Definitions of Quantum Gates

This notebook is based on

- M. Nielsen and I. Chuang, Quantum Computation and Quantum Information, Cambridge University Press, 2010 (tenth anniversary edition)
- **2.** Bittner, Eric R. Quantum Dynamics: Applications in Biological and Materials Systems. CRC Press, 2009.
- **3.** arXiv:1403.7050 [quant-ph]

Single-qubit pseudo spin operators: $\mathbf{1}$, $\hat{\sigma}_x$, $\hat{\sigma}_y$, $\hat{\sigma}_z$

In[*]:= {id, σx, σy, σz} = Table[SparseArray[PauliMatrix[i]], {i, 0, 3}];

Projectors onto $|0\rangle$ and $|1\rangle$ states

$$ln[\theta]:= proj0 = \frac{id + \sigma z}{2}$$
 // SparseArray;
 $proj1 = \frac{id - \sigma z}{2}$ // SparseArray;

Operator \hat{a} acting on the k^{th} qubit of a set of n qubits:

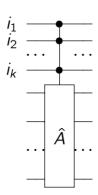
$$1 \otimes 1 \otimes \cdots \otimes 1 \otimes \hat{a} \otimes 1 \otimes \cdots \otimes 1 \otimes 1$$

 $ln[*]:= op[n_Integer, k_Integer, a_] /; 1 \le k \le n \land Dimensions[a] == \{2, 2\} := KroneckerProduct[IdentityMatrix[2^{k-1}, SparseArray], a, IdentityMatrix[2^{n-k}, SparseArray]]$

Controlled operator

In a set of n qubits, the n-qubit operator $\mathtt{ctrl}[n,\lambda,A]$ acts like the operator \hat{A} if all qubits in the list $\lambda = \{i_1, i_2, \ldots, i_k\}$ are in the

 $|1\rangle$ state, and has no action (acts like the identity operator, 1 on n qubits) if any of the qubits in the list λ are in the $|0\rangle$ state:



$$\mathsf{CTRL} = \left[\bigotimes_{j=1}^k |1\rangle\langle 1|^{(i_j)} \right] \cdot \hat{A} + \left[\mathbb{1} - \bigotimes_{j=1}^k |1\rangle\langle 1|^{(i_j)} \right] \cdot \mathbb{1} = \mathbb{1} + \left[\bigotimes_{j=1}^k |1\rangle\langle 1|^{(i_j)} \right] \cdot (\hat{A} - \mathbb{1})$$

In[@]:= ctrl[n_Integer, λ_ /; VectorQ[λ, IntegerQ], A_] /; $\left(\mathsf{Unequal} @@\ \lambda \right) \land \mathsf{Min}[\lambda] \ge 1 \land \mathsf{Max}[\lambda] \le \mathsf{n} \land \mathsf{Dimensions}[\mathsf{A}] == \left\{ 2^\mathsf{n},\ 2^\mathsf{n} \right\} := \left\{ 2^\mathsf{n},\$ IdentityMatrix[2ⁿ, SparseArray] + Apply[Dot, op[n, #, proj1] & $/@\lambda$]. (A - IdentityMatrix[2^n , SparseArray])

Single-Qubit Gates

Pauli gates

Pauli gates acting on the k^{th} qubit of a set of n qubits:

- The Pauli-X gate _____ acts like $\hat{\sigma}_x$ on the desired qubit, and has no effect on all other qubits.
- $In[\cdot]:= xGate[n_Integer, k_Integer] /; 1 \le k \le n := op[n, k, \sigma x]$
 - The Pauli-Y gate $\overline{\gamma}$ acts like $\hat{\sigma}_y$ on the desired qubit, and has no effect on all other qubits.
- $ln[\cdot]:= yGate[n_Integer, k_Integer] /; 1 \le k \le n := op[n, k, \sigma y]$
 - The Pauli-Z gate Z acts like $\hat{\sigma}_z$ on the desired qubit, and has no effect on all other qubits.
- $ln[\cdot]:= zGate[n_Integer, k_Integer] /; 1 \le k \le n := op[n, k, \sigma z]$

Pauli rotations

$$\hat{R}_{x}(\boldsymbol{\phi}) \propto \boldsymbol{e}^{-i\boldsymbol{\phi}\,\hat{\sigma}_{x}/2}$$

$$ln[*]:= rxGate[n_Integer, k_Integer, \phi_] /; 1 \le k \le n := op[n, k, \frac{1 + e^{i \phi}}{2} id + \frac{1 - e^{i \phi}}{2} \sigma x]$$

$$\hat{R}_{v}(\boldsymbol{\phi}) \propto \boldsymbol{e}^{-i \, \boldsymbol{\phi} \, \hat{\sigma}_{v}/2}$$

$$ln[*]:= ryGate[n_Integer, k_Integer, \phi_] /; 1 \le k \le n := op[n, k, \frac{1 + e^{i \phi}}{2} id + \frac{1 - e^{i \phi}}{2} \sigma y]$$

$$\hat{R}_{z}(\phi) \propto e^{-i\phi \hat{\sigma}_{z}/2}$$

$$ln[\cdot]:=$$
 rzGate[n_Integer, k_Integer, ϕ _] /; $1 \le k \le n := op[n, k, \frac{1 + e^{i \cdot \phi}}{2} id + \frac{1 - e^{i \cdot \phi}}{2} \sigma z]$

Hadamard gate

The Hadamard gate H acts like $\frac{\hat{\sigma}_x + \hat{\sigma}_z}{\sqrt{2}}$ on the desired qubit, and has no effect on all other qubits.

$$\textit{ln[*]}:= \text{hadamardGate[n_Integer, k_Integer] /; } 1 \leq k \leq n := op[n, k, \frac{\sigma x + \sigma z}{\sqrt{2}}]$$

Two-qubit gates

SWAP gate

Exchanges the state of qubits *j* and *k* in a set of *n* qubits

$$\begin{array}{ccc}
j & \xrightarrow{} \\
k & \xrightarrow{} \\
SWAP^{(jk)} = (\mathbb{1}^{(j)} \otimes \mathbb{1}^{(k)} + \hat{\sigma}_x^{(j)} \otimes \hat{\sigma}_x^{(k)} + \hat{\sigma}_y^{(j)} \otimes \hat{\sigma}_y^{(k)} + \hat{\sigma}_z^{(j)} \otimes \hat{\sigma}_z^{(k)})
\end{array}$$

The matrix representation of a two - qubit SWAP takes on the familiar form

In[@]:= swapGate[2, {1, 2}] // MatrixForm

Out[• 1//MatrixForm=

Square root of the SWAP gate

In[*]:= sqrtSwapGate[2, {1, 2}] // MatrixForm

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \frac{1}{2} + \frac{\dot{1}}{2} & \frac{1}{2} - \frac{\dot{1}}{2} & 0 \\ 0 & \frac{1}{2} - \frac{\dot{1}}{2} & \frac{1}{2} + \frac{\dot{1}}{2} & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

CNOT gate

In a set of n qubits, if qubit j is in state $|0\rangle$ then there is no effect on qubit k, whereas if qubit j is in state $|1\rangle$ then the NOT operator $\hat{\sigma}_x$ acts on qubit k.

$$\frac{j}{k} \longrightarrow \text{CNOT}^{(jk)} = |0\rangle\langle 0|^{(j)} \otimes \mathbb{1}^{(k)} + |1\rangle\langle 1|^{(j)} \otimes \sigma_x^{(k)}$$

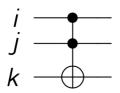
CNOT is simply the CTRL operator with a single element in the list $\lambda = \{j\}$ and a single-qubit $\hat{A} = \hat{\sigma}_x^{(k)}$ operator.

```
ln[e]:= cnotGate[n_Integer, j_Integer \rightarrow k_Integer] /; 1 \le j \le n \land 1 \le k \le n \land j \ne k :=
         ctrl[n, {j}, op[n, k, σx]]
  In[⊕]:= MatrixForm@cnotGate[2, 1 → 2]
Out[ • ]//MatrixForm=
         1 0 0 0
         0 1 0 0
         0 0 0 1
```

Three-qubit gates

CCNOT (Toffoli) gate

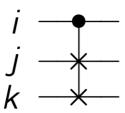
In a set of n qubits, if both qubits i and j are in state 1, then the NOT operator $\hat{\sigma}_x$ acts on qubit



```
ln[\cdot]:= ccnotGate[n_Integer, {i_Integer, j_Integer}] \rightarrow k_Integer] /;
         1 \le i \le n \land 1 \le j \le n \land 1 \le k \le n \land Unequal[i, j, k] := ctrl[n, \{i, j\}, op[n, k, \sigmax]]
```

Controlled SWAP (Fredkin) gate

In a set of *n* qubits, if qubit *i* is in state 0 then there is no effect on qubits *j* and *k*, whereas if qubit *i* is in state 1 then the state of qubits *j* and *k* is exchanged,



```
ln[⊕]:= cswapGate[n_Integer, i_Integer → {j_Integer, k_Integer}] /;
        1 \le i \le n \land 1 \le j \le n \land 1 \le k \le n \land Unequal[i, j, k] := ctrl[n, \{i\}, swapGate[n, \{j, k\}]]
```

A Simple Circuit

As a simple example, we study the quantum circuit,

qubit 1:
$$|0\rangle$$
 H qubit 2: $|0\rangle$

The operator for the circuit,

$$ln[*]:=$$
 s = cnotGate[2, 1 \rightarrow 2].hadamardGate[2, 1];
MatrixForm[s]

$$\begin{pmatrix} \frac{1}{\sqrt{2}} & 0 & \frac{1}{\sqrt{2}} & 0 \\ 0 & \frac{1}{\sqrt{2}} & 0 & \frac{1}{\sqrt{2}} \\ 0 & \frac{1}{\sqrt{2}} & 0 & -\frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & 0 & -\frac{1}{\sqrt{2}} & 0 \end{pmatrix}$$

the basis set is $\mathcal{B}_2 = \{ \mid 00 \rangle, \mid 01 \rangle, \mid 10 \rangle, \mid 11 \rangle \}$

 $ln[\cdot]:= bas[n_Integer/; n \ge 1] := Tuples[{0, 1}, n]$ bas[2]

Out[\circ]= {{0,0}, {0,1}, {1,0}, {1,1}}

operate on the $|\psi_{in}\rangle = |0\rangle \otimes |0\rangle = |00\rangle$ input state:

 $ln[\bullet] := \psi in = \{1, 0, 0, 0\};$ ψ out = s. ψ in; MatrixForm@ψout

Out[@]//MatrixForm

$$\begin{pmatrix} \frac{1}{\sqrt{2}} \\ 0 \\ 0 \\ \frac{1}{\sqrt{2}} \end{pmatrix}$$

 \mathcal{B}_2 shows that we have $|\psi_{\text{out}}\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$, and projective measurements on the two qubits, give $\frac{1}{2}$ probability of finding the classical result $|00\rangle$ and $\frac{1}{2}$ probability of finding $|11\rangle$, whereas the bit combinations $|10\rangle$ and $|01\rangle$ never occur

In[*]:= Abs[
$$\psi$$
out]²
Out[*]= $\left\{\frac{1}{2}, 0, 0, \frac{1}{2}\right\}$

Quantum Fourier Transform

The quantum circuit for a QFT is assembled from blocks that connect the *i*th qubit to the qubits $\{i+1, i+2, ..., n\}$ via controlled phase gates:

$$\begin{array}{ll} \textit{In[e]} = & \text{qftBlock[n_Integer, i_Integer] /; } 1 \leq i \leq n := \\ & \text{Apply[Dot, Table[ctrl[n, {j}, rzGate[n, i, 2\pi/2^{j+1-i}]], {j, n, i+1, -1}]].} \\ & \text{hadamardGate[n, i]} \end{array}$$

assemble the quantum Fourier transformation by connecting all qubits:

- 1. Connect all qubits through the above QFTblock operation
- 2. Swap the order of the qubits

```
n_{[n]} = \text{qft}[n_{\text{Integer}}] / ; n \ge 1 := \text{Apply}[Dot, Table[swapGate[n, \{i, n+1-i\}], \{i, 1, \frac{n}{2}\}]].
        Apply[Dot, Table[qftBlock[n, i], {i, n, 1, -1}]]
```

QFT consists of a polynomial number of gates:

- n Hadamard gates
- $\frac{n(n-1)}{2}$ controlled rotations
- \blacksquare $\left|\frac{n}{2}\right|$ swap gates

simple formula for the resulting matrix representation:

```
ln[*] = \frac{1}{2^{n/2}} \text{ Table} \left[ e^{2\pi i j k/2^n}, \{j, 0, 2^n - 1\}, \{k, 0, 2^n - 1\} \right], \{n, 6\} \right] / / 
        FullSimplify
Out[*]= {True, True, True, True, True, True}
```

Quantum phase estimation

estimate the phase of this:

(see ReducedDensityMatrix.nb)

```
ln[\circ]:= W = \{1\};
     u[\phi_{-}] = \{\{e^{2\pi i \phi}\}\};
     check that we set up the problem correctly:
ln[@]:= \{u[\phi].w === e^{2\pi i \phi} w, Norm[w] == 1\}
Out[•]= {True, True}
      unit operator in the space of the system U:
In[*]:= u0 = IdentityMatrix[Length[w], SparseArray];
      U operator attached to n qubits, controlled by the i<sup>th</sup> qubit:
ln[\bullet]:= ctrlU[n_Integer, i_Integer, \phi_] /; 1 \le i \le n :=
       KroneckerProduct[op[n, i, proj0], u0] + KroneckerProduct[op[n, i, proj1], u[\phi]]
      estimate to t digits of precision:
ln[\circ] := t = 4;
      \epsilon[\phi_{-}] = KroneckerProduct[ConjugateTranspose[qft[t]], u0].
           Apply Dot, Table [\text{ctrlU}[t, i, 2^{t-i}\phi], \{i, t, 1, -1\}]].
           KroneckerProduct[Apply[Dot, Table[hadamardGate[t, i], {i, t}]], u0].
           Flatten[KroneckerProduct[SparseArray[1 → 1, 2<sup>t</sup>], w]] // Normal;
      probabilities for measuring the different basis states: trace out the SUT and look at the diagonal
      elements of the reduced density matrix
```

```
ln[\psi] = rdm[\psi ABC_? VectorQ, \{dA_Integer/; dA \ge 1, dB_Integer/; dB \ge 1,
          dC_Integer /; dC ≥ 1}] /; Length[\psi ABC] == dA dB dC :=
      With[\{P = ArrayReshape[\psi ABC, \{dA, dB, dC\}]\},
       Flatten[Transpose[P, {1, 3, 2}].ConjugateTranspose[P], {{1, 2}, {3, 4}}]]
     traceout[\psi_?VectorQ, d_Integer /; d \geq 1] /; Divisible[Length[\psi], d] :=
      rdm[\psi, \{1, d, \frac{Length[\psi]}{d}\}]
    traceout[\psi_?VectorQ, d_Integer /; d \le -1] /; Divisible[Length[\psi], -d] :=
      rdm[\psi, \{\frac{Length[\psi]}{d}, -d, 1\}]
```

 $log_{i} = prob[\phi_?NumericQ] := Re[Diagonal[traceout[e[N[\phi]], -Length[w]]]]$

When ϕ is an integer multiple of 2^{-t}, only one basis state has 100% probability of occurring. The ith basis state corresponds to a measurement of $\phi = \frac{i-1}{2^t}$:

 $ln[\cdot]:=$ Table[prob[ϕ], $\{\phi, 0, 1, 2^{-t}\}$] // Chop // MatrixForm

Out[]//MatrixForm=

When ϕ is not an integer multiple of 2^{-t}, all basis states can occur in measurement:

```
In[*]:= Round[prob[0.2], 0.001]
Out[\circ] = \{0.004, 0.008, 0.025, 0.876, 0.055, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.011, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.005, 0.00
                                                                                        0.003, 0.002, 0.002, 0.001, 0.001, 0.001, 0.002, 0.002, 0.003}
```

The mean measurement is a bad estimator (doesn't converge as $t \to \infty$). The most likely measurement is a good estimator.

