Time Complexities of Three Sorting Algorithms

Selection Sort

First, select the smallest element

Then, select the smallest element among the remaining elements

Then, select the smallest element among the remaining elements

... ...

Select the smallest element among elements that have not yet been processed every time.

8 6 2 3 1 5 7 4

1 6 2 3 8 5 7 4

1 2 6 3 8 5 7 4

1 2 6 3 8 5 7 4

```
public class SelectionSort{
  // Our algorithm class does not allow any instances
  private SelectionSort(){}
  public static void sort(Comparable[] arr){
     int n = arr.length;
     for( int i = 0; i < n; i ++ ){
       // Find the index of the smallest value in the interval [i, n]
        int minIndex = i;
       for(int i = i + 1; i < n; i + +)
          // Use compareTo method to compare the size of two Comparable objects
          if( arr[i].compareTo( arr[minIndex] ) < 0 )</pre>
             minIndex = j;
        swap( arr , i , minIndex);
  private static void swap(Object[] arr, int i, int j) {
     Object t = arr[i];
     arr[i] = arr[i];
     arr[i] = t;
```

Sorted: arr[0, i) Out of order: arr[i, 0)

Insertion Sort

8 6 2 3 1 5 7 4

6 8 2 3 1 5 7 4

6 8 2 3 1 5 7 4

6 2 8 3 1 5 7 4

2 6 8 3 1 5 7 4

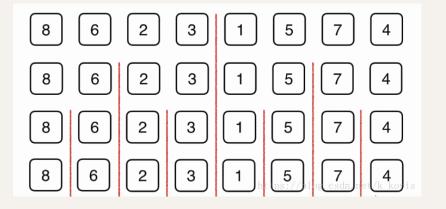
Sorted: arr[0, i)
Out of order: arr[i, 0)



.

```
public class InsertionSort{
 // Our algorithm class does not allow any instances
  private InsertionSort(){}
// Use InsertionSort to sort the entire arr array
  public static void sort(Comparable[] arr){
     int n = arr.length;
     for (int i = 0; i < n; i++) {
                                                 1. simplify
        Comparable e = arr[i];
        int j = i;
        for(; j > 0 && arr[j-1].compareTo(e) > 0; j--){
          arr[j] = arr[j-1];
        arr[j] = e;
                    2. optimization
```

Merge Sort

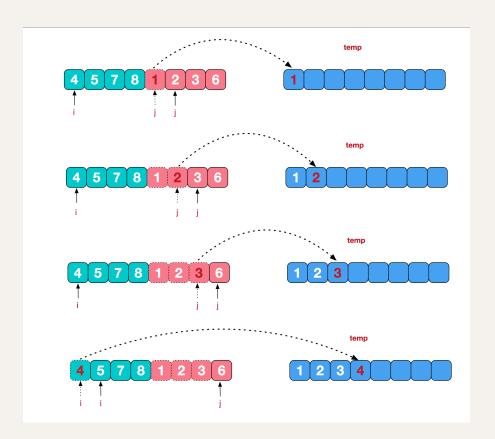


In mergeSort(E[] arr, I, r):

Step 1: Sort arr[l, mid]

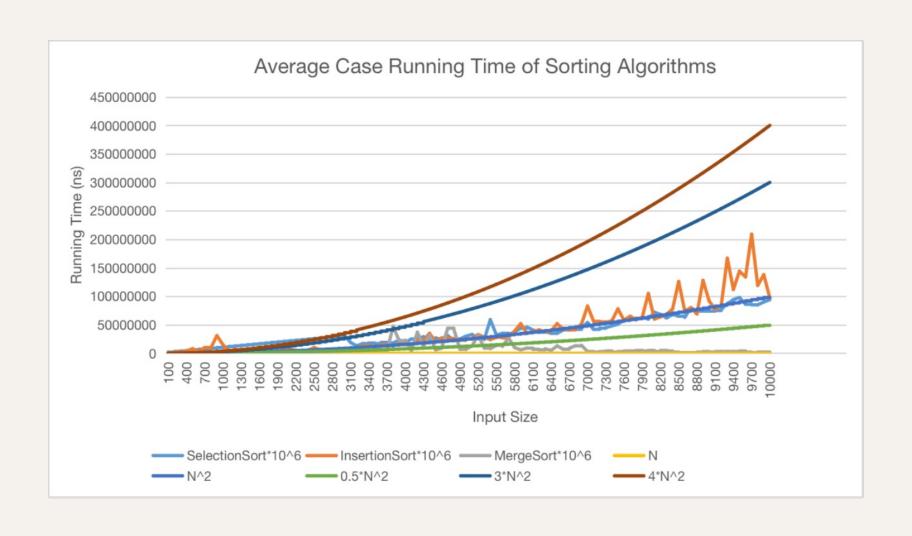
Step 2: Sort arr[mid+1, r]

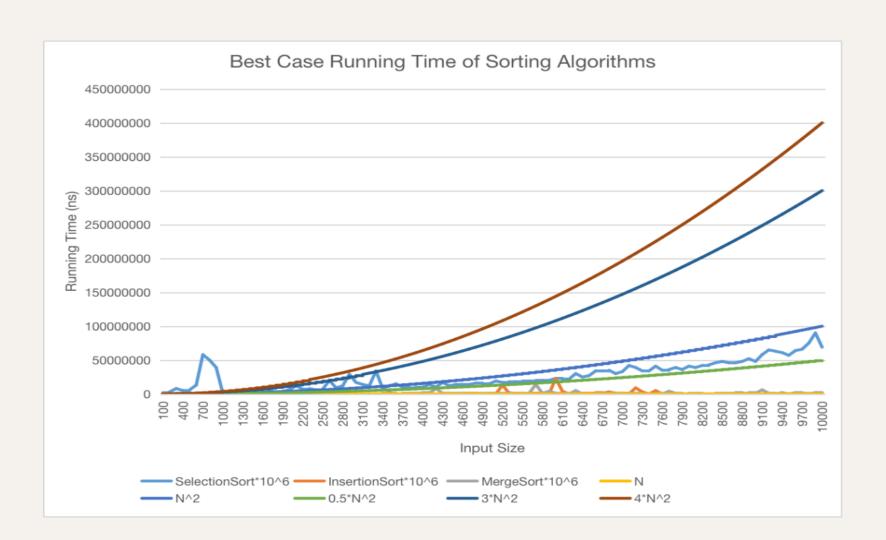
Step 3: Merge sorted arr[l, mid] and arr[mid+1, r]

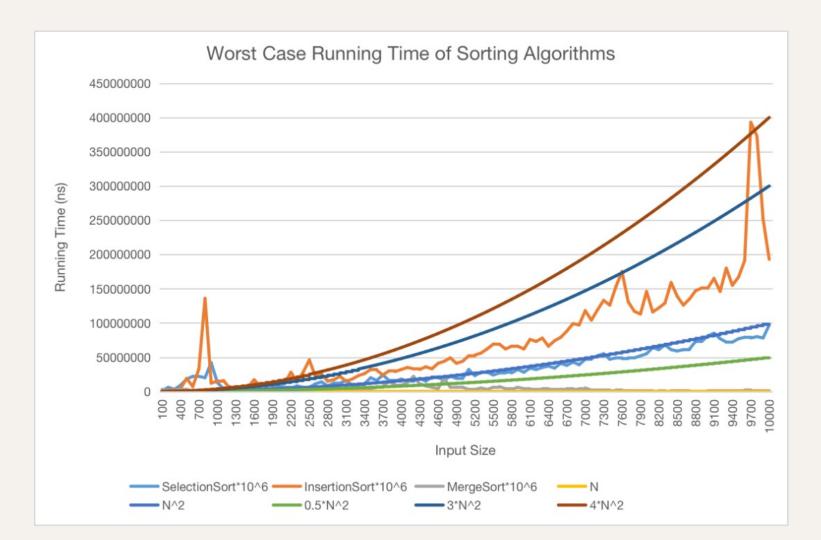


```
// Combine the two parts of arr[l...mid] and arr[mid+1...r]
  private static void merge(Comparable[] arr, int I, int mid, int r) {
     Comparable[] aux = Arrays.copyOfRange(arr, I, r+1);
  // Initialization, i points to the left half of the starting index position I; j points to the right half of the starting
index position mid+1
    int i = I, j = mid+1;
    for( int k = I; k \le r; k ++ ){
       if( i > mid ){ // If the left half of the elements have all been processed
          arr[k] = aux[j-l]; j ++;
       else if(i > r){ // If the right half of the elements have all been processed
          arr[k] = aux[i-l]; i ++;
       else if( aux[i-l].compareTo(aux[j-l]) < 0 ){ // The element pointed to by the left half <the element
pointed by the right half
          arr[k] = aux[i-l]; i ++;
       else{ // The element referred to in the left half >= the element referred to in the right half
          arr[k] = aux[j-l]; j ++;
```

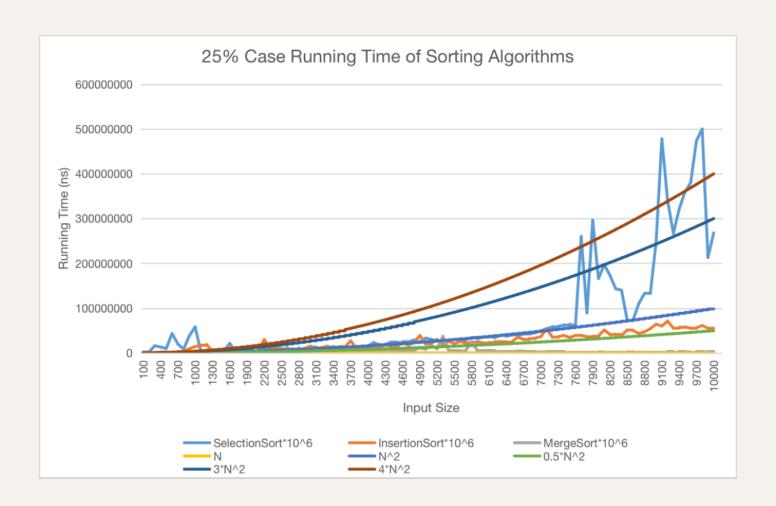
```
private static void sort(Comparable[] arr, int I, int r) {
 // Optimization 2: For small-scale arrays, use insertion sort
  if( r - 1 \le 15 ){
     InsertionSort.sort(arr, I, r);
     return:
  int mid = (1+r)/2;
  sort(arr, I, mid);
  sort(arr, mid + 1, r);
// Optimization 1: For the case of arr[mid] <= arr[mid+1], do not merge
  // It is very effective for nearly ordered arrays, but for general cases, there is a certain performance loss
  if( arr[mid].compareTo(arr[mid+1]) > 0)
     merge(arr, I, mid, r);
public static void sort(Comparable[] arr){
  int n = arr.length;
                                                             Optimization 1: Best case: O(n)!!!
  sort(arr, 0, n-1);
```

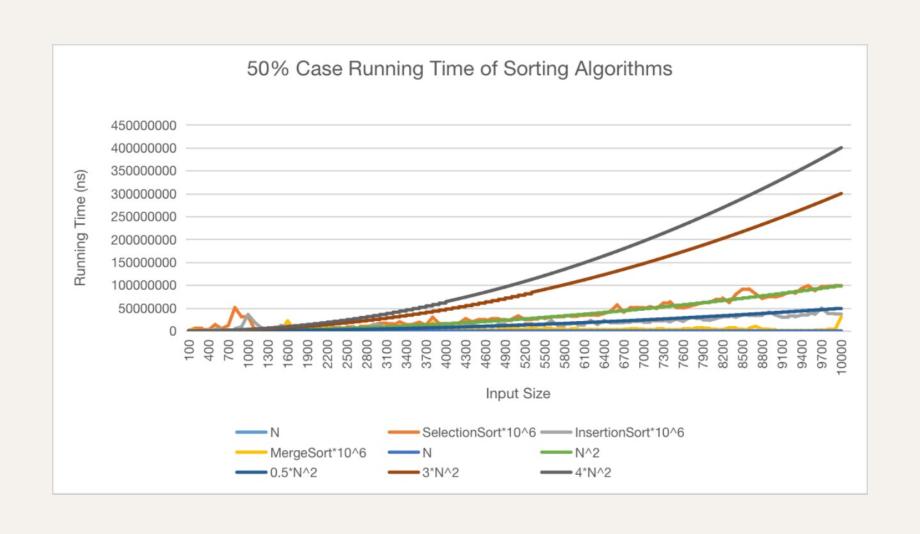


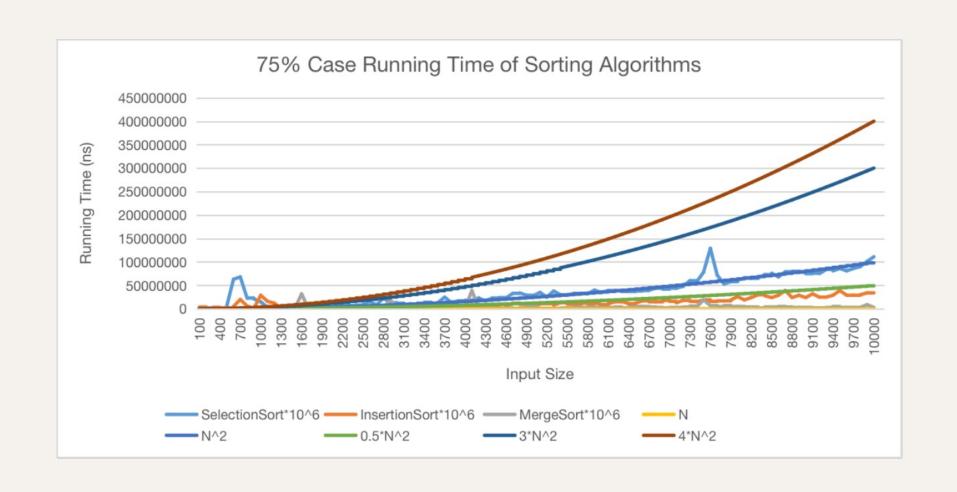




```
// Generate an almost ordered array
// First generate a completely ordered array containing [0...n-1], then randomly exchange swapTimes pairs of data
// swapTimes defines the degree of disorder of the array:
// When swapTimes == 0, the array is completely ordered
// The larger the swapTimes, the more disorder the array tends to be
public static Integer[] generateNearlyOrderedArray(int n, int swapTimes){
    Integer[] arr = new Integer[n];
    for( int \underline{i} = 0; \underline{i} < n; \underline{i} ++)
         arr[<u>i</u>] = new <u>Integer(i)</u>;
    for( int \underline{i} = 0 ; \underline{i} < swapTimes ; \underline{i} ++ ){
         int a = (int)(Math.random() * n);
         int b = (int)(Math.random() * n);
         int t = arr[a];
         arr[a] = arr[b];
         arr[b] = t;
    return arr;
```







Questions?

Thanks!