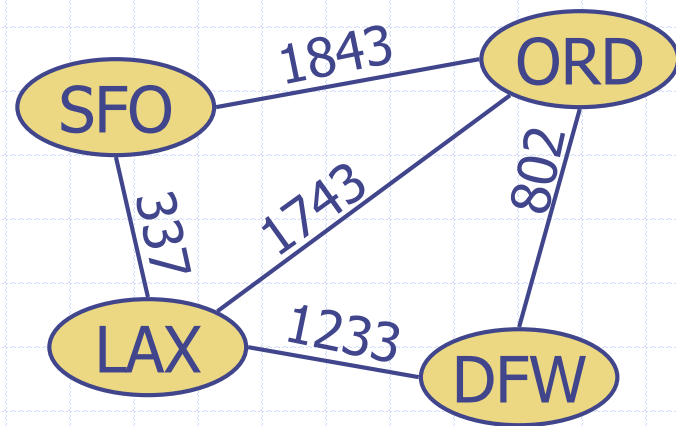
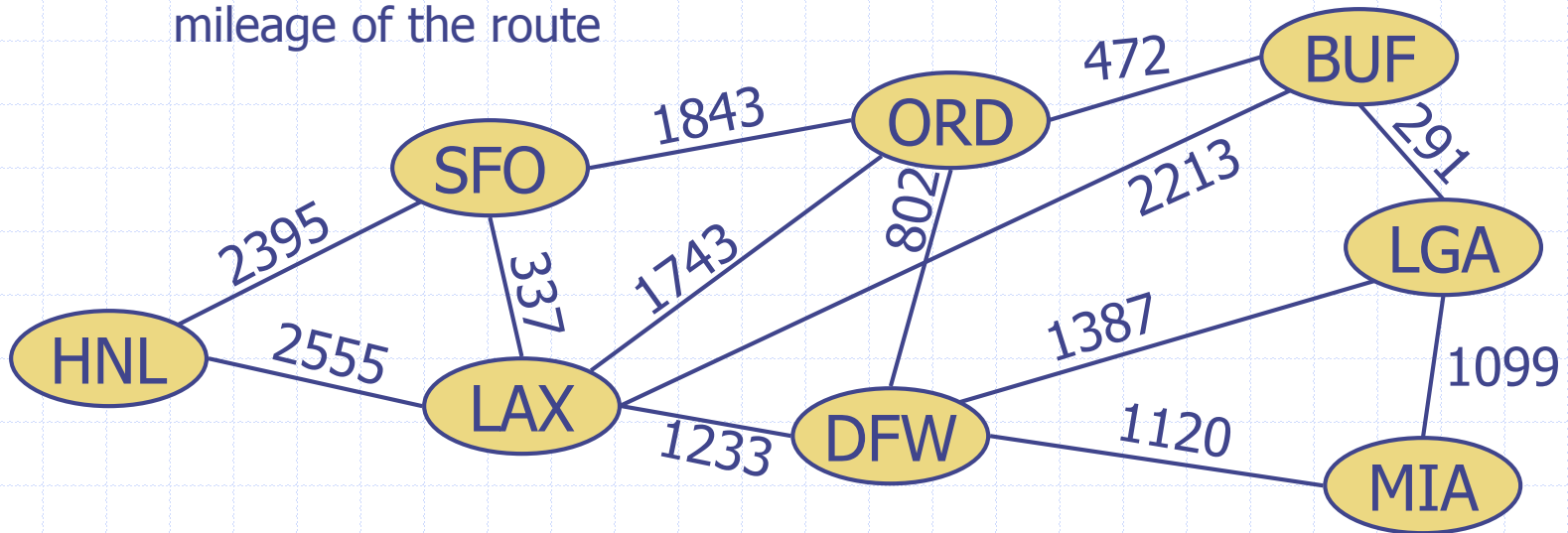


# Graphs



# Graphs

- A graph is a pair  $(V, E)$ , where
  - $V$  is a set of nodes, called **vertices**
  - $E$  is a collection of pairs of vertices, called **edges**
  - Vertices and edges are positions and store elements
- Example:
  - A vertex represents an airport and stores the three-letter airport code
  - An edge represents a flight route between two airports and stores the mileage of the route



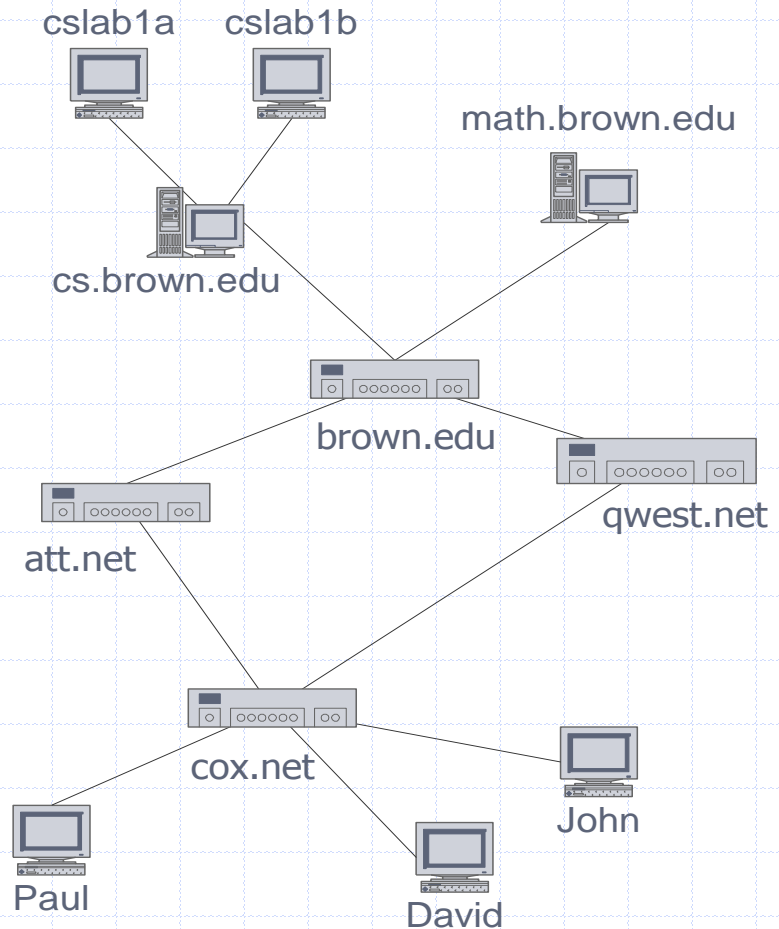
# Edge Types

- Directed edge
  - ordered pair of vertices  $(u,v)$
  - first vertex  $u$  is the origin
  - second vertex  $v$  is the destination
  - e.g., a flight
- Undirected edge
  - unordered pair of vertices  $(u,v)$
  - e.g., a flight route
- Directed graph
  - all the edges are directed
  - e.g., route network
- Undirected graph
  - all the edges are undirected
  - e.g., flight network



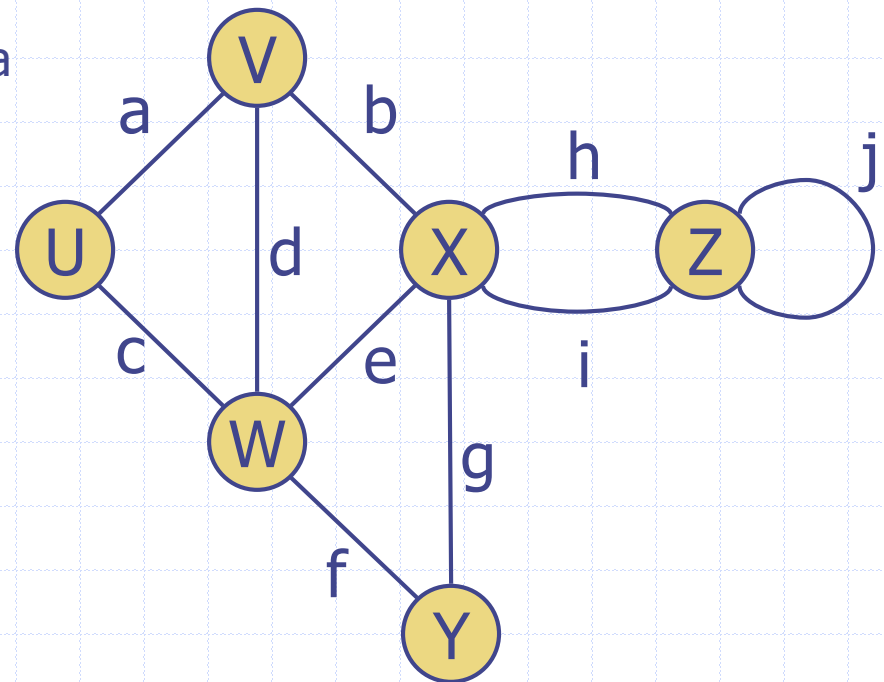
# Applications

- ❑ Electronic circuits
  - Printed circuit board
  - Integrated circuit
- ❑ Transportation networks
  - Highway network
  - Flight network
- ❑ Computer networks
  - Local area network
  - Internet
  - Web
- ❑ Databases
  - Entity-relationship diagram



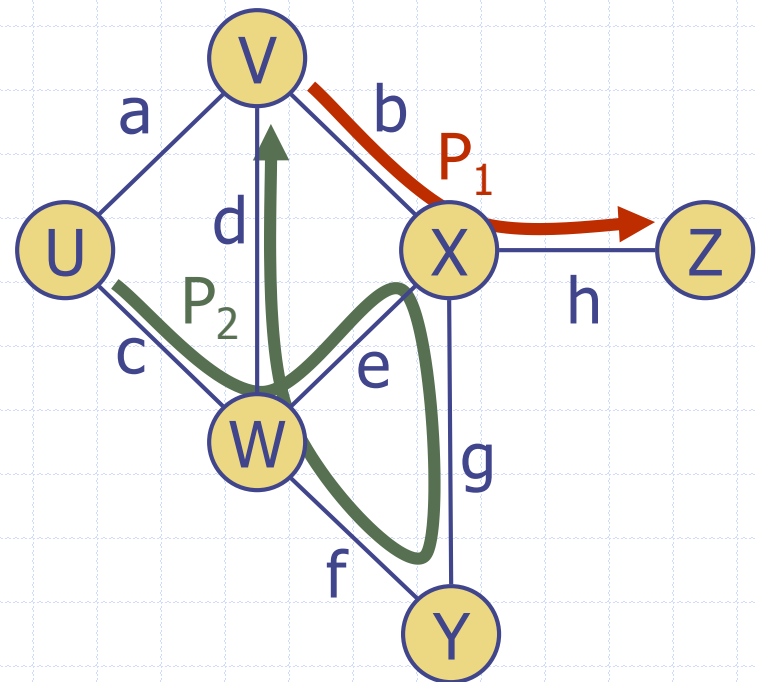
# Terminology

- End vertices (or endpoints) of an edge
  - U and V are the endpoints of a
- Edges incident on a vertex
  - a, d, and b are incident on V
- Adjacent vertices
  - U and V are adjacent
- Degree of a vertex
  - X has degree 5
- Parallel edges
  - h and i are parallel edges
- Self-loop
  - j is a self-loop



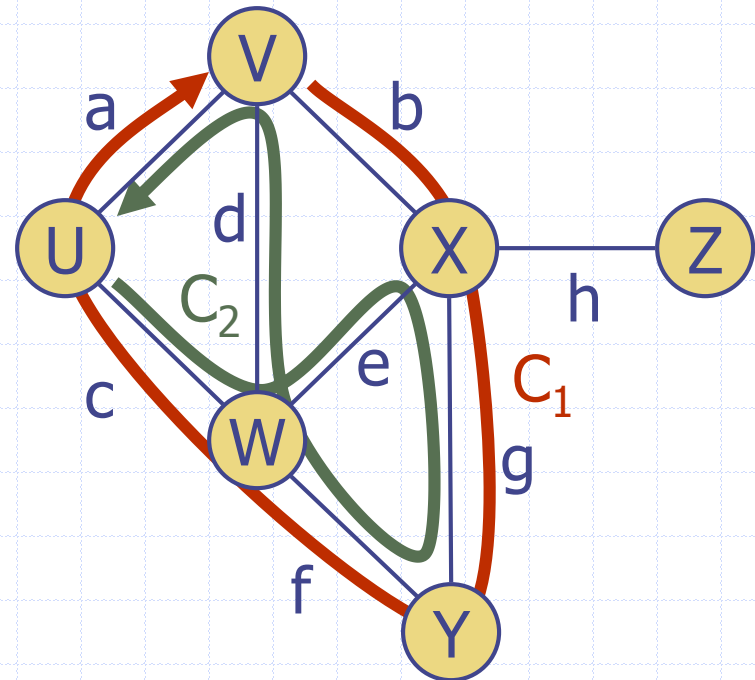
# Terminology (cont.)

- Path
  - sequence of alternating vertices and edges
  - begins with a vertex
  - ends with a vertex
  - each edge is preceded and followed by its endpoints
- Simple path
  - path such that all its vertices and edges are distinct
- Examples
  - $P_1 = (V, b, X, h, Z)$  is a simple path
  - $P_2 = (U, c, W, e, X, g, Y, f, W, d, V)$  is a path that is not simple



# Terminology (cont.)

- Cycle
  - circular sequence of alternating vertices and edges
  - each edge is preceded and followed by its endpoints
- Simple cycle
  - cycle such that all its vertices and edges are distinct
- Examples
  - $C_1 = (V, b, X, g, Y, f, W, c, U, a, V)$  is a simple cycle
  - $C_2 = (U, c, W, e, X, g, Y, f, W, d, V, a, U)$  is a cycle that is not simple



# Properties

## Property 1

$$\sum_v \deg(v) = 2m$$

Proof: each edge is counted twice

## Property 2

In an undirected graph with no self-loops and no multiple edges

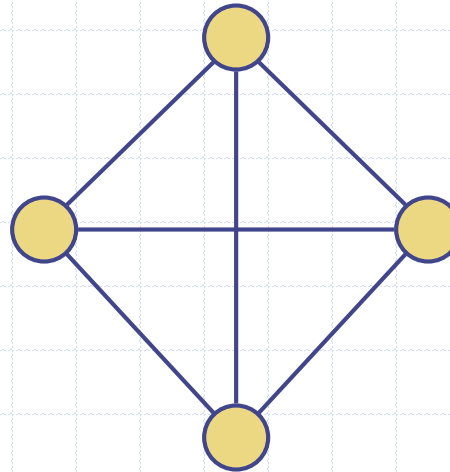
$$m \leq n(n-1)/2$$

Proof: each vertex has degree at most  $(n-1)$

What is the bound for a directed graph?

## Notation

$n$	number of vertices
$m$	number of edges
$\deg(v)$	degree of vertex $v$



## Example

- $n = 4$
- $m = 6$
- $\deg(v) = 3$

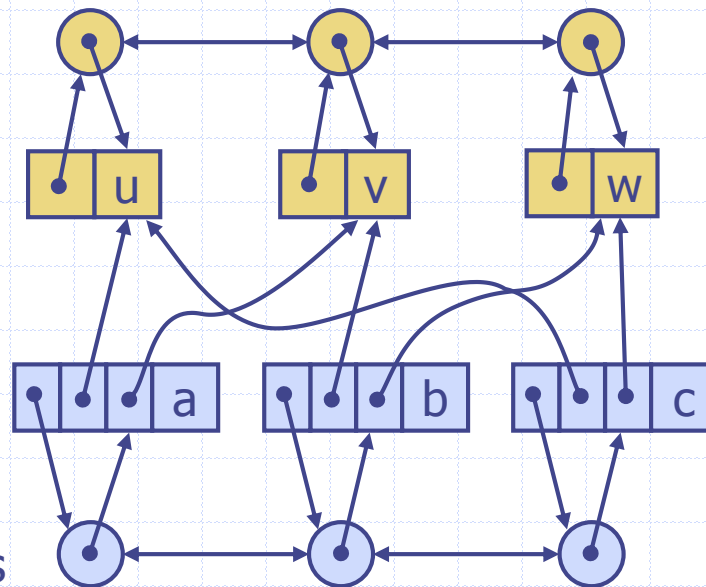
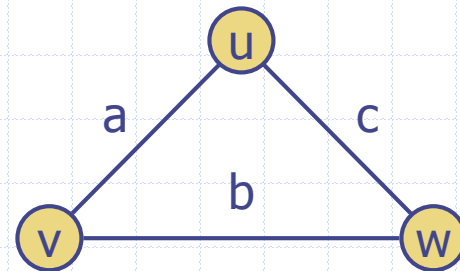


# Main Methods of the Graph ADT

- Vertices and edges
  - are positions
  - store elements
  - are objects themselves
- Accessor methods
  - `e.endVertices()`: a list of the two endvertices of `e`
  - `u.isAdjacentTo(v)`: true iff `u` and `v` are adjacent
  - `*v`: reference to element associated with vertex `v`
  - `*e`: reference to element associated with edge `e`
- Update methods
  - `insertVertex(o)`: insert a vertex storing element `o`
  - `insertEdge(v, w, o)`: insert an edge `(v,w)` storing element `o`
  - `eraseVertex(v)`: remove vertex `v` (and its incident edges)
  - `eraseEdge(e)`: remove edge `e`
- Iterable collection methods
  - `incidentEdges(v)`: list of edges incident to `v`
  - `vertices()`: list of all vertices in the graph
  - `edges()`: list of all edges in the graph

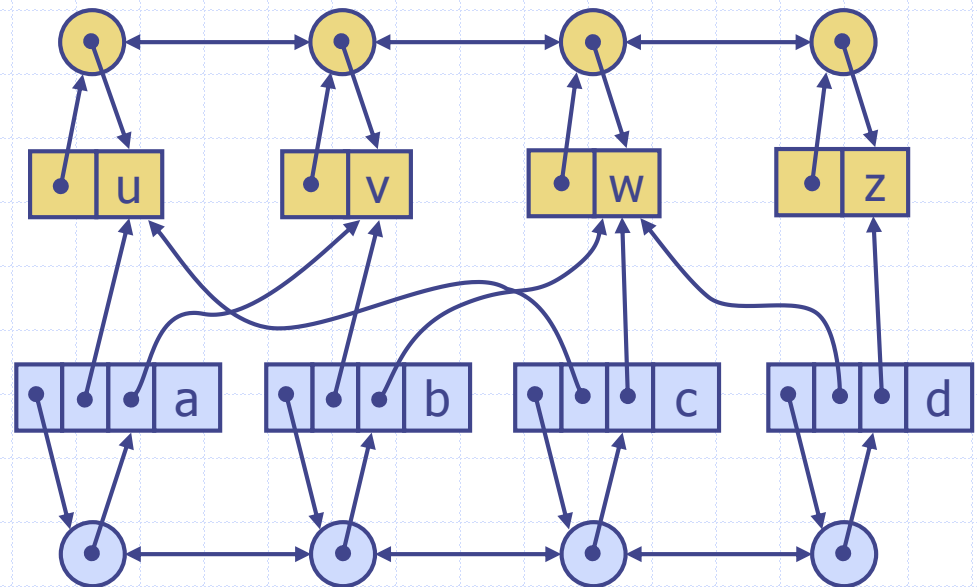
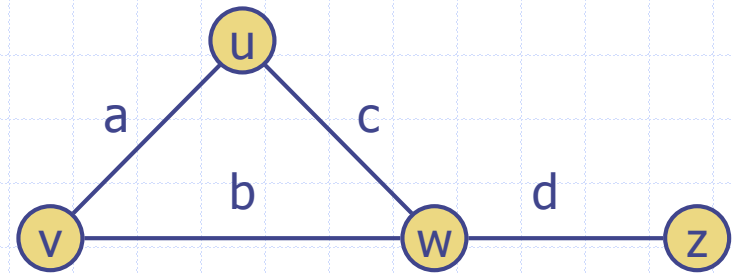
# Edge List Structure

- Vertex object
  - element
  - reference to position in vertex sequence
- Edge object
  - element
  - origin vertex object
  - destination vertex object
  - reference to position in edge sequence
- Vertex sequence
  - sequence of vertex objects
- Edge sequence
  - sequence of edge objects



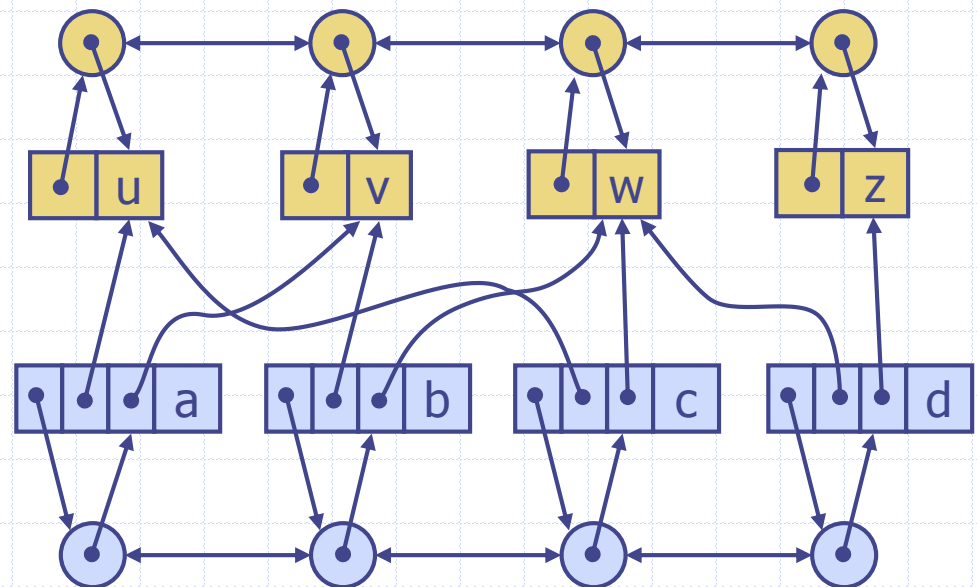
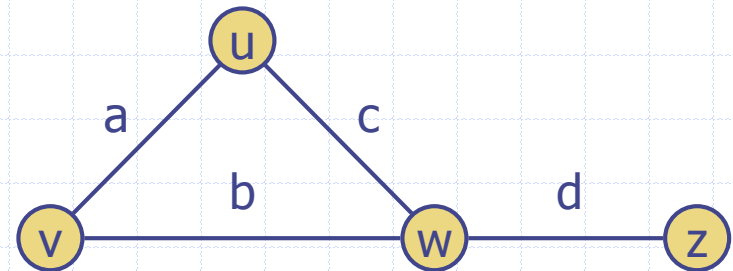
# Edge List Structure

- Adding vertices:
  - Input: vertex data.
  - Create vertex object.
  - Add node to vertex list head/tail.
  - Update links.
  - $O(1)$  runtime.
- Adding edges:
  - Input: two vertex objects, edge data.
  - Create edge object with pointers to vertices.
  - Add node to edge list.
  - Update links.
  - $O(1)$  runtime.



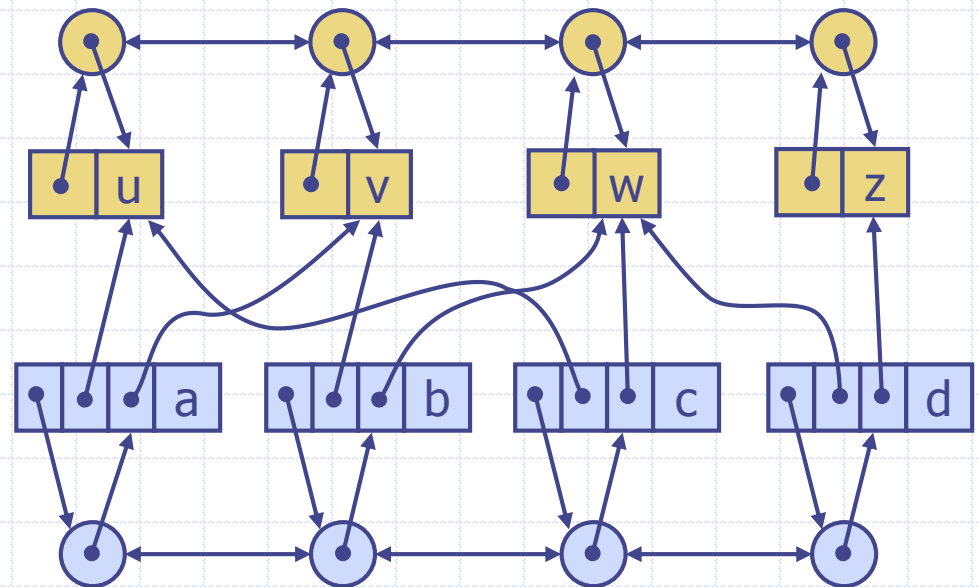
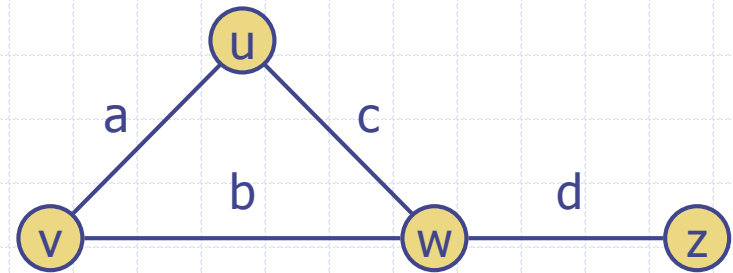
# Edge List Structure

- Determine if two vertices are adjacent:
  - Input: two vertex objects.
  - From beginning of edge sequence, look for edge.
  - $O(m)$  runtime.
- Find all adjacent edges:
  - Input: vertex object.
  - From beginning of edge sequence, look at each edge for vertex.
  - $O(m)$  runtime.



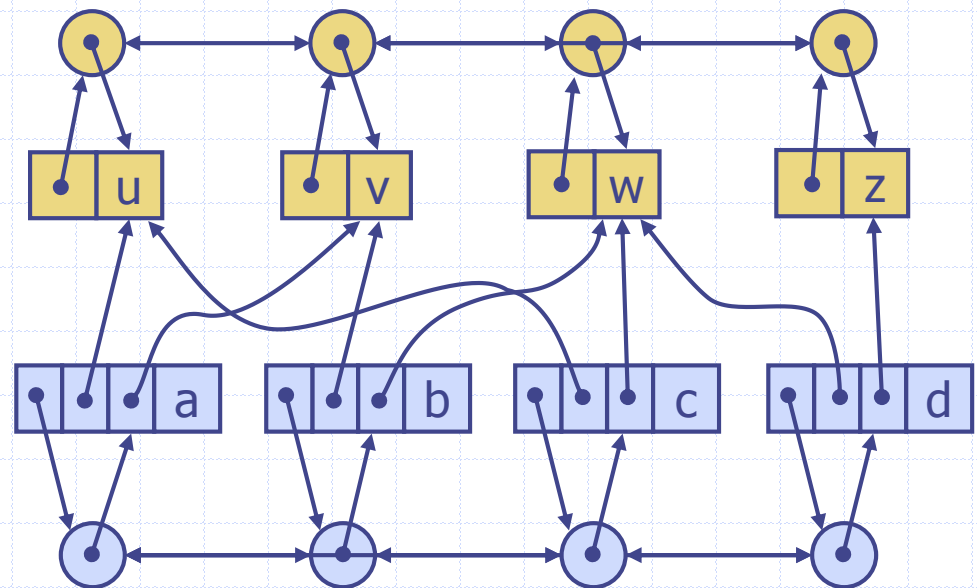
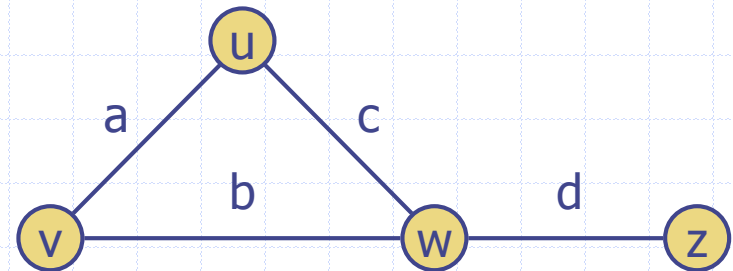
# Edge List Structure

- Remove an edge:
  - Input: edge object.
  - Remove edge from edge sequence.
  - $O(1)$  runtime.
- E.g., remove edge a.



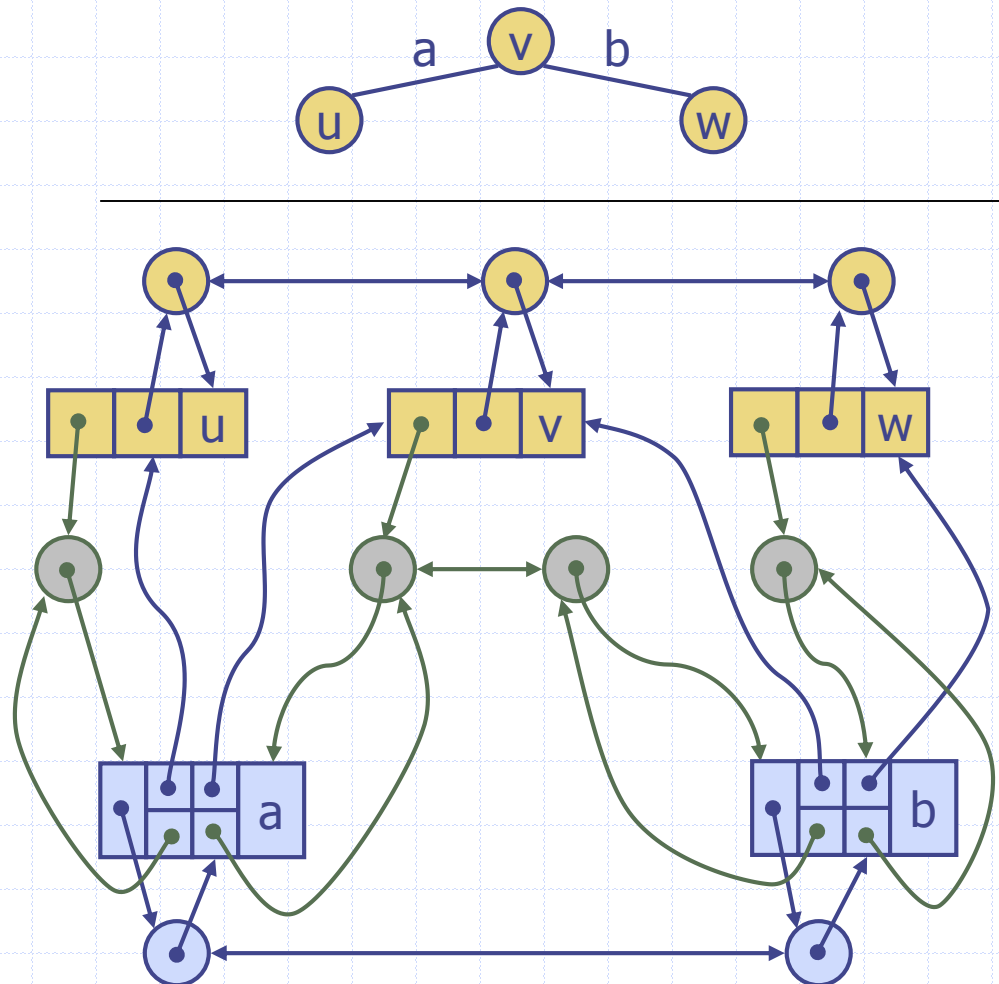
# Edge List Structure

- Remove a vertex:
  - Input: vertex object.
  - From beginning of edge sequence, find edges containing vertex and remove edge.
  - Remove vertex from vertex sequence.
  - $O(m) + O(1)$  runtime.
- E.g., remove vertex  $w$ .



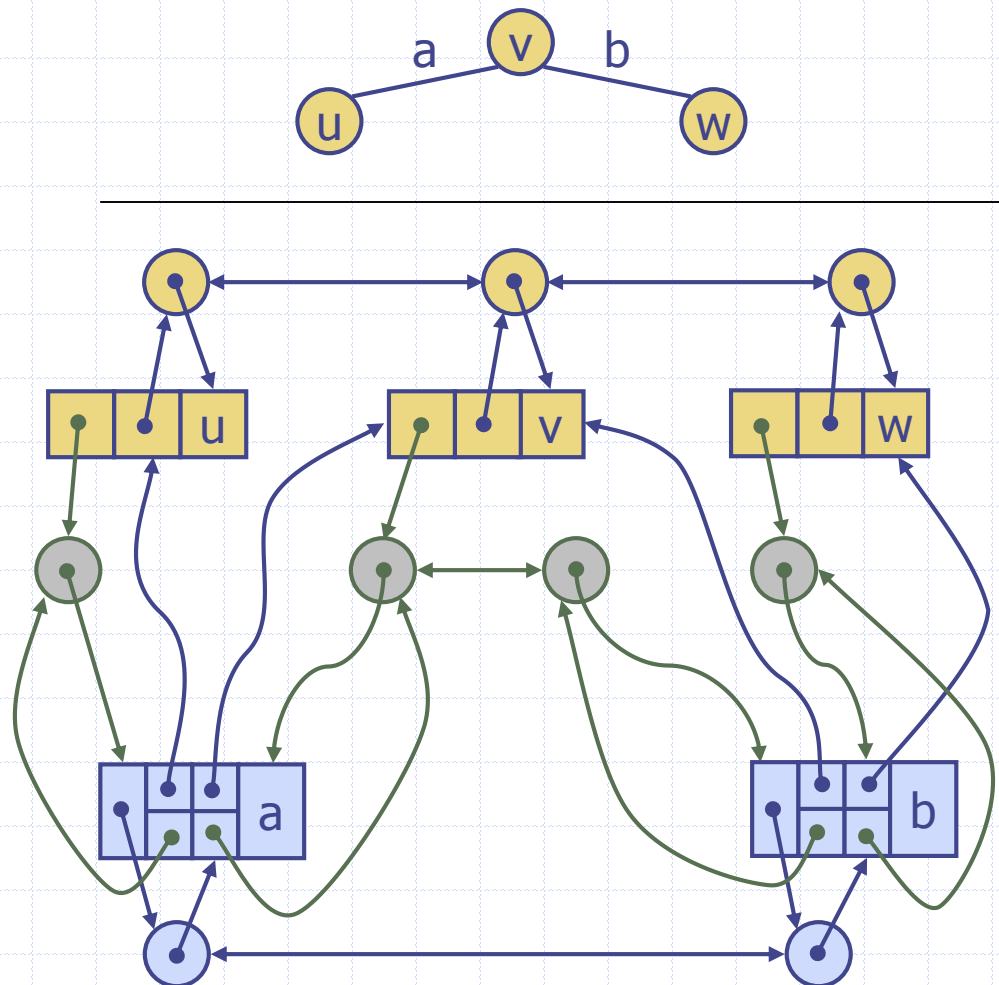
# Adjacency List Structure

- Edge list structure
- Incidence sequence for each vertex
  - sequence of references to edge objects of incident edges
- Augmented edge objects
  - references to associated positions in incidence sequences of end vertices



# Adjacency List Structure

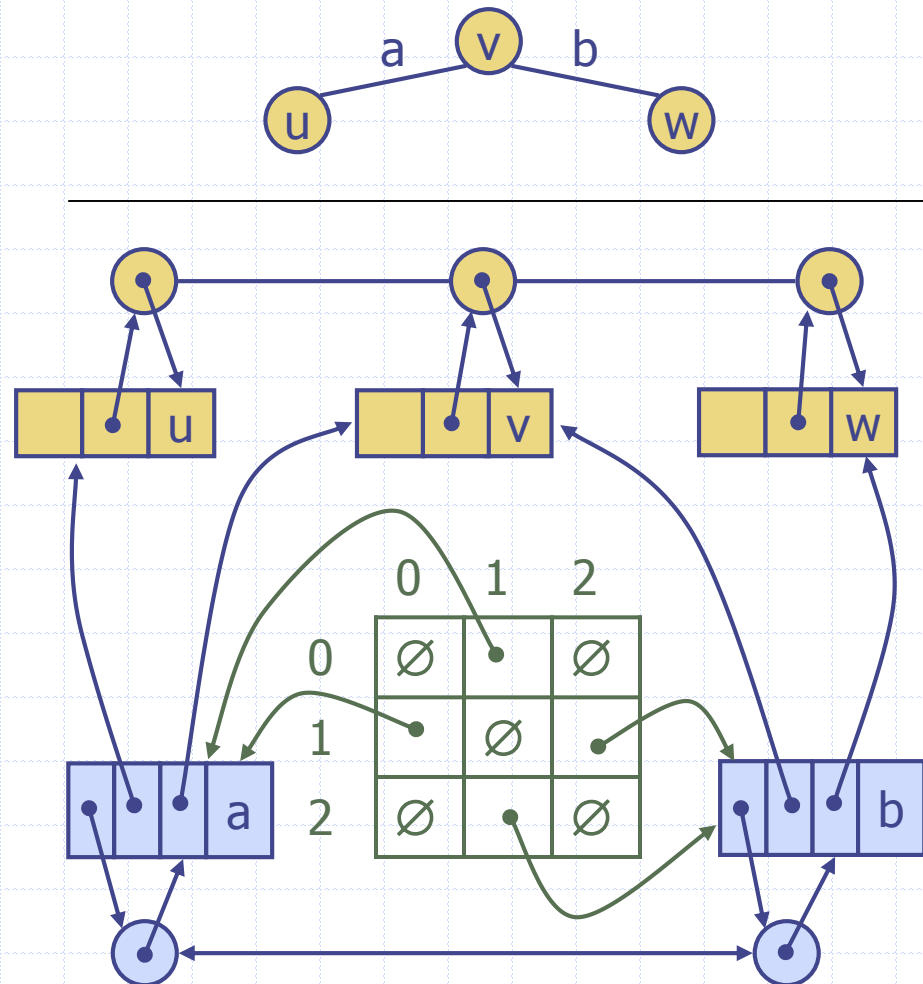
- ❑ Add vertex/edge:
  - Still  $O(1)$  runtime.
- ❑ Determine if two vertices are adjacent:
  - Traverse corresponding edge lists.
  - $\min(\deg(u), \deg(v))$  runtime.
- ❑ Find all adjacent edges:
  - $O(1)$  runtime.
- ❑ Remove edge:
  - $O(1)$  runtime.
- ❑ Remove vertex:
  - Traverse corresponding edge list.
  - $O(\deg(v))$  runtime.





# Adjacency Matrix Structure

- ❑ Edge list structure
- ❑ Augmented vertex objects
  - Integer key (index) associated with vertex
- ❑ 2D-array adjacency array
  - Reference to edge object for adjacent vertices
  - Null for non adjacent vertices
- ❑ The “old fashioned” version just has 0 for no edge and 1 for edge



# Performance

<ul style="list-style-type: none"><li>▪ <math>n</math> vertices, <math>m</math> edges</li><li>▪ no parallel edges</li><li>▪ no self-loops</li></ul>	Edge List	Adjacency List	Adjacency Matrix
Space	$n + m$	$n + m$	$n^2$
<code>v.incidentEdges()</code>	$m$	$\deg(v)$	$n$
<code>u.isAdjacentTo(v)</code>	$m$	$\min(\deg(v), \deg(w))$	1
<code>insertVertex(o)</code>	1	1	$n^2$
<code>insertEdge(v, w, o)</code>	1	1	1
<code>eraseVertex(v)</code>	$m$	$\deg(v)$	$n^2$
<code>eraseEdge(e)</code>	1	1	1