# PCC: Practical Code Completion Tool with Fuzzy Matching and Precise Synthesis on Statistical Language Modelling

Yixiao Yang Tsinghua University, TNLIST, KLISS, Beijing, China yangyixiaofirst@163.com

## **ABSTRACT**

As open source has become a trend, more and more good projects can be found on github. In countless projects, much of the code can be similar. But in most cases, the existing code cannot be directly copied and pasted. Besides, searching for similar code is difficult, especially for structurally similar one.

We present PCC<sup>1</sup>, a plugin of eclipse IDE which analyzes the code being written, uses the sequence fuzzy matching algorithm to find similar contexts, predicts the most likely sequence based on the pre-trained statistical language model and automatically synthesizes the code from the sequence.

PCC is an intelligent code writer which synthesizes a whole statement which may include all syntax elements of Java in real time. PCC takes the context of currently being edited Java file into consideration and ensures that there are no compilation errors in the generated code.

The correctness of the completed code is 95% among 100 test cases. The time cost of 80% of the completion process is less than 2 seconds, which is effective and user-friendly. In the worst case, the time cost is still less than 4 seconds. To best of our knowledge, it is the first time to apply fuzzy matching algorithm to match and predict the code patterns on the statistical language model. It is also the first time to make a practical tool to automatically write the code with no compilation errors.

# **CCS Concepts**

ullet Software and its engineering eta Search-based software engineering;

# **Keywords**

Code Completion, Code Prediction, Synthesis, Automatically Code Writing, Statistical Language Modelling

# 1. INTRODUCTION

Since 2012, researchers have tried to apply different machine learning algorithms to software engineering. When applying algorithms of natural languages to source codes, things become more complicated. Unlike natural languages, source codes are closely related. Minor modifications may cause compilation errors. This brings challenges to automatic code generations.

PCC synthesizes codes based on patterns which is different from Test Driven Synthesis technique which is based on

genetic algorithms. There have been multiple code synthesis and auto-fix tools [18][10][13] driven by genetic algorithms. These tools can generate accurate code based on oracle test cases. However, there is still a long way to achieve the goal of automatic code writing. Although the integer expression synthesis and the Boolean expression synthesis have been developed for a long time, the problem synthesizing with external method invocations is still difficult. There would be too many possible situations. The genetic algorithm is slow for the problem of a large size.

To narrow the whole possible search space, applying machine learning algorithms to big-code to learn how to write codes has become an effective method.

We apply n-gram model to train the source codes of Java. We predict the codes which are of highest probabilities based on the trained n-gram model.

## 2. RELATED WORK

The first one to apply natural language processing algorithms to source codes was Hindle [7], he used the n-gram [2] model to train the source codes of Java. The next work was done by Allamanis [1], he trained 100 times larger model than Hindle and made a topic visualization tool to show the frequently used codes of the software. Nguyen [11] refined the tokens, used the n-gram topic [17] model, performed the experiments similarly. Nguyen [9], Raychev [14] used the statistical language model to do researches on API usages. The work of Raychev can predict the API, but more generic codes such as if-block, for-block, variable-declaration and so on can not be predicted. The other works can hardly predict and synthesize the codes which can actually be used.

The reason is that their models are too scattered. For example, "String str=a.toString()" will be split into 8 to-kens: String#str#=#a#.#toString#(#). Every word split by # is a token. If a 5-gram model is available, when inferring the next token from the above example, what will be inferred? Now the longest context is the last 4 tokens: .#toString#(#). Even with the longest context, any token behind statements such as "String str=a.toString()", "String any=anything.toString()", "print(sth.toString())" and so on will show up. Information before ".toString()" is lost. According to our experiments, if the model is large, about 100 different tokens with similar probabilities will be obtained. Choice has become a problem. What is more, different people have their own name conventions. Some people like to declare variable like this: "String ssssss1 = a.toString()".

The model cannot be compact either. Due to the database being searched by the exact key, if taking the statement "String str=a.toString()" as one token, nothing will be in-

<sup>&</sup>lt;sup>1</sup>Hhhaha

ferred from the statement "String str=b.toString()" although there is only one character different between two tokens.

Therefore, we propose a novel code completion tool. An intermediate representation of Java codes is introduced and the fuzzy matching algorithm is used.

#### **ARCHITECTURE**

Figure 1 shows the architecture and the execution flow of PCC.

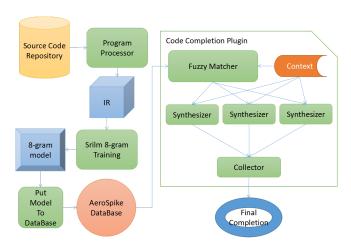


Figure 1: PCC architecture

Source Code Repository contains Java source files crawled from github and other websites. Program Processor translates Java source files to IR. IR is a special form which meets the condition: one operation one token. The basic idea is that if a statement is too long, we split it into several tokens. IR is carefully designed to ensure that these separated tokens can be rebuilt to the original statement. Variable names are eliminated in IR. Due to one token contains only one operation, it is easy for IR to do sequence fuzzy matching. As discussed in Related Work, IR cannot be too scattered nor too compact. After trying for 2 months, a suitable kind of IR was finally found.

```
List<String> list = new LinkedList<String>();
Iterator<String> itr = list.iterator();
VD@List VH@=@HO MI@LinkedList(new);
VD@Iterator VH@=@HO MI@iterator(@C0?0);
```

Figure 2: Simple translation to IR

Firgure 2 gives a concrete idea about the translation and the IR. VD@ means a type declared. VH@ means a variable declared. @HO means that some operations are after this token. MI@ means the method invocation.

Notice that, In Figure 2, the strange symbol C0?0 refer to the variable "list" which is of type List < String >.

Figure 3 shows the translation rule of the variable name. Figure 4 shows more examples about the IR translation.

The grammar of the IR is available in [21]. The program which can parse the IR is available in [22].

Then the IR is trained by Srilm[15] tool. If the Srilm[10] tool is told to train 3-gram models, the 2-gram model and 1-gram model will also be trained. So, if the 8-gram model is trained, the 1-gram to 7-gram models are also be trained.

Show which scope variable is declared. 0 refer to current scope, 1 refers to the parent scope and so on CUSU C means not If there are more than 1 variable of the same type in declared in Field. nearest scope, this integer shows the number of the F means declared

variables of the same type behind where this

Figure 3: Variables to IR

variable is declared.

in Field of a class.

```
int a = 0. b = a + 1:
                                                 int i=0;
boolean c = false:
                                                 new Runnable() {
while (c == true)
                                                 @Override
                                                 public void run() {
         a = 100:
                                                          System.out.println(i);
VD@int VH@=0 VH@=@HO IxE@@C0?0+1;
VD@boolean VH@=false;
                                                 VD@int VH@=0;
IxE@@C0?0==true
                                                 MI@Runnable(new)
WS@while#@PE A@@C0?1=100;
                                                 DH@AnonymousDeclaration;
IxE@ means it is an infix expression.
                                                 MD@void()run
WS@ means it is a while block.
                                                 MI@println(out.System.@C1?0):
@PE means some codes pre-existed should be
put here.
                                                 variable i is in the parent scope of
A@ means it is an assignment.
                                                 anonymous class Runnable, so the IR
a has two represntations. As program evolves,
                                                 form of i is C1?0.
variables may have more and more variables
declared of the same type behind their positions.
```

Figure 4: More examples to IR

In order to preserve the information as much as possible, the 8-gram model is trained.

After the 8-gram model been trained, it is put into Aero Spike[5] which is a mature distributed key-value database.

When users are editing Java files in eclipse and pressing the shortcut keys to invoke the content assistance, the PCC code completion plugin will extract the context of the currently edited Java file, use the program processor to parse the context to IR, put the IR to the Fuzzy Matcher. The possible sequences returned by Fuzzy Matcher will be put into multiple Synthesizers to generate the final completions.

Every synthesizer runs in a separated thread. The Collector will collect all the results of all threads and append the results to the GUI of the eclipse.

## N-GRAM MODEL

The training process computes and stores all the conditional probabilities:

$$P(s_m|s_{m-1}...s_{m-(n-1)}) = \frac{c(s_m...s_{m-(n-1)})}{c(s_{m-1}...s_{m-(n-1)})}$$
 where  $c(\cdot)$  is the count of the given sequence of tokens,  $s_i$ 

represents the token,  $m \in \mathbb{Z}$ .

After training, given a sequence of tokens  $s_1...s_{n-1}$  and an available n-gram model, the probability of next token is:

$$NP(s_{next}) = P(s_{next}|s_1...s_{n-1}) * F(s_1...s_{n-1})$$
  
$$F(s_1...s_{n-1}) = \prod_{m=1}^{n-1} P(s_m|s_{m-1}...s_{m-(n-1)})$$

# **FUZZY MATCHER**

In most cases, The codes mined are not exactly same as the codes written by users. That is why the fuzzy matching algorithm is needed here. The fuzzy matching algorithm is combined with the genetic algorithm [8] and the longest common subsequence LCS algorithm [3].

Figure 5 gives a concrete idea about how this algorithm is running. For the sake of simplicity, the complex IR tokens are replaced by the English alphabets.

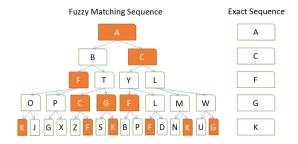


Figure 5: Fuzzy Matching Algorithm

Assume that the already written codes are A, C, F, G, and K. The fuzzy algorithm starts at A, infers the next generation of A based on 2-gram model, chooses two tokens with the highest probabilities computed by P(token|A). Assume that the values of P(B|A) and P(C|A) are largest. So B and C are inferred. The longest common subsequences between the two newly generated sequences: AB, AC and the already written codes: ACFGK are computed. Then, the fuzzy algorithm does the similar things to each token of the new generation: B, C, but this time, 3-gram model is used because the length of the context is 2. Assume that the values of P(F|AB), P(T|AB), P(Y|AC) and P(L|AC) are largest, so, in the next generation, F, T, Y, and L are inferred. The corresponding longest common sequences are recalculated. Keep doing the algorithm until the length of the newly generated sequence is more than the length of the already written codes. In realization, 4 to 8 tokens will be inferred from a single sequence instead of 2. Through the selection of the most likely next generation, the overall search space has been reduced.

The most similar sequences will be retained which are the sequences with orange nodes on its path in Figure 5.

#### 6. PRECISE SYNTHESIZER

PCC puts every sequence returned by the Fuzzy Matcher to an unique Precise Synthesizer. Based on the given sequence, the Precise Synthesizer keeps using the 8-gram model to infer and synthesize the next token until the the sequence of tokens can form a statement.

Algorithm 1 shows the framework of the synthesizer. The function SynthesizeCode is a recursive function which is based on depth-first search. SynthesizeCode has a parameter which represents the current sequence of tokens. The function FindBound finds the tokens which are a whole. For example, given the IR tokens: mi() @PE&&C0?0, the bound of token: @PE&&C0?0 is mi() because mi()&&C0?0 is a whole. The function Synthesize keeps merging the last two tokens from the second argument back to the first argument, and involves the type checking to ensure that there are no compilation errors. If the first argument is null, the tokens from the newly inferred token back to the start of the sequence will all be merged. Figure 6 gives a concrete idea about synthesizing one sequence.

The internal code completion in eclipse could infer the specifications based on literals. This functionality is integrated into PCC. For example, the internal code comple-

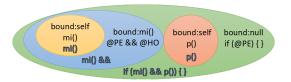


Figure 6: One Sequence Synthesis Example

tion could translate String to java.lang.String or translate a.subStr to  $a.substring(int\ beginIndex)$ .

Because the symbol of the variable name such as C0?0, F0?1 does not contain the information about which type this symbol represents. Therefore, all the types must be tried to synthesize the code. For example, assume two variables declared:  $int\ a=0; boolean\ b=false;$ , when synthesizing the code: C0?0<1, a<1 and b<1 are tried and b<1 is eliminated due to the type conflict.

```
Algorithm 1: SynthesizeCode(contextsequence)
```

```
Require: context sequence! = null
next generations \leftarrow InferTokens(context sequence)
for token: next generations do
bound token \leftarrow Find Bound(token)
success \leftarrow Synthesize(bound token, token)
if !success then
return
end if
if CanFormStatement(context sequence + token)
then
Append Result(Synthesize(null, token))
return
end if
SynthesizeCode(context sequence + token)
end for
```

# 7. IMPLEMENTATION

The source of PCC is in [19]. PCC is implemented in Java. PCC has 28555 lines of code. The Antlr4 [12] library is used to parse the special IR. The model is stored in memory. The AeroSpike database is configured to use memory as the storage device. The AeroSpike database is installed on a Linux machine. The eclipse with PCC installed is running on a Windows machine. The computers used in experiments all have 16G memory and a 2-core i7 CPU.

# 8. PREPARATION OF PCC

Before using PCC, an AeroSpike database which is properly configured must be available.

The program of translating Java to the IR is in [25]. The tutorial is in [26]. The bash script of training the IR to the 8-gram model is in [23]. The tutorial is in [24]. The program of putting the 8-gram model into the Aero Spike database is in [27]. The corresponding tutorial is in [28]. A crawler which can crawl projects from github is in [20]. The format of the files to be downloaded is zip. A script to unzip these zipped files is available in [29].

# 9. USE OF PCC

The installation way [16] of the PCC tool is the same as the common plug-in [6] of eclipse [4]. When the AeroSpike databases filled with the 8-gram model are available, users need to set the IP address of one machine of the AeroSpike clusters on the preference page of their own eclipse. Figure 7 shows the GUI of the preference page.

When everything is ready, the only thing to do is to press the hot key of content assistance of eclipse which is usually Alt+/ to invoke the internal content assistance plugin and the PCC. Figure 8 shows the screenshots of the running eclipse with the PCC installed. The proposals prefixed by the apple icon are generated by PCC. The lower the position, the higher the priority.

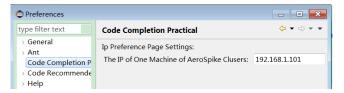


Figure 7: GUI of IP preference setting

```
🕖 fo.java 🛭
        private JTable table;
        private JTextField newFileField;
         protected String getClassName(Object o) {
             String classString = o.getClass().getName();
             int dotIndex = classString.lastIndexOf(".");
               toarray - convert collection to array
                🗏 try_catch - try catch block
                try_finally - try finally block
              while - iterate with enumeration
                🗏 while - iterate with iterator
                while - while loop with condition
                int i1 = classString.lastIndexOf("String Content");
                int i1 = classString.lastIndexOf(File.separator);
                java.lang.String jlaSgs = classString.substring(0,dotIn
                ∫java.lang.String jlaSgs = classString.substring(dotInd∈
                Treturn classString.substring(dotIndex+1);
                                     Press 'Alt+/' to show Template Proposals
```

Figure 8: Running Screen Shot

## 10. EMPIRICAL RESULTS

By learning different projects from github, we found that if two projects deal with different areas, only a small part is similar. Therefore, to simplify our experiments, we only choose projects which deal with similar areas.

Due to the strong type checking of PCC, the project must import the necessary jar packages, otherwise wrong codes may be obtained. In order to reduce the workload of configuring the jar packages, we choose the projects that do not require or require little external jar packages.

After screening by the two conditions just mentioned, 2187 Java files are selected to do the experiments. Due the small size of the model, we put the model into memory.

If the last five proposals returned by PCC contain the right answer, we think the completion is right.

We choose 10 different functions including lambda expressions in Java8. For each function, we randomly choose 10 positions. At each selected position, we invoke the content

assistance. We compare the proposals returned by PPC with the actual statement following the position where we invoked the content assistance. The 100 test cases consist of these 100 positions.

The experiments show that the accuracy is 98%. Only two test cases are not passed. One failed test case is due to too many possible choices when inferring from 7-gram. Another is due to the failure of choosing the right variable from three variables which all pass the type checks.

In order to evaluate the performance of the fuzzy matching algorithm, we randomly add and delete one or two statements of the context of the 100 test cases, and compare the proposals returned by PPC with the actual statement.

The experiments show that the accuracy is greater than or equal to 95%. If we randomly add some statements to the context of the 100 test cases, there is no influence on the accuracy. However, if we randomly delete some statements of the context of the 100 test cases, the accuracy decrease to 95%. The reason for the decline in the correct rate is that the genetic algorithm fails to infer some vital tokens in the process of producing the next generation.

#### 11. LIMITATION AND FUTURE WORK

PCC automatically writes the codes if the completion context is similar to some codes in the database. The computation of the similarity is mainly based on comparing the string contents of each token. If two programs are similar in logic but not similar in structure, PCC will take them as two different programs. What is more, if two programs are structurally similar, but their method names are totally different, PCC will not generate the right code either. When the model becomes more and more large, the genetic algorithm may fail to infer some vital tokens, which leads to the bad results. The future work will mainly focus on solving the problems just mentioned.

# 12. CONCLUSIONS

We present PCC, a plugin of eclipse IDE, to automatically complete the code when users are writing Java codes in eclipse. PCC does not influence the existing functionality of eclipse.

We make a careful analysis of the advantages and disadvantages of the existing work, put forward an IR in which one token represents one operation. Different name conventions of variables are translated into an unified form.

The fuzzy matching algorithm is applied to cope with possible changes in codes. The synthesizer involves runtime type-checking to ensure there is no compilation errors except some unresolved method invocations.

PCC is a real-time tool with quick responses. It is suitable to train some targeted sets according to different users and different projects. Users may find that many codes do not need to be written by themselves.

#### 13. REFERENCES

- M. Allamanis and C. A. Sutton. Mining source code repositories at massive scale using language modeling. In Proceedings of the 10th Working Conference on Mining Software Repositories, MSR '13, San Francisco, CA, USA, May 18-19, 2013, pages 207-216, 2013.
- [2] W. B. Cavnar, J. M. Trenkle, et al. N-gram-based text categorization. Ann Arbor MI, 48113(2):161–175, 1994.
- [3] C. E. R. R. L. S. C. Cormen, Thomas H. (EDT)/ Leiserson. *Introduction to Algorithms 3rd Edition*. Mit Press, LCS algorithm, 2009.
- [4] I. Corporation. Eclipse help platform. http://help.eclipse.org/mars/index.jsp, 1996.
- [5] J. Dillon. Aerospike documentation. http://www.aerospike.com/docs/, 2016.
- [6] D. R. Eric Clayberg. Eclipse Plug-ins (4th Edition). Addison-Wesley Professional, Introduction, 2012.
- [7] A. Hindle, E. T. Barr, Z. Su, M. Gabel, and P. T. Devanbu. On the naturalness of software. In 34th International Conference on Software Engineering, ICSE 2012, June 2-9, 2012, Zurich, Switzerland, pages 837–847, 2012.
- [8] M. Melanie. An introduction to genetic algorithms. www.boente.eti.br/fuzzy/ebook-fuzzy-mitchell.pdf, 1996.
- [9] A. T. Nguyen and T. N. Nguyen. Graph-based statistical language model for code. In 37th IEEE/ACM International Conference on Software Engineering, ICSE 2015, Florence, Italy, May 16-24, 2015, Volume 1, pages 858-868, 2015.
- [10] H. D. T. Nguyen, D. Qi, A. Roychoudhury, and S. Chandra. Semfix: program repair via semantic analysis. In Proceedings of the 2013 International Conference on Software Engineering, pages 772–781. IEEE Press, 2013.
- [11] T. T. Nguyen, A. T. Nguyen, H. A. Nguyen, and T. N. Nguyen. A statistical semantic language model for source code. In Joint Meeting of the European Software Engineering Conference and the ACM SIGSOFT Symposium on the Foundations of Software Engineering, ESEC/FSE'13, Saint Petersburg, Russian Federation, August 18-26, 2013, pages 532-542, 2013.
- [12] T. Parr. Antlr4. http://www.antlr.org/, 2016.
- [13] D. Perelman, S. Gulwani, D. Grossman, and P. Provost. Test-driven synthesis. In ACM SIGPLAN Notices, volume 49, pages 408–418. ACM, 2014.
- [14] V. Raychev, M. T. Vechev, and E. Yahav. Code completion with statistical language models. In ACM SIGPLAN Conference on Programming Language Design and Implementation, PLDI '14, Edinburgh, United Kingdom - June 09 - 11, 2014, page 44, 2014.
- [15] A. Stolcke, J. Zheng, W. Wang, and V. Abrash. Srilm at sixteen: Update and outlook. In Proceedings of IEEE Automatic Speech Recognition and Understanding Workshop, page 5, 2011.
- $[16] \begin{tabular}{l} we nukb. Eclipse plugin installation. \\ http://www.venukb.com/2006/08/20/ \\ install-eclipse-plugins-the-easy-way/, 2006. \\ \end{tabular}$
- [17] X. Wang, A. McCallum, and X. Wei. Topical n-grams: Phrase and topic discovery, with an application to

- information retrieval. In *Data Mining, 2007. ICDM 2007. Seventh IEEE International Conference on*, pages 697–702. IEEE, 2007.
- [18] W. Weimer, T. Nguyen, C. Le Goues, and S. Forrest. Automatically finding patches using genetic programming. In 31st International Conference on Software Engineering, ICSE 2009, May 16-24, 2009, Vancouver, Canada, Proceedings, pages 364-374, 2009.
- [19] Y. Yang. Codecompletion. https: //github.com/yangyixiaof/CodeCompletionPlugin/ tree/master/cn.yyx.research.codecompleteplugin, 2016.
- [20] Y. Yang. The github crawler. https://github.com/ yangyixiaof/gitcrawler/tree/master/gitcrawler, 2016.
- [21] Y. Yang. Ir grammar. https://github.com/ yangyixiaof/specialparse/blob/master/specialparse/ src/main/java/cn/yyx/parse/szparse8java/Java8.g4, 2016.
- [22] Y. Yang. Ir parser. https://github.com/yangyixiaof/ specialparse/tree/master/specialparse, 2016.
- [23] Y. Yang. N-gram training script. https://github.com/yangyixiaof/gitcrawler/tree/ master/programprocessor/ngram-train-scripts, 2016.
- [24] Y. Yang. N-gram training script tutorial. https://github.com/yangyixiaof/gitcrawler/blob/master/programprocessor/ngram-train-scripts/train-script-tutorial, 2016.
- [25] Y. Yang. Ppc program processor. https://github.com/ yangyixiaof/gitcrawler/tree/master/programprocessor, 2016.
- [26] Y. Yang. Ppc program processor tutorial. https://github.com/yangyixiaof/gitcrawler/blob/master/README.md, 2016.
- [27] Y. Yang. The program which puts the model into the aerospike database. https://github.com/yangyixiaof/ModelIntoAeroSpike/ tree/master/PutModelIntoAeroSpike, 2016.
- [28] Y. Yang. The tutorial of the program which puts the model into the aerospike database. https://github.com/yangyixiaof/ModelIntoAeroSpike/ blob/master/README.md, 2016.
- [29] Y. Yang. The unzip script. https://github.com/yangyixiaof/gitcrawler/tree/master/gitcrawler/UnzipScript, 2016.