## CSC240 Notes

## Week 11

## 0.1 Languages

**Definition.** A language over  $\Sigma$  is a subset of  $\Sigma^*$ .

i.e.  $L \subseteq \Sigma^*$  is a language of  $\Sigma$ .

**Definition.** (Kleene Star \*)

$$L^* = \bigcup_{k \in \mathbb{N}} L^k$$

- $L^+ = \bigcup_{k \in \mathbb{Z}^+} L^k$  Notice that  $L^* = L^+ \cup \{\lambda\}$ , so  $L^+ = L^*$  iff  $\lambda \in L$ .

  Let  $L_0, L_1 \subseteq \Sigma^*$  be languages. Then  $L_0 \cup L_1, L_0 \cap L_1$  are languages.
- The complement of a language  $L \subseteq \Sigma^*$ ,  $\bar{L} = \Sigma^* L$ , is also a language

## 0.2**Regular Expressions**

Regular Expressions are a concise way to describe some language.

**Definition.** Given a finite alphabet  $\Sigma$ ,  $R_{\Sigma}$  denotes the set of regular expressions over  $\Sigma$ .  $R_{\Sigma}$  is the smallest set of strings such that

- Base case:  $\emptyset \in R_{\Sigma}, \lambda \in R_{\Sigma}$ , for every  $a \in \Sigma, a \in R_{\Sigma}$ .
- Constructor cases: For every  $r, r' \in R_{\Sigma}$ ,

$$(r+r') \in R_{\Sigma}, (r \cdot r') \in R_{\Sigma}, r^* \in R_{\Sigma}.$$

Given a regular expression