq 0$$
, as A is of tull

column rank. So 
$$\begin{pmatrix} I & A \\ A^T & D \end{pmatrix}$$
 is non-singular.

If we pick 
$$y \in \text{Null}(A^T)$$
 and  $y \neq 0$ ,  
then  $y^T y + 2y^T Ax = \|y\|^2 > 0$ .

If we pick x such that 
$$Ax \neq 0$$
 and  $y = -Ax$ ,  
then  $y^Ty + 2y^TAx = -(Ax)^T(Ax) < 0$ .

Hence, 
$$\begin{pmatrix} I & A \\ A^T & o \end{pmatrix}$$
 is symmetric indefinite.

Let 
$$A = U \Sigma V^*$$
 the SVD of  $A$ .

U and  $V$  are unitary,  $\Sigma = \begin{pmatrix} \sigma_{i_1} \sigma_{i_2} \\ \sigma_{i_3} \end{pmatrix}_{men}$ ,  $\sigma_i \circ \sigma_i \circ$ 

 $= \frac{\sqrt{2} + 1}{\sqrt{2} + 1} \sqrt{2} \cos \left( \sqrt{2} \right)$ 

3. The QR algorithm can also be used in theoretical analysis. Let  $\mathbf{A} = [\mathbf{a}_1 \ \mathbf{a}_2 \ \dots \ \mathbf{a}_n] \in \mathbb{R}^{n \times n}$  with column vectors  $\boldsymbol{a}_i$  for  $i = 1, \dots, n$ . Prove

$$|\det(\boldsymbol{A})| \leq \prod_{i=1}^n \|\boldsymbol{a}_i\|_2,$$

and the upper bound is attained if and only if  $A^TA$  is diagonal or  $||a_i||_2 = 0$  for some i. (Hint: Consider the Gram-Schmidt procedure.)

upper bound is attained if and only if 
$$A^*A$$
 is diagonal or  $||a_i||_2 = 0$  for some  $i$ . (Hint: the Gram-Schmidt procedure.)

Consider  $QR$  via  $Gram-Schmidt$  procedure.  $A=QR$  where  $Q=(Q_1,Q_2,\cdots Q_n)$  is orthogonal and  $R$  is upper triangular.

$$||det A|| = ||det Q|| ||det R|| = \prod_{i=1}^n |r_{ii}|$$

$$= \prod_{i=1}^{n} | q_{i}^{T} \alpha_{i} |$$

$$\leq \prod_{i=1}^{n} ||q_{i}||_{2} ||\alpha_{i}||_{2}$$

$$= \ \frac{1}{|I|} \| \alpha_i \|_2$$

the upper bound is attained

$$\iff \overline{\prod}_{i=1}^{n} | q_i^{\mathsf{T}} \alpha_i | = \overline{\prod}_{i=1}^{n} | q_i | |_2 | | | | |\alpha_i | |_2$$
 (\*)

$$\Leftrightarrow$$
  $\alpha_i = \alpha_i q_i$   $i=1,2,...,n$ 

$$\iff A = BD \quad D = \begin{pmatrix} a_1 \\ a_2 \end{pmatrix}$$

4. Let the initial vector  $x_0 = 0$ . When the GMRES is applied to solve Ax = b, where

$$m{A} = egin{bmatrix} 0 & & & & 1 \ 1 & 0 & & & \ & 1 & 0 & & \ & & \ddots & \ddots & \ & & & 1 & 0 \ \end{pmatrix}, \qquad m{b} = egin{bmatrix} 1 \ 0 \ dots \ 0 \ \end{bmatrix},$$

what is the convergence rate?

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$$X_0 = 0$$
,  $r_0 = b - Ax_0 = b$ 

$$\begin{cases}
k = span \{ r_0, Ar_0, ..., A^{k-1}r_0 \} \\
= span \{ e_1, Ae_1, ..., A^{k-1}e_1 \} \\
= span \{ e_1, e_2, ..., e_k \}
\end{cases}$$

$$\begin{cases}
x_k = argmin || Ax - b||_2 \\
x \in X_0 + X_k
\end{cases}$$

$$|| \Gamma_k ||_2 = \min_{y \in R^k} || (b - Ax_0) - \frac{k}{2} y_j A^j r_0 ||_2
\end{cases}$$

$$= \min_{y \in R^k} || (\frac{b}{b} - Ax_0) - \frac{k}{2} y_j A^j r_0 ||_2$$

For k=1,2,...,n-1,  $||\Gamma_k||_2=1$  and when k=n,  $||\Gamma_n||_2=0$ , then the convergence rate is  $\frac{||\Gamma_{k-1}||_2}{||\Gamma_{k-1}||_2}=1$ , and it will converge

at k=n.

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