

Information and coding theory HW#1

Evgeny Marshakov

Problem 1

We know that

$$I(X; Z|Y) = H(X|Y) - H(X|Y, Z) = H(X, Y) - H(Y) - H(X, Y|Z) + H(Y|Z)$$

So we have that

$$I(X; Z|Y) + I(Y; Z) = H(X, Y) - H(Y) - H(X, Y|Z) + H(Y|Z) + H(Y) - H(Y|Z) = H(X, Y) - H(X, Y|Z) = I(X, Y; Z)$$

Problem 2

Using the chain rules for $H(X^n)$ and $H(X_i^-)$ we obtain

$$H(X^n) = \sum_{j=1}^n H(X_j | X_{j-1}, \dots, X_1)$$

$$H(X_i^-) = \sum_{j=1}^n H(X_j | X_i^{j-1}) = \sum_{j=1}^n H(X_j | X_{j-1}, \dots, X_{i+1}, X_{i-1}, \dots, X_1)$$

Obviously, we have the following inequalities

$$\begin{cases} H(X_j | X_{j-1}, \dots, X_1) \leq H(X_j | X_i^{j-1}), & i < j \\ H(X_j | X_{j-1}, \dots, X_1) = H(X_j | X_i^{j-1}), & i > j \end{cases}$$

We see that there is no term corresponding to i in $H(X_{\bar{i}})$. So using these inequalities we can easily find that

$$\sum_{j \neq i} H(X_j | X_{j-1}, \dots, X_1) \leq \sum_{j=1}^n H(X_j | X_{\bar{i}}^{j-1})$$

or

$$H(X^n) - H(X_i | X_{i-1}, \dots, X_1) \leq H(X_{\bar{i}})$$

Summing over i we obtain

$$\sum_{i=1}^n [H(X^n) - H(X_i | X_{i-1}, \dots, X_1)] \leq \sum_{i=1}^n H(X_{\bar{i}}) \Leftrightarrow (n-1)H(X^n) \leq \sum_{i=1}^n H(X_{\bar{i}})$$

Problem 3

1. Let us consider the following probabilities

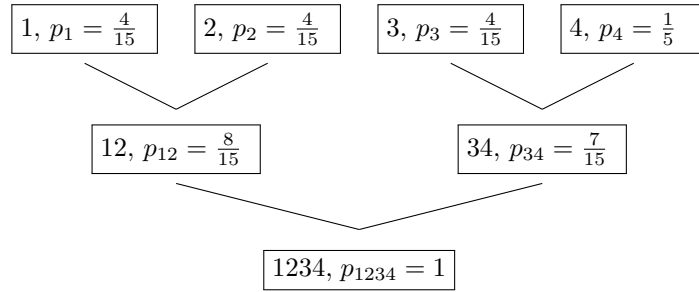
$$p_i = \mathbb{P}(X = i) = \frac{\#i \text{ in the string}}{\#all \text{ elements in the string}}$$

More precisely,

$$p_1 = p_2 = p_3 = \frac{4}{15}, p_4 = \frac{1}{5}$$

2. $H(X) = -3 \cdot \frac{4}{15} \log_2 \frac{4}{15} - \frac{1}{5} \log_2 \frac{1}{5} \approx 1.98$ bits.

- 3.



So the Huffman code for these probabilities is the following

1	00
2	01
3	10
4	11

It can be easily seen that the average number of bits is 2.

4. Encoded file F

$$F \rightarrow 000010101010111101110101010110000011011011001000010000111001$$

The average number of bits in the file F is also 2.

Problem 4

1. We want to find the capacity of the Z channel. Let p_0 be the probability $\mathbb{P}(X = 0)$. Set

$$q_0 = \mathbb{P}(Y = 0) = p_0 + (1 - p_0)\delta$$

$$q_1 = \mathbb{P}(Y = 1) = (1 - p_0)(1 - \delta)$$

Then we have

$$H(Y) = -q_0 \log_2(q_0) - q_1 \log_2(q_1) = -[p_0 + (1 - p_0)\delta] \log_2[p_0 + (1 - p_0)\delta] - [(1 - p_0)(1 - \delta)] \log_2[(1 - p_0)(1 - \delta)]$$

$$H(Y|X) = p_0 H(Y|X = 0) + (1 - p_0) H(Y|X = 1) = (1 - p_0)(-\delta \log_2 \delta - (1 - \delta) \log_2(1 - \delta)) =: (1 - p_0) \mathbf{H}(\delta)$$

Since $H(Y|X = 0) = -1 \log_2 1 = 0$. So the the mutual information is the following

$$\begin{aligned} I(X; Y)(p_0) &= H(Y) - H(Y|X) = -[p_0 + (1 - p_0)\delta] \log_2[p_0 + (1 - p_0)\delta] - [(1 - p_0)(1 - \delta)] \log_2[(1 - p_0)(1 - \delta)] - \\ &\quad -(1 - p_0) \mathbf{H}(\delta) = \mathbf{H}((1 - p_0)(1 - \delta)) - (1 - p_0) \mathbf{H}(\delta) \end{aligned}$$

Here $\mathbf{H}(\delta) = -\delta \log_2 \delta - (1 - \delta) \log_2 (1 - \delta)$

After calculating the derivative and putting it to zero we obtain

$$\begin{aligned} \frac{dI(X; Y)}{dp_0} &= -(1 - \delta) \log_2 \left[\frac{p_0 + (1 - p_0)\delta}{(1 - p_0)(1 - \delta)} \right] + \mathbf{H}(\delta) = 0 \Leftrightarrow \\ &\Leftrightarrow p_0 = 1 - \frac{1}{(1 - \delta)(1 + 2^{H(\delta)/(1 - \delta)})} \end{aligned}$$

Substituting the result into the capacity formula we obtain

$$C(\delta) = \mathbf{H} \left(\frac{1}{1 + 2^{H(\delta)/(1 - \delta)}} \right) - \frac{1}{(1 - \delta)(1 + 2^{H(\delta)/(1 - \delta)})} \mathbf{H}(\delta) = \log_2 \left[1 + 2^{-\frac{\mathbf{H}(\delta)}{1 - \delta}} \right]$$

2. Firstly, we use l'Hopital's rule

$$\lim_{\delta \rightarrow 1} \left(\frac{\frac{1}{1 - \delta}}{1 + 2^{H(\delta)/(1 - \delta)}} \right) = \lim_{\delta \rightarrow 1} \frac{\left(\frac{1}{1 - \delta} \right)'}{\left(1 + 2^{H(\delta)/(1 - \delta)} \right)'} = \lim_{\delta \rightarrow 1} \frac{1}{\ln 2} \frac{\frac{1}{(1 - \delta)^2}}{2^{H(\delta)/(1 - \delta)} \cdot \left(\frac{H(\delta)}{1 - \delta} \right)'} = \quad (1)$$

$$= \frac{1}{\ln 2} \lim_{\delta \rightarrow 1} \frac{\frac{1}{(1 - \delta)^2}}{2^{H(\delta)/(1 - \delta)} \cdot \left(-\frac{\log_2 \delta}{(1 - \delta)^2} \right)} = -\frac{1}{\ln 2} \lim_{\delta \rightarrow 1} \frac{1}{\log_2 \delta \cdot 2^{-\frac{\delta \log_2 \delta}{1 - \delta} - \log_2 (1 - \delta)}} = \quad (2)$$

$$= -\frac{1}{\ln 2} \lim_{\delta \rightarrow 1} \frac{1}{\frac{\log_2 \delta}{1 - \delta} \cdot 2^{-\frac{\delta \log_2 \delta}{1 - \delta}}} \quad (3)$$

So we need to find the following limit

$$\lim_{\delta \rightarrow 1} \frac{\log_2 \delta}{1 - \delta} = -\log_2 e \cdot \lim_{x \rightarrow 0} \frac{1 + x}{x} = -\log_2 e$$

because it is zamechatelniy limit. Substituting this result to (3) we obtain

$$-\frac{1}{\ln 2} \frac{1}{(-\log_2 e) \cdot e} = \frac{1}{e}$$

So we find that the capacity achieving distribution is the following

$$\mathbb{P}(X = 0) = p_0 = 1 - \frac{1}{e}$$

$$\mathbb{P}(X = 1) = 1 - p_0 = \frac{1}{e}$$

And the distribution is the following

$$C(1) = \lim_{\delta \rightarrow 1} \log_2 \left[1 + 2^{-\frac{H(\delta)}{1-\delta}} \right] = 0$$

Problem 5

Let p_0 be the probability $\mathbb{P}(X = 0)$. Set

$$q_0 = \mathbb{P}(Y = 0) = (1 - \epsilon)(p_0(1 - p) + p(1 - p_0))$$

$$q_1 = \mathbb{P}(Y = 1) = (1 - \epsilon)(p_0 p + (1 - p_0)(1 - p))$$

$$q_* = \mathbb{P}(Y = *) = \epsilon$$

It can be easily seen that this is symmetric channel so the capacity of this channel achieves with the uniform distribution of X , i.e. $p_0 = \frac{1}{2}$, hence

$$q_0 = q_1 = \frac{1}{2}(1 - \epsilon)$$

TO find the capacity we need to calculate $H(Y)$ and $H(Y|X)$.

$$H(Y) = -q_0 \log_2 q_0 - q_1 \log_2 q_1 - q_* \log_2 q_* = -(1 - \epsilon) \log_2 \left[\frac{1 - \epsilon}{2} \right] - \epsilon \log_2 \epsilon$$

$$H(Y|X) = -(1 - p)(1 - \epsilon) \log_2 [(1 - p)(1 - \epsilon)] - p(1 - \epsilon) \log_2 (p(1 - \epsilon)) - \epsilon \log_2 \epsilon$$

Hence, the capacity is the following

$$C = (1 - \epsilon)(1 - \log_2(1 - \epsilon)) + (1 - p) \log_2(1 - p) + (1 - p) \log_2(1 - \epsilon) + p \log_2(p) + p \log_2(1 - \epsilon) =$$

$$(1 - \epsilon)(1 - H(p))$$

Problem 6

We have the optimization problem

$$\text{maximize} \quad \sum_{i=1}^{499} \frac{1}{2} \log_2 \left(1 + \frac{P_i}{100} \right) + \frac{1}{2} \log_2 \left(1 + \frac{P_{500}}{25} \right) \quad (4)$$

$$\text{subject to} \quad \sum_i P_i = 50 \quad (5)$$

We know that the solution if this problem has the following form

$$P_i = (\nu - N_i)^+$$

where ν is a constant that satisfies

$$\sum_{i=1}^{500} (\nu - N_i)^+ = 50 \quad (6)$$

In this case it is obvious that if $\nu \geq 100$ equality (6) doesn't hold. If $\nu < 100$ then the expression (6) has the form

$$(\nu - 25)^+ = 50$$

So $\nu = 75$ and we have the following solution

$$P_1 = \dots = P_{499} = 0$$

$$P_{500} = 50$$

Hence, the capacity is the following

$$C = \frac{1}{2} \log_2 \left(1 + \frac{50}{25} \right) = \frac{1}{2} \log_2 3 \approx 0.79 \text{bits}$$

Problem 7

Our capacity region is determined by the following inequalities

$$R_1 \leq I(X_1; Y|X_2) \tag{7}$$

$$R_2 \leq I(X_2; Y|X_1) \tag{8}$$

$$R_1 + R_2 \leq I(X_1, X_2; Y) \tag{9}$$

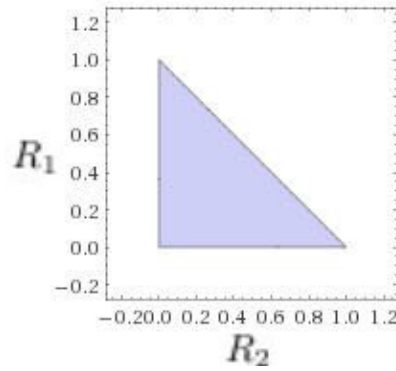
It can be easily seen that

$$I(X_1; Y|X_2) = H(X_1|X_2) - H(X_1|Y, X_2) \leq H(X_1) \leq 1$$

and, obviously, the equality holds when $\mathbb{P}(X_1 = 0) = \mathbb{P}(X_2 = 0) = \frac{1}{2}$. The same is true for (8). Also we can find that

$$R_1 + R_2 \leq I(X_1, X_2; Y) \leq H(Y) = 1$$

So we have the following plot



Problem 8

We have the following channel

X_1	X_2	Y
-1	-1	-2
1	-1	0
-1	1	0
1	1	2

It is obvious that we can consider symmetric case

$$\mathbb{P}(X_1 = 1) = \mathbb{P}(X_2 = 1) = p_1$$

Firstly, let us find $I(X_1; Y|X_2)$.

$$I(X_1; Y|X_2) \leq 1 \quad (10)$$

$$I(X_2; Y|X_1) \leq 1 \quad (11)$$

As in the previous exercise, and the equalities holds when $p_1 = \frac{1}{2}$. Now let us find $I(X_1, Y)$

$$I(X_1, Y) = H(X) - H(X|Y) = p_0 \log_2 \frac{1}{p_0} + (1 - p_0) \log_2 \left(\frac{1}{1 - p_0} \right) - \mathbb{P}(Y = -2)H(X_1|Y = -2) - \quad (12)$$

$$\mathbb{P}(Y = 2)H(X_1|Y = 2) - \mathbb{P}(Y = 0)H(X_1|Y = 0) = \quad (13)$$

$$p_0 \log_2 \frac{1}{p_0} + (1 - p_0) \log_2 \left(\frac{1}{1 - p_0} \right) - 4p_0(1 - p_0) \left(p_0(1 - p_0) \log_2 \frac{1}{p_0(1 - p_0)} \right) \quad (14)$$

It can be easily calculated that

$$\frac{dI(X_1; Y)}{dp_0} = 0 \Leftrightarrow p_0 = \frac{1}{2}$$

And for $p_0 = \frac{1}{2}$

$$I(X_1; Y) \left[\frac{1}{2} \right] = \frac{1}{2}$$

We have the same for $I(X_2; Y)$. Finally, the $\max I(X_1, X_2, Y) = \frac{3}{2}$ and we have the following picture

